



King Saud University

College of Computer and Information Sciences

Department of Computer Science

Data Structures CSC 212**Final Exam Solution - Fall 2018**

Date: 15/12/2018

Duration: 3 hours

Guidelines: No calculators or any other electronic devices are allowed in this exam.

Student ID:

Name:

Section:

Instructor:

1	2.1	2.2	3.1	3.2	4	5	6	7	8	Total

Question 1.....16 (Part (a) and (b): 1pt / question. Part (c): 2pts for 1, 2, 3 and 1pt for 4) points

(a) Choose the correct frequency for every line as well as the total O of the following code:

```

1  int A = 0;
2  for (int i = 1; i <= n; i++)
3      for (int j = 0; j < i; j++)
4          A++;

```

- Line 1: (A) 0 (B) 1 (C) 2 (D) n (E) A
- Line 2: (A) A (B) i (C) $i + 1$ (D) n (E) $n + 1$
- Line 3: (A) n^2 (B) $n(n + 1)/2$ (C) $n(n + 1)/2 + 1$ (D) $(n^2 + 3n)/2$ (E) $n(n - 1)/2 - 1$
- Line 4: (A) A^2 (B) n^2 (C) $(n^2 + 3n)/2$ (D) $n^2(n + 1)/2 + 1$ (E) $n(n + 1)/2$
- Tightest Total O : (A) n (B) n^2 (C) n^3 (D) n^4 (E) None

(b) Choose the correct frequency for every line as well as the total O of the following code:

```

1  int i = 1;
2  while (i < n) {
3      i++;
4      if (i > 7) break;
5  }

```

- Line 1: (A) 1 (B) 0 (C) i (D) n (E) $n + 1$
- Line 2: (A) 8 (B) 7 (C) n (D) $n - 1$ (E) $n + 1$
- Lines 3 (and similarly 4): (A) n (B) $n - 1$ (C) 6 (D) 7 (E) 8
- Tightest Total O : (A) 1 (B) n (C) $\log(n)$ (D) n^2 (E) 2^n

(c) Choose the correct answer:

- $n^7 + n^4 + n^2 + \log n$ is : (A) $O(n^2)$ (B) $O(n^4)$ (C) $O(n^7)$ (D) $O(\log(n))$ (E) None
- $2^n + n!$ is : (A) $O(n^2)$ (B) $O(2^n)$ (C) $O(n!)$ (D) $O(n^n)$ (E) None
- $n + \log n^3 + 6$ is : (A) $O(n)$ (B) $O(\log n^3)$ (C) $O(n \log n)$ (D) $O(n^3)$ (E) None
- The time complexity of inserting an element in a heap of n elements is:
(A) $O(n^2)$ (B) $O(n)$ (C) $O(2^n)$ (D) $O(\log(n))$ (E) None

Question 2 10 (0.5pt / question) points

- (a) Given the interfaces `Map` and `LocNot` below, write the method `int nbNots(Map<String, Queue<LocNot>> ind, String w, double t1, double g1, double t2, double g2)` which takes as input an index map, where the key is the word and data is a queue containing all notifications where the word appears. The method returns the number of notifications containing the word `w` and located within the rectangle having bottom left corner at `(t1, g1)` and upper right corner at `(t2, g2)`.

<pre>public interface Map<K extends Comparable< K>, T> { boolean empty(); boolean full(); T retrieve(); void update(T e); boolean find(K key); boolean insert(K key, T data); boolean remove(K key); }</pre>	<pre>public interface LocNot { double getLat(); // Latitude double getLng(); // Longitude int getMaxNbRepeats(); int getNbRepeats(); String getText(); }</pre>
--	--

Complete the code below by choosing the correct answer:

```
1  int nbNots(Map<String, Queue<LocNot>> ind, String w, double t1, double g1, double t2,
2      double g2) {
3      if (...)
4          return 0;
5      int cpt = 0;
6      Queue<LocNot> q = ...;
7      ... {
8          LocNot not = ...;
9          ...;
10         double t = ...;
11         double g = ...;
12         if (...)
13             ...;
14     }
15     ...
16 }
```

- Line 2:

- ☐ (A) `if (ind.find(w))`
☒ (B) `if (!ind.find(w))`
☐ (C) `if (ind.find(w) == null)`
☐ (D) `if (ind.retrieve(w) == null)`
☐ (E) None

- Line 5:

- ☐ (A) `Queue<LocNot> q = ind.find(w);`
☐ (B) `Queue<LocNot> q = ind.remove(w);`
☐ (C) `Queue<LocNot> q = ind.retrieve(w);`
☒ (D) `Queue<LocNot> q = ind.retrieve();`
☐ (E) None

- Line 6:

- ☐ (A) `while (!q.empty()){`
☐ (B) `for (int i = 0; i <= q.length(); i++){`
☐ (C) `while (!q.last()){`
☒ (D) `for (int i = 0; i < q.length(); i++){`
☐ (E) None

- Line 7:

- ☒ (A) `LocNot not = q.serve();`
☐ (B) `LocNot not = q.head.data;`
☐ (C) `LocNot not = q.retrieve();`
☐ (D) `LocNot not = q.pop();`
☐ (E) None

- Line 8:

☐ (A) `q.push(not);`
☐ (B) `q.serve();`
☐ (C) `q.insert(not);`
☐ (D) `q.enqueue();`
☒ (E) **None**

- Line 9:

☒ (A) `double t = not.getLat();`
☐ (B) `double t = q.retrieve().getLat();`
☐ (C) `double t = q.serve().getLat();`
☐ (D) `double t = q.pop().getLat();`
☐ (E) **None**

- Line 10:

☐ (A) `double g = q.serve().getLng();`
☐ (B) `double g = q.retrieve().getLng();`
☒ (C) `double g = not.getLng();`
☐ (D) `double g = q.pop().getLng();`
☐ (E) **None**

- Line 11:

☒ (A) `if (t1<=t && t<=t2 && g1<=g && g<=g2)`
☐ (B) `if (t1<=t && t>=t2 && g1<=g && g>=g2)`
☐ (C) `if (t1<=t && t<=t2 || g1<=g && g<=g2)`
☐ (D) `if (t1<=t && t<=t2 && g1>=g && g>=g2)`
☐ (E) **None**

- Line 12:

☐ (A) `return cpt;`
☐ (B) `{cpt++; break;}`
☒ (C) `cpt++;`
☐ (D) `break;`
☐ (E) **None**

- Line 14:

☒ (A) `return cpt;`
☐ (B) `return q.length()- cpt;`
☐ (C) `return q.length()+ cpt;`
☐ (D) `return q.length();`
☐ (E) **None**

- (b) Write the method `Stack<LocNot> copyNots(Map<String, Stack<LocNot>> ind, String w)` which takes as input an index map, where the key is the word and data is a stack containing all notifications where the word appears. The method returns **a copy** of the stack of notifications where the word `w` appears. If `w` does not exist, an empty stack is returned.

Complete the code below by choosing the correct answer:

```

1 Stack<LocNot> copyNots(Map<String, Stack<LocNot>> ind, String w) {
2     Stack<LocNot> rs = new LinkedStack<LocNot>();
3     if (...)
4         return ...;
5     Stack<LocNot> ts = ...;
6     Stack<LocNot> st = ...;
7     while (...) {
8         ...;
9     }
10    while (...) {
11        LocNot not = ...;
12        ...;
13        ...;
14    }
15    return rs;
16 }

```

- Line 3:

- (A) `if (ind.find(w))`
- (B) `if (!ind.find(w))`
- (C) `if (ind.retrieve(w) == null)`
- (D) `if (ind.find(w) == null)`
- (E) None

- Line 4:

- (A) `return ind.retrieve();`
- (B) `return null;`
- (C) `return rs.empty();`
- (D) `return rs;`
- (E) None

- Line 5:

- (A) `Stack<LocNot> ts = null;`
- (B) `Stack<LocNot> ts = new LinkedStack<LocNot>();`
- (C) `Stack<LocNot> ts = new LinkedStack<String>();`
- (D) `Stack<LocNot> ts = new Stack<LocNot>();`
- (E) None

- Line 6:

- (A) `Stack<LocNot> st = ind.serve();`
- (B) `Stack<LocNot> st = ind.retrieve();`
- (C) `Stack<LocNot> st = ind.pop();`
- (D) `Stack<LocNot> st = ind.retrieve(w);`
- (E) None

- Line 7:

- (A) `while (!st.empty()){`
- (B) `while (!st.last()){`
- (C) `while (st.empty()){`
- (D) `while (st.length() > 0){`

- (E) None

- Line 8:

- (A) `ts.pop(st.pop());`
- (B) `st.push(st.push());`
- (C) `ts.enqueue(st.serve());`
- (D) `ts.push(st.pop());`
- (E) None

- Line 10:

- (A) `while (!ts.empty()){`
- (B) `while (!ts.last()){`
- (C) `while (!st.empty()){`
- (D) `while (ts.empty()){`
- (E) None

- Line 11:

- (A) `LocNot not = ts.pop();`
- (B) `LocNot not = rs.pop();`
- (C) `LocNot not = ts.push();`
- (D) `LocNot not = st.pop();`
- (E) None

- Line 12:

- (A) `st.push(st.pop());`
- (B) `st.push(not);`
- (C) `st.push(rs.pop());`
- (D) `st.push(ts.pop());`
- (E) None

- Line 13:

- (A) `rs.push(ts.pop());`
- (B) `rs.push(st.pop());`
- (C) `rs.push(not);`
- (D) `ts.push(rs.pop());`
- (E) None

Question 3 10 (Part (a) 1.5pts / question. Part (b) 1pt / question) points

- (a) Write the method `public boolean isBal()`, member of the BT class, which returns true if the BT is a balanced, and false otherwise. A BT is balanced if for each node, the absolute difference in height of its two subtrees is at most 1. Assume you have a method called `private int height(BTNode<T> p)` that

returns the height the sub-tree `p`. The method `isBal()` makes a call to the recursive method **private** `boolean isBalRec(BTNode<T> p)`. Choose the correct option to complete the code of these methods:

```

1 public boolean isBal() {
2     ...
3 }
4 private boolean isBalRec(BTNode<T> p) {
5     ...
6     ...
7     ...
8 }

```

1. Line 2:

- ☐ (A) `return isBalRec(root.left) || isBalRec(root.right);`
- ☐ (B) `return isBalRec(root.left) && isBalRec(root.right);`
- ☐ (C) `return !isBalRec(root.left) && !isBalRec(root.right);`
- ☒ (D) `return isBalRec(root);`
- ☐ (E) None

2. Line 5:

- ☐ (A) `if (p == null) return false;`
- ☒ (B) `if (p == null) return true;`
- ☐ (C) `if (p != null) return true;`
- ☐ (D) `if (p != null) return false;`
- ☐ (E) None

3. Line 6:

- ☐ (A) `if (Math.abs(height(p.right) - height(p.left)) >= 1) return true;`

☒ (B) `if (Math.abs(height(p.right) - height(p.left)) >= 2) return false;`

☐ (C) `if (Math.abs(height(p.right) - height(p.left)) <= 2) return false;`

☐ (D) `if (Math.abs(height(p.right) - height(p.left)) != 0) return false;`

☐ (E) None

4. Line 7:

☐ (A) `return !isBalRec(p.left) && !isBalRec(p.right);`

☐ (B) `return isBalRec(p.left) + isBalRec(p.right);`

☒ (C) `return isBalRec(p.left) && isBalRec(p.right);`

☐ (D) `return isBalRec(p.left) || isBalRec(p.right);`

☐ (E) None

(b) Consider the function `f` below, member of `DoubleLinkedList`:

```

public void f(int n) {
    Node<T> p = head;
    for(int i = 0; i < n; i++) {
        if (p.next != null)
            p = p.next;
    }
    p.previous.next = p.next;
    if (p.next != null)
        p.next.previous = p.previous;
    p.next = head;
    p.next.previous = p;
    p.previous = null;
    head = p;
}

```

Choose the correct result in each of the following cases:

1. The list 1: A, B, C, D, E , after calling $1.f(3)$, 1 becomes:

(A) B, C, D, E (B) D, A, B, C, E (C) E, B, C, D (D) A, D, E, B, C (E) None
2. The list 1: A, B, C, D, E , after calling $1.f(1)$, 1 becomes:

(A) A, B, C (B) E, A, B, C, D (C) B, C, D, E, A (D) B, A, C, D, E (E) None
3. The list 1: A, B, C, D, E , after calling $1.f(5)$, 1 becomes:

(A) A (B) E, A, B, C, D (C) A, B, C, D, E (D) E, A, B, C, D (E) None
4. The list 1: A, B, C, D, E , after calling $1.f(2)$, 1 becomes:

(A) *empty* (B) E, D, C, B, A (C) C, A, B, D, E (D) E, C, D (E) None

Question 4 14 (2pts / question) points

(a) Consider the following heap represented as an array: 2, 7, 5, 8, 20, 10, 12. Choose the correct answer for every operation (all operations are done on the above heap).

1. Heap after inserting 1: (A) **1, 2, 5, 7, 20, 10, 12, 8** (B) 1, 2, 5, 7, 20, 10, 8, 12 (C) 2, 5, 7, 20, 10, 12, 8, 1 (D) 2, 5, 7, 20, 10, 12, 1, 8 (E) None
2. Heap after inserting 3 then 4: (A) 2, 3, 4, 5, 20, 10, 12, 8, 7 (B) **2, 3, 5, 4, 20, 10, 12, 8, 7** (C) 2, 3, 4, 5, 8, 7, 20, 10, 12 (D) 2, 3, 4, 5, 8, 10, 12, 7, 20 (E) None
3. Heap after inserting 11 then deleting one key: (A) 11, 2, 7, 5, 8, 20, 10, 12 (B) 5, 3, 4, 20, 10, 12, 8, 7 (C) 5, 7, 11, 8, 20, 10, 12 (D) 5, 7, 10, 8, 20, 12, 11 (E) **None**
4. Heap after deleting two keys: (A) 2, 7, 5, 8, 20 (B) 2, 5, 7, 20, 8 (C) **7, 8, 10, 12, 20** (D) 7, 10, 8, 12, 20 (E) None

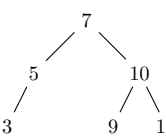
(b) Suppose we have two heaps (5, 9, 6) and (7, 8, 10) represented as arrays and a key 12, what will be the resultant heap after merging them? (A) 12, 5, 9, 6, 7, 8, 10 (B) 5, 9, 6, 12, 7, 8, 10 (C) **5, 6, 7, 9, 12, 8, 10** (D) 5, 9, 6, 7, 8, 10, 12 (E) None

(c) What is the result of a bottom-up min-heap construction of the array 5, 11, 2, 7, 16, 15, 4? (A) **2, 7, 4, 11, 16, 15, 5** (B) 5, 11, 2, 7, 16, 15, 4 (C) 2, 7, 5, 11, 16, 15, 4 (D) 2, 4, 7, 11, 16, 15, 5 (E) None

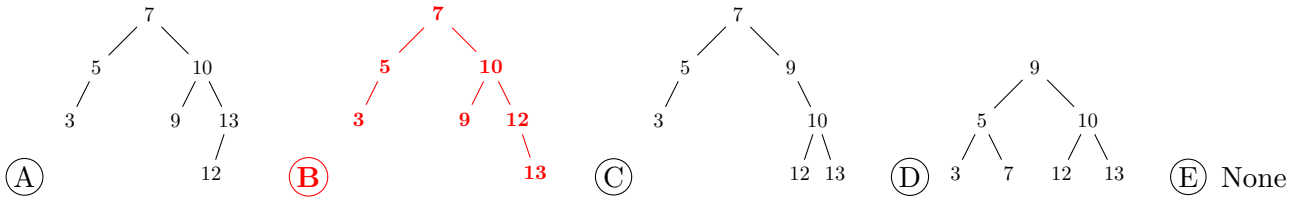
(d) What is the height of a heap containing 10 elements? (A) **3** (B) 10 (C) **4** (D) 5 (E) None.

Question 5 14 (2pts / question) points

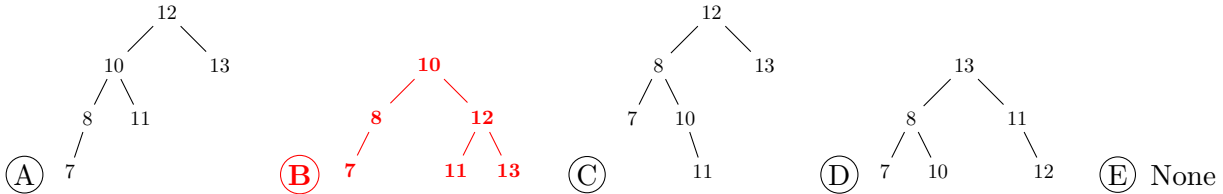
Choose the correct result in each of the following cases (follow the the convention of replacing with the smallest key in the right sub-tree when necessary):



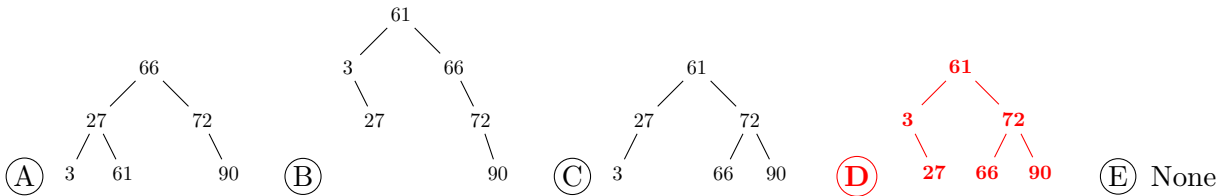
1. After inserting the key 13 in the AVL tree, the tree becomes:



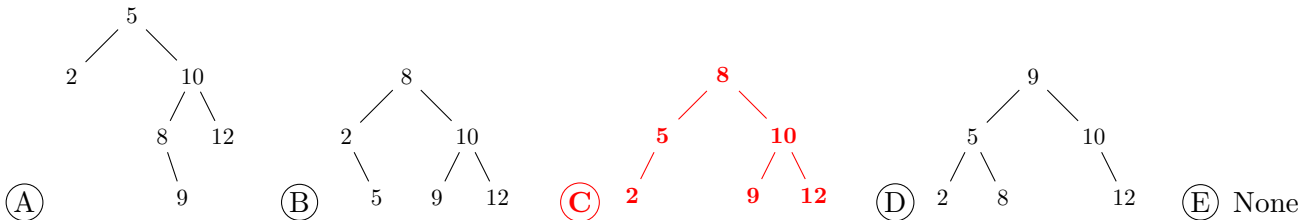
2. After inserting the key 7 in the AVL tree, the tree becomes:



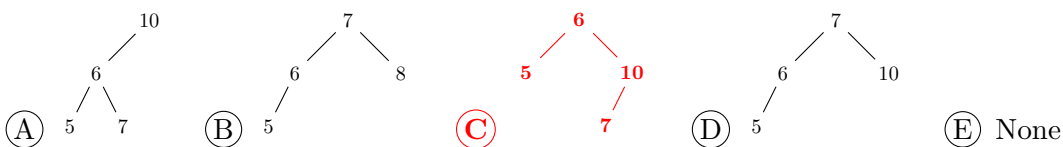
3. After inserting the key 90 in the AVL tree, the tree becomes:



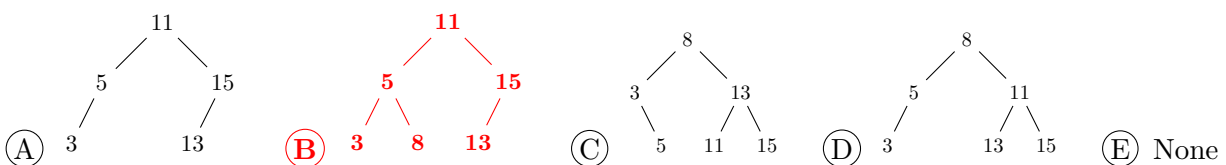
4. After inserting the key 9 in the AVL tree, the tree becomes:

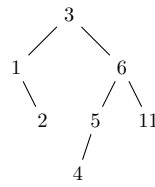


5. After deleting the key 9 from the AVL tree, the tree becomes:

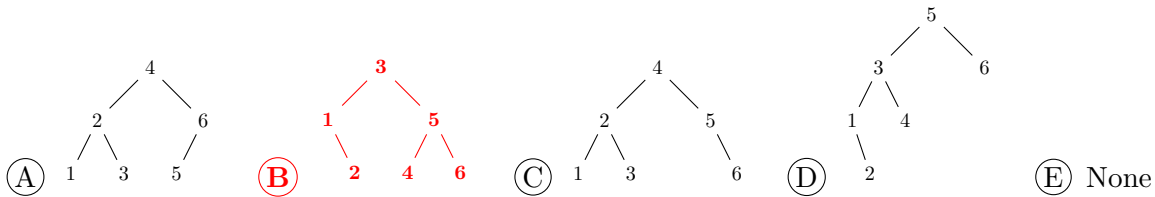


6. After deleting the key 7 from the AVL tree, the tree becomes:





7. After deleting the key 11 from the AVL , the tree becomes:



Question 6 14 ((45 - number of mistakes) / 45 * 14 then rounded up to the nearest 0.5) points

Use the hash function $H(key) = key \% 9$ to store the sequence of keys 21, 15, 18, 12, 27, 30, 35, 19, 10 in a hash table of size 9. Use the following collision resolution strategies:

1. Linear rehashing (c=1). Fill in the following table:

Key	21	15	18	12	27	30	35	19	10
Position	3	6	0	4	1	5	8	2	7
Number of probes	1	1	1	2	2	3	1	2	7

2. External chaining. Fill in the following table:

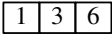
Key	21	15	18	12	27	30	35	19	10
Index of the list	3	6	0	3	0	3	8	1	1

3. Coalesced chaining with cell size 3 and address region size 7 (you must change the hash function to $H(key) = key \% 7$.) Fill in the following table (put -1 if there is no next element):

Key	21	15	18	12	27	30	35	19	10
Position	0	1	4	5	6	2	9	8	3
Index of next element	9	-1	-1	8	-1	-1	-1	-1	-1

Question 7 14 (2pts / question) points

Choose the correct result in each of the following cases (when possible, always borrow and transfer to the left):

1. After inserting the key 2 in the B+ tree , the **root** of tree becomes:

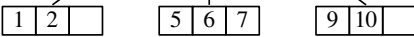
- (A)  (B)  (C)  (D)  (E) None

2. After inserting the key 4 in the B+ tree , the **root** of the tree becomes:

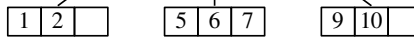
- (A)  (B)  (C)  (D)  (E) None





3. After inserting the key 8 in the B+ tree , the **root** of the tree becomes:

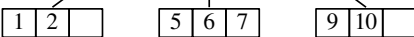
- (A)  (B)  (C)  (D)  (E) None


4. After deleting the key 2 from the B+ tree , the **root** of the tree becomes:

- (A)  (B)  (C)  (D)  (E) None

5. After deleting the key 10 from the B+ tree , the **root** of the tree becomes:

- (A)  (B)  (C)  (D)  (E) None

6. After deleting the key 6 from the B+ tree , the **root** of the tree becomes:

- (A)  (B)  (C)  (D)  (E) None

7. The leaves of a B+ tree of order 5 can contain the following number of data elements:

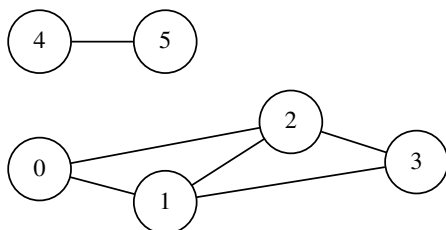
- (A) 2 to 5 elements (B) **3 to 5 elements** (C) 4 to 5 elements (D) 0 to 5 elements (E) None

Question 8 8 points

(1) 3pts: -0.5pts for every incorrect edge (missing or extra); (2) 3pts: -0.25pts for every incorrect edge (missing or extra) and round up to the nearest 0.5; (3) 1pt; (4) 1pt.

1. Given the following adjacency list, draw the graph it represents.

0	→ 1 → 2
1	→ 0 → 2 → 3
2	→ 0 → 1 → 3
3	→ 1 → 2
4	→ 5
5	→ 4



2. Give the adjacency matrix representation of the graph.

	0	1	2	3	4	5
0		1	1			
1	1		1	1		
2	1	1		1		
3		1	1			
4						1
5					1	

3. This graph is connected: **[False]**

4. This graph has a cycle: **[True]**

ADT Queue Specification

- enqueue (Type e): **requires:** Queue Q is not full. **input:** Type e. **results:** Element e is added to the queue at its tail. **output:** none.
- serve (Type e): **requires:** Queue Q is not empty. **input:** none. **results:** the element at the head of Q is removed and its value assigned to e. **output:** Type e.
- length (int length): **requires:** none. **input:** none. **results:** The number of elements in the Queue Q is returned. **output:** length.
- full (boolean flag): **requires:** none. **input:** none. **results:** If Q is full then flag is set to true, otherwise flag is set to false. **output:** flag.

ADT Stack Specification

- push (Type e): **requires:** Stack S is not full. **input:** Type e. **results:** Element e is added to the stack as its most recently added elements. **output:** none.
- pop (Type e): **requires:** Stack S is not empty. **input:** none. **results:** the most recently arrived element in S is removed and its value assigned to e. **output:** Type e.
- empty (boolean flag): **requires:** none. **input:** none. **results:** If Stack S is empty then flag is true, otherwise false. **output:** flag.
- full (boolean flag): **requires:** none. **input:** none. **results:** If S is full then Full is true, otherwise Full is false. **output:** flag.