

Different Turing machine classes have different power; which means each class accept different set of languages.
☐ True ☒ False

QUESTION 2

Since b^* is regular language, then it is Turing Machine acceptable one.
☒ True ☐ False

QUESTION 3

Multi-tape Turing machine does not have the same power as a single-tape Turing machine.
☐ True ☒ False

QUESTION 4

Turing Machine with multiple track tape have different read/write head for each track.
☐ True ☒ False

QUESTION 5

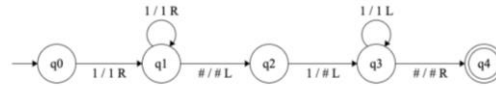
For every context-free language there exists a Turing Machine accepts that language.
☒ True ☐ False

QUESTION 6

$L = \{uv^R \mid u \in \Sigma^*\}$ is Turing Machine acceptable language.
☒ True ☐ False

QUESTION 7

Consider the following Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_4, \delta)$ where $Q = \{q_0, q_1, q_2, q_3, q_4\}$, $\Sigma = \{1\}$, and $\Gamma = \{1, \#\}$.

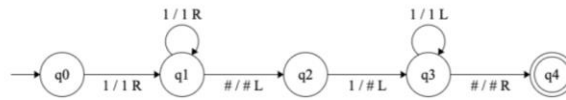


Let $w = 11111$, the initial configuration is:

- ☐ a. $q_4 11111$
- ☐ b. $11111 q_0$
- ☒ c. $q_0 11111$
- ☐ d. $q_1 11111$

QUESTION 9

Recall **Question 7** Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_4, \delta)$ where $Q = \{q_0, q_1, q_2, q_3, q_4\}$, $\Sigma = \{1\}$, and $\Gamma = \{1, \#\}$.



M computes $f(x)$ which:

- ☐ a. Removes the first 1 of x where $x \in L = \{1w \mid w \in \{1\}^*\}$
- ☒ b. Removes the last 1 of x where $x \in L = \{w1 \mid w \in \{1\}^*\}$
- ☐ c. Inserts 1 to the last of x where $x \in L = \{w \mid w \in \{1\}^*\}$
- ☐ d. Inserts 1 to the first of x where $x \in L = \{w \mid w \in \{1\}^*\}$

QUESTION 8

Recall **Question 7** Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_4, \delta)$ where $Q = \{q_0, q_1, q_2, q_3, q_4\}$, $\Sigma = \{1\}$, and $\Gamma = \{1, \#\}$.

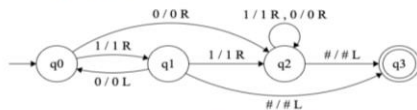


Let $w = 11111$, the final configuration is:

- ☐ a. $q_4 111111$
- ☒ b. $q_4 1111$
- ☐ c. $11111 q_4$
- ☐ d. $q_0 1111$

QUESTION 10

Consider the following Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_3, \delta)$ where $Q = \{q_0, q_1, q_2, q_3\}$, $\Sigma = \{0, 1\}$, and $\Gamma = \{0, 1, \#\}$.

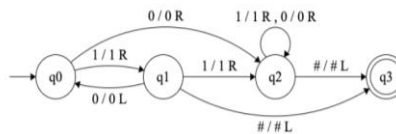


Suppose the current configuration is $1 q_1 0100$. The next configuration will be:

- ☒ a. $q_0 10100$
- ☐ b. None
- ☐ c. $10 q_2 100$
- ☐ d. $10 q_0 100$

QUESTION 11

Recall **Question 10** Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_3, \delta)$ where $Q = \{q_0, q_1, q_2, q_3\}$, $\Sigma = \{0, 1\}$, and $\Gamma = \{0, 1, \#\}$.

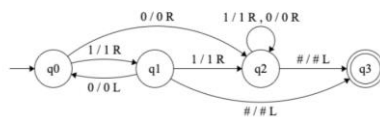


Is the string 10100 accepted by M ?

- ☐ Yes
- ☒ No

QUESTION 12

Recall **Question 10** Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_3, \delta)$ where $Q = \{q_0, q_1, q_2, q_3\}$, $\Sigma = \{0, 1\}$, and $\Gamma = \{0, 1, \#\}$.



Is the string 11100 accepted by M ?

- ☒ Yes
- ☐ No

QUESTION 13

Recall **Question 10** Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_3, \delta)$ where $Q = \{q_0, q_1, q_2, q_3\}$, $\Sigma = \{0, 1\}$, and $\Gamma = \{0, 1, \#\}$.

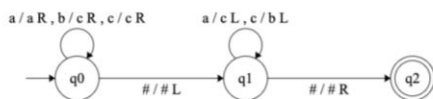


The language accepted by M is:

- ☐ a. None
- ☐ b. $L = \{w \in \{0, 1\}^* \mid w \text{ does not end with } 10\}$
- ☒ c. $L = \{w \in \{0, 1\}^* \mid w \text{ does not start with } 10\}$
- ☐ d. $L = \{w \in \{0, 1\}^* \mid w \text{ does not start with } 1\}$

QUESTION 14

Consider the following Turing machine $M: M = (Q, \Sigma, \Gamma, \#, q_0, q_2, \delta)$ where $Q = \{q_0, q_1, q_2\}$, $\Sigma = \{a, b, c\}$, and $\Gamma = \{a, b, c, \#\}$.



Let $q_0 a a a c$ be the initial configuration, the final configuration is:

- ☐ a. $q_2 a a a b$
- ☐ b. $q_2 c c c c$
- ☒ c. $q_2 c c c b$
- ☐ d. None