

# King Saud University

College of Computer and Information Sciences

Department of Computer Science

### **Data Structures CSC 212**

### Final Exam - Fall 2017

Date: 30/12/2017 Duration: 3 hours

#### **Guidelines**

•No calculators or any other electronic devices are allowed in this exam.

•Use a pencil in choice questions.

Student ID: Name:										
Section	:					Instru	ictor:			
1	2.1	2.2	3.1	3.2	4	5	6	7	8	Total

(a) Choose the correct frequency for every line as well as the total O of the following code:

```
1    sum = 1;
2    for (i = 1; i <= n; i++) {
3        sum+= i;
4        for (j = i; j >= 2; j--)
5        sum--;}
```

- 1. Line 1:  $(\widehat{A})$  1  $(\widehat{B})$  2  $(\widehat{C})$  3  $(\widehat{D})$  n  $(\widehat{E})$  2n
- 2. Line 2: (A) n (B) n+1 (C) n-1 (D) n+2 (E) n-2
- 3. Line 3: (A) n (B) n+1 (C) n-1 (D) n+2 (E) n-2
- 4. Line 4: (A)  $n^2$  (B) n(n-1)/2 (C) (2n+1)/2 (D) (2n-1)/2 (E) n(n+1)/2
- 5. Line 5: (A)  $n^2$  (B) n(n-1)/2 (C) (2n+1)/2 (D) (2n-1)/2 (E) n(n+1)/2
- 6. Total O: (A) 1 (B) n (C)  $n^2$  (D)  $n \log(n)$  (E)  $n^3$

(b) Choose the correct frequency for every line as well as the total O of the following code:

```
count = 0;
for (i = 1; i < n+1; i++)
count ++;
for (j = 0; j <= count; j++)
k = j+1;</pre>
```

- 1. Line 1: (A) 0 (B) 1 (C) 2 (D) n (E)  $n^2$
- 2. Line 2: (A) n (B) n+1 (C) n-1 (D) n+2 (E) n-2
- 3. Line 3: (A) n (B) n+1 (C) n-1 (D) n+2 (E) n-2
- 4. Line 4: (A) count + 2 (B) n + 1 (C) n 1 (D) n + 2 (E) n(n + 1)/2
- 5. Line 5: (A) count + 1 (B) n + 1 (C) n 1 (D) n + 2 (E) n(n 1)/2
- 6. Total O: (A) 1 (B) n (C)  $n^2$  (D)  $n \log(n)$  (E)  $n^3$

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(c)	Choose the correct answer:
	1. $n^2 + n \log n^4$ is : (A) $O(n)$ (B) $O(n^2)$ (C) $O(n \log(n))$ (D) $O(n^4)$ (E) None
	2. $n^2 + 1000n$ is : (A) $O(n)$ (B) $O(n^2)$ (C) $O(n \log(n))$ (D) $O(nn^2)$ (E) None
	3. $n^4 \log n + n!$ is: (A) $O(n!)$ (B) $O(n^4)$ (C) $O(n^5)$ (D) $O(\log(n))$ (E) None
	4. Algorithm A is $O(n)$ , and Algorithm B is $O(2n)$ . Given the same input:
	(A) A always finishes before B. (B) B always finishes before A. (C) A and B finish at the same
	time. (D) B requires double the time taken by A. (E) None
Que	stion 2
(a)	Given a queue of time intervals represented as pairs of integers, write the method <b>public static</b> Queues
	Pair <integer, integer="">&gt; getIntervals(Queue<pair<integer, integer="">&gt; q, int start, int end), which returns</pair<integer,></integer,>
	all intervals of ${\tt q}$ intersecting the interval [start, end]. The input ${\tt q}$ must not change. Assume that all
	intervals in q are valid (that is $first \le second$ ), non-overlapping and ordered in chronological order
	The class Pair is given below.
	<pre>public class Pair<u, v=""> {   public U first;   public V second;   Pair(U first, V second) {     this.first = first;     this.second = second;   } }</u,></pre>

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(b) Consider a stack of decreasing time intervals, that is, starting from the top, each interval contains the next. Write the method public static Pair<Integer, Integer> smallest(Stack<Pair<Integer, Integer>> st, int t), which returns the smallest interval containing t if it exists, null otherwise. Assume that all intervals in st are valid (that is first <= second).</p>

**Example 0.1.** If st:  $\{[0,8], [1,6], [1,5], [2,4]\}$ , then smallest(st, 1) returns [1,5], smallest(st, 3) returns [2,4], smallest(st, 9) returns null.

Complete the code below by choosing the correct answer:

```
public static Pair < Integer , Integer > smallest(Stack < Pair < Integer , Integer >> st , int t) {
1
2
3
      Pair < Integer , Integr > itm = null;
4
      while (!st.empty()) {
5
        Pair < Integer , Integr > it = st.pop();
6
        if (...)
7
8
           itm = it;
9
        else
10
11
12
      while (...) {
13
14
15
      return itm; }
```

- 1. Line 2:
  - A Queue<Pair<Integer, Integer>> r = new
    LinkedQueue<Pair<Integer, Integer>>();
  - B Stack<Integer> r = new LinkedStack
    Integer>();
  - C List<Pair<Integer, Integer>> r = new
  - LinkedList<Pair<Integer, Integer>>();
  - (D) Stack<Pair<Integer, Integer>> r = new
    LinkedStack<Pair<Integer, Integer>>();
  - (E) None
- 2. Line 6:
  - (A) r.push(it.first);
  - (B) r.insert(it);
  - (C) r.enqueue(it);
  - (D) r.push(it);
  - (E) None
- 3. Line 7:
  - (A) if (it.first < t && t <= it.second)
  - $\stackrel{\textstyle oxed{(B)}}{}$  if (it.first <= t && t <= it.second)
  - (C) if (it.first < t || it.second > t)

- $\stackrel{\textstyle \frown}{\rm (D)}$  if (it.first <= t && it.second <= t)
- (E) None
- 4. Line 10:
  - (A) r.serve();
  - (B) break;
  - (C) r.pop();
  - $(\widehat{D})$  r.findNext();
  - (E) None
- 5. Line 12:
  - (A) while (r.empty()){
  - (B) while (!r.empty()){
  - (C) while (r.pop()!= null){
  - (D) while (r.length()!= 0){
  - (E) None
- 6. Line 13:
  - (A) st.push(r.pop());
  - (B) st.push(r.serve());
  - C st.push(r.retrieve()); r.findNext();
  - (D) st.push(r.push());
  - (E) None

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(a) The method private BTNode<T> mirrorCopy(BTNode<T> t) creates recursively a mirror copy of the subtree

t. Choose the correct option to complete the code of this method:

- 1. Line 2:
  - (A) if (t.left == null || t.right == null)
  - (B) if (t.left == null && t.right == null)
  - (C) if (t == null)
  - (D) if (root != null)
  - (E) None
- 2. Line 3:
  - (A) return null;
  - (B) return root;
  - (C) return mirrorCopy(root);
  - (D) return mirrorCopy(t);
  - (E) None
- 3. Line 4:
  - (A) BTNode<T> p = new BTNode<T>(t.data);
  - (B) BTNode<T> p = new BTNode<T>(root);
  - (C) BTNode<T> p = new BTNode<T>(t);
  - (D) BTNode<T> p = new BTNode<T>(root.data);
  - (E) None

- 4. Line 5:
  - (A) p.right = mirrorCopy(t.left);
  - (B) t.left = mirrorCopy(t.left);
  - (C) p.right = mirrorCopy(t.right);
  - (D) t.left = mirrorCopy(t.right);
  - (E) None
- 5. Line 6:
  - (A) t.right = mirrorCopy(t.left);
  - (B) p.left = mirrorCopy(t.left);
  - (C) p.left = mirrorCopy(t.right);
  - (D) t.right = mirrorCopy(t.right);
  - (E) None
- 6. Line 7:
  - (A) return p;
  - (B) return mirrorCopy(t);
  - (C) mirrorCopy(t.left); mirrorCopy(t.right);
  - (D) return t;
  - (E) None
- (b) Consider the function f below, member of DoubleLinkedList:

```
public void f(int n) {
   Node < T > p = head; Node < T > q = null;
   for (int i = 0; i < n; i++) {
      q = p;
      p = p.next;
   }
   if (p != null) {
      p.previous = null;
      while (p.next != null)
           p = p.next;
      p.next = head;
      head = q.next;
      q.next = null;
   }
}</pre>
```

Choose the correct result in each of the following cases:

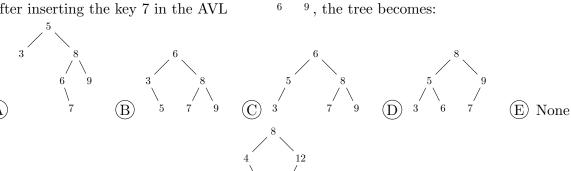
1. The list 1: $A, B, C, D, E$ , after calling 1.f(2), 1 becomes:
(A) $A, B$ (B) $D, E, A, B, C$ (C) $C, D, E, A, B$ (D) $A, D, E, B, C$ (E) None
2. The list 1: $A, B, C, D, E$ , after calling 1.f(0), 1 becomes:
(A) $empty$ (B) $A, B, C, D, E$ (C) $B, C, D, E, A$ (D) $B, C, D, E$ (E) None
3. The list 1: $A, B, C, D, E$ , after calling 1.f(5), 1 becomes:
(A) empty (B) $A, B, C, D, E$ (C) $E, D, C, B, A$ (D) $A, D, E, B, C$ (E) None
4. The list 1: $A, B, C, D, E$ , after calling 1.f(1), 1 becomes:
A $A$ $B$ $E,A,B,C,D$ $C$ $C,D,E,A,B$ $D$ $B,C,D,E,A$ $E$ None
Question 4
(a) Consider the following heap represented as an array: 3, 7, 9, 13, 8, 11. Choose the correct answer for
every operation (all operations are done on the above heap).
1. Heap after inserting 5:
A 3,7,9,13,8,11,5 $B$ 3,5,7,13,8,11,9 $C$ 3,7,9,13,8,5,11 $D$ 5,7,3,13,8,11,9 $E$ None
2. Heap after inserting 10:
A 3,7,10,13,8,11,9 $B$ 3,7,9,13,8,10,11 $C$ 3,7,9,13,8,11,10 $D$ 3,7,9,10,8,11,13 $E$ None
3. Heap after inserting 2:
4. Heap after deleting one key:
A 7,13,9,11,8 $B$ 7,8,9,13,11 $C$ 9,7,11,13,8 $D$ 7,9,8,11,13 $E$ None
5. Heap after deleting two keys:
(A) $7,13,9,11$ (B) $8,11,9,13$ (C) $7,8,9,13$ (D) $13,9,8,11$ (E) None
(b) What is the result of a bottom-up min-heap construction of the following array: 2,4,6,3,5,1?
(A) $1,2,3,5,4,6$ (B) $2,1,3,4,5,6$ (C) $1,3,2,4,5,6$ (D) $1,3,2,5,6,4$ (E) None.
(c) Choose the correct answer:
1. What is the height of a heap of size $k$ ?
$\bigcirc$
2. Bottom-up heap construction is:
$\bigcirc A$ $O(n)$ $\bigcirc B$ $O(\log n)$ $\bigcirc C$ $O(n\log n)$ $\bigcirc D$ $O(n^2)$ $\bigcirc E$ None.
3. The enqueue operation in a heap priority queue is:
$\textcircled{A} \ O(1)  \textcircled{B} \ O(\log n)  \textcircled{C} \ O(n)  \textcircled{D} \ O(n\log n)  \textcircled{E} \ \text{None}.$

Question  $5 \dots 5$ 

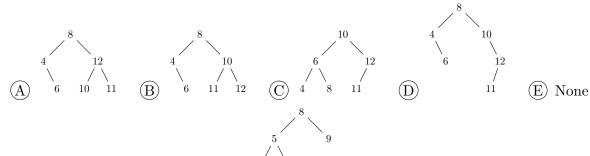
Choose the correct result in each of the following cases:



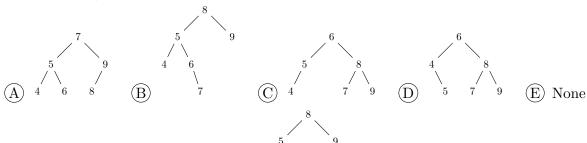
1. After inserting the key 7 in the AVL



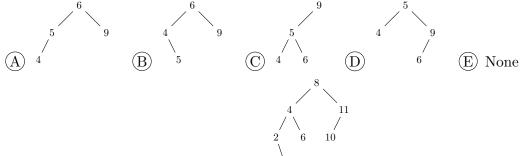
2. After inserting the key 11 in the AVL , the tree becomes:



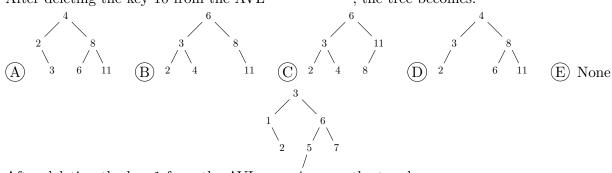
3. After inserting the key 7 in the AVL  $^4$ , the tree becomes:



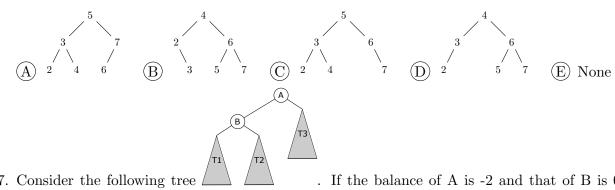
4. After deleting the key 8 from the AVL  $^4$ , the tree becomes:



5. After deleting the key 10 from the AVL , the tree becomes:



6. After deleting the key 1 from the AVL , the tree becomes:



7. Consider the following tree \_\_\_\_\_ . If the balance of A is -2 and that of B is 0, then after performing a single right rotation at A, then:

(a) The balance of A becomes:

(b) The balance of B becomes:

Use the hash function H(key) = key%11 to store the sequence of keys 16, 14, 27, 5, 21, 43, 10, 38, 19, 18, 20 in the hash table. Use the following collision resolution strategies:

1. Linear rehashing (c=1). Fill in the following table:

Key	16	14	27	5	21	43	10	38	19	18	20
Position											
Number of probes											

2. External chaining. Fill in the following table:

Key	16	14	27	5	21	43	10	38	19	18	20
List position											

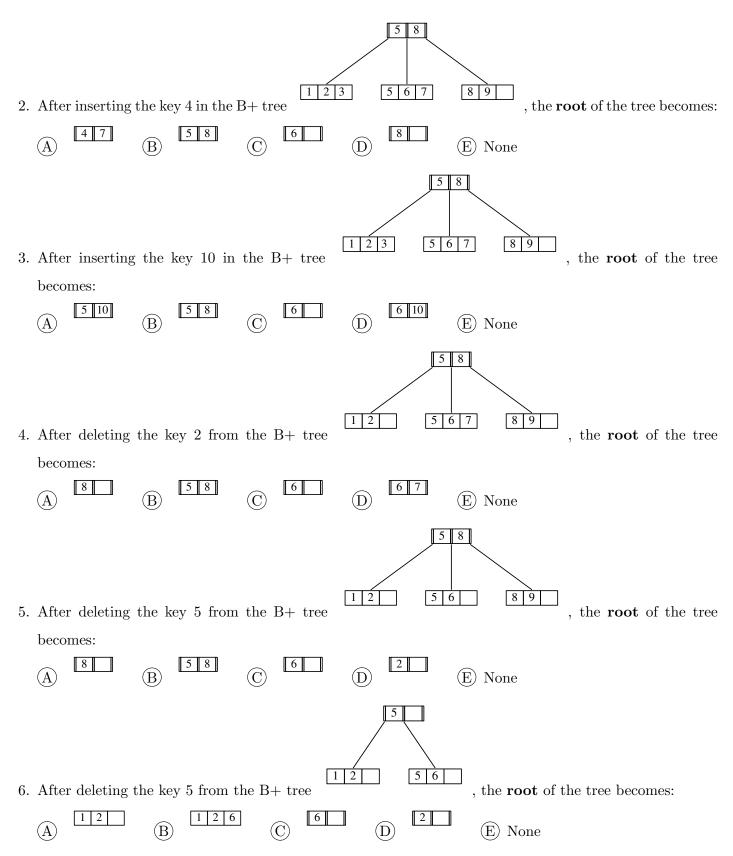
3. Coalesced chaining with cellar size 2 (do not change the hash function). Fill in the following table (put -1 if there is no next element):

Key	16	14	27	5	21	43	10	38	19	18	20
Position											
Next											

Choose the correct result in each of the following cases:

1. After inserting the key 6 in the B+ tree , the **root** of tree becomes:

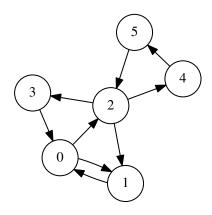
(A) (B) (C) (C) (D) (T) (E) None



7. A B+ tree of order 3 containing n keys has a height that is:

 $\bigcirc$  A O(n)  $\bigcirc$  B O(n/3)  $\bigcirc$  C  $O(n^3)$   $\bigcirc$  D  $O(\log n)$   $\bigcirc$  None

Consider the following graph.



1.	Give the	e adjacency	matrix of the graph.
	•••••		

2.	Give the adjacency list representation of the
	graph.
3.	What is the number of edges in the subgraph
	containing the nodes $\{1, 2, 3\}$ .
4.	What is the maximum number of edges in a di-
	rected graph with $n$ nodes (loops, edges from a
	node to itself, are not allowed)?

# **ADT Queue Specification**

- enqueue (Type e): **requires**: Queue Q is not full. **input**: Type e. **results**: Element e is added to the queue at its tail. **output**: none.
- serve (Type e): **requires**: Queue Q is not empty. **input**: none. **results**: the element at the head of Q is removed and its value assigned to e. **output**: Type e.
- length (int length): **requires**: none. **input**: none. **results**: The number of elements in the Queue Q is returned. **output**: length.
- full (boolean flag): **requires**: none. **input**: none. **results**: If Q is full then flag is set to true, otherwise flag is set to false. **output**: flag.

## **ADT Stack Specification**

- push(Type e): requires: Stack S is not full.

  input: Type e. results: Element e is added to
  the stack as its most recently added elements.

  output: none.
- pop(Type e): requires: Stack S is not empty.

  input: results: the most recently arrived element in S is removed and its value assigned to
  e. output: Type e.
- empty(boolean flag): **requires**: none. **input**: none. **results**: If Stack S is empty then flag is true, otherwise false. **output**: flag.
- full(boolean flag): **requires**: none. **input**: none. **results**: If S is full then Full is true, otherwise Full is false. **output**: flag.