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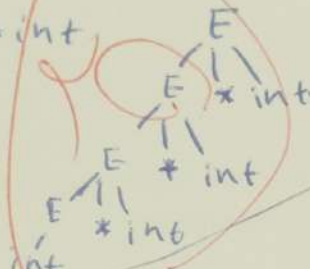
Q1) Consider the following CFG (12 marks)

$S \rightarrow \text{int} * E \mid \text{int} + E$   
 $E \rightarrow \text{int}$

A) Write two left most derivation for the string  $\text{int} * \text{int} + \text{int} * \text{int}$  (2 marks)

First Derivation:

$E \rightarrow E * \text{int} \mid E + \text{int}$   
 $E \rightarrow \text{int}$

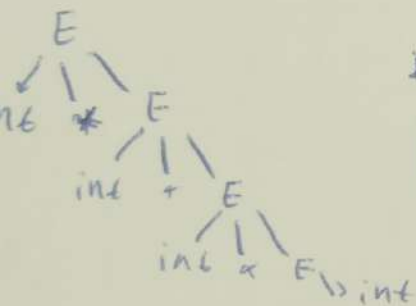


Second Derivation:

$E \rightarrow E * \text{int}$   
 $E \rightarrow E + \text{int}$



B) Since the string  $\text{int} * \text{int} + \text{int} * \text{int}$  has more than one leftmost derivations, can we conclude that the grammar is ambiguous. (2 marks)



~~the grammar~~  
the string can do only in one parse tree  
so is unambiguous

C) Which operation has the highest precedence? Why? (2 marks)

is  $*$ , because, highest ~~priority~~ priority then  $+$   
let's

D) Is the associativity of the operation  $*$  right or left associativity? (2 marks)

right associativity

- E) Rewrite the above grammar so that you reverse the \* operation has higher precedence than the + operation and both operations have right associativity. (4 marks)

$$\begin{aligned} E &\rightarrow \text{int } E' \\ E' &\rightarrow *E' \mid \epsilon \\ X &\rightarrow +E' \mid \epsilon \end{aligned}$$

Q2. Consider the following CFG (please notice that it is slightly different than the grammar in the previous question) (12 marks)

$$S \rightarrow S * E \mid E + S$$

$$E \rightarrow \text{int} \mid \epsilon$$

- a) Find the first and follow set for each terminal and non-terminal symbol. (4 marks)

Symbol	First Set	Follow Set
S	{int, $\epsilon$ }	{*, \$}
E	{int, $\epsilon$ }	{+, *, \$}
int	int	{+, *, \$}
*	*	{int, *, \$}
+	+	{int, *, \$}

- b) Draw the LL(1) parsing Table (3 marks)

	int	*	+	\$
S	$E + S$	$\epsilon$	$\epsilon$	$\epsilon$
E	int	$\epsilon$	$\epsilon$	$\epsilon$

- c) Is the grammar LL(1) grammar? Why? or why not? (2 mark)

Yes, this is LL(1) grammar since you not have

conflict in the parsing table

- d) If your answer to the above question was no, then rewrite the grammar so that the grammar is LL(1). (3 marks)

*(Handwritten scribble)*

*eliminate left recursion*

Q3) Consider the CFG

(16 marks)

$S \rightarrow S^* E \mid E + S$

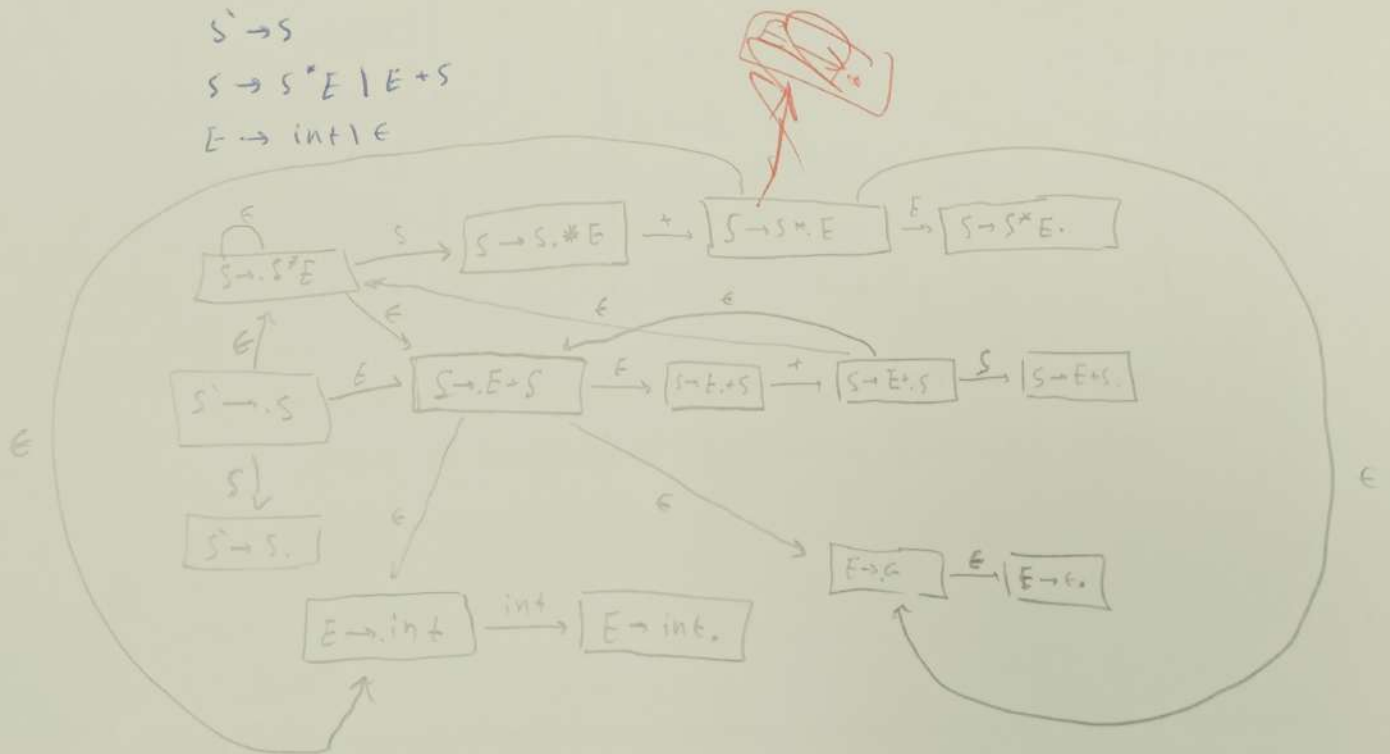
$E \rightarrow \text{int} \mid \epsilon$

- a) Design an NFA that recognizes the viable prefixes for the above grammar. (5 marks)

$S' \rightarrow S$

$S \rightarrow S^* E \mid E + S$

$E \rightarrow \text{int} \mid \epsilon$



...an def  $f(x,y) = e$  the activation follows:

- b) Use the following LR(0) DFA to parse the string  $(id, id + id) \$$  using the SLR parsing algorithm. Please use a suitable table to illustrate your answer. (5 marks)



$(id, id + id) \$$	1	shift
$(L, id, id + id) \$$	3	shift
$(id, ., id + id) \$$	2	reduce $S \rightarrow id$
$(S, ., id + id) \$$	7	reduce $L \rightarrow S$
$(L, ., id + id) \$$	5	shift
$(L, ., id + id) \$$	8	shift
$(L, id, . + id) \$$	2	reduce $S \rightarrow id$
$(L, S, . + id) \$$	9	reduce $L \rightarrow L, S$

$(L, + id) \$$	9	shift
$(L, +, id) \$$	8	<del>shift</del> Reject
$(L, + id, .) \$$	9	<del>shift</del>
$(L, + id, .) \$$	6	reduce
$(L, + id, .) \$$	6	<del>shift</del>
$S, \$$	4	$S \rightarrow (L, id)$
$S, \$$	4	Shift
$S, \$$	4	final state

- c) Is the CFG represented using the above DFA an SLR grammar? Why? (2 marks)

easy to implement and ~~to test the input~~

- d) Give two examples of problems that cannot be detected by both the lexical analyzer and the parser. (4 marks)

1.

2.