



KING SAUD UNIVERSITY  
COLLEGE OF COMPUTER & INFORMATION SCIENCES  
DEPT OF COMPUTER SCIENCE

CSC311 Algorithms Design and Analysis  
Second Semester 1444 AH  
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## Tutorial #3

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### Problem 1

For each algorithm listed below,

give a recurrence that describes its worst-case running time, and give its worst-case running time using  $\theta$ -notation.

You need **not** justify your answers.

(a) Merge sort      Recurrence:  $T(n) = 2T(n/2) + O(n)$   
Worst-case:  $\Theta(n \log n)$

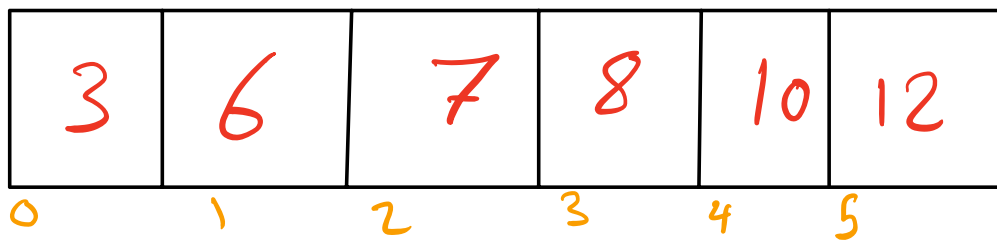
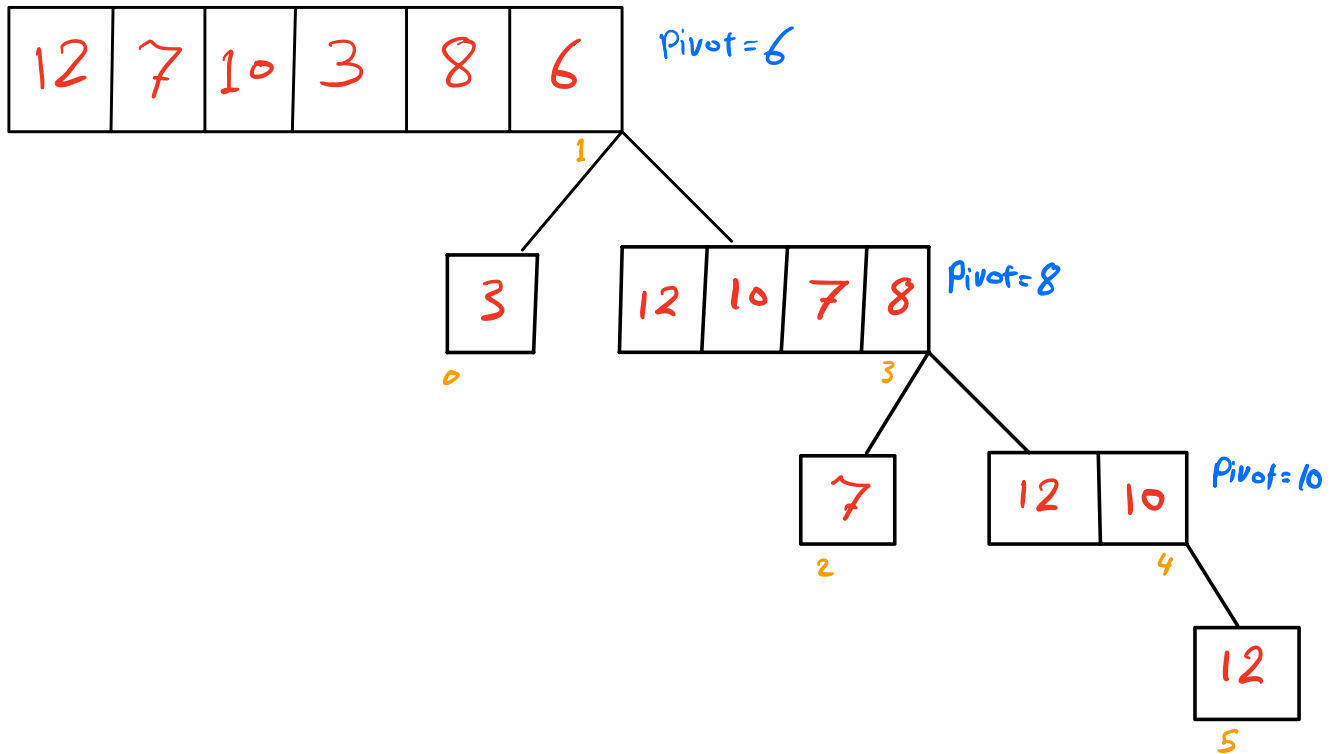
(b) Insertion sort      Recurrence:  $T(n) = T(n-1) + O(n)$   
Worst-case:  $\Theta(n^2)$

(c) Merge's algorithm      No Recursive.

(d) Binary Search      Recurrence:  $T(n) = T(n/2) + O(1)$   
Worst-case:  $\Theta(\log n)$

## Problem 2

Do Quicksort on the input numbers 12, 7, 10, 3, 8, 6. Illustrate the execution using a binary tree where each node is an array or subarray when calling the PARTITION procedure. The two children of a node are the subarrays immediately after partitioning their parent array.



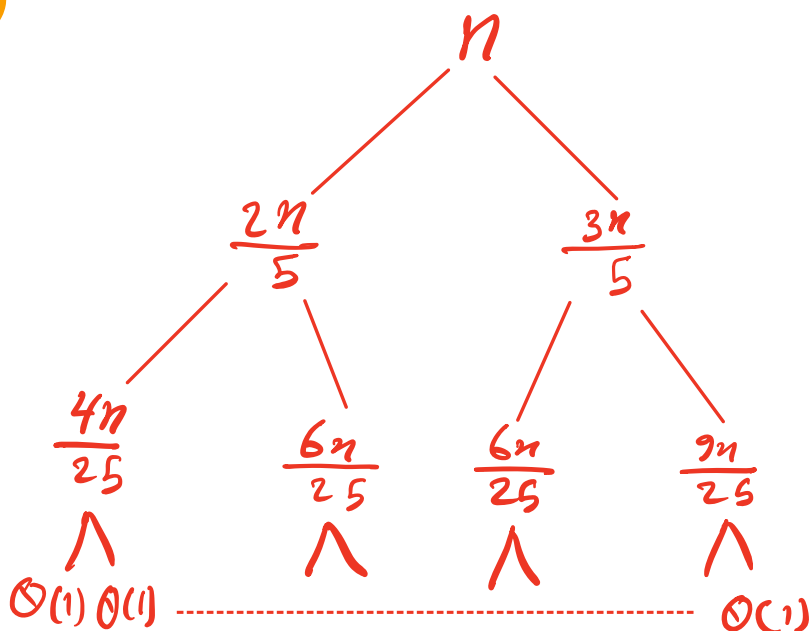
### Problem 3

Consider a variation of MergeSort which divides the list of elements into two lists of size  $2/5$  and  $3/5$ , recursively at each step, instead of dividing it into halves. The Merge procedure does not change.

- Give a recurrence relation for this algorithm
- Draw a recursion tree for the algorithm
- Using the recursion tree, explain how the worst case upper bound is  $O(n \log n)$ .

a)  $T(n) = T\left(\frac{2n}{5}\right) + T\left(\frac{3n}{5}\right) + n$

b)



c)

$h$ : height

stop P

$$n = \left(\frac{5}{3}\right)^k$$

$$k = \log_{\frac{5}{3}} n$$

$$k = h$$

$$T(n) \leq n \cdot \log_{\frac{5}{3}} n$$

$$\leq c \cdot g(n)$$

$$\leq n \cdot \log_{\frac{5}{3}} n$$

$$c = 1$$

$$n_0 = \frac{5}{3}$$

$$O(n \log_{\frac{5}{3}} n)$$