

King Saud University

College of Computer and Information Sciences

Department of Computer Science

Data Structures CSC 212

Final Exam - Fall 2017

Date: 30/12/2017 Duration: 3 hours

Guidelines

•No calculators or any other electronic devices are allowed in this exam.

•Use a pencil in choice questions.

Student ID: Name:											
Section: Instructor:											
1	2.1	2.2	3.1	3.2	4	5	6	7	8	Total	

(a) Choose the correct frequency for every line as well as the total O of the following code:

```
1  sum = 1;
2  for (i = 1; i <= n; i++) {
3    sum+= i;
4    for (j = i; j >= 2; j--)
5    sum--;}
```

- 1. Line 1: (\widehat{A}) 1 (\widehat{B}) 2 (\widehat{C}) 3 (\widehat{D}) n (\widehat{E}) 2n
- 2. Line 2: (A) n (B) n+1 (C) n-1 (D) n+2 (E) n-2
- 3. Line 3: (A) n (B) n+1 (C) n-1 (D) n+2 (E) n-2
- 4. Line 4: (A) n^2 (B) n(n-1)/2 (C) (2n+1)/2 (D) (2n-1)/2 (E) n(n+1)/2
- 5. Line 5: (A) n^2 (B) n(n-1)/2 (C) (2n+1)/2 (D) (2n-1)/2 (E) n(n+1)/2
- 6. Total O: (A) 1 (B) n (C) n^2 (D) $n \log(n)$ (E) n^3

(b) Choose the correct frequency for every line as well as the total O of the following code:

```
1 count = 0;
2 for (i = 1; i < n+1; i++)
3   count ++;
4 for (j = 0; j <= count; j++)
5   k = j+1;</pre>
```

- 1. Line 1: (A) 0 (B) 1 (C) 2 (D) n (E) n^2
- 2. Line 2: (A) n (B) n+1 (C) n-1 (D) n+2 (E) n-2
- 3. Line 3: (\widehat{A}) n (\widehat{B}) n+1 (\widehat{C}) n-1 (\widehat{D}) n+2 (\widehat{E}) n-2
- 4. Line 4: (A) count + 2 (B) n + 1 (C) n 1 (D) n + 2 (E) n(n + 1)/2
- 5. Line 5: (A) count + 1 (B) n + 1 (C) n 1 (D) n + 2 (E) n(n 1)/2
- 6. Total O: (A) 1 (B) n (C) n^2 (D) $n \log(n)$ (E) n^3

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(c)	Choose the correct answer:
	1. $n^2 + n \log n^4$ is: (A) $O(n)$ (B) $O(n^2)$ (C) $O(n \log(n))$ (D) $O(n^4)$ (E) None
	2. $n^2 + 1000n$ is : (A) $O(n)$ (B) $O(n^2)$ (C) $O(n\log(n))$ (D) $O(nn^2)$ (E) None
	3. $n^4 \log n + n!$ is : (A) $O(n!)$ (B) $O(n^4)$ (C) $O(n^5)$ (D) $O(\log(n))$ (E) None
	4. Algorithm A is $O(n)$, and Algorithm B is $O(2n)$. Given the same input:
	(A) A always finishes before B. (B) B always finishes before A. (C) A and B finish at the same
	time. (D) B requires double the time taken by A. (E) None
3	stion 2
(a)	Given a queue of time intervals represented as pairs of integers, write the method public static Queue
	Pair <integer, integer="">> getIntervals(Queue<pair<integer, integer="">> q, int start, int end), which returns</pair<integer,></integer,>
	all intervals of q intersecting the interval [start, end]. The input q must not change. Assume that al
	intervals in q are valid (that is first <= second), non-overlapping and ordered in chronological order
	The class Pair is given below.
	<pre>public class Pair<u, v=""> { public U first; public V second; Pair(U first, V second) { this.first = first; this.second = second; } }</u,></pre>

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(b) Consider a stack of decreasing time intervals, that is, starting from the top, each interval contains the next. Write the method public static Pair<Integer, Integer> smallest(Stack<Pair<Integer, Integer>> st, int t), which returns the smallest interval containing t if it exists, null otherwise. Assume that all intervals in st are valid (that is first <= second).

Example 0.1. If st : $\{[0,8],[1,6],[1,5],[2,4]\}$, then smallest(st, 1) returns [1,5], smallest(st, 3) returns [2,4], smallest(st, 9) returns null.

Complete the code below by choosing the correct answer:

```
public static Pair<Integer, Integer>smallest(Stack<Pair<Integer, Integer>> st, int t){
2
      Pair < Integer , Integr > itm = null;
3
4
      while (!st.empty()) {
        Pair < Integer , Integr > it = st.pop();
5
6
        if (...)
7
8
          itm = it;
9
        else
10
11
12
      while (...) {
13
14
15
      return itm; }
```

- 1. Line 2:
 - A Queue<Pair<Integer, Integer>> r = new
 LinkedQueue<Pair<Integer, Integer>>();
 - B Stack<Integer> r = new LinkedStack<
 Integer>();
 - C List<Pair<Integer, Integer>> r = new
 LinkedList<Pair<Integer, Integer>>();
 - D Stack<Pair<Integer, Integer>> r = new
 LinkedStack<Pair<Integer, Integer>>();
 - (E) None
- 2. Line 6:
 - (A) r.push(it.first);
 - (B) r.insert(it);
 - (C) r.enqueue(it);
 - (D) r.push(it);
 - (E) None
- 3. Line 7:
 - (A) if (it.first < t && t <= it.second)
 - $\stackrel{\textstyle igoreal}{igorplus}$ if (it.first <= t && t <= it.second)
 - (C) if (it.first < t || it.second > t)

- $\stackrel{\textstyle \frown}{
 m (D)}$ if (it.first <= t && it.second <= t)
- (E) None
- 4. Line 10:
 - (A) r.serve();
 - B break;
 - (C) r.pop();
 - (D) r.findNext();
 - (E) None
- 5. Line 12:
 - (A) while (r.empty()){
 - (B) while (!r.empty()){
 - (C) while (r.pop()!= null){
 - (D) while (r.length()!= 0){
 - (E) None
- 6. Line 13:
 - (A) st.push(r.pop());
 - (B) st.push(r.serve());
 - (C) st.push(r.retrieve()); r.findNext();
 - (D) st.push(r.push());
 - (E) None

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(a) The method **private** BTNode<T> mirrorCopy(BTNode<T> t) creates **recursively** a mirror copy of the subtree

```
t. Choose the correct option to complete the code of this method:
```

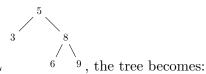
- 1. Line 2:
 - (A) if (t.left == null || t.right == null)
 - (B) if (t.left == null && t.right == null)
 - (C) if (t == null)
 - (D) if (root != null)
 - (E) None
- 2. Line 3:
 - (A) return null;
 - (B) return root;
 - (C) return mirrorCopy(root);
 - (D) return mirrorCopy(t);
 - (E) None
- 3. Line 4:
 - (A) BTNode<T> p = new BTNode<T>(t.data);
 - (B) BTNode<T> p = new BTNode<T>(root);
 - (C) BTNode<T> p = new BTNode<T>(t);
 - (D) BTNode<T> p = new BTNode<T>(root.data);
 - (E) None

- 4. Line 5:
 - (A) p.right = mirrorCopy(t.left);
 - (B) t.left = mirrorCopy(t.left);
 - (C) p.right = mirrorCopy(t.right);
 - (D) t.left = mirrorCopy(t.right);
 - (E) None
- 5. Line 6:
 - (A) t.right = mirrorCopy(t.left);
 - (B) p.left = mirrorCopy(t.left);
 - (C) p.left = mirrorCopy(t.right);
 - (D) t.right = mirrorCopy(t.right);
 - (E) None
- 6. Line 7:
 - (A) return p;
 - (B) return mirrorCopy(t);
 - (C) mirrorCopy(t.left); mirrorCopy(t.right);
 - (D) return t;
 - (E) None
- (b) Consider the function f below, member of DoubleLinkedList:

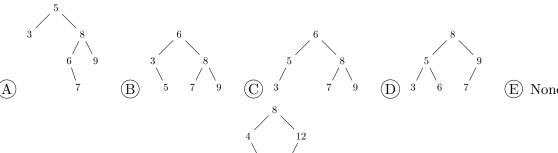
```
public void f(int n) {
   Node < T > p = head; Node < T > q = null;
   for (int i = 0; i < n; i++) {
      q = p;
      p = p.next;
   }
   if (p != null) {
      p.previous = null;
      while (p.next != null)
           p = p.next;
      p.next = head;
      head = q.next;
      q.next = null;
   }
}</pre>
```

Choose the correct result in each of the following cases:
1. The list 1: A, B, C, D, E , after calling 1.f(2), 1 becomes:
2. The list 1: A, B, C, D, E , after calling 1.f(0), 1 becomes:
3. The list 1: A, B, C, D, E , after calling 1.f(5), 1 becomes:
4. The list 1: A, B, C, D, E , after calling 1.f(1), 1 becomes:
(A) A (B) E, A, B, C, D (C) C, D, E, A, B (D) B, C, D, E, A (E) None
Question 4
(a) Consider the following heap represented as an array: 3, 7, 9, 13, 8, 11. Choose the correct answer for
every operation (all operations are done on the above heap).
1. Heap after inserting 5:
2. Heap after inserting 10:
3. Heap after inserting 2:
4. Heap after deleting one key:
igoplus 7,13,9,11,8 igoplus 7,8,9,13,11 igoplus 9,7,11,13,8 igoplus 7,9,8,11,13 igoplus None
5. Heap after deleting two keys:
\bigcirc 7,13,9,11 \bigcirc 8,11,9,13 \bigcirc 7,8,9,13 \bigcirc 13,9,8,11 \bigcirc None
(b) What is the result of a bottom-up min-heap construction of the following array: 2,4,6,3,5,1?
A 1,2,3,5,4,6 B 2,1,3,4,5,6 C 1,3,2,4,5,6 D 1,3,2,5,6,4 E None.
(c) Choose the correct answer:
1. What is the height of a heap of size k ?
$\textcircled{A} \log \log k \textcircled{B} \ k/2 \textcircled{C} \ k \log k \textcircled{D} \log k \textcircled{E} $ None.
2. Bottom-up heap construction is:
$igotimes_{O}O(n)$ $igotimes_{O}O(\log n)$ $igotimes_{O}O(n\log n)$ $igotimes_{O}O(n^2)$ $igotimes_{O}O(n^2)$ $igotimes_{O}O(n^2)$ $igotimes_{O}O(n^2)$
3. The enqueue operation in a heap priority queue is:
$ig(A) \ O(1) ig(B) \ O(\log n) ig(C) \ O(n) ig(D) \ O(n \log n) ig(E) \ ext{None.}$

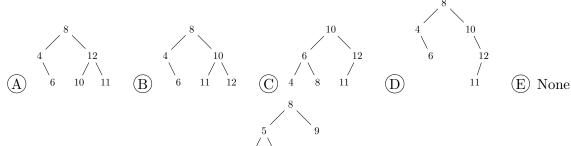
Choose the correct result in each of the following cases:



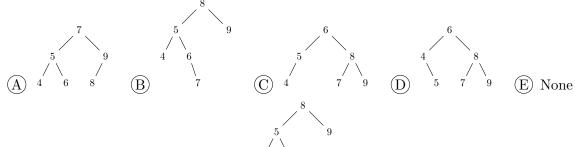
1. After inserting the key 7 in the AVL



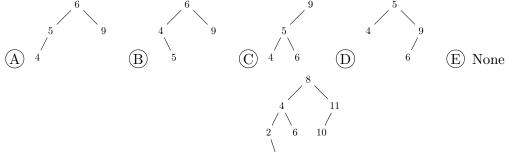
2. After inserting the key 11 in the AVL $^{\circ}6$ $^{\circ}10$, the tree becomes:



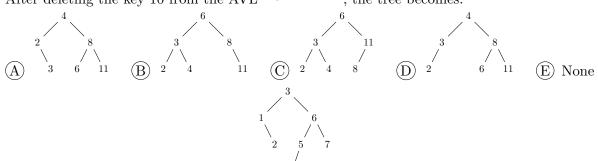
3. After inserting the key 7 in the AVL $_{4}^{4}$ $_{6}^{6}$, the tree becomes:



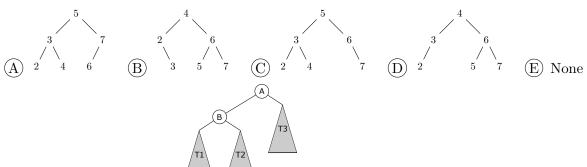
4. After deleting the key 8 from the AVL 4 6, the tree becomes:



5. After deleting the key 10 from the AVL $\stackrel{\backslash}{3}$, the tree becomes:



6. After deleting the key 1 from the AVL 4, the tree becomes:



7. Consider the following tree _____ . If the balance of A is -2 and that of B is 0, then after performing a single right rotation at A, then:

(a) The balance of A becomes:	
-------------------------------	--

1. Linear rehashing (c=1). Fill in the following table:

Key	16	14	27	5	21	43	10	38	19	18	20
Position											
Number of probes											

2. External chaining. Fill in the following table:

Key	16	14	27	5	21	43	10	38	19	18	20
List position											

3. Coalesced chaining with cellar size 2 (do not change the hash function). Fill in the following table (put -1 if there is no next element):

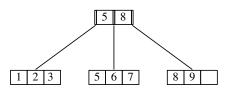
Key	16	14	27	5	21	43	10	38	19	18	20
Position											
Next											

1. After inserting the key 6 in the B+ tree

Choose the correct result in each of the following cases:

3 | 5 | 7, the **root** of tree becomes:

(A) (B) (C) (6) (D) (T) (E) None



2. After inserting the key 4 in the B+ tree

, the ${f root}$ of the tree becomes:

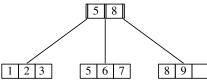


(B) 5 8



(D) [8]

(E) None



3. After inserting the key 10 in the B+ tree becomes:



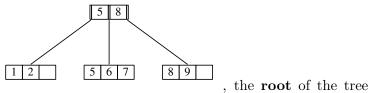
(A) 5 10

(B) 5 8



(D) 6 10

(E) None



4. After deleting the key 2 from the B+ tree becomes:

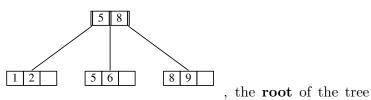




(C) 6



(E) None



5. After deleting the key 5 from the B+ tree becomes:

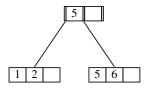




 $\widehat{\mathbf{C}}$



(E) None



6. After deleting the key 5 from the B+ tree

, the **root** of the tree becomes:



 $\widehat{\mathbf{B}}$ 12



(D) 2

(E) None

7. A B+ tree of order 3 containing n keys has a height that is:

$$\widehat{(A)} O(n)$$

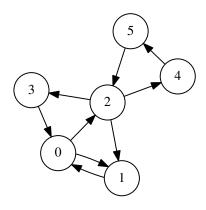
$$\bigcirc$$
 $O(n/3)$

$$\bigcirc$$
 $O(n^3)$

$$\bigcirc$$
 $O(\log n)$

(E) None

Consider the following graph.



1.	Give	the	adja	cency	matrix	of the g	raph.
			• • • • •				
			• • • • •				
			• • • • •				

۷.	Give	tne	aaja	acei	ісу	nst	rep	rese	enta	tion	OI	TT.	ıe
	graph	1.											
				• • •									
										• • • •			•
3.	What						_		n th	ie si	ıbgı	ap	h
	conta	Ì				,	,	,					
				• • •	• • •					• • • •			•
4.	What	is t	he n	ıaxi	imu	m n	umb	er o	of ec	lges	in a	a d	i-

ADT Queue Specification

- enqueue (Type e): **requires**: Queue Q is not full. **input**: Type e. **results**: Element e is added to the queue at its tail. **output**: none.
- serve (Type e): **requires**: Queue Q is not empty. **input**: none. **results**: the element at the head of Q is removed and its value assigned to e. **output**: Type e.
- length (int length): **requires**: none. **input**: none. **results**: The number of elements in the Queue Q is returned. **output**: length.
- full (boolean flag): **requires**: none. **input**: none. **results**: If Q is full then flag is set to true, otherwise flag is set to false. **output**: flag.

ADT Stack Specification

node to itself, are not allowed)?

• push(Type e): requires: Stack S is not full. input: Type e. results: Element e is added to the stack as its most recently added elements. output: none.

rected graph with n nodes (loops, edges from a

- pop(Type e): requires: Stack S is not empty.
 input: results: the most recently arrived element in S is removed and its value assigned to
 e. output: Type e.
- empty(boolean flag): **requires**: none. **input**: none. **results**: If Stack S is empty then flag is true, otherwise false. **output**: flag.
- full(boolean flag): **requires**: none. **input**: none. **results**: If S is full then Full is true, otherwise Full is false. **output**: flag.