

King Saud University, department of Computer Science			Name:		
Theory of Computation - CSC 339		S	Student ID:		
Midterm exam – 25/2/2021 Section ID (or time):			Section ID (or time):		
Question 1					.6 points
For each of the following statements, circle either True or False. (No explanation needed.)					
a) If $S1 = \{a,b,c\}$ and $S2 = \{S1\}$ . Then, $S1$ and $S2$ are equivalent.				True	<u>False</u>
b) If $L_1$ is regular and $L_2$ is regular. Then, $L_1 \cup L_2$ is regular.				<u>True</u>	False
c) If $L_1$ is regular and $L_2$ is regular. Then, $L_1 \circ L_2$ is not regular.				True	<b>False</b>
d) The class of regular languages is closed under the (kleene) star operation.				<u>True</u>	False
e) Every NFA can be converted into an equivalent DFA.				— <u>True</u>	False
f) There exists some regular expression that cannot be converted into DFA.				True	<b>False</b>
g) Any regular expression can be converted into an equivalent NFA, and vice versa.				<u>True</u>	False
h) Any context-free grammar (CFG) can be converted into an equivalent NFA.				True	<b>False</b>
i) Let $L = \{a^ib^k \mid i \neq k\}$ . L is a regular language.				True	<b>False</b>
j) If string $w \in B$ passes the three conditions of the pumping lemma, then B must be regular.				True	<b>False</b>
k) All regular languages are context-free.				<u>True</u>	False
1) A CFG is said to be ambiguous if there are different parse trees for different strings.				True	<b>False</b>
Ouestion 2					5 points
& mesorom =					.e penies
For each of the following noted.	ng questions, <u>select exact</u>	t <u>ly one option</u> . As.	sume the alphabet $\Sigma$ is $\{0,1\}$	unless oti	herwise
1) Which of the following	ing languages is regular?				
(a) $\{0^i1^i \mid i \ge 0\}$	(b) $\{0^i1^i \mid i=2\}$	(c) $\{0^i1^i0^i \mid i \geq 0\}$	(d) None of the menti	oned.	
2) Which of the following ends with 0}	ing regular expressions re	epresent the langu	$age L = \{w \mid w \text{ of an even length}\}$	ngth and	starts and
(a) $0 \Sigma^* \Sigma^* 0^*$	(b) 0 Σ Σ*0	$\underline{(c)\ 0(\Sigma\Sigma)*0}$	(d) All of the mention	ied.	

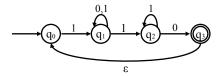
- 3) Consider a language L represented by the regular expression  $\Sigma*011*$ . Which of the following strings belongs to this language?
- (a) 0001
- (b) 000111
- (c) 1101
- (d) All mentioned strings  $\in L$
- 4) When using the pumping lemma to prove non-regularity, we select a string w, and fragment it into three parts such that w = xyz. Which of these parts cannot be an empty string?
- (a) x

**(b) y** 

(c) z

- (d) All of the mentioned
- 5) What regular expression is equivalent to the NFA on the right?
- (a)  $1*(0 \cup 1)*110$
- (b)  $1(0 \cup 1)*11*0$
- (c)  $(1(0 \cup 1)*11*0)+$

(d)  $(1(0 \cup 1)*11*0)*$ 



Answer each of the following questions about NFA  $N_1$ .

a) (1.5 point) Assume  $\Sigma = \{a,b\}$ . What are Q,  $q_0$ , and F for  $N_1$ ?

$$Q = \{q_0, q_1, q_2, q_3\}$$

$$q_0 = q_0$$

$$F = \{q_2, q_3\}$$

 $N_1$   $q_0$  a  $q_1$  b  $q_2$  b  $g_3$  b

b) (2 point) Provide 2 example strings (of any length) that  $N_I$  will accept, and provide 2 other example strings that  $N_I$  will not accept.

Not accepted

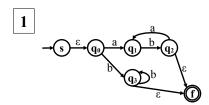
$$wI = ab$$

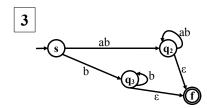
$$w3 = aa$$

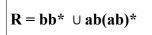
$$w2 = b$$

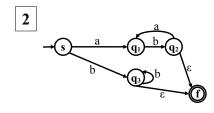
$$w4 = ba$$

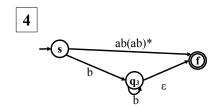
c) (3.5 points) Convert  $N_1$  into an equivalent regular expression. Show your work step by step, and write the final expression inside the designated box. You may use the back of this page for extra space.

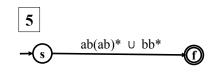












Answer the following questions about language  $L_1 = \{0^n 1^{2n} \mid n \ge 0\}$ .

a) (1 point) Informally describe the language  $L_1$  (i.e., what type of strings does it consist of?).

 $L_1$  consists of all strings that have a number of 0's followed by twice the number of 1's. The empty string is also a member of  $L_1$ .

b) (2 points) List down the three conditions of the pumping lemma. And use the pumping lemma to prove that  $L_1$  is not regular, and clearly state the length of the string you choose.

The three conditions of the pumping lemma are:

- (1)  $xy^iz \in L_I$  for each  $i \ge 0$
- (2) |v| > 0, and
- (3)  $|xy| \le p$ .

## Proof:

We assume  $L_I$  is regular. Let  $s = 0^p 1^{2p}$ , where p is the pumping length. Then, |s| = 3p. Condition (3) of the pumping lemma tells us that  $|xy| \le p$ . By condition (2), we know that y can consist of only 0's. Let |y| = p. Now, if we pump the string such that |y| = p+1, then the new pumped string  $sI = 0^{p+1} 1^{2p}$ . Condition (1) tells us the new (pumped) string should  $\in L_I$ , but sI is not in the right form, and therefore  $sI \notin L_I$ . We have a contradiction; thus,  $L_I$  is not regular.

a) (3 points) Design a CFG  $G_1$  that generates the same language that NFA  $N_1$  (in question 3) recognizes.

 $A \rightarrow aB \mid bD$ 

 $B \rightarrow bC$ 

 $C \rightarrow aB \mid \epsilon$ 

 $D \rightarrow bD \mid \epsilon$ 

b) (*I point*) Provide the formal definition for  $G_I$  (without the rules set). (hint: it should consist of: V,  $\Sigma$ , and S)

$$V = \{A, B, C, D\}$$
  

$$\Sigma = \{a,b, \epsilon\}$$
  

$$S = A$$

c) (2 <u>bonus</u> points) Choose 1 string that is accepted by  $N_1$  (from your answer to question 3-b), and use grammar  $G_1$  to draw the parse tree for that string.

