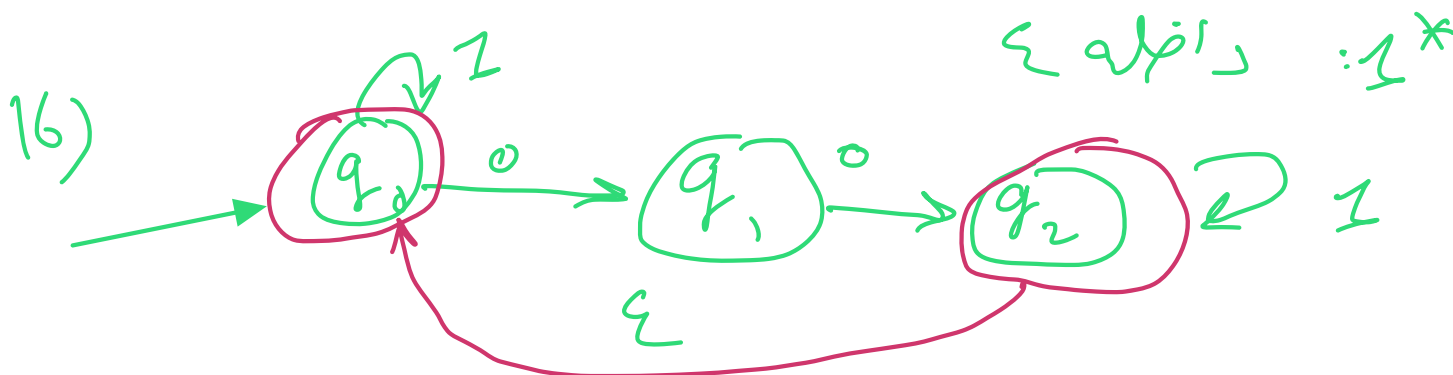
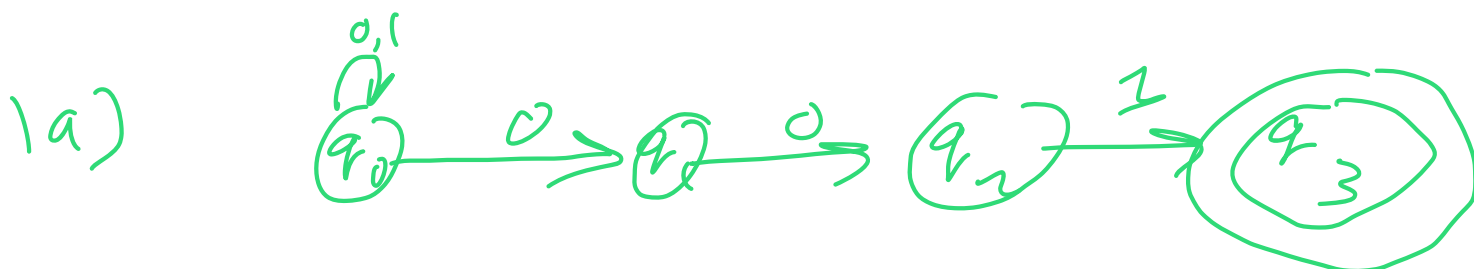


Theory of Computation (CSC 339) – Fall 2023

Instructor: Prof. M.B. Menai

Tutorial 3: Nondeterministic Finite Automata

- Give state diagrams of NFAs with the specified number of states recognizing each of the following languages. In all parts, the alphabet is $\Sigma = \{0, 1\}$.
 - The language $\{w | w \text{ ends with } 001\}$ with four states.
 - The language $1^*(001^*)^*$ with three states.
 - The language ϵ with one state.
 - The language 0^* with one state.
 - The language $\{w | w \text{ is any string except } 11 \text{ and } 111\}$.
- Convert the NFAs obtained in the previous question to DFAs.
- (3 points) Give the state diagrams of NFAs recognizing the concatenation and the star of the following languages. The alphabet is $\Sigma = \{0, 1\}$.
 - The language $\{w | w \text{ ends with } 001\}$; The language $1^*(001^*)^*$.
 - The language ϵ with one state; The language 0^* .
- Convert the following regular expressions to NFAs. The alphabet is $\Sigma = \{0, 1\}$.
 - $(0 + 1)^*000(0 + 1)^*$.
 - $((00)^*(11) + 01)^*$.
 - \emptyset^*



$$(001^*)^* = \{\epsilon, 00, 001, 001001, \dots\}$$

$$L = \{ \underbrace{1^*}_{\epsilon} \underbrace{(001^*)^*}_{\epsilon} \}$$

مثال خطية بسيطة جداً :

$(1)^*$: ممكن يكون ϵ

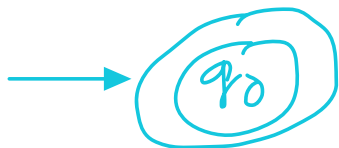
$(001^*)^*$: ممكن اي عدد اقل القوس يكون ϵ

فيس ~~String~~ : $\epsilon \epsilon$ يعني لازم q_0

~~accept state~~

1c)

$$L = \{\epsilon\}$$



it is an NFA

1d)

$$L = \{0^*\}$$

$\rightarrow \textcircled{g_0} \mathbb{Z}_0$