PARALLEL PROCESSING

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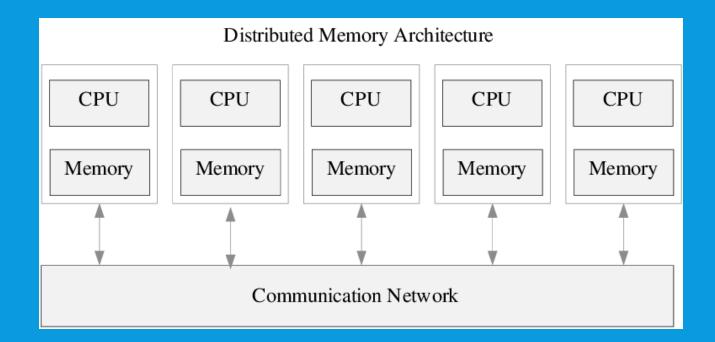


- Shared memory machines have faster communication time than distributed memory machine but it is less scalable in terms of number of processors.
- Distributed memory machines are more scalable in number of processors.

 But with more communication time.
- Hence the speed of the communication network need to be fast enough to close this gap.

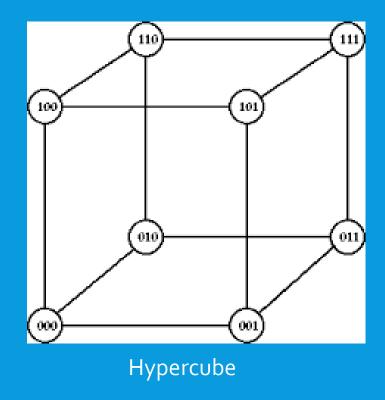


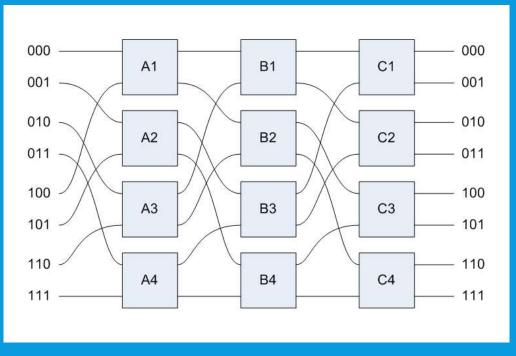
- Distributed memory machines can not communicate through shared memory.
- Exchanging data and communications are done using communication network.





Interconnection networks (examples)







Supercomputers with interconnection network





- In distributed memory machines cooperation among processes implies the data exchange using message passing (communication overhead).
- The overall execution time is consequently a sum of <u>computation</u> and <u>communication time</u>.
- Algorithms with only local communication between neighboring processors are faster and more scalable than the algorithms with the global communication among all processors.



- MPI is Message Passing Interface.
- MPI is introduced from the practical point of view, with a set of basic operations that enable implementation of parallel programs.
- The MPI library interface is a specification, not an implementation.
- The MPI is not a language, and all MPI operations are expressed as functions, subroutines, or methods, according to the appropriate language bindings for C and Fortran.



- MPI standard includes:
 - Process creation and management.
 - Language bindings for C and Fortran.
 - Point-to-point and collective communications.
 - Group and communicator concepts.



- MPI program consists of autonomous processes that are able to execute their own code in the sense of Multiple Instruction Multiple Data (MIMD) paradigm.
- The processes communicate via calls to MPI communication operations independently of operating system.
- Any MPI program should have operations to initialize execution environment and to control starting and terminating procedures of all generated processes.



- MPI processes can be collected into **groups** of specific *size* that can communicate in its own environment where each message sent in a context must be received only in the same context.
- A process group and context together form an MPI communicator.
- A process is identified by its <u>rank</u> in the group associated with a communicator.



- MPI has default communicator MPI_COMM_WORLD
- The group associated with this communicator includes all initial processes with default context.
- MPI operations return a status of the execution success; in C routines it is returned as the value of the function.
- The MPI standard *does not* specify how **st dout** from different nodes should be collected for printing at the originating process.



- MPI_SEND (buf, count, datatype, dest, tag, comm)
- The basic MPI communication is characterized by two fundamental MPI operations MPI_SEND and MPI_RECV.
- MPI_SEND is blocking call it will not complete until there is matching MPI_RECV on the other receiving process.
- The MPI_RECV will empty the input send buffer buf of matching MPI_SEND.
- MPI_SEND will return when the message data has been delivered to the communication system and the send buffer buf of the sender process source can be reused.



- MPI_SEND (buf, count, datatype, dest, tag, comm)
 - **buf** is a pointer to the send buffer
 - count is the number of data items
 - datatype is the type of data items (more about data types in MPI)
 - dest the rank of receiver process (similar to thread number)
 - *tag* a message tag to distinguish between different messages for the same receiver process
 - comm the communicator.



- MPI_RECV (buf, count, datatype, source, tag, comm, status)
 - **buf** is a pointer to the output buffer
 - count is the number of data items to be received
 - datatype is the type of data items
 - *source* the rank of sending process
 - *tag* a message tag
 - comm the communicator.
 - **status** contains further information in case of error



| MPI datatype | C equivalent |
|------------------------|------------------------|
| MPI_SHORT | short int |
| MPI_INT | int |
| MPI_LONG | long int |
| MPI_LONG_LONG | long long int |
| MPI_UNSIGNED_CHAR | unsigned char |
| MPI_UNSIGNED_SHORT | unsigned short int |
| MPI_UNSIGNED | unsigned int |
| MPI_UNSIGNED_LONG | unsigned long int |
| MPI_UNSIGNED_LONG_LONG | unsigned long long int |
| MPI_FLOAT | float |
| MPI_DOUBLE | double |
| MPI_LONG_DOUBLE | long double |
| MPI_BYTE | char |



Wildcard may be used at the destination side

MPI_RECV (buf, count, datatype, <u>source</u>, <u>tag</u>, comm, status)

- Source MPI_ANY_SOURCE
- Tag MPI_ANY_TAG



- MPI_Init initializes the MPI execution environment.
- MPI_Comm_size(MPI_COMM_WORLD, &size) returns size, which is the number of started processes.
- MPI_Comm_rank(MPI_COMM_WORLD, &rank) that returns rank or the ID of each process.
- MPI_Finalize exits the MPI



```
#include <stdio.h>
#include "mpi.h"
int main( argc, argv)
int argc;
char **argv;
 int rank, size;
 MPI_Init( & argc, & argv );
 MPI_Comm_size( MPI_COMM_WORLD, &size );
 MPI_Comm_rank( MPI_COMM_WORLD, &rank );
  printf( "Hello world from process %d of %d\n", rank, size );
 MPI_Finalize();
 return o;
```



% mpicc -o helloworld helloworld.c

% mpirun -np 4 helloworld

Hello world from process 0 of 4

Hello world from process 3 of 4

Hello world from process 1 of 4

Hello world from process 2 of 4

%



```
int world_rank;
MPI_Comm_rank(MPI_COMM_WORLD, &world_rank);
int world_size;
MPI_Comm_size(MPI_COMM_WORLD, &world_size);
int number;
if (world_rank == o) {
 number = -1;
 MPI_Send(&number, 1, MPI_INT, 1, 0, MPI_COMM_WORLD);
 } else if (world_rank == 1) {
        MPI_Recv(&number, 1, MPI_INT, 0, 0, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
        printf("Process 1 received number %d from process o\n", number);
```



MPI_Send(&number, 1, MPI_INT, 1, 0, MPI_COMM_WORLD);

- **Buf** &number
- count 1
- datatype MPI_INT
- dest 1
- tag o
- comm MPI_COMM_WORLD.



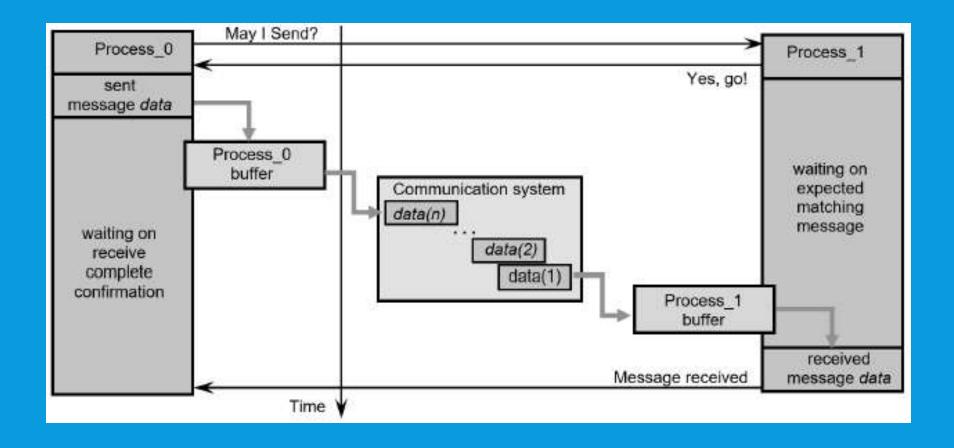
mpirun -n 2 ./send_recv

Process 1 received number 1

Process 1 received number -1 from process 0



How it works





MPI_RECV (buf, count, datatype, source, tag, comm, status)

- The number of received data items of <u>datatype</u> must be equal of fewer as specified by <u>count</u>. (count >= o).
- Receiving more data items results in an error and the <u>status</u> value return more information.



 MPI_SENDRECV (sendbuf, sendcount, sendtype, dest, sendtag, recvbuf, recvcount, recvtype, source, recvtag, comm, status)





 MPI_SENDRECV (sendbuf, sendcount, sendtype, dest, sendtag, recvbuf, recvcount, recvtype, source, recvtag, comm, status)

• MPI_SENDRECV combines a sending of message to destination process **dest** and a receiving of another message from process **source**, in a single call in sender and receiver process.



Measuring time

```
double start, finish;
start = MPI_Wtime ();
... //MPI program segment to be measured
finish = MPI_Wtime ();
printf ("Elapsed time is %f\n", finish - start);
```



- The elapsed time (wall-clock) between two points in an MPI program can be measured by using operation MPI_WTIME ().
- The time measured in seconds.



MPI_BARRIER (comm)

- This operation is used to synchronize the execution of a group of processes specified within the communicator comm.
- When a process reach this operation it waits until all other processes reach the MPI_BARRIER.
- The barrier is a simple way of separating two phases of a computation to ensure that messages generated in different phases do not interfere.



MPI_BARRIER (comm)

