

Data Structures in C

Outline

- Linked Lists
- Binary Trees
- **Stacks**
- Queues
- Hash Tables

Linked Lists

- Linked List: A dynamic data structure that consists of a sequence of nodes
 - o each element contains a link or more to the <u>next</u> node(s) in the sequence
 - Linked lists can be singly or doubly linked, linear or circular.
- Every node has a payload and a link to the next node in the list
- ❖ The start (head) of the list is maintained in a separate variable
- **End of the list is indicated by NULL (sentinel).**

```
Example:
struct Node{
   void* data;
   struct Node* next;
};
struct LinkedList {
   struct Node* head;
};

Next pointer field of 2<sup>nd</sup> node
   struct Node* head;
};

Information field of second node

**Today **T
```

Linked Lists: Operations

```
typedef struct Node Node;
Node* new node(void*);
typedef struct LinkedList LinkedList;
LinkedList* new linked list();
void insert at front(LinkedList*, void*);
void insert at back(LinkedList*, void*);
void* remove from front(LinkedList*);
void* remove from back(LinkedList*);
int size(LinkedList*);
int is empty(LinkedList*);
```

```
struct Node{
  void* data;
                Node* next;
};
Node* new node(void* data) {
Node* n=(Node*)
          calloc(1, sizeof(Node));
 n->data = data;
 return n;
struct LinkedList {
Node* head;
};
LinkedList* new linked list() {
 LinkedList* ll=(LinkedList*)
    calloc(1, sizeof(LinkedList));
  return 11;
```

Linked Lists: Operations

```
! Iterating:
    O for (p=head; p!=NULL; p=p->next) /* do something */
    O for (p=head; p->next !=NULL; p=p->next) /* do something */
    O for (p=head; p->next->next !=NULL; p=p->next) /* do something */
❖ int size(LinkedList* ll){
     int result = 0;
     Node* p = ll->head;
     while (p) {
       p=p->next; result++;
     return result;
   int is empty(LinkedList* 11) {
     return !ll->head;
```

Linked Lists: Operations - insert

```
void insert_at_front(LinkedList* ll, void* data) {
  Node* n = new_node(data);
  if (!n) return;
  n->next = ll->head;
  ll->head = n;
}
```

```
void insert_at_back(LinkedList* ll, void* data) {
  Node* n = new_node(data);
  if (!n) return;
  Node* p = ll->head;
  if (!p) ll->head = n;
  else {
    while (p->next) p=p->next;
    p->next = n;
  }
}
```

Linked Lists: Operations - insert

```
void insert after nth(LinkedList* ll, void* data, int n) {
 Node* nn = new node(data);
  if (!nn) return;
  int i=0;
 Node* p = 11->head;
  if (!p) ll->head = nn;
  else {
    while (p->next \&\& i < n) {
     p = p-next; i++;
    nn->next = p->next;
    p->next = nn;
```

Linked Lists: Operations - insert

```
void insert in order(LinkedList* 11, void* data, int(*comp)(void*, void*)){
 Node* n = new node(data);
  if (!n) return;
 Node* p = 11->head;
  if (!p \mid | comp(data, p->data)<0)
   n->next = p;
    11->head = n;
  else {
    while (p->next && comp(data, p->next->data)>0) p=p->next;
    n->next = p->next;
    p->next = n;
```

Linked Lists: Operations - remove

```
void* remove from front(LinkedList*ll) {
 void* result;
 Node* p = ll->head;
  if (!p) return NULL;
  result = p->data;
  ll->head = p->next;
  free(p);
  return result;
```

```
void* remove from back(LinkedList*ll) {
   void* result;
   Node* p = ll->head;
   if (!p) return NULL;
   if (!(p->next)) {
     result = p->data;
     ll->head = NULL;
     free(p);
   else {
     while (p->next->next) p=p->next;
     result = p->next->data;
     free(p->next);
     p->next = NULL;
   return result:
```

Linked List vs Arrays - operations

❖ Time complexity:

0	Operation	Linked List	Array
	Indexing	O(n)	O(1)
	Insert at front	O(1)	O(n)
	Insert at back	O(n)	O(1)
	Remove from front	O(1)	O(n)
	Remove from back	O(n)	O(1)

Other aspects:

Aspect

 Extensibility
 Shifting
 Random access
 Sequential access
 Memory use

Linked List dynamic size not required inefficient slow efficient

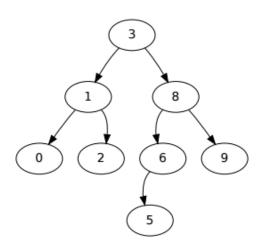
Array
fixed size: expansion is costly
some operations (discuss)
efficient
fast (discuss)
inefficient for large arrays and few data

Binary Trees

- ❖ Binary Tree: dynamic data structure where <u>each node</u> has <u>at most</u> two children
- ❖ A binary search tree is a binary tree with ordering among its children
 - o all elements in the left subtree are assumed to be "less" than the root element
 - o and all elements in the right subtree are assumed to be "greater" than the root element

***** Example:

```
struct tnode{
  void* data; /* payload */
  struct tnode* left;
  struct tnode* right;
};
struct tree{
  struct tnode root;
}
```



Binary Trees

- **The operation on trees can be framed as recursive operations.**
 - Traversal (printing, searching):
 - pre-order: root, left subtree, right subtree
 - inorder: left subtree, root, right subtree
 - post-order: right subtree, right subtree, root

♦ Add node:

```
struct tnode* addnode(struct tnode* root, int data){
  if (root==NULL) { /* termination condition */
    /* allocate node and return new root */
  }
  else if (data < root->data) /* recursive call */
    return addnode(root->left, data);
  else
    return addnode(root->right, data);
}
```

Stack

- A structure that stores data with restricted insertion and removal:
 - o insertion occurs from the top exclusively: push
 - o removal occurs from the top exclusively: pop
- typedef struct Stack Stack;
 Stack* new_stack(int size);
 void* pop(Stack* q);
 void push(Stack* q, void* data);
- * may provide void* top (void); to read last (top) element without removing it

- Stores in an array buffer (static or dynamic allocation)
- * insert and remove done at end of array; need to track end

```
Stack* new_stack(int size) {
   Stack* result = (Stack*)calloc(1,sizeof(Stack));
   result->capacity = size;
   result->buffer = (void**)calloc(size, sizeof(void*));
   return result;
}
```

```
void push(Stack* s, void* data){
  if (s->top < s->capacity)
    s->buffer[s->top++] = data;
}
```

```
void* pop(Stack* s) {
  if (s->top > 0)
    return s->buffer[--(s->top)];
  else return NULL;
}
```

```
struct Stack{
  int capacity;
  void** buffer;
  int top;
};
```

Stack as a Linked List

ll.h ll.c stack.h stack ll.c

- Stores in a linked list (dynamic allocation)
- * "Top" is now at front of linked list (no need to track)

```
struct Stack{
  LinkedList* buffer;
};
```

```
Stack* new_stack(int size) { /* size is not needed */
   Stack* result = (Stack*)calloc(1,sizeof(Stack));
   result->buffer = new_linked_list();
   return result;
}
```

```
void push(Stack* s, void* data) {
  insert_at_front(s->buffer, data);
}
void* pop(Stack* s) {
  return remove_from_front(s->buffer);
}
```

Queue

• Opposite of stack:

- o first in: enqueue
- o first out: dequeue
- Read and write from opposite ends of list

❖ Important for:

- UIs (event/message queues)
- o networking (Tx, Rx packet queues)
- 0

❖ Imposes an ordering on elements

```
typedef struct Queue Queue;
Queue* new_queue(int size);
void* dequeue(Queue* q);
void enqueue(Queue* q, void* data);
```

Queue as an Array

- queue.h queue ar.c test1.c
- ❖ Stores in an array buffer (static or dynamic allocation);
- Elements added to rear, removed from front

```
o need to keep track of front and rear: int front=0, rear=0;
```

```
o or, track the front and number of elements: int front=0, count=0;
Queue* new_queue(int size){
  Queue* result = (Queue*)calloc(1,sizeof(Queue));
  result->capacity = size;
  result->buffer = (void**)calloc(size,sizeof(void*));
```

```
struct Queue{
  int capacity;
  void** buffer;
  int front;
  int count;
};
```

```
void enqueue(Queue* q, void* data){
  if (q->count < q->capacity){
    q->buffer[q->front+q->count] = data;
    q->count++;
  }
}
```

return result;

```
void* dequeue(Queue* q) {
  if (q->count > 0) {
    q->count--;
    return q->buffer[q->front++];
  }
  else return NULL;
}
```

Queue as an Array

queue.h queue arr.c test1.c

- Let us try a queue of capacity 4:
 - o enqueue a, enqueue b, enqueue c, enqueue d
 - o queue is now full.
 - o dequeue, enqueue e: where should it go?
- Solution: use a circular (or ring) buffer
 - 'e' would go in the beginning of the array
- * Need to modify enqueue and dequeue:

```
a b c d front rear
```

```
void* dequeue(Queue* q) {
  void* result = NULL;
  if (q->count > 0) {
    q->count--;
    result=q->buffer[q->front++];
    if (q->front == q->capacity)
        q->front = 0;
  }
  return result;
}
```

Queue as a Linked List 11.h 11.c queue.h queue 11.c

❖ Stores in a linked list (dynamic allocation)

void enqueue (Queue* q, void* data) {
 insert at back(q->buffer, data);

```
Queue* new_queue(int size){
    /* size is not needed*/
    Queue* result = (Queue*)calloc(1,sizeof(Queue));
    result->buffer = new_linked_list();
    return result;
}
```

```
void* dequeue(Queue* q) {
  return remove_from_front(q->buffer);
}
```

struct Queue{

LinkedList* buffer;

Example: Postfix Evaluator

- ❖ Stacks and queues allow us to design a simple expression evaluator
- Prefix, infix, postfix notation:
 - o operator before, between, and after operands, respectively
 - o Infix
 - \blacksquare A + B
 - A * B C
 - \blacksquare (A + B) * (C D)
 - Prefix
 - + A B
 - -*ABC
 - * + A B C D
 - Postfix
 - AB+
 - AB*C
 - \blacksquare AB+CD-*
 - Infix more natural to write, postfix easier to evaluate

Example: Postfix Evaluator

```
float pf eval(char* exp) {
 Stack* S = new stack(0);
 while (*exp) {
   if (isdigit(*exp) || *exp=='.'){
      float* num; num = (float*)malloc(sizeof(float));
      sscanf(exp, "%f", num);
      push(S, num);
     while(isdigit(*exp) || *exp=='.') exp++; exp--;
   else if (*exp!= ' ') {
      float num1 = *(float*)pop(S), num2 = *(float*)pop(S);
      float* num; num = (float*)malloc(sizeof(float));
      switch (*exp){
       case '+': *num = num1+num2; break;
       case '-': *num = num1-num2; break;
       case '*': *num = num1*num2; break;
       case '/': *num = num1/num2; break;
     push(S, num);
   exp++;
 return *(float*)pop(S);
```

- ❖ Hash tables (hashmaps) an efficient data structure for storing dynamic data.
- commonly implemented as an array of linked lists (hash tables with chaining).
- **Each** data item is associated with a key that determines its location.
 - Hash functions are used to generate an evenly distributed hash value.
 - A hash collision is said to occur when two items have the same hash value.
 - Items with the same hash keys are chained
 - \circ Retrieving an item is O(1) operation.
- * Hash function: map its input into a finite range: hash value, hash code.
 - The hash value should ideally have uniform distribution. why?
 - Other uses of hash functions: cryptography, caches (computers/internet), bloom filters etc.
 - Hash function types:
 - Division type
 - Multiplication type
 - Other ways to avoid collision: linear probing, double hashing.

```
struct Pair{
     char* key;
     void* data;
  struct HashTable{
     LinkedList* buckets;
     int capacity;
   };
unsigned long int hash(char* key);
   HashTable* new hashtable(int size);
   int is empty ht(HashTable* ht);
   int length(HashTable* ht);
   void insert(HashTable* ht, Pair* p);
   void* remove(char* key);
   void* retrieve(char* key);
```

```
HashTable* new hashtable(int size){
  HashTable* result = (HashTable*)calloc(1, sizeof(HashTable));
  result->size = size;
  result->buckets = (LinkedList*)calloc(size, sizeof(LinkedList));
  return result;
int is empty ht(HashTable* ht){
  int i;
 for (i=0; i < ht->size; i++)
   if (!is empty(&(ht->buckets[i])))
     return 0:
  return 1:
int length(HashTable* ht){
  int i, result=0;
  for (i=0; i < ht->size; i++)
   result += size(&(ht->buckets[i]));
  return result;
```

```
unsigned long int hash(char* key) {
  /* any good hashing algorithm */
  const int MULTIPLIER = 31;
  unsigned long int hashval = 0;
  while (*key)
   hashval = hashval * multiplier + *key++;
  return hashval;
void insert(HashTable* ht, Pair* p) {
  if (retrieve(ht, p->key)) return NULL;
  int index = hash(p->key) % ht->capacity;
  insert at front(&(ht->buckets[index]));
void* remove(char* key) {
  /* since we do not have a ready supporting ll function we'll go low level */
void* retrieve(char* key) {
  /* since we do not have a ready supporting ll function we'll go low level */
```