

Identify regular and non-regular languages

Question 1.1 [2.5 pts]

Answer with TRUE or FALSE next to each of the following statements:

Statement	True	False
1. There are some languages that can be designed by finite automata but not by regular expression.		<input checked="" type="checkbox"/>
2. If L is a regular language and F is a finite language, then $L \cup F$ is a regular language.	<input checked="" type="checkbox"/>	
3. Define $EVEN(w)$, for a finite string w , to be the string consisting of the symbols of w in <u>even-numbered positions</u> . For example, $EVEN(1011010) = 011$. If L is a regular language, then $\{EVEN(w); w \in L\}$ must be regular.		<input checked="" type="checkbox"/>
4. If L is a regular language, then $\{ww^R; w \in L\}$ must be a regular language.		<input checked="" type="checkbox"/>
5. If $L_1 \cup L_2$ is a non-regular language, then both L_1, L_2 must be non-regular languages.	<input checked="" type="checkbox"/>	

Question 1.2 [2.5 pts]

Select the correct answer:

1. Which of the following languages are not regular:

A. $L = \{(01)^n 0^k \mid n > k, k \geq 0\}$ ☒

B. $L = \{c^n b^k a^{n+k} \mid n \geq 0, k \geq 0\}$ ☒

C. $L = \{0^n 1^k \mid n \neq k\}$ ☒

a) A and B only

c) A and C only

b) B and C only

d) A, B and C

2. Which of the following languages is regular:

a) $L = \{a^i b^i \mid i \geq 0\}$ ☒

b) $L = \{a^i b^i \mid i \geq 1\}$

c) $L = \{a^i b^i \mid 0 < i < 5\}$ ☒

d) None of the above. ☒

3. The reverse of $(0+1)^*$ will be:

a) \emptyset ☒

b) ϵ

c) $(0+1)^*$

d) $(0+1)$

4. $L = \{w c x : w, x \in \{a, b\}^* \text{ and the number of } a\text{'s in } w \text{ is equal to the number of } b\text{'s in } x\}$ is a non-regular language. For example, $w = abababcb \in L$. We use pumping lemma to prove that the language L is non-regular. Fill in blanks to complete the proof.

A. Pick a string $w \in L$ and length $|w| \geq m$, such that m is the critical length.

$w = a^m c b^m = xyz, |xy| \leq m, |y| \geq 1$.

B. $y = a^k, 1 \leq k \leq m$.

C. From pumping lemma, $xy^i z \in L, i = 0, 1, 2, \dots$, we choose $i = 2$. Thus the string after pumping is $a^{m+k} c b^m \notin L$, which is a contradiction.

Part 2: Produce computing-based solutions using regular expressions, and context free grammar

Question 2: [7 pts]

Question 2.1 Construct a context-free grammar that generates $L = \{w \in \{a, c\}^* \mid w \text{ contains at least 3 } c\text{'s}\}$

$S \rightarrow AcAcAcA$
 $A \rightarrow aA \mid cA \mid \lambda$

Question 2.2 Construct a context-free grammar that generates $L = \{0^n 1^n 0^m 1^m \mid n, m \geq 0\}$ and $\Sigma = \{0, 1\}$

$S \rightarrow 0S1A \mid \lambda$
 $A \rightarrow 0A1 \mid \lambda$

Question 2.3 Which of the following context-free grammar productions generates words of balanced brackets. An example for the generated string is $(())()$.

A. $P = \{S \rightarrow \lambda, S \rightarrow TS, T \rightarrow (T)\}$

B. $P = \{S \rightarrow \lambda, S \rightarrow SS, S \rightarrow ()\}$

C. $P = \{S \rightarrow \lambda, S \rightarrow (T)S, T \rightarrow ()\}$

$(T)S \rightarrow (())S \rightarrow ((()))$

D. $P = \{S \rightarrow \lambda, S \rightarrow (S)S\}$

Question 2.4 Construct regular expressions representing the following languages. [4 pts]

(1) The language over the alphabet $\{0, 1\}$ that doesn't contain the substring 110

$(0 + 10)^* 1^*$

(2) The language $\{w \in \Sigma^* \mid w \text{ ends with a double letter}\}$ over the alphabet $\{a, b\}$. A double letter over this alphabet is aa or bb.

$(a+b)^* (aa+bb)$

(3) $L = \{x \in \{a, b\}^* \mid x \text{ ends with } a \text{ and does not contain the substring } bb\}$

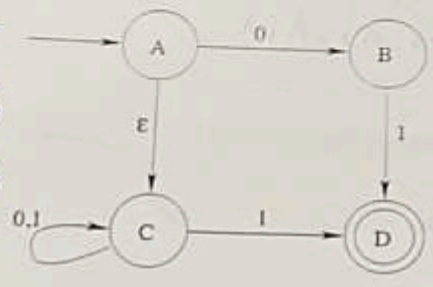
$(b + \lambda)(a + ab)^* a$

- (4) Find a regular expression corresponding to the language of all strings over the alphabet $\{a, b\}$ that do not end with ab .
- A. $(a+b)^*(a+bb)^*$ ✓
 B. (ab^*+bb)
 C. $(a+ab)^*(a+bb)$
 D. None of the above

Part 3: Design different machine models (DFA, NFA, PDA, TM)

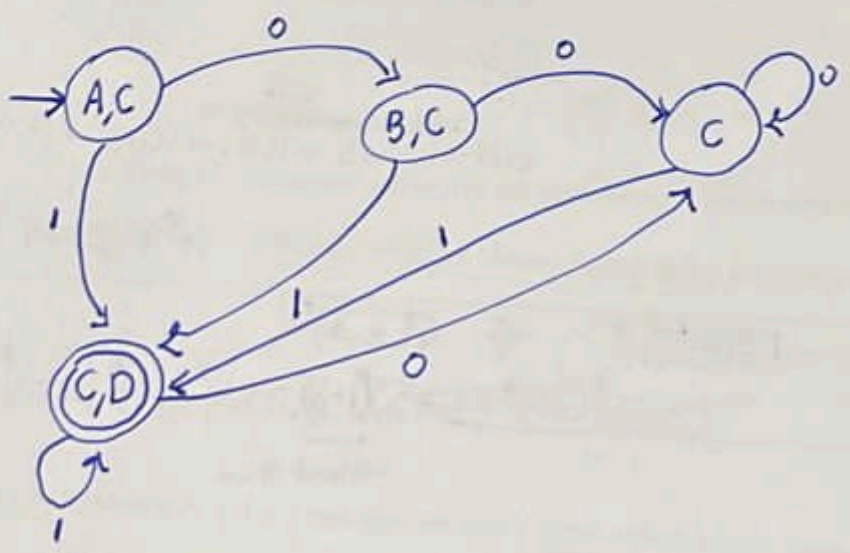
Question 3.1: [5 pts]
 Consider the following NFA.

state	0	1
$\{A, C\}$	$\{B, C\}$	$\{D, C\}$
$\{B, C\}$	$\{C\}$	$\{D, C\}$
$\{D, C\}$	$\{C\}$	$\{D, C\}$
$\{C\}$	$\{C\}$	$\{D, C\}$



$\delta(A, 1) = \{A, C\}$
 $\delta(B, 1) = \{D\}$
 $\delta(C, 1) = \{C, D\}$
 $\delta(D, 1) = \{D\}$

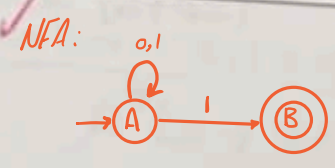
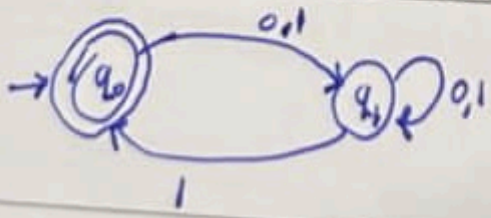
1. Convert this NFA into an equivalent DFA. Your answer should be the state diagram of a DFA. [3 pts]



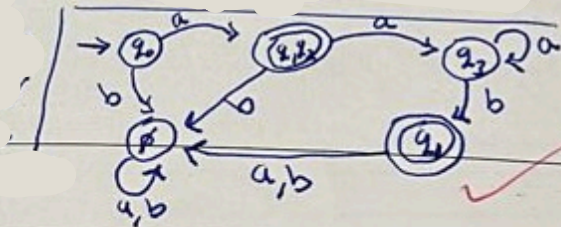
2. What is the language recognize by this DFA. Express your answer as a regular expression. [1 pt]

$(01 + (0+1)^*1)$

3. Construct an NFA with two states that recognizes the same language. [1 pt]



Question 3.2: Construct a DFA that recognizes the regular expression $(a + aa^*b)$. [3 pts]



Part 4: Evaluate the language accepted by a machine, a regular expression, and a context free grammar

Question 4.1: [2 pts]

Describe the language generated by the following Grammar productions.

Grammar	Corresponding $L(G)$
(1) $S \rightarrow ABC$ $A \rightarrow 0A1 \mid \epsilon$ $B \rightarrow 1B \mid 1$ $C \rightarrow 1C0 \mid \epsilon$	$L = \{0^n 1^n 1^m 1^n 0^n \mid n, m \geq 0 \text{ and } m \geq 1\}$
(2) $S \rightarrow 0Y1 \mid 1Y0$ $Y \rightarrow 0Y \mid 1Y \mid \epsilon$	$L = \{\text{start with 0 (every string start with 0 and ends with 1) or (every string start with 1 and ends with 0)}\}$

Question 4.3 : [3 pts]

1. Is the grammar below ambiguous?

$S \rightarrow aS \mid aSbS \mid c$

A. Yes ☒

B. No ☒

acub

2. Is the grammar below ambiguous?

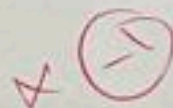
$S \rightarrow T \mid U$

$T \rightarrow aTbT \mid c$

$U \rightarrow aS \mid aTbU$

A. Yes ✓

B. No



3. Are CFGs provided in (1) and (2) equivalent?

A. Yes ✓

B. No

End of Exam

