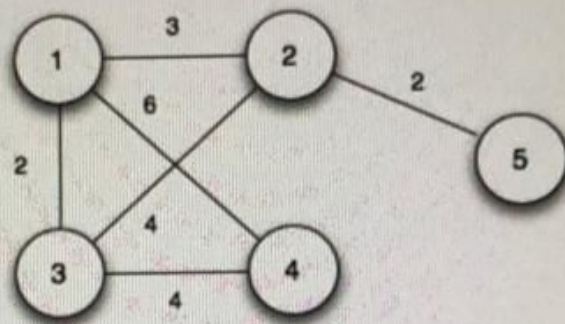


### Question 1

Given the following undirected graph:



We apply Kruskal's algorithm to construct a MST. Follow the steps of the algorithm by filling in the blanks:

a) We define  $E_T$  to maintain the set of edges composing a minimum spanning tree, which is initially empty.

b) we pick edge {1,3}, then we pick edge {2,5}, then we pick edge {1,2} and then we pick edge {3,4}.

c) The algorithm stops when the number of picked edges = 4.

d) The resultant MST will have a weight equal to 11.

⚠ Moving to another question will save this response.

1 points

### Question 2

Both dynamic programming and greedy method will solve optimization problems. In the space provided answer the following two questions:

- a) In your own words, what is an optimization problem?
- b) explain to me how does DP is different than the greedy method, when solving optimization problems.

**B** **T** **I** Arial 3 (12pt) **T** **≡** **≡** **↓** **↺** **↻**

a) The optimization problem is finding the best solution of all possible solution, such as: minimum or maximum .

b) In DP it solves the all sub-problems, then choose the best of all of the solved. Unlike greedy, first will have decision of each subproblem, then we solve the result of sub-problem.

Given the following Alphabet and frequency table, use the Huffman coding algorithm to construct a coding tree by filling in the blanks in each step.

Alphabet: A, B, C, D, E, F, G, O

Frequency table:

| A   | B  | C   | D   | E   | F   | G   | O   |
|-----|----|-----|-----|-----|-----|-----|-----|
| 12% | 8% | 19% | 25% | 43% | 15% | 13% | 31% |

Fill in the blanks in each of the following:

- After the algorithm terminates, the resultant code tree will have a root node with the value .

- The code for letter A: .

- The code for letter C: .

- The code for letter D: .

- The code for letter E: .

$G = 1110$



Let  $W = 10$  and

| $i$   | 1  | 2  | 3  | 4  |
|-------|----|----|----|----|
| $v_i$ | 10 | 40 | 30 | 50 |
| $w_i$ | 5  | 4  | 6  | 3  |

| $V[i, w]$ | 0 | 1 | 2 | 3  | 4  | 5  | 6  | 7  | 8  | 9  | 10 |
|-----------|---|---|---|----|----|----|----|----|----|----|----|
| $i = 0$   | 0 | 0 | 0 | 0  | 0  | 0  | 0  | 0  | 0  | 0  | 0  |
| 1         | 0 | 0 | 0 | 0  | 0  | 10 | 10 | 10 | 10 | 10 | 10 |
| 2         | 0 | 0 | 0 | 0  | 40 | 40 | 40 | 40 | 40 | 50 | 50 |
| 3         | 0 | 0 | 0 | 0  | 40 | 40 | 40 | 40 | 40 | 50 | 70 |
| 4         | 0 | 0 | 0 | 50 | 50 | 50 | 50 | 90 | 90 | 90 | 90 |

According to the solution table, the maximum item value that we can achieve **90**. By reconstructing the solution, we know that the following items **{2,4}** are included in the solution. Carefully, fill in the following blanks:

- The answer to the problem appears at cell  $\langle 4, 10 \rangle$  in the matrix.
- The cell in the table that lets me pick item 4 is cell at location:
- The cell in the table that lets me pick item 3 is cell at location:

Question 5

1 points

Solve the following instance of the 0/1 Knapsack Problem using Greedy method. The greedy choice we are going to use is: pick the item with the maximum value.

The maximum allowed weight is  $W = 10$ .

|        |   |   |   |   |   |   |
|--------|---|---|---|---|---|---|
| Item   | 1 | 2 | 3 | 4 | 5 | 6 |
| Value  | 6 | 4 | 5 | 3 | 9 | 7 |
| Weight | 4 | 2 | 3 | 1 | 6 | 4 |

Fill in the blanks in the following:

- the maximum value we can get is: .
- the items we pick: {  }.

Question 6

Solve the following instance of the 0/1 Knapsack Problem using Dynamic programming.

The maximum allowed weight is  $W = 10$ .

|        |   |   |   |   |   |   |
|--------|---|---|---|---|---|---|
| Item   | 1 | 2 | 3 | 4 | 5 | 6 |
| Price  | 6 | 4 | 5 | 3 | 9 | 7 |
| Weight | 4 | 2 | 3 | 1 | 6 | 4 |

Fill in the blanks in the following table:

| 0 |   |   | 1 |  |  | 2 |  |  | 3 |  |  | 4 |  |  | 5 |  |  | 6  |  |  | 0  |
|---|---|---|---|--|--|---|--|--|---|--|--|---|--|--|---|--|--|----|--|--|----|
| 0 | 0 | 0 |   |  |  | 0 |  |  | 0 |  |  | 0 |  |  | 0 |  |  | 0  |  |  | 0  |
| 0 | 0 | 0 |   |  |  |   |  |  | 0 |  |  | 6 |  |  | 6 |  |  | 6  |  |  | 6  |
| 1 | 0 | 0 |   |  |  | 0 |  |  |   |  |  |   |  |  |   |  |  |    |  |  |    |
| 2 | 0 | 0 |   |  |  | 4 |  |  | 4 |  |  | 6 |  |  | 6 |  |  | 10 |  |  | 10 |
| 3 | 0 | 0 |   |  |  | 4 |  |  | 5 |  |  | 6 |  |  | 9 |  |  | 10 |  |  | 11 |
| 4 | 0 | 3 |   |  |  | 4 |  |  | 7 |  |  | 8 |  |  | 9 |  |  | 12 |  |  | 13 |
| 5 | 0 | 3 |   |  |  | 4 |  |  | 7 |  |  | 8 |  |  | 9 |  |  | 12 |  |  | 13 |

| 4 | 5  | 6  | 7  | 8  | 9  | 10 |
|---|----|----|----|----|----|----|
| 0 | 0  | 0  | 0  | 0  | 0  | 0  |
| 6 | 6  | 6  | 6  | 6  | 6  | 6  |
| 6 | 10 | 10 | 10 | 10 | 10 | 10 |
| 9 | 10 | 11 | 11 | 15 | 15 | 15 |
| 9 | 12 | 13 | 14 | 14 | 18 | 18 |
| 9 | 12 | 13 | 14 | 14 | 18 | 18 |
| 9 | 12 | 13 | 14 | 14 | 18 | 18 |