

Department of Computer Science
Data Structures (CSC212)
Final Exam (2nd Semester 1428-29)
Date: 8/6/1429H

Time: 3 hours

Marks: 100

Question 1 (30 marks)

- (a) What is the step count of the 'for' statement: $\text{for (int } i = 0; i \leq n; i++) \text{ } x++;$ $n+2$ $n+1$
- (b) What is the step count of the statement S1 in the following code segment:

for (int i = 1; i < n; i++) for (int j = i; j < n; j++) S1;

Note : $\sum_{i=1}^n i = \frac{n(n+1)}{2}$.

- (c) What is the growth rate of the following functions in terms of the big-O notation:
- (i) $f(n) = 10$ (ii) $f(n) = (2n^3 + n^2 + 1) \log(n^2 + 1)$.
- (d) Why is remove operation in singly linked list $O(n)$?
- (e) Give one advantage of having sentinel nodes in lists?
- (f) What type of structure do trees have? درکبیت پسائی
- (g) Give a formula for the number of nodes in a full binary tree in terms of its height, h .
- (h) Give a formula for the minimum number of nodes in a complete binary tree in terms of its height, h .
- (i) What is the *worst-case* and the *best-case* time complexity of the FindKey operation in BSTs?
- (j) Draw an expression tree for the expression: $(a*b/2+3*(c-1))$.
- (k) What is the minimum number of keys a non-leaf node (except the root) can have in a B+-tree of order M ?
- (l) What is the minimum and maximum number of data elements a root can have in a B+-tree?
- (m) What is the time complexity of HeapSort algorithm?
- (n) What is $(\log_2 100)$ equal to?

Question 2. (15 marks)

- (a) Draw a graphical representation of a doubly-linked list.
- (b) Implement **insert** and **remove** operations for a doubly-linked list, according to the specification of ADT List.

Question 3. (15 marks)

- (a) Convert the following array of integers into a min-heap and show the resulting heap as an array and as a tree. The array is {12, 14, 3, 16, 8, 7, 10, 17, 5, 11, 9, 6, 13, 15, 4}. Use SiftDown operation.
- (b) From the heap you obtained in part (a) delete the minimum element (deleteMin) twice and show the result after each deletion as an array.

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Q1-

- (a) What is the step count of the 'for' statement :
for(int i = 0 ; i <= n ; i++)
x++;

Answer:

$$f(n) = \begin{matrix} n+2+n+1 \\ 2n + 3 \end{matrix}$$

- (b) What is the step count of the statement S1 in the following code segment:

```
for(int i = 1 ; i < n ; i++)  
    for(int j = i ; j < n ; j++)  
        S1;
```

Answer:

Suppose n = 5.

first 'for' has steps : n
second 'for' :

i	j	j < n	S1
1	1	1	0
2	2	2	1
3	3	3	2
4	4	4	3
5			

step count for S1 : 0 + 1 + 2 + 3 + + n-2

$$\sum_{i=1}^N i = 1+2+3+\dots+(n-2)+(n-1)+(n) = \frac{n(n+1)}{2}$$

$$\begin{aligned}
&= \frac{n(n+1)}{2} - (n-1) - (n) = \frac{n(n+1)}{2} - n + 1 - n \\
&= \frac{n(n+1)}{2} - 2n + 1 = \frac{n^2 + n}{2} - 2n + 1 \\
&= 2n^2 - 2n - 2
\end{aligned}$$

(c)

(i) $f(n) = 10$
 $O(1)$

(ii)

$$f(n) = (2n^3 + n^2 + 1) \log(n^2 + 1)$$

$$2n^3 + n^2 + 1 < 2n^3 + n^3 + n^3 \text{ where } n_0 \geq 1$$

$$< 4n^3 \text{ where } n_0 \geq 1$$

$$O(n^3) \text{ where } c=4 \text{ and } n_0 \geq 1$$

$$\log(n^2 + 1) \leq \log(n^2 + n^2) \text{ where } n_0 \geq 1$$

$$\leq \log(2n^2)$$

$$\leq \log(2) + \log(n^2)$$

$$\leq \log(n) + 2 \log(n)$$

$$\leq 3 \log(n)$$

$$O(\log(n)) \text{ where } c=3 \text{ and } n_0 \geq 1$$

$$O(n^3) O(\log(n))$$

$$O(n^3 \log(n)) \text{ where } C=12 \text{ and } n_0 \geq 1$$

(d) Because you to do ~~shift~~ ^{go through} for all nodes ~~before~~ ^{before} current, you do the shift with loop.

(e) Advantage of Sentinel nodes Linked list : It simplifies code because all nodes will have previous and next node, you will not have any special case.

(f) Hierarchical.

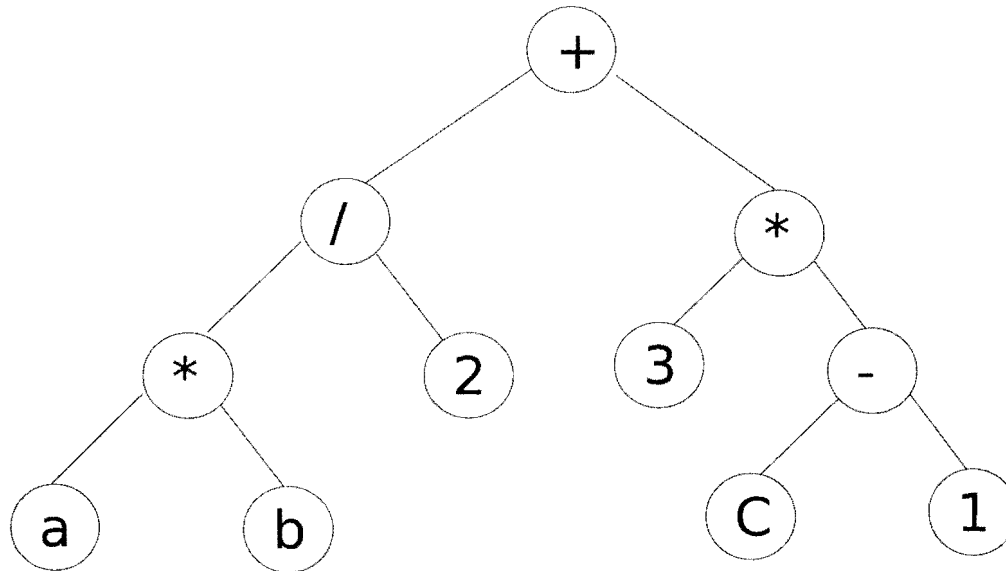
(g) Max. Number of nodes in a full binary tree = $2^h - 1$

(h) ~~Best case $O(1)$ Worst case $O(\log(n))$~~

Best case $O(\log(n))$ Worst case $O(n)$.

(i) Min Number of nodes in a complete binary tree
 $= 2^{h-1}$.

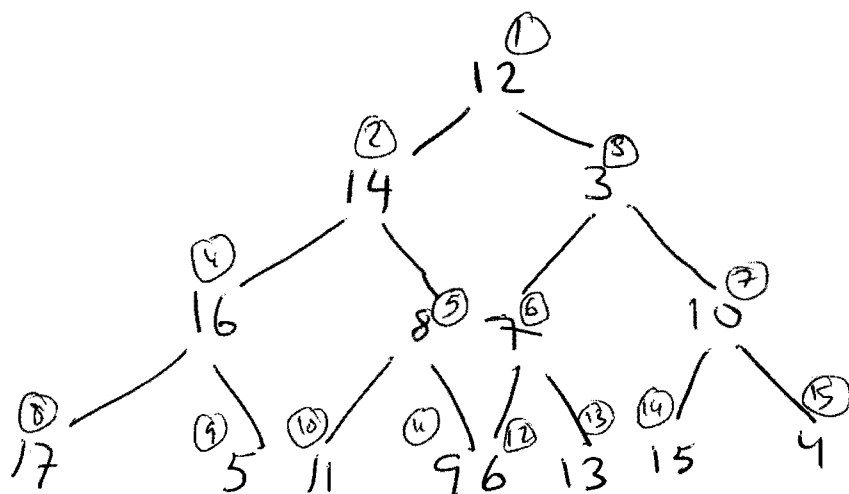
(j) $(a*b/2 + 3*(c-1))$
 $((a*b)/2) + (3*(c-1))$



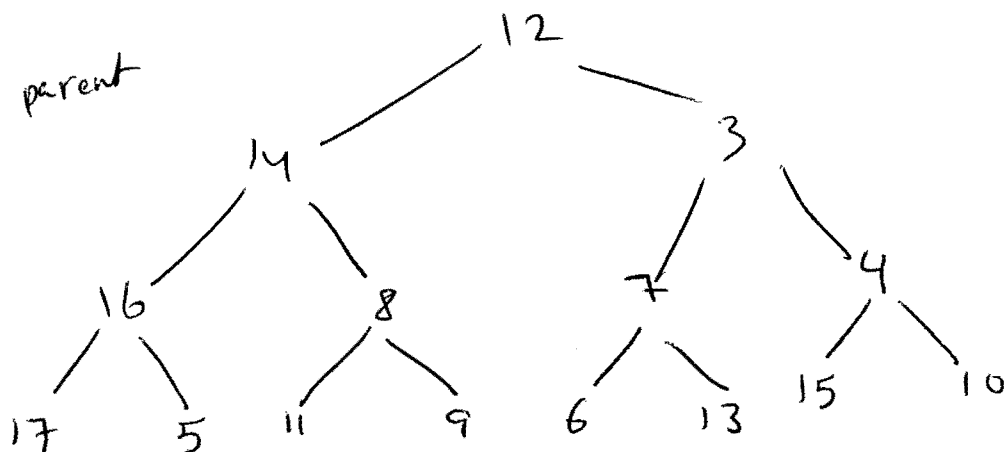
- (k) Minimum number of keys in a leaf node except root
 with order $m = (m/2 - 1)$.
 (l) Minimum = 1 Maximum = order m .
 (m) Time complexity of heapsort = $O(n \log_2(n))$.
 (n) $\log_2(100) = \log_2(10^2) = 2 \log_2(10)$

Q3

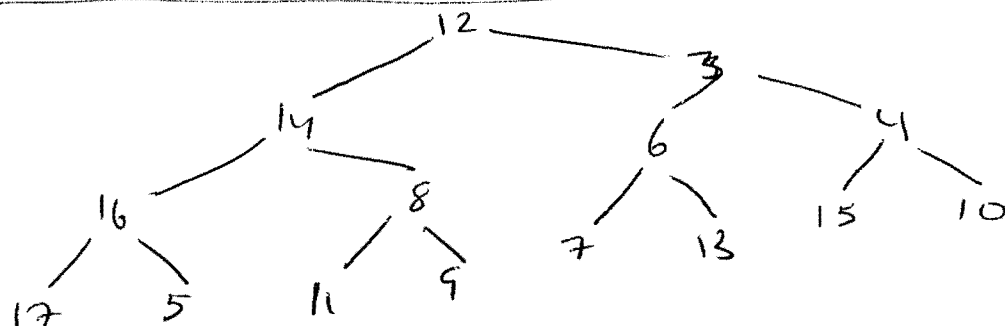
(a)



Start with last parent
index 7

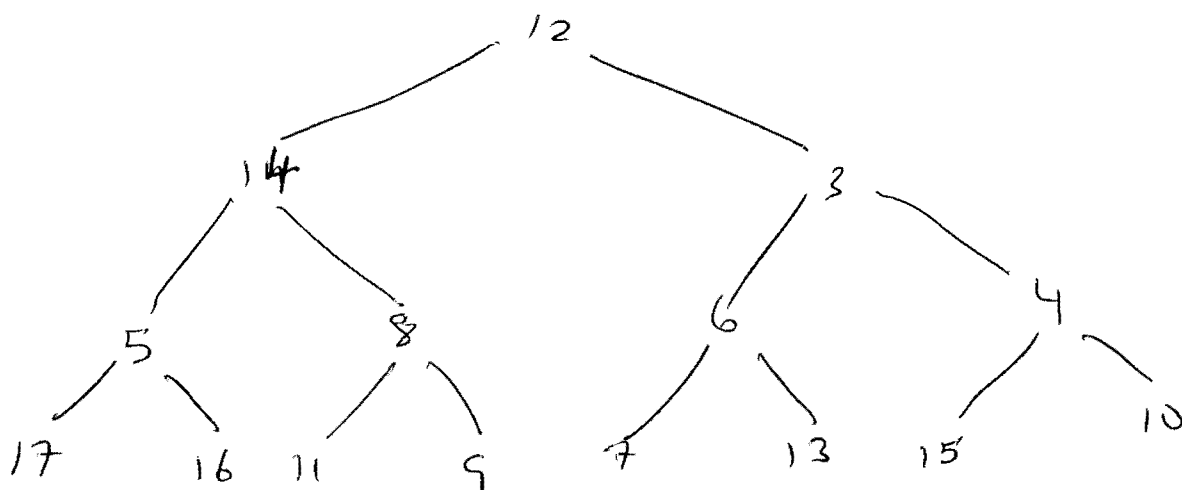


index 6



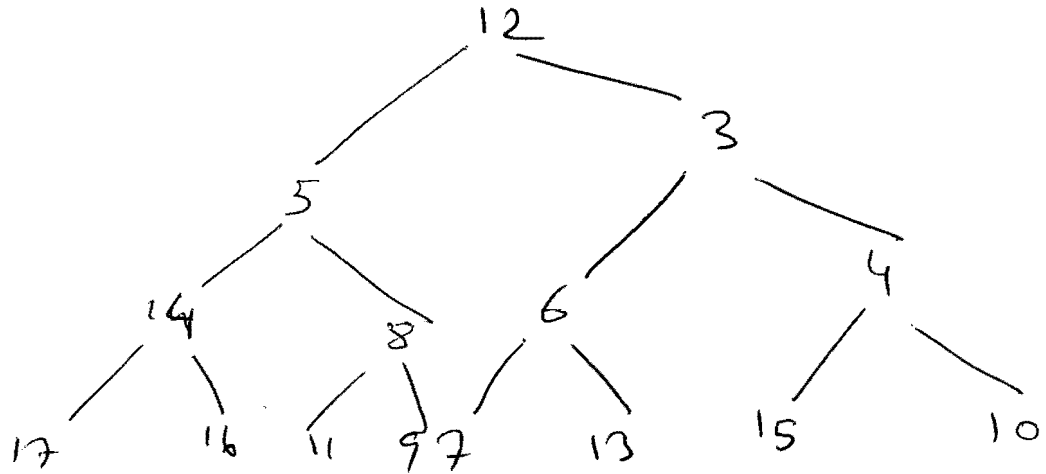
index 5 (10)

index 4

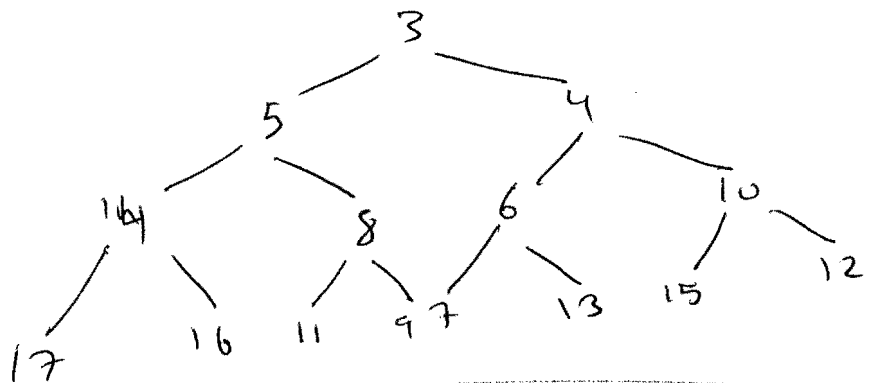
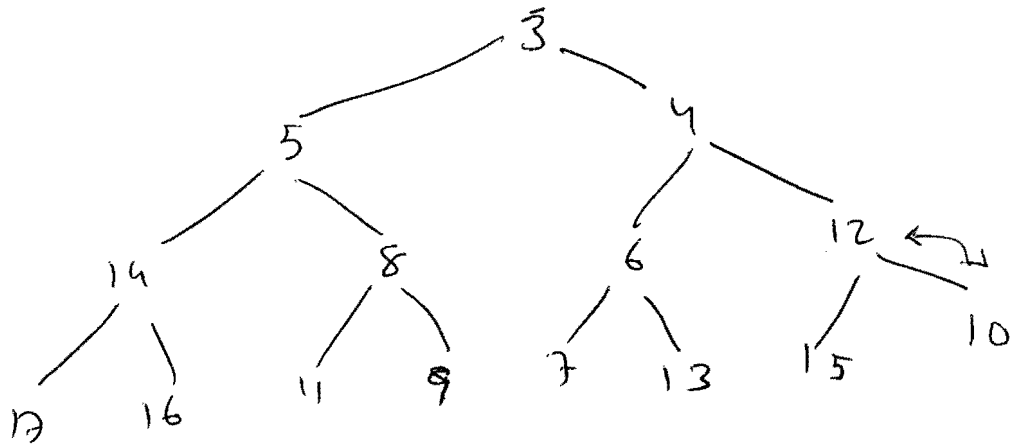
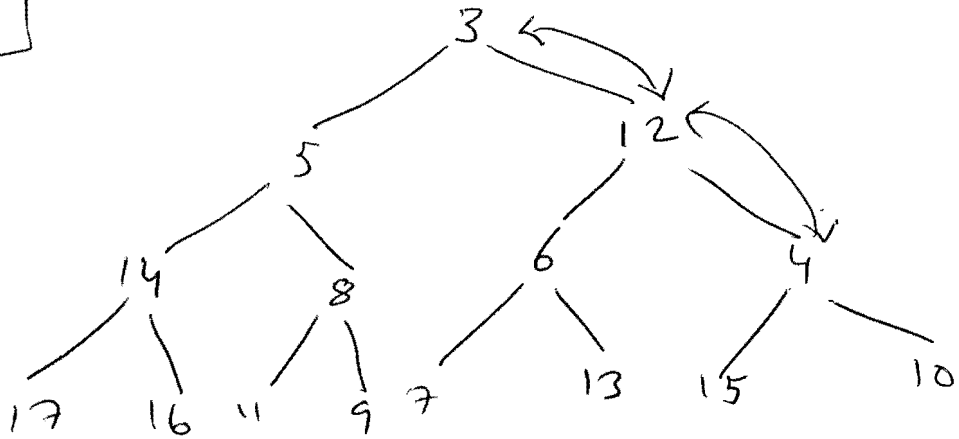


index 3 ok

index 2



index 1



as an array :

3	5	4	14	8	6	10	17	16	11	9	7	13	15	12	
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

⑥ delete first time:

3	5	4	14	8	6	10	17	16	11	9	7	13	15	12
1	2	3	4	5	6	7	8	9	10	11	12	13	14	15

12	5	4	14	8	6	10	17	16	11	9	7	13	15	
1	2	3	4	5	6	7	8	9	10	11	12	13	14	

Sift down

Index = 1 Left child in = 2 Right child in = 3

Swap 12 with 4.

4	5	12	14	8	6	10	17	16	11	9	7	13	15
1	2	3	4	5	6	7	8	9	10	11	12	13	14

Index = 3 Left child = 6 Right child = 7

Swap 12 with 6.

4	5	6	14	8	12	10	17	16	11	9	7	13	15
1	2	3	4	5	6	7	8	9	10	11	12	13	14

Index = 6 Left child = 12 Right child = 13

Swap 12 with 7

4	5	6	14	8	7	10	17	16	11	9	12	13	15
1	2	3	4	5	6	7	8	9	10	11	12	13	14

delete
second

15	5	6	14	8	7	10	17	16	11	9	12	13	15
1	2	3	4	5	6	7	8	9	10	11	12	13	14

Index = 1 Left child = 2 Right child = 3

Swap 15 with 5.

5	15	6	14	8	7	10	17	16	11	9	12	13
1	2	3	4	5	6	7	8	9	10	11	12	13

Index 2 Left child = 4 Right child = 5

Swap 15 with 8.

5	8	6	14	15	7	10	17	16	11	9	12	13
1	2	3	4	5	6	7	8	9	10	11	12	13

Index = 5 Left child = 10 Right child = 11

Swap 15 with 9,

5	8	6	14	9	7	10	17	16	11	15	12	13
1	2	3	4	5	6	7	8	9	10	11	12	13

