

(a) Choose the correct frequency for every line as well as the total O of the following code:

```
1 int A = 0;
2 for (int i = 1; i <= n; i++)
3 for (int j = 0; j < i; j++)
4 A++;</pre>
```

- 1. Line 1: (A) 0 (B) 1 (C) 2 (D) n (E) A
- 2. Line 2: (A) A (B) i (C) i+1 (D) n (E) n+1
- 3. Line 3: (A) n^2 (B) n(n+1)/2 (C) n(n+1)/2+1 (D) $(n^2+3n)/2$ (E) n(n-1)/2-1
- 4. Line 4: (A) A^2 (B) n^2 (C) $(n^2 + 3n)/2$ (D) $n^2(n+1)/2 + 1$ (E) n(n+1)/2
- 5. Tightest Total $O: (\widehat{A}) n (\widehat{B}) n^2 (\widehat{C}) n^3 (\widehat{D}) n^4 (\widehat{E})$ None

(b) Choose the correct frequency for every line as well as the total O of the following code:

```
1 int i = 1;
2 while (i < n) {
3    i++;
4    if (i > 7) break;
5
```

- 1. Line 1: \bigcirc A 1 \bigcirc B 0 \bigcirc C i \bigcirc D n \bigcirc E n+1
- 2. Line 2: A 8 B 7 C n D n-1 E n+1
- 3. Lines 3 (and similarly 4): (A) n (B) n-1 (C) (G) (D) (G) (D) (D) (D)
- 4. Tightest Total O: (A) 1 (B) n (C) $\log(n)$ (D) n^2 (E) 2^n

(c) Choose the correct answer:

- $1. \ n^7 + n^4 + n^2 + \log n \text{ is :} \quad \textcircled{A} \ O(n^2) \quad \textcircled{B} \ O(n^4) \quad \textcircled{C} \ O(n^7) \quad \textcircled{D} \ O(\log(n)) \quad \textcircled{E} \ \text{None}$
- 2. $2^n + n!$ is : (A) $O(n^2)$ (B) $O(2^n)$ (C) O(n!) (D) $O(n^n)$ (E) None
- 3. $n + \log n^3 + 6$ is : (A) O(n) (B) $O(\log n^3)$ (C) $O(n \log n)$ (D) $O(n^3)$ (E) None
- 4. The time complexity of inserting an element in a heap of n elements is:
 - $\bigcirc A O(n^2)$ $\bigcirc B O(n)$ $\bigcirc C O(2^n)$ $\bigcirc D O(\log(n))$ $\bigcirc E$ None

CSC 212 Page 2 of 10

(a) Given the interfaces Map and LocNot below, write the method int nbNots(Map<String, Queue<LocNot>> ind, String w, double t1, double g1, double t2, double g2) which takes as input an index map, where the key is the word and data is a queue containing all notifications where the word appears. The method returns the number of notifications containing the word w and located within the rectangle having bottom left corner at (t1, g1) and upper right corner at (t2, g2).

```
public interface Map<K extends Comparable<</pre>
                                              public interface LocNot {
   K>, T> {
                                                double getLat(); // Latitude
                                                double getLng(); // Longitutde
  boolean empty();
                                                int getMaxNbRepeats();
  boolean full();
  T retrieve();
                                                int getNbRepeats();
  void update(T e);
                                                String getText();
  boolean find(K key);
  boolean insert(K key, T data);
  boolean remove(K key);
}
```

Complete the code below by choosing the correct answer:

```
1
    int nbNots(Map<String, Queue<LocNot>> ind, String w, double t1, double g1, double t2,
        double g2) {
      if (...)
 2
 3
        return 0;
 4
      int cpt = 0;
      Queue < LocNot > q = ...;
 5
 6
      ... {
 7
        LocNot not = ...;
 8
        . . . ;
 9
        double t = ...;
10
        double g = \dots;
        if (...)
11
12
           . . . ;
13
      }
14
15
```

- Line 2:
 - (A) if (ind.find(w))
 - (B) if (!ind.find(w))
 - (C) if (ind.find(w)== null)
 - \bigcirc if (ind.retrieve(w)== null)
 - (E) None
- Line 5:
 - (A) Queue<LocNot> q = ind.find(w);
 - (B) Queue < Loc Not > q = ind.remove(w);
 - Queue<LocNot> q = ind.retrieve(w);
 - \bigcirc Queue<LocNot> q = ind.retrieve();
 - (E) None

- Line 6:
 - (A) while (!q.empty()){
 - (B) for (int i = 0; i <= q.length(); i++){
 - (C) while (!q.last()){
 - (D) for (int i = 0; i < q.length(); i++){
 - (E) None
- Line 7:
 - (A) LocNot not = q.serve();
 - (B) LocNot not = q.head.data;
 - C LocNot not = q.retrieve();
 - (D) LocNot not = q.pop();
 - (E) None

CSC 212 Page 3 of 10

```
• Line 8:
```

```
(A) q.push(not);
```

- (B) q.serve();
- (C) q.insert(not);
- (D) q.enqueue();
- (E) None
- Line 9:
 - (A) double t = not.getLat();
 - (B) double t = q.retrieve().getLat();
 - (C) double t = q.serve().getLat();
 - (D) double t = q.pop().getLat();
 - (E) None
- Line 10:
 - (A) double g = q.serve().getLng();
 - (B) double g = q.retrieve().getLng();
 - (C) double g = not.getLng();
 - (D) double g = q.pop().getLng();
 - (E) None

- Line 11:
 - (A) if (t1<=t && t<=t2 && g1<=g && g<=g2)
 - (B) if (t1<=t && t>=t2 && g1<=g && g>=g2)
 - (C) if (t1<=t && t<=t2 || g1<=g && g<=g2)
 - (D) if (t1<=t && t<=t2 && g1>=g && g>=g2)
 - (E) None
- Line 12:
 - (A) return cpt;
 - (B) {cpt++; break;}
 - (C) cpt++;
 - (D) break;
 - (E) None
- Line 14:
 - (A) return cpt;
 - (B) return q.length()- cpt;
 - (C) return q.length()+ cpt;
 - (D) return q.length();
 - (E) None
- (b) Write the method Stack<LocNot> copyNots(Map<String, Stack<LocNot>> ind, String w) which takes as input an index map, where the key is the word and data is a stack containing all notifications where the word appears. The method returns a copy of the stack of notifications where the word w appears. If w does not exists, an empty stack is returned.

Complete the code below by choosing the correct answer:

```
1
    Stack < Loc Not > copyNots (Map < String, Stack < Loc Not >> ind, String w) {
 2
      Stack<LocNot> rs = new LinkedStack<LocNot>();
 3
      if (...)
 4
        return ...;
 5
      Stack < LocNot > ts = ...;
      Stack < LocNot > st = ...;
 6
 7
      while (...) {
 8
9
10
      while (...) {
11
        LocNot not = ...;
12
         . . . ;
13
         . . . ;
14
      }
15
      return rs;
16
```

```
• Line 3:
  (A) if (ind.find(w))
  (B) if (!ind.find(w))
  (C) if (ind.retrieve(w) == null)
  (\widehat{D}) if (ind.find(w)== null)
  (E) None
• Line 4:
  (A) return ind.retrieve();
  (B) return null;
  (C) return rs.empty();
  (D) return rs;
  (E) None
• Line 5:
  (A) Stack<LocNot> ts = null;
  (B) Stack<LocNot> ts = new LinkedStack<LocNot
  >();
  (C) Stack<LocNot> ts = new LinkedStack<String
  >();
  (\mathrm{D}) Stack<LocNot> ts = new Stack<LocNot>();
  (E) None
• Line 6:
  (A) Stack<LocNot> st = ind.serve();
  (B) Stack<LocNot> st = ind.retrieve();
  (C) Stack<LocNot> st = ind.pop();
  (D) Stack<LocNot> st = ind.retrieve(w);
  (E) None
• Line 7:
  (A) while (!st.empty()){
  (B) while (!st.last()){
  (C) while (st.empty()){
  (D) while (st.length()> 0){
```

```
(E) None
```

• Line 8:

```
(A) ts.pop(st.pop());
```

- (E) None
- Line 10:

```
(A) while (!ts.empty()){
```

- (E) None
- Line 11:

```
(A) LocNot not = ts.pop();
```

```
(B) LocNot not = rs.pop();
```

$$\bigcirc$$
 LocNot not = st.pop();

- (E) None
- Line 12:

```
(A) st.push(st.pop());
```

```
(B) st.push(not);
```

```
(C) st.push(rs.pop());
```

```
(D) st.push(ts.pop());
```

- (E) None
- Line 13:
 - (A) rs.push(ts.pop());
 - (B) rs.push(st.pop());
 - (C) rs.push(not);
 - (D) ts.push(rs.pop());
 - (E) None

(a) Write the method public boolean isBal(), member of the BT class, which returns true if the BT is a balanced, and false otherwise. A BT is balanced if for each node, the absolute difference in height of its two subtrees is at most 1. Assume you have a method called private int height(BTNode<T> p) that

returns the height the sub-tree p. The method isBal() makes a call to the recursive method private boolean isBalRec(BTNode<T> p). Choose the correct option to complete the code of these methods:

```
public boolean isBal() {
    ...
}

private boolean isBalRec(BTNode<T> p) {
    ...

    ...
}

...
}
```

```
1. Line 2:
```

```
A return isBalRec(root.left)|| isBalRec(
root.right);
```

- B return isBalRec(root.left)&& isBalRec(root.right);
- C return !isBalRec(root.left)&& !isBalRec(root.right);
- (D) return isBalRec(root);
- (E) None
- 2. Line 5:
 - (A) if (p == null)return false;
 - (B) if (p == null)return true;
 - (C) if (p != null)return true;
 - (D) if (p != null)return false;
 - (E) None
- 3. Line 6:
 - A if (Math.abs(height(p.right)- height(p.
 left))>= 1)return true;

```
B if (Math.abs(height(p.right) - height(p.
left))>= 2)return false;
```

```
C if (Math.abs(height(p.right) - height(p.
left)) <= 2) return false;</pre>
```

- D if (Math.abs(height(p.right)- height(p.
 left))!= 0)return false;
- (E) None
- 4. Line 7:
 - (A) return !isBalRec(p.left)&& !isBalRec(p.
 right);
 - B return isBalRec(p.left)+isBalRec(p.right)
 - C return isBalRec(p.left)&& isBalRec(p. right):
 - D return isBalRec(p.left)|| isBalRec(p.
 right);
 - (E) None

(b) Consider the function f below, member of DoubleLinkedList:

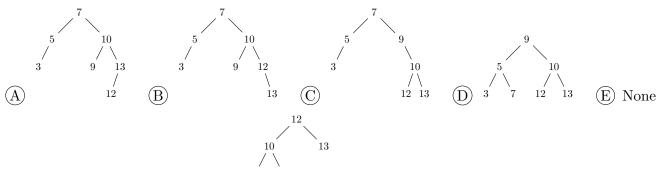
```
public void f(int n) {
  Node <T> p = head;
  for(int i = 0; i < n; i++) {
    if (p.next != null)
      p = p.next;
}

p.previous.next = p.next;
if (p.next != null)
    p.next.previous = p.previous;
p.next = head;
p.next.previous = p;
p.previous = null;
head = p;
}</pre>
```

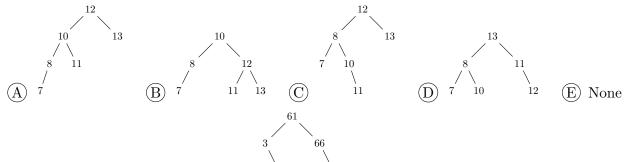
Choose the correct result in each of the following cases:

1. The list 1: A, B, C, D, E, after calling 1.f(3), 1 becomes:

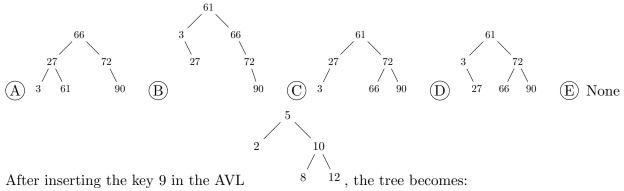
2. The list 1: A, B, C, D, E , after calling 1.f(1), 1 becomes:
3. The list 1: A, B, C, D, E , after calling 1.f(5), 1 becomes:
$igotimes_{A} A igotimes_{A} E, A, B, C, D igotimes_{A} C, D, E igotimes_{B} E, A, B, C, D igotimes_{B} None$
4. The list 1: A, B, C, D, E , after calling 1.f(2), 1 becomes:
Question 4
(a) Consider the following heap represented as an array: 2, 7, 5, 8, 20, 10, 12. Choose the correct answer
for every operation (all operations are done on the above heap).
1. Heap after inserting 1: (A) 1, 2, 5, 7, 20, 10, 12, 8 (B) 1, 2, 5, 7, 20, 10, 8, 12 (C) 2, 5, 7, 20,
10, 12, 8, 1
2. Heap after inserting 3 then 4: (A) 2, 3, 4, 5, 20, 10, 12, 8, 7 (B) 2, 3, 5, 4, 20, 10, 12, 8, 7
\bigcirc 2, 3, 4, 5, 8, 7, 20, 10, 12 \bigcirc D 2, 3, 4, 5, 8, 10, 12, 7, 20 \bigcirc E None
3. Heap after inserting 11 then deleting one key: (A) 11, 2, 7, 5, 8, 20, 10, 12 (B) 5, 3, 4, 20, 10,
$12, 8, 7$ \bigcirc \bigcirc $5, 7, 11, 8, 20, 10, 12$ \bigcirc \bigcirc $5, 7, 10, 8, 20, 12, 11$ \bigcirc None
4. Heap after deleting two keys: (A) 2, 7, 5, 8, 20 (B) 2, 5, 7, 20, 8 (C) 7, 8, 10, 12, 20
① $7, 10, 8, 12, 20$ ② None
(b) Suppose we have two heaps (5, 9, 6) and (7, 8, 10) represented as arrays and a key 12, what will be
the resultant heap after merging them? \textcircled{A} 12, 5, 9, 6, 7, 8, 10 \textcircled{B} 5, 9, 6, 12, 7, 8, 10 \textcircled{C} 5, 6,
$7, 9, 12, 8, 10$ \bigcirc $5, 9, 6, 7, 8, 10, 12 \bigcirc None$
(c) What is the result of a bottom-up min-heap construction of the array 5, 11, 2, 7, 16, 15, 4? (A) 2, 7,
4, 11, 16, 15, 5
© None
(d) What is the height of a heap containing 10 elements? \textcircled{A} 3 \textcircled{B} 10 \textcircled{C} 4 \textcircled{D} 5 \textcircled{E} None.
Question 5
Choose the correct result in each of the following cases (follow the the convention of replacing
with the smallest key in the right sub-tree when necessary):
1. After inserting the key 13 in the AVL 3 5 10 , the tree becomes:
, the tree becomes.



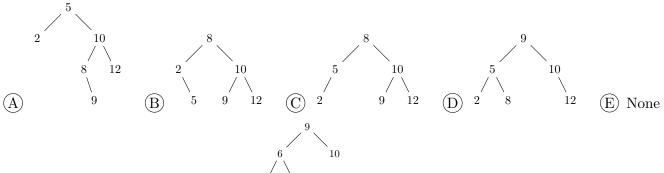
2. After inserting the key 7 in the AVL ⁸ , the tree becomes:



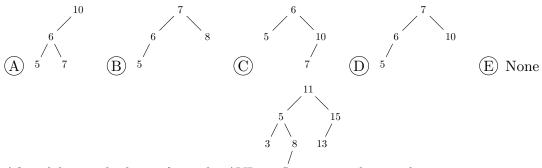
 $\overline{}^{72}$, the tree becomes: 3. After inserting the key 90 in the AVL



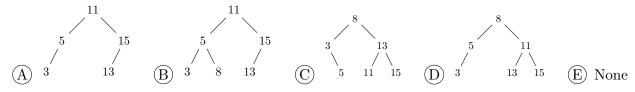
4. After inserting the key 9 in the AVL



5. After deleting the key 9 from the AVL 5 , the tree becomes:



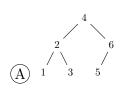
6. After deleting the key 7 from the AVL $\,$, the tree becomes:

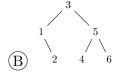


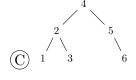


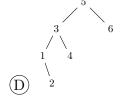
7. After deleting the key 11 from the AVL

, the tree becomes:









(E) None

Use the hash function H(key) = key%9 to store the sequence of keys 21, 15, 18, 12, 27, 30, 35, 19, 10 in a hash table of size 9. Use the following collision resolution strategies:

1. Linear rehashing (c=1). Fill in the following table:

Key	21	15	18	12	27	30	35	19	10
Position									
Number of probes									

2. External chaining. Fill in the following table:

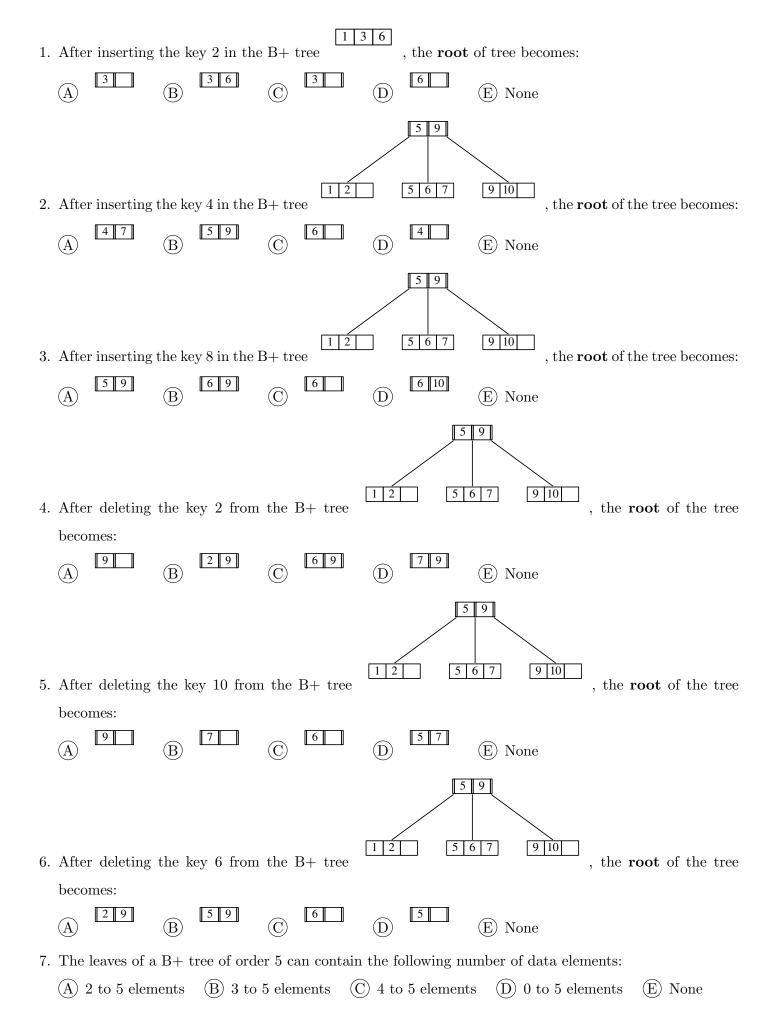
Key	21	15	18	12	27	30	35	19	10
Index of the list									

3. Coalesced chaining with cellar size 3 and address region size 7 (you must change the hash function to H(key) = key%7.) Fill in the following table (put -1 if there is no next element):

Key	21	15	18	12	27	30	35	19	10
Position									
Index of next element									

Question 7	points
------------	--------

Choose the correct result in each of the following cases (when possible, always borrow and transfer to the left):



1. Given the following adjacency list, draw the	
graph it represents.	
$\boxed{0} \rightarrow 1 \rightarrow 2$	2. Give the adjacency matrix representation of the
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	graph.
$\boxed{2} \rightarrow 0 \rightarrow 1 \rightarrow 3$	
$3 \rightarrow 1 \rightarrow 2$	
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	
$5 \rightarrow 4$	

ADT Queue Specification

• enqueue (Type e): requires: Queue Q is not full. input: Type e. results: Element e is added to the queue at its tail. output: none.

- serve (Type e): **requires**: Queue Q is not empty. **in- put**: none. **results**: the element at the head of Q is
 removed and its value assigned to e. **output**: Type e.
- length (int length): requires: none. input: none. results: The number of elements in the Queue Q is returned. output: length.
- full (boolean flag): requires: none. input: none. results: If Q is full then flag is set to true, otherwise flag is set to false. output: flag.

ADT Stack Specification

• push(Type e): requires: Stack S is not full. input: Type e. results: Element e is added to the stack as its most recently added elements. output: none.

3. This graph is connected: [True / False]

4. This graph has a cycle: [True / False]

- pop(Type e): requires: Stack S is not empty. input: results: the most recently arrived element in S is removed and its value assigned to e. output: Type e.
- empty(boolean flag): requires: none. input: none. results: If Stack S is empty then flag is true, otherwise false. output: flag.
- full(boolean flag): requires: none. input: none. results: If S is full then Full is true, otherwise Full is false. output: flag.