

What is CPU utilization ?

Keep CPU busy as possible

What is CPU-I/O burst cycle ?

Process execution consists of a cycle of CPU execution and I/O wait
That means the CPU either work or wait for I/O

What is CPU scheduler (short-term scheduler) ?

Selects process from the ready queue to be execute

When the CU scheduler will take place (work)?

1. Switches from running to waiting state (I/O interrupt) **nonpreemptive**
2. Switches from running to ready state (timer interrupt) **preemptive**
3. Switches from waiting to ready(higher priority process arrive to ready queue) **preemptive**
4. Terminates **nonpreemptive**

What is the preemptive?

- **nonpreemptive** : once CPU given to the process it cannot be preempted until completes its CPU burst
- **preemptive** : if a new process arrives with CPU burst length less than remaining time of current executing process, preempt. This scheme is know as the Shortest-Remaining-Time-First (SRTF)

What is dispatcher ?

Dispatcher module gives control of the CPU to the process selected by the short-term scheduler

What is the task of dispatcher ?

- switching context
- switching to user mode
- jumping to the proper location in the user program to restart that program

What is the dispatcher latency ?

time it takes for the dispatcher to stop one process and start another running

What is the scheduling criteria ?

- **CPU utilization** : keep the CPU as busy as possible
- **Throughput** : # of processes that complete their execution per time unit
- **Turnaround time** : amount of time to execute a particular process
- **Waiting time** : amount of time a process has been waiting in the ready queue
- **Response time** : amount of time it takes from when a request was submitted until the first response is produced, not output (for time-sharing environment)

What we must keep for scheduling criteria ?

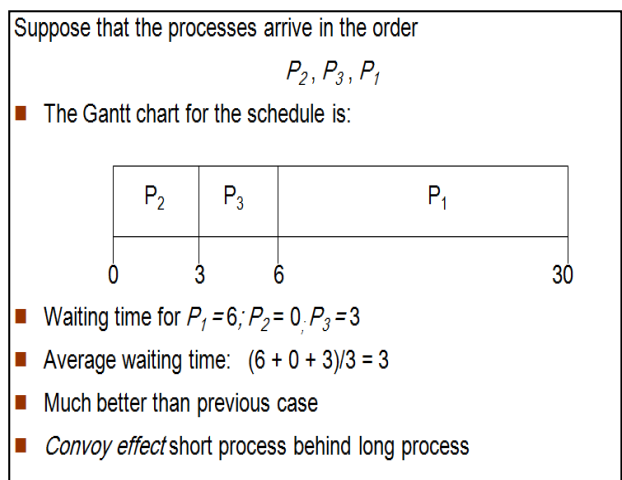
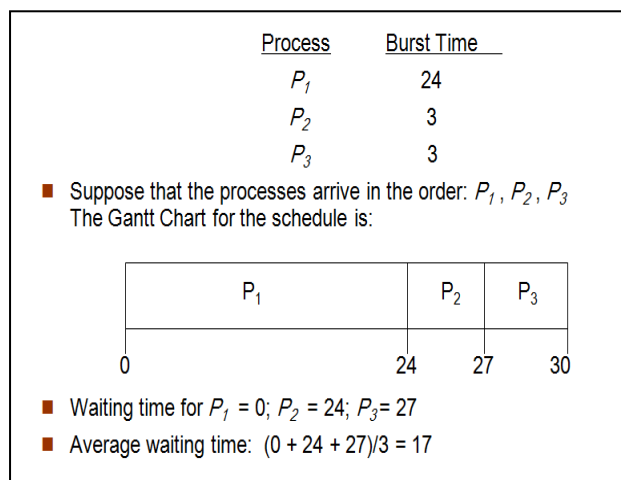
- **Max** CPU utilization
- **Max** throughput
- **Min** turnaround time
- **Min** waiting time
- **Min** response time

What is the types of scheduler algorithms ?

- 1- First come, first served (FCFS)
- 2- Shortest Job First (SJF)
- 3- Priority
- 4- Round robin (RR)
- 5- Multilevel queue

Explain FCFS scheduling ?

- **Advantages** : very simple in implementation
- **disadvantages** : when process with high burst time come first the other processes will wait for long time to start execute.

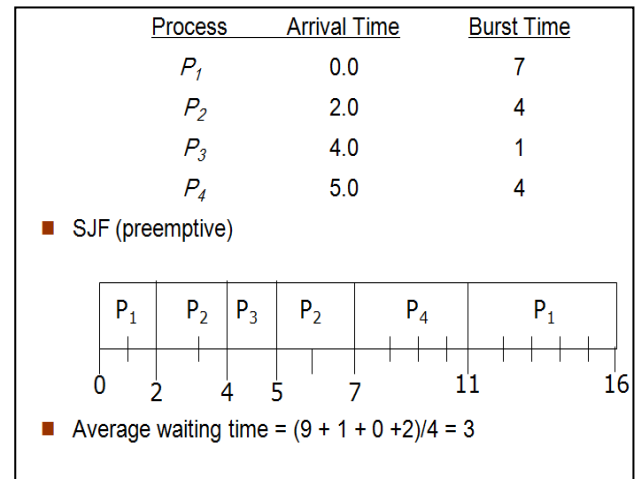
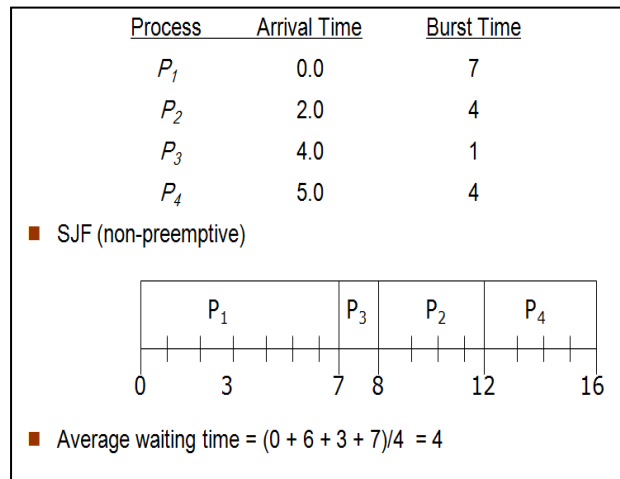


explain SJR scheduling ?

Associate with each process the length of its next CPU burst. Use these lengths to schedule the process with the shortest time

Advantages : SJF is optimal . gives minimum average waiting time for a given set of processes

Disadvantages : The difficulty is knowing the length of the next CPU request



explain the previous preemptive example ?

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	17
P1	7	6	5	5	5	5	5	5	5	5	5	5	4	3	2	1	0
P2	---	---	4	3	2	2	1	0									
P3	---	---	---	---	1	0											
P4	---	---	---	---	---	4	4	4	3	2	1	0					

How to Determining Length of Next CPU Burst ?

1. t_n = actual length of n^{th} CPU burst
2. τ_{n+1} = predicted value for the next CPU burst
3. $\alpha, 0 \leq \alpha \leq 1$
4. Define : $\tau_{n+1} = \alpha t_n + (1 - \alpha)\tau_n$.

- If we assume that : **t0 = 6** and **T0 = 10**

CPU burst (t_i)	6	4	6	4	13	13	13	...	
"guess" (τ_i)	10	8	6	6	5	9	11	12	...

Explain the priority scheduling ?

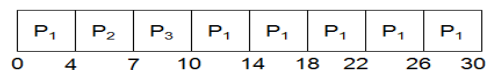
- A priority number (integer) is associated with each process
- The CPU is allocated to the process with the highest priority (smallest)
- SJF is a priority scheduling where priority is the predicted next CPU burst time
- **Problem** \equiv **Starvation** – low priority processes may never execute
- **Solution** \equiv **Aging** –as time progresses increase the priority of the process (over head)
- **Disadvantage** : not suitable to time sharing

Explain round robin scheduling ?

- Each process gets a small unit of CPU time (*time quantum*), usually 10-100 milliseconds. After this time has elapsed, the process is preempted and added to the end of the ready queue.
- If there are n processes in the ready queue and the time quantum is q , then each process gets $1/n$ of the CPU time in chunks of at most q time units at once. No process waits more than $(n-1)q$ time units.
- **Advantage** : not suitable to time sharing
- **Disadvantage** : higher average turnaround than SJF but better response
- **Performance**
 - If q large \Rightarrow FIFO
 - If q small \Rightarrow overhead

Process	Burst Time
P_1	24
P_2	3
P_3	3

■ The Gantt chart is:



■ Typically, higher average turnaround than SJF, but better response

Are there a relation between time quantum and the average of turnaround ?

NO

Explain multilevel scheduling ?

- Ready queue is partitioned into separate queues:
 - foreground (interactive)
 - background (batch)
- Each queue has its own scheduling algorithm
 - foreground – RR
 - background – FCFS

What is types of scheduling between queues in multilevel scheduling ?

- **Fixed priority scheduling:** (i.e., serve all from foreground then from

background). Possibility of starvation.

- **Time slice** : each queue gets a certain amount of CPU time which it can schedule amongst its processes; i.e., 80% to foreground in RR and 20% to background in FCFS

Explain multilevel feedback queue ?

A process can move between the various queues

What is the parameters that Multilevel-feedback-queue scheduler defined by them ?

- number of queues
- scheduling algorithms for each queue
- method used to determine when to upgrade a process
- method used to determine when to demote a process
- method used to determine which queue a process will enter when that process needs service

Give example of multilevel feedback queue ?

- Three queues:
 - Q_0 – RR with time quantum 8 milliseconds
 - Q_1 – RR time quantum 16 milliseconds
 - Q_2 – FCFS
- Scheduling
 - A new job enters queue Q_0 which is served RR. When it gains CPU, job receives 8 milliseconds. If it does not finish in 8 milliseconds, job is moved to queue Q_1 .
 - At Q_1 job is again served RR and receives 16 additional milliseconds. If it still does not complete, it is preempted and moved to queue Q_2 .

