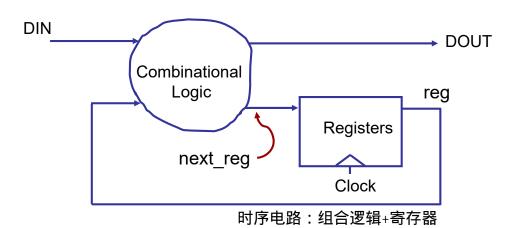
Chapter 6: Modeling at the RT Level

 A register transfer level (RTL) design consists of a set of registers connected by combinational logic.



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6.1 Combinational circuit

- A combinational circuit, by definition, is a circuit whose output, after the initial transient period, is a function of current input.
- It has no internal state and therefore is "memoryless" about the past event (or past inputs).

```
signal A, B, Cin, Cout : bit;
...
process (A, B, Cin) is
begin
   Cout <= (A and B) or ((A xor B) and Cin);
end;</pre>
```

To describe a combinational circuit

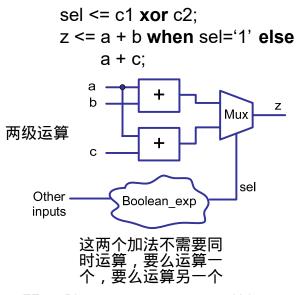
- The variables or signals in the process must not have initial values.
- A signal or a variable must be assigned a value before being referenced.
- The arithmetic operators (such as +, -, *, etc), relational operators (such as <, >, =, etc), and logic operators (such as and, or, not, etc) can be used in an expression.

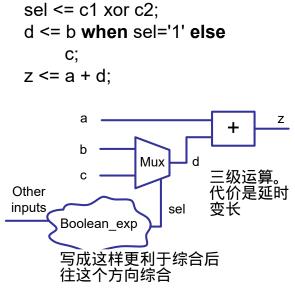
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Operator sharing

 One way to reduce the overall size of synthesized hardware is to identify the resources that can be used by different operations. This is know as resource sharing.





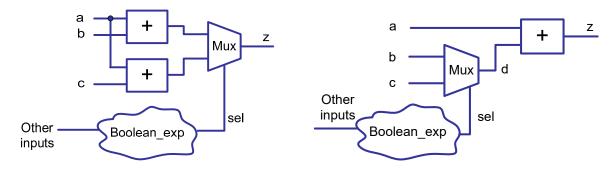
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- Performing resource sharing normally introduces some overhead and may penalize performance.
 - In the above examples, assume T_{adder} , T_{mux} , $T_{boolean}$,
 - For the circuit not sharing the adders:

$$T = \max (T_{adder}, T_{boolean}) + T_{mux}$$

• For the circuit sharing the adders:

$$T = T_{adder} + T_{boolean} + T_{mux}$$



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Shaping the circuit

 Using VHDL code, it is possible to outline the general shape of the circuit.

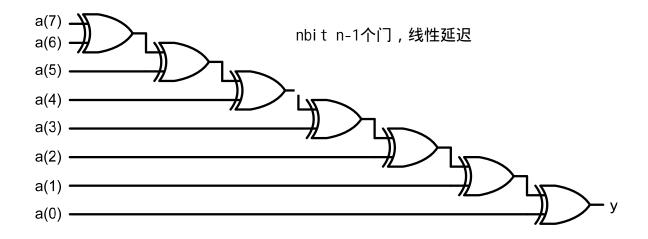
Reduced-xor circuit

$$y = a_7 \oplus a_6 \oplus a_5 \oplus a_4 \oplus a_3 \oplus a_2 \oplus a_1 \oplus a_0$$

奇偶校验
signal a: std_logic_vector (7 downto 0);
signal y: std_logic;
...
 $y \le a(7) xor a(6) xor a(5) xor a(4) xor a(3)$
 $xor a(2) xor a(1) xor a(0);$

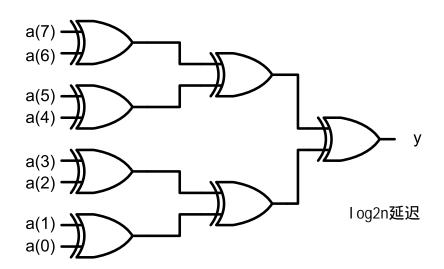
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$y \le ((((((((a_7 \oplus a_6) \oplus a_5) \oplus a_4) \oplus a_3) \oplus a_2) \oplus a_1) \oplus a_0);$



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$$y \le ((a_7 \oplus a_6) \oplus (a_5 \oplus a_4)) \oplus ((a_3 \oplus a_2) \oplus (a_1 \oplus a_0));$$



Example: Combinational adder-based multiplier

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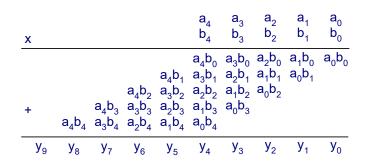
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The algorithm includes three tasks:

Multiply the digits of the multiplier (b4, b3, b2, b1 and b0) by the multiplicand A = (a4, a3, a2, a1, a0) one at a time to obtain b4*A, b3*A, b2*A, b1*A and b0*A.

bi * A =
$$(a4 \cdot bi, a3 \cdot bi, a2 \cdot bi, a1 \cdot bi, a0 \cdot bi)$$

- Shift bi * A to left by i position.
- Add the shifted bi * A terms to obtain the final product.



Initial description of an adder-based multiplier

```
\mathbf{a}_0
                                                                                               b_1
                                                                                          b_2
                                                                                                     b_0
                                                                                b_{\lambda}
                                                                                     b_3
library IEEE;
                                                                               a_4b_0 \ a_3b_0 \ a_2b_0 \ a_1b_0
use ieee.std logic 1164.all;
                                                                          a_4b_1 \ a_3b_1 \ a_2b_1 \ a_1b_1 \ a_0b_1
use ieee.std logic arith.all;
                                                                     a_4b_2 \ a_3b_2 \ a_2b_2 \ a_1b_2 \ a_0b_2
                                                                a_4b_3 \ a_3b_3 \ a_2b_3 \ a_1b_3 \ a_0b_3
use ieee.std logic unsigned.all;
                                                           a_4b_4 \ a_3b_4 \ a_2b_4 \ a_1b_4 \ a_0b_4
                                                       y<sub>9</sub>
                                                            y<sub>8</sub>
                                                                 y<sub>7</sub>
                                                                      y<sub>6</sub>
                                                                           y<sub>5</sub>
                                                                                y_4
                                                                                     y_3
                                                                                          y_2
                                                                                               y_1
                                                                                                     y_0
entity mult5 is
port (a, b : in std logic vector(4 downto 0);
        y: out std logic vector(9 downto 0));
end entity mult5:
architecture comb1 arch of mult5 is
constant WIDTH: integer := 5;
signal au, bv0, bv1, bv2, bv3, bv4: std logic vector(WIDTH-1 downto 0);
signal p0, p1, p2, p3, p4, prod: std logic vector(2*WIDTH-1 downto 0);
                                                                   product延长到10位
```

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```
Array aggregate:
                                             a VHDL construct to assign a value
                                             to an object of array data type.
begin
                                             v <= "1011";
    au <= a:
                                             v <= ('1', '0', '1', '1');
    bv0 \le (others => b(0)):
                                             v <= (3=>'1', 2=>'0', 1=>'1', 0=>'1');
                                     same
    bv1 \le (others => b(1));
                                             v \le (3|1|0=>'1', 2=>'0');
    bv2 \le (others => b(2)):
                                             v \le (2=>'0', others=>'1');
    bv3 \le (others => b(3));
                                             v \le (others = >'0');
    bv4 \le (others => b(4));
    p0 \le "00000" \& (bv0 and au);
    p1 <= "0000" & (bv1 and au) & '0';
                                               需要四个十
    p2 <= "000" & (bv2 and au) & "00";
                                               位加法器
    p3 <= "00" & (bv3 and au) & "000"; 共四轮赋值
    p4 <= '0' & (bv4 and au) & "0000";
    prod \le ((p0+p1)+(p2+p3))+p4;
                                                                          b_3
                                                                              b_2
                                                                                   b_1
                                                                                       b_0
    y \le prod;
                                                                     a_4b_0 \ a_3b_0 \ a_2b_0 \ a_1b_0 \ a_0b_0
end architecture comb1 arch;
                                                                a_4b_1 \ a_3b_1 \ a_2b_1 \ a_1b_1 \ a_0b_1
                                                            a_4b_2 \ a_3b_2 \ a_2b_2 \ a_1b_2 \ a_0b_2
                                                        a_4b_3 \ a_3b_3 \ a_2b_3 \ a_1b_3
                                                                a_1b_4 a_0b_4
                                                                      y_4
                                                                          y_3
                                                                              y_2
                                                                                       y_0
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                                                                                       178
```

```
X
                                               a_4b_0 a_3b_0 a_2b_0 a_1b_0 a_0b_0
                                      pp0_5 pp0_4 pp0_3 pp0_2 pp0_1 pp0_0 partial product pp0
                              + a_4b_1 a_3b_1 a_2b_1 a_1b_1 a_0b_1
                            pp1<sub>5</sub> pp1<sub>4</sub> pp1<sub>3</sub> pp1<sub>2</sub> pp1<sub>1</sub> pp1<sub>0</sub>
                                                                                          partial product pp1
                     + a_4b_2 a_3b_2 a_2b_2 a_1b_2 a_0b_2
                    pp2<sub>5</sub> pp2<sub>4</sub> pp2<sub>3</sub> pp2<sub>2</sub> pp2<sub>1</sub> pp2<sub>0</sub>
                                                                                           partial product pp2
                a_4b_3 a_3b_3 a_2b_3 a_1b_3 a_0b_3
            pp3<sub>5</sub> pp3<sub>4</sub> pp3<sub>3</sub> pp3<sub>2</sub> pp3<sub>1</sub> pp3<sub>0</sub>
                                                                                          partial product pp3
          a_4b_4 a_3b_4 a_2b_4 a_1b_4 a_0b_4
    \mathsf{pp4}_5 \ \mathsf{pp4}_4 \ \mathsf{pp4}_3 \ \mathsf{pp4}_2 \ \mathsf{pp4}_1 \ \mathsf{pp4}_0
                                                                                          partial product pp4
     y_9 y_8 y_7 y_6 y_5 y_4 y_3 y_2 y_1 y_0
```

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More efficient description of an adder-based multiplier

```
a<sub>0</sub> multiplicand
                                                                                                              b<sub>0</sub> multiplier
                                                                                    a_4b_0 \ a_3b_0 \ a_2b_0 \ a_1b_0 \ a_0b_0
                                                                               pp0_5 pp0_4 pp0_3 pp0_2 pp0_1 pp0_0 partial product pp0
                                                                               a_4b_1 a_3b_1 a_2b_1 a_1b_1 a_0b_1
begin
                                                                         pp1<sub>5</sub> pp1<sub>4</sub> pp1<sub>3</sub> pp1<sub>2</sub> pp1<sub>1</sub> pp1<sub>0</sub>
                                                                                                                 partial product pp1
                                                                          a_4b_2 \ a_3b_2 \ a_2b_2 \ a_1b_2 \ a_0b_2
      au <= a:
                                                                   pp2<sub>5</sub> pp2<sub>4</sub> pp2<sub>3</sub> pp2<sub>2</sub> pp2<sub>1</sub> pp2<sub>0</sub>
                                                                                                                 partial product pp2
      bv0 \le (others => b(0));
                                                                    a_4b_3 \ a_3b_3 \ a_2b_3 \ a_1b_3 \ a_0b_3
      bv1 \le (others => b(1)):
                                                             pp3<sub>5</sub> pp3<sub>4</sub> pp3<sub>3</sub> pp3<sub>2</sub> pp3<sub>1</sub> pp3<sub>0</sub>
                                                                                                                 partial product pp3
                                                             a_4b_4 a_3b_4 a_2b_4 a_1b_4 a_0b_4
      bv2 \le (others => b(2)):
                                                        pp4<sub>5</sub> pp4<sub>4</sub> pp4<sub>3</sub> pp4<sub>2</sub> pp4<sub>1</sub> pp4<sub>0</sub>
                                                                                                                 partial product pp4
      bv3 \le (others => b(3));
                                                                                            y_3 y_2 y_1
                                                                                                            y_0
      bv4 \le (others => b(4));
      pp0 <= '0' & (bv0 and au);
      pp1 \le ('0' \& pp0(WIDTH downto 1)) + ('0' \& (bv1 and au));
      pp2 \le (0' \& pp1(WIDTH downto 1)) + (0' \& (bv2 and au));
      pp3 <= ('0' & pp2(WIDTH downto \frac{1}{1})+ ('0' & (bv3 and au));
      pp4 \le ('0' \& pp3(WIDTH downto 1)) + ('0' \& (bv4 and au));
      prod \le pp4 \& pp3(0) \& pp2(0) \& pp1(0) \& pp0(0);
      y \le prod;
end architecture comb2 arch;
                                                                     只需要4个六位
                                                                     加法器,面积减
                                                                     小,功耗降低
```

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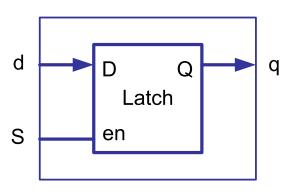
6.2 Sequential circuit

- A sequential circuit is a circuit that has an internal state, or memory.
- Its output is a function of current input as well as the internal state. Thus the output is affected by current input values as well as past input values.
- A synchronous sequential circuit, in which all memory elements are controlled by a global synchronizing signal, greatly simplifies the design process and is the most important design methodology.
- Flip-flops and latches are two commonly used one-bit memory devices.

6.2.1 Latch

A latch is a level-sensitive memory device.

```
signal S, d, q: bit
.....
process (S, d) is
begin
    if (S='1') then
        q <= d;
    end if;
end process;</pre>
```



Truth table S q²
0 q
1 d

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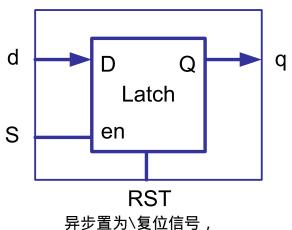
- In general, latches are synthesized from incompletely specified conditional expressions in a combinational description.
- Latch inferences occur normally with if statements or case statements.
- To avoid having a latch inferred, assign a value to the signal under all conditions.

```
signal S, d, q: bit
.....
process (S, d) is
begin
   if (S='1') then q <= d;
   else q <= '0';
   end if;
end process;</pre>
```

asynchronous reset or preset

 An asynchronous reset (or preset) will change the output of a latch to 0 (or 1) immediately.

```
signal S, RST, d, q: bit
.....
process (S, RST, d) is
begin
    if (RST = '1') then
        q <= '0';
    elsif (S='1') then
        q <= d;
    end if;
end process;</pre>
```



カン量が、後位にう、 同步电路中允许存在异步复位

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6.2.2 Flip-Flops (f/f) 边缘触发

- A flip-flop is an edge-triggered memory device.
- To detect the rising edge (or falling edge), or the event occurred for a signal, we can make use of the attribute of a signal.

```
signal CLK : bit;
.....

CLK'event true if CLK changes its value.

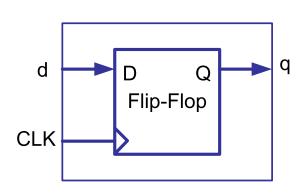
CLK'event and CLK = '1' true for the CLK rising edge

CLK'event and CLK = '0' true for the CLK falling edge
```

 The event attribute on a signal is the most commonly used edge-detecting mechanism. It operates on a signal and returns a Boolean value. The result is true if the signal shows a change in value.

An example of a simple flip-flop

 An edge triggered flip-flop will be generated from a VHDL description if a signal assignment is executed on the rising (or falling) edge of another signal.

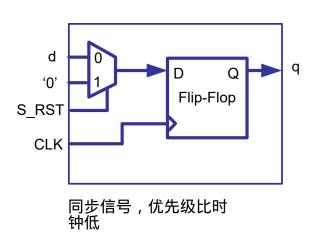


```
entity dff is
port (d, CLK: in bit; q: out bit);
end entity dff;
architecture behavior of dff is
begin
process (CLK) is 敏感列表只需要有clk,其他
begin 的都不需要
if (CLK'event and CLK='1') then
q <= d;
end if;
end process;
end architecture behavior;
```

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Synchronous sets and resets

 Synchronous inputs set (preset) or reset (clear) the output of flip-flops when they are asserted. The assignment will only take effect while the clock edge is active.



```
signal CLK, d, q, S_RST: bit;
.....

process (CLK) is

begin

if (CLK'event and CLK='1') then

if (S_RST = '1') then

cq <= '0';

else

q <= d;

end if;

end process;
```

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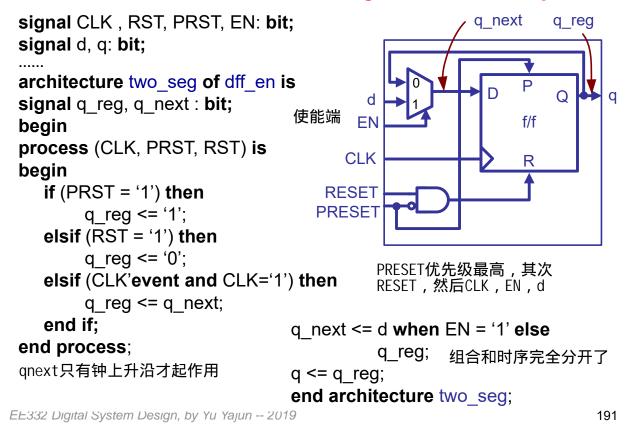
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Asynchronous sets and resets

 Asynchronous inputs set (preset) or reset (clear) the output of flip-flops whenever they are asserted independent of the clock.

```
signal CLK, A RST, d, q: bit;
                                                                  q
                                        d
                                                 D
process (CLK, A RST) is
                                                 Flip-Flop
begin
  if (A RST = '1') then
                                     CLK
       q \le '0';
  elsif (CLK'event and CLK='1') then
                                                   A RST
       q \le d;
                             异步逻辑,复位优先级比时钟信号高
  end if:
end process;
```

A f/f with more than one asynchronous input



6.2.3 VHDL templates for sequential circuits

- An RTL circuit can be described in two segments:
 - A synchronous section updates the register information at the rising edge of the clock.
 - q req <= q next;
 - A combinational section describes combinational logics, for example, update q_next;

A synchronous section

or

A synchronous section with asynchronous inputs

.

A combinational section

if (CLK'event AND CLK='1') then

 $q_reg \le q_next;$

-- CLK is the clock input to reg q

end if;

```
if (async_sig = '1') then
q_reg <= '0';
```

-- active high asynchronous reset

elsif (CLK's event AND CLK='1') then

 $q_reg \le q_next;$

-- CLK is the clock input to reg Q

end if;

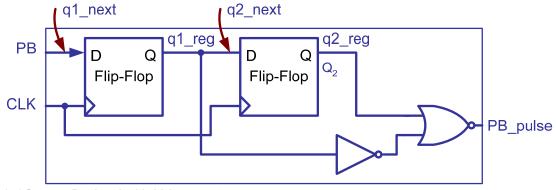
```
q_next <= expression;
-- other combinational logics.</pre>
```

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```
An example
```

```
entity PULSER is
    port (CLK, PB : in bit ;
        PB_pulse : out bit);
end PULSER;
architecture BHV of PULSER is
    signal q1_reg, q2_reg,
        q1_next, q2_next : bit;
begin
```

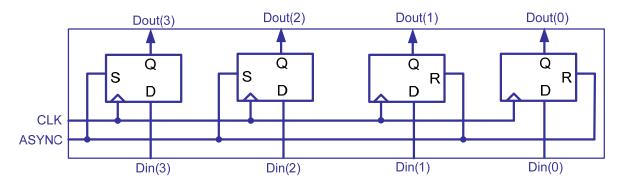
```
process (CLK) is
begin
   if (CLK'event and CLK='1') then
        q1_reg <= q1_next;
        q2_reg <= q2_next;
   end if;
end process;
   q1_next <= PB;
   q2_next <= q1_reg;
   PB_pulse <= (not q1_reg) nor q2_reg;
end architecture BHV;</pre>
```



6.2.4 Registers

-- 4-bit simple register

```
signal CLK , ASYNC : bit;
singal Din, Dout :
   bit_vector (3 down to 0);
```



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-- 4-bit serial-in and serial-out shift register

```
signal CLK ,d, q : bit;
   architecture two seg of shift register is
   signal r reg, r next: bit vector (3 downto 0);
   begin
   process (CLK) is
   begin
       if (CLK'event and CLK='1') then
                                                 r \text{ next} \le d \& r \text{ reg}(3 \text{ downto } 1);
             r reg <= r next;
                                                 q \leq r \operatorname{reg}(0);
       end if:
                                                 end architecture two seg;
   end process;
                                                               r_next(1)
                  r_reg(2)
              r_next(2)
                                                                  r_reg(1)
                                                                              r next(0)
    r_reg(3)
r_next(3)
                                                                                   r_reg(0)
                                                                                    q
    d
                    Q
                               D
                                       Q
            D
                                                  D
                                                          Q
                                                                     D
                f/f
                                   f/f
                                                      f/f
                                                                        f/f
 CLK
```

6.2.5 Synchronous counter

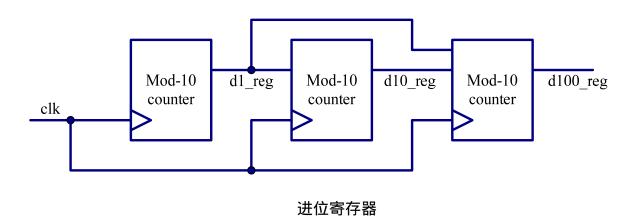
-- 4-bit synchronous counter library IEEE; use ieee.std_logic_1164.all; use ieee.std_logic_arith.all; use ieee.std_logic_unsigned.all; entity counter is port (CLK, RESET: in std_logic; load, Count, UpDown: in std_logic; Dataln: in std_logic_vector(3 downto 0); DataOut: out std_logic_vector(3 downto 0)); end entity counter; architecture two-seg of counter is signal Reg, Reg_Next: std_logic_vector (3 downto 0);

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```
begin
process (CLK, RESET) is
begin
if RESET = '1' then
   Reg <="0000";
elsif CLK'event and CLK='1' then
   Reg <= Reg Next;
end if;
end process;
-- next-state logic
Reg_next <= DataIn when load = '1' else
   (Reg+1) when (Count='1' and UpDown = '1') else
   (Reg-1) when (Count='1' and UpDown ='0') else
   Reg:
-- Output logic
DataOut <= Reg;
end architecture two-seg;
```

Decimal counter

- A decimal counter circulates the patterns in binary-coded decimal (BCD) format.
- The BCD code use 4 bits to represent a decimal number.



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Three-digit decimal counter using conditional concurrent statements

```
library IEEE;
use ieee.std_logic_arith.all;
use ieee.std_logic_unsigned.all;
entity decimal_counter is
port (CLK, RESET: in std_logic;
    d1, d10, d100: out std_logic_vector(3 downto 0));
end entity decimal_ounter;
architecture concurrent_arch of decimal_ounter is
signal d1_reg, d10_reg, d100_reg: std_logic_vector (3 downto 0);
signal d1_next, d10_next, d100_next: std_logic_vector (3 downto 0);
begin
```

```
-- register
process (CLK, RESET) is
begin
if RESET = '1' then
   d1 reg <="0000";
   d10 reg <="0000";
   d100 reg <="0000";
elsif CLK'event and CLK='1' then
   d1 req <= d1 next;
   d10 \text{ reg} \leq d10 \text{ next};
   d100 reg <= d100 next;
end if:
end process;
-- next-state logic
d1 next <= "0000" when d1 reg = 9 else d1 reg+1;
d10 next <= "0000" when (d1 reg = 9 and d10 reg = 9) else
             d10 reg+1 when d1 reg = 9 else d10 reg;
d100 next <= "0000" when (d1 reg=9 and d10 reg=9 and d100 reg=9) else
             d100 reg+1 when (d1 reg=9 and d10 reg=9) else d100 reg;
-- Output logic
d1 <= d1 reg; d10 <= d10 reg; d100 <= d100 reg;
end architecture concurrent arch;
```

Three-digit decimal counter using a nested if statement

```
architecture if arch of decimal ounter is
   signal d1 reg, d10 reg, d100 reg; std logic vector (3 downto 0);
  signal d1 next, d10 next, d100 next: std logic vector (3 downto 0);
begin
process (CLK, RESET) is
begin
if RESET = '1' then
  d1 reg <="0000":
  d10_reg <="0000";
  d100 reg <="0000";
elsif CLK'event and CLK='1' then
  d1 req <= d1 next;
  d10 reg <= d10 next;
  d100 reg <= d100 next;
end if;
end process:
```

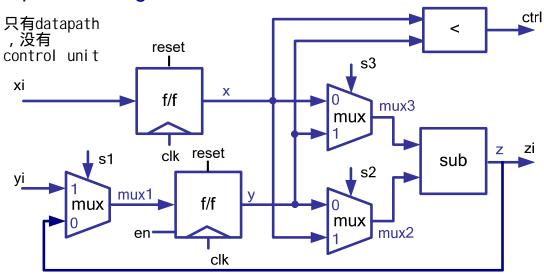
```
-- next-state logic
process (d1_reg, d10_reg, d100_reg)敏感列表不需要next
begin
    d10 next <= d10_reg;
                                用i f的话必须要有
    d100 next <= d100_reg; default value
    if d1 reg \neq 9 then
        d1 \text{ next} = d1 \text{ reg+1};
    else -- reach 9
        d1 next = "0000";
        if d10 \text{ reg} = 9 \text{ then}
              d10 \text{ next} \le d10 \text{ reg} + 1;
        else -- reach 99
              d10 next <= "0000";
              if d100 \text{ reg} = 9 \text{ then}
                  d100 next <= d100_reg + 1;
              else -- reach 999
                  d100 next <= "0000";
              end if:
        end if;
    end if;
end process;
-- Output logic
d1 <= d1 reg; d10 <= d10 reg; d100 <= d100 reg;
end architecture if arch;
```

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6.3 Netlist of RTL components

 A data path usually consists of a netlist of RTL components such as function units, multiplexers, comparators, registers, etc.



```
signal clk, en, s1, s2, s3 : std logic;
signal xi, yi, zi : std logic vector(7 downto 0);
signal ctrl: boolean;
architecture two seg of datapath is
signal x, y, z, x_next, y_next, mux1, mux2, mux3 : std_logic_vector (7 downto 0);
begin
process (clk) is
begin
if reset ='1' then
  x \le "00000000";
  y <= "00000000";
                                               mux1 <= z when s1 = '0' else
elsif (clk'event and clk='1') then -- registers
                                                        yi,
                                                                 -- multiplexer
                                               mux2 <= y when s2 = '0' else
  x \le x \text{ next}
                                                         x; -- multiplexer
  y \le y_next;
end if;
                                               mux3 <= x when s3 = '0' else
end process;
                                                                  -- multiplexer
ctrl \le (x \le y); -- comparator
                                               z <= mux3 - mux2; -- subtracter
x next <= xi;
                                               zi \le z;
y next <= mux1 when en = '1' else
                                               end architecture two seg;
          у;
```

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6.4 Test benches for sequential system

- All synchronous system require a system clock signal.
- A reset signal is required. The reset signal is asserted at power on to place the sequential system in its initial state.

6.4.1 Generating a system clock

50% duty cycle clock

```
clock_gen: process
constant period : time := 100 ns;
begin
  clk <= '0';
  wait for period/2;
  clk <= '1';
  wait for period/2;
end process;</pre>
```

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6.4.2 Generating the system reset

- The reset signal typically
 - starts in its asserted state at power on,
 - remains in that state for a specified period of time, then
 - changes to its unasserted state, and
 - remains there for as long as power continues to be applied to the system.

- The duration of the assertion of the reset signal is specified as
 - either a fixed time

```
reset <= '1', '0' after 160 ns;
```

 or some multiple of the clock's period and is synchronized to the system clock

```
reset_process : process
begin
reset <= '1';
for i in 1 to 2 loop
wait until clk = '1';
end loop;
reset <= '0';
wait;
end process;
```

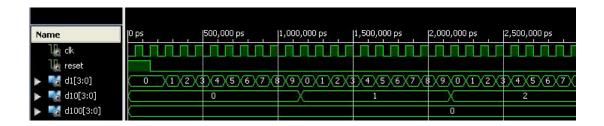
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6.4.3 Synchronizing stimulus generation and monitoring

```
monitor: process
   constant n : integer := 1000;
   variable number: integer range 0 to 999:=0;
begin
   wait until reset <= '0';
   wait for 1 ns;
   for i in 0 to n loop
      number := to integer(unsigned(d100))*100+to integer(unsigned(d10))*10
                   + to integer(unsigned(d1));
      assert number = i mod n
      report "count of " & integer'image(i mod n) & " failed"
      severity error;
      wait until clk = '1';
      wait for 1 ns; 完成赋值的
   end loop;
                    时间
   wait;
end process;
```

Waveforms of clk and reset

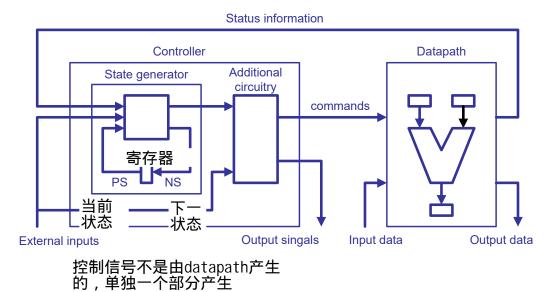


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```
use ieee.numeric_std.all;
signal x : std_logic_vector(3 downto 0); -- vector with element std_logic
signal y : unsigned(3 downto 0); -- vector with element std_logic
signal z : integer range 0 to 15;
conversion between std_logic_vector, signed, unsigned
x <= y; -- illegal assignment, type conflict
y <= x; -- illegal assignment, type conflict
x <= std_logic_vector(y); -- legal assignment
y <= unsigned(x); -- legal assignment
conversion between signed, unsigned, integer
z <= to_integer(y); -- legal assignment
z <= to_integer(unsigned(x)); -- legal assignment
y <= to_unsigned(z, 4); -- legal assignment
x <= std_logic_vector(to_unsigned(z, 4)); -- legal assignment</pre>
```

Chapter 7: Modeling at the FSMD Level

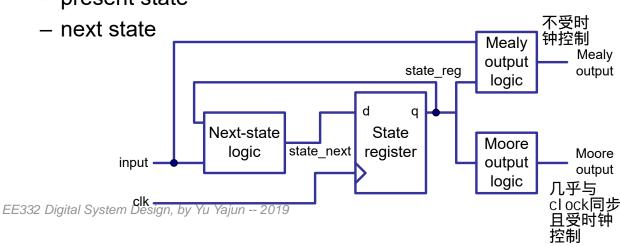
 A digital design is conceptually divided into two parts – a controller and a datapath.



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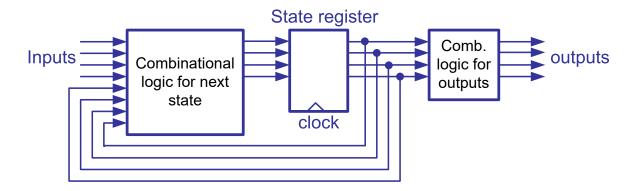
- A sequential circuit which is implemented in a fixed number of possible states is called a finite state machine (FSM).
- It contains five elements:
 - symbolic state
 - input signal
 - output signal
 - present state

- Two types of FSM:
 - Moore machines
 - Mealy machines



7.1 Moore machine

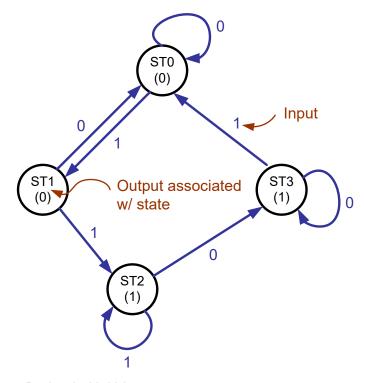
• In the Moore modal of sequential circuits, the outputs are the functions of the present state only.



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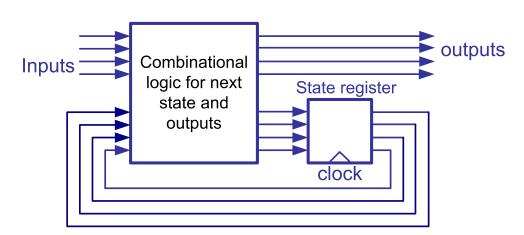
A state transition diagram of a Moore machine



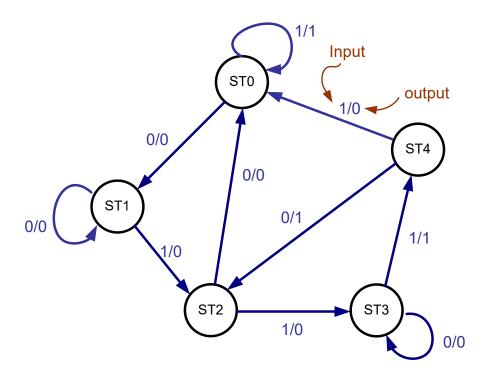
```
-- Moore machine
                                          case State is
                                            when ST0 =>
entity MOORE is
                                              O <= '0' :
port (Clk, RST, I: in std logic;
                                              if ( I ='0') then Next State <= ST0;</pre>
     O: out std logic);
                                              else
                                                            Next State <= ST1;
end entity MOORE;
                                              end if;
architecture two_seg_arch of MOORE is when ST1 =>
type state_type is (ST0, ST1, ST2, ST3);
                                              O <= '0' :
signal State, Next State: state type;
                                             if ( I = '0' ) then Next state <= ST0;</pre>
begin
                                              else
                                                              Next State <= ST2;
 clk proc: process (CLK, RST) is
                                              end if:
 begin
                                            when ST2 =>
      if (RST = '1') then
                                              O <= '1';
         State <= ST0:
                                              if ( I ='0') then Next State <= ST3;
     elsif (Clk'event and Clk = '1') then
                                                            Next State <= ST2;
                                              else
         State <= Next State:
                                              end if;
     end if;
                                            when ST3 =>
 end process clk proc;
                                              O <= '1';
 comb proc: process (State, I) is
                                             if ( I = '0' ) then Next state <= ST3;</pre>
                                                              Next State <= ST0;
 begin
                                              end if:
                                          end case; end process comb proc;
                                          end architecture two seg arch;
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                                                                                  217
```

7.2 Mealy machine

 In the Mealy modal, the outputs are the functions of both the present state and current inputs.



A state transition diagram of a Mealy machine



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-- two segments coding style

```
library IEEE;
use IEEE.STD_LOGIC_1164.all
entity MEALY is
port (Clk, RST, I: in std logic;
     O: out std logic);
end entity MEALY;
architecture two_seg_arch of MEALY is
type state type is (ST0, ST1, ST2, ST3, ST4);
signal State, Next_State : state_type;
begin
  clk_proc: process (CLK, RST) is
  begin
      if (RST = '1') then
         State <= ST0;
     elsif (Clk'event and Clk = '1') then
         State <= Next State;
     end if;
  end process clk_proc;
```

```
comb proc: process (State, I) is
                                     when ST2 =>
begin
                                       if ( I ='0') then O <= '0'; Next State <= ST0;
case State is
                                       else O <= '0'; Next State <= ST3;
 when ST0 =>
                                       end if;
    if ( I = '0') then
                                     when ST3 =>
        O <= '0';
                                       if ( I = '0') then O <= '0'; Next State <= ST3;
        Next State <= ST1;
                                       else O <= '1'; Next State <= ST4;
    else
                                       end if:
        O <= '1':
                                     when ST4 =>
        Next State <= ST0;
                                       if ( I ='0') then O <= '1'; Next State <= ST2;
    end if;
                                       else O <= '0'; Next State <= ST0;
 when ST1 =>
                                       end if:
    if ( | ='0') then
                                   end case:
        O <= '0';
                                   end process comb proc;
        Next State <= ST1;
                                   end architecture two seg arch;
    else
        O <= '0';
        Next State <= ST2;
    end if;
```

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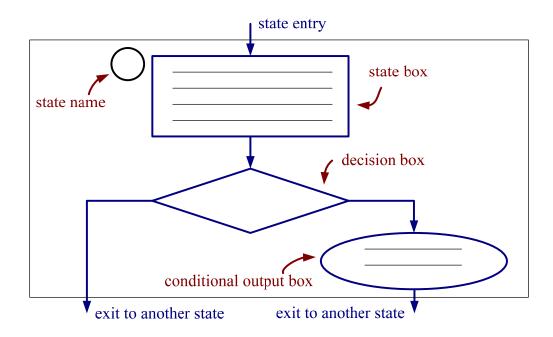
7.3 An FSM with a datapath (FSMD)

- A traditional FSM
 - cannot represent storage elements (register) except the state registers.
 - works well for a design with a few to several hundred states.
- An FSM with a datapath (FSMD) is an extension of a traditional FSM.
 - storage and signals can be declared.
 - Within a state expression, comparison, arithmetic or logic operations on these signals can be performed.

Algorithm state machine (ASM) chart

- The behavior of a FSMD can be represented as a flowchart-like description – algorithm state machine (ASM) chart.
- ASM chart is constructed from ASM blocks;
- An ASM block consists of three basic elements:
 - the state box
 - the decision box
 - the conditional output box.

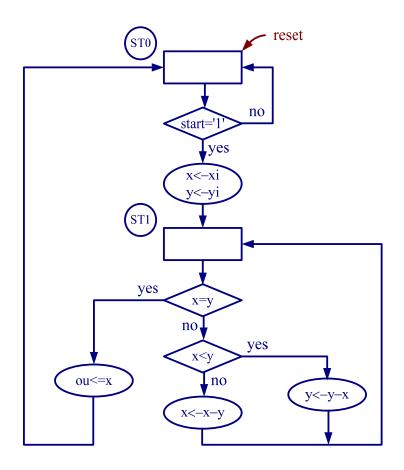
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Example:

find the greatest common divisor of two eight-bit numbers xi and yi

```
x = xi;
y= yi;
St1: If x=y then
ou=x;
Else {
if x> y then
x = x-y;
Else y= y-x;
Go to st1;
}
```



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-- GCD calculator

```
library IEEE;
use IEEE.STD_LOGIC_1164.ALL;
use IEEE.STD_LOGIC_ARITH.ALL;
use IEEE.STD_LOGIC_UNSIGNED.ALL;
entity GCD is
    Port ( xi, yi : in std_logic_vector(7 downto 0);
        clk, reset, start : in std_logic;
        ou : out std_logic_vector (7 downto 0));
end entity GCD;
architecture FSMD of GCD is
    signal x, y, x_next, y_next: std_logic_vector(7 downto 0);
    type S_Type is (ST0, ST1);
    signal state, next state : S Type;
```

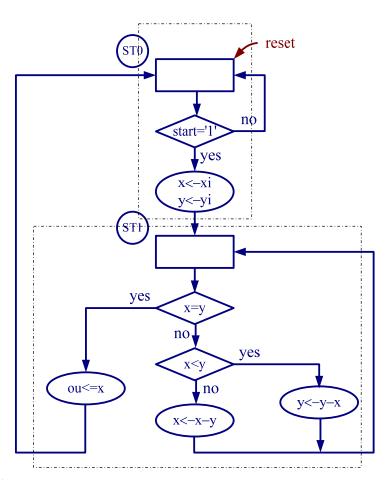
```
begin
                                                  case state is
 clkproc: process (clk, reset) is
                                                    when ST0 =>
   begin
                                                         if (start ='1') then
      if (reset = '1') then
                                                                                   整的if
                                                             x next <= xi;
         state <= ST0;
                                                             y next <= yi;
         x <= 0;
                                                             next state <= ST1;</pre>
         v <= 0;
                                                                                  角有缺
                                                         end if: -- (start = '1')
      elsif (clk'event and clk = '1') then
                                                    when ST1 =>
         state <= next state;
                                                         if (x = y) then
         x \le x \text{ next};
                                                             ou \le x;
         y \le y_next;
                                                         else
      end if;
                                                             if (x < y) then
   end process clkproc;
                                                                  y_next \le y - x;
 operproc: process (state, xi, yi, x, y, start) is
                                                             else -- (x > y)
                                                                  x next \le x - y;
   ou <= (others => '0'); 缺省值
                                                             end if:
   x next \le x;
                                                             next state <= ST1;</pre>
   y next \le y;
                                                         end if;
   next_state <= ST0;</pre>
                                                  end case; -- State
                                               end process operproc;
                                               end architecture FSMD;
```

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Example:

find the greatest common divisor of two eight-bit numbers xi and yi



Rules to Construct ASM Chart:

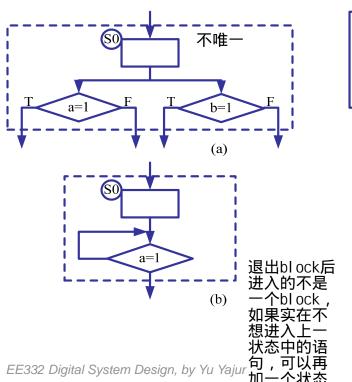
- For a given input combination, there is one unique exit path from the current AMS block.
- The exit path of an ASM block must always lead to a state box. The state box can be the state box of the current ASM block or a state box of another ASM block.

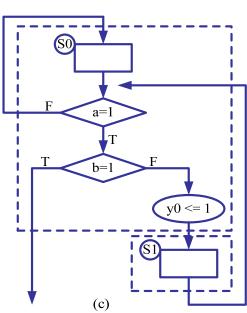
不能退出状态框 后直接进入判断 等组合逻辑,必 须还是进入一个 状态框

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Common errors in ASM Chart Construction





Example: FSMD design of a repetitive-addition multiplier

 Consider a multiplier with a_in and b_in, and with output r_out. The repetitive-addition algorithm can be formalized in the following pseudo-code:

```
if (a_in =0 or b_in =0) then{
    r = 0;}
    else{
        a = a_in; n = b_in; r = 0;
    op:     r = r + a;
        n = n - 1;
        if (n = 0) then {goto stop;}
        else {goto op;}
    }
stop:    r_out = r;
```

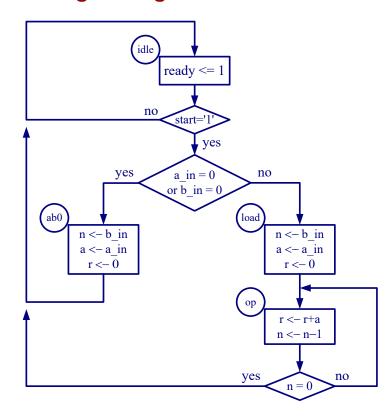
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Step 1: Defining the input and output signals

- Input signals:
 - a_in and b_in: input operands. 8-bit signals with std_logic_vector data type and interpreted as unsigned integers
 - start: command. The multiplier starts operation when the start signal is activated.
 - clk: system clock;
 - reset: asynchronous reset signal for system initialization.
- Output signals
 - r_out: the product. 16-bit signals.
 - ready: external status signal. It is asserted when the multiplication circuit is idle and ready to accept new inputs.

Step 2: Converting the algorithm to an ASM chart



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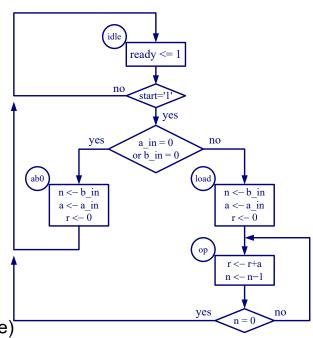
Step 3: Constructing the FSMD

- Basic data path can be constructed as follows:
 - 1. List all possible RT operations in the ASM chart.
 - 2. Group RT operations according to their destination registers.
 - 3. Derive the circuit for each group RT operation.
 - 4. Add the necessary circuits to generate the status signals.
- 3.1 The circuit require 3 registers, to store signals r, n, and a respectively.

3.2. The RT operations:

这里不用写reset, 因为reset是独立于 RT的一个异步时钟操 作

- RT operation with the r register:
 - r <- r (in the idle state)
 - r <- 0 (in the load and ab0 state)
 - r <- r + a (in the op state)
- RT operation with the n register:
 - n <- n (in the idle state)
 - n <- b_in (in the load and ab0 state)
 - n <− n − 1 (in the op state)
- RT operation with the a register:
 - a <- a (in the idle and op state)
 - a <- a_in (in the load and ab0 state)

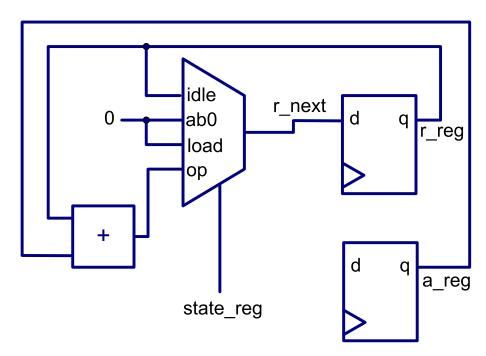


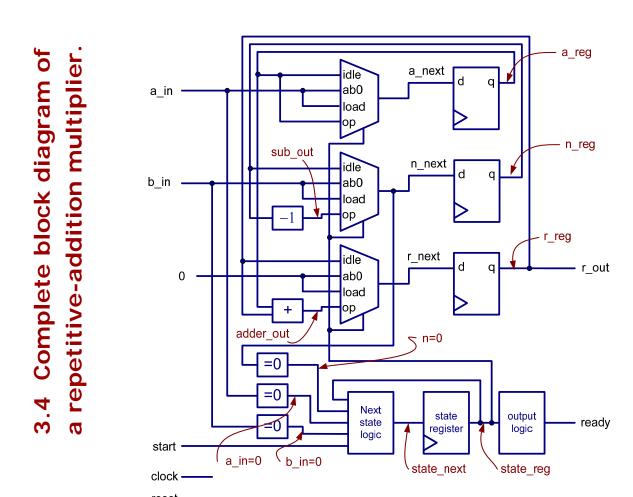
可以同时有两种类型的 状态机在一个程序中, ready是moore,r是 mealy

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3.3 the conceptual diagram of the circuit associated with the r register





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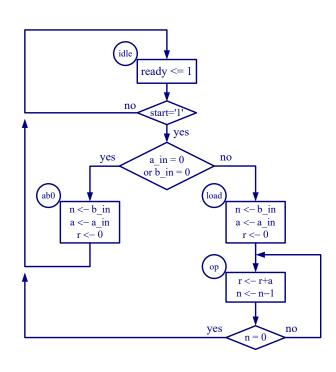
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Step 4: VHDL descriptions of FSMD

```
library IEEE;
use ieee.std logic 1164.all;
use ieee.std logic arith.all;
use ieee.std_logic_unsigned.all;
entity seq mult is
port (CLK, RESET, start: in std logic;
    a in, b in: in std logic vector(7 downto 0);
    ready: out std logic;
    r: out std_logic_vector(15 downto 0));
end entity seq mult;
architecture seg arch of seg mult is
constant WIDTH: integer:=8;
type state_type is (idle, ab0, load, op);
signal state reg, state next : state type;
signal a reg, a next, n reg, n next: std logic vector (WIDTH-1 downto 0);
signal r reg, r next : std logic vector (2*WIDTH-1 downto 0);
```

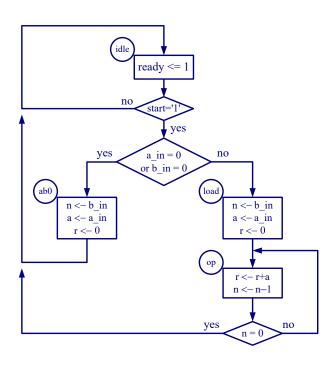
begin

```
-- state and data registers
process (CLK, RESET) is
begin
if RESET = '1' then
    state_reg <= idle;
    a_reg <= "00000000";
    n_reg <="00000000";
    r_reg <= x" 0000";
elsif CLK'event and CLK='1' then
    state_reg <= state_next;
    a_reg <= a_next;
    n_reg <= n_next;
    r_reg <= r_next;
end if;
end process;</pre>
```

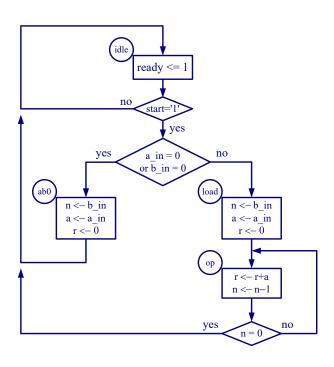


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```
-- combinational circuit
process (start, state_reg, a_reg, n_reg,
  r_reg, a_in, b_in, n_next) is
begin
    -- default value
    a next <= a reg;
    n next <= n reg;
    r_next <= r_reg;
    ready <= '0';
    case state reg is
       when idle =>
           if start = '1' then
             if (a in = "00000000" or
                b in = "00000000") then
                 state next <= ab0;
              else
                 state next <= load;
             end if:
          else
              state next <= idle;
          end if;
          ready <= '1';
```



```
when ab0 =>
           a next <= a in;
           n next <= b in;
           r = x"0000";
           state next <= idle;
        when load =>
           a next <= a in;
           n next <= b in;
           r \text{ next} \le x"0000":
           state next <= op;
        when op =>
           n \text{ next} \leq n \text{ reg} - 1;
           r next <= ("00000000" & a reg)
                      + r reg;
           if (n next = "00000000") then
              state_next <= idle;
           else
              state next <= op;
           end if:
       end case;
    end process;
    r \le r \text{ reg};
end architecture seg_arch;
```

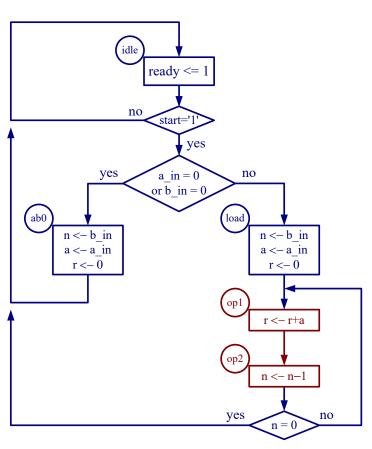


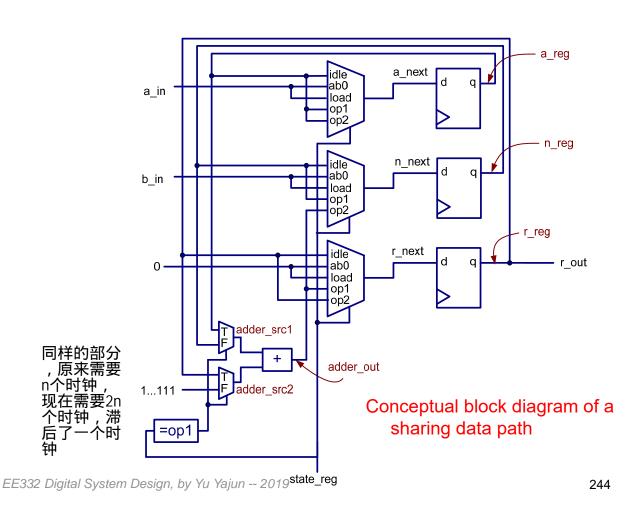
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Resource sharing via FSMD example of repetitive-addition multiplier

- Many RT operations perform the same or similar function.
- Some function unit can be shared as long as these operations are scheduled in different states.
- the 16-bit adder and 8-bit decrementor are shared in the following example.







sharing on a repetitive-addition multiplier

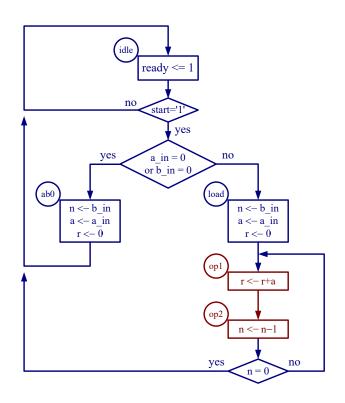
```
architecture sharing_arch of seq_mult is
constant WIDTH : integer :=8;
type state_type is (idle, ab0, load, op1, op2);
signal state_reg, state_next : state_type;
signal a_reg, a_next, n_reg, n_next : std_logic_vector (WIDTH-1 downto 0);
signal r_reg, r_next : std_logic_vector (2*WIDTH-1 downto 0);
signal adder_scr1, adder_scr2: std_logic_vector (2*WIDTH-1 downto 0);
signal adder_out: std_logic_vector (2*WIDTH-1 downto 0);
```

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begin

```
-- state and data registers
process (CLK, RESET) is
begin
if RESET = '1' then
    state_reg <= idle;
    a_reg <= "00000000";
    n_reg <="00000000";
    r_reg <= x" 0000";
elsif CLK'event and CLK='1' then
    state_reg <= state_next;
    a_reg <= a_next;
    n_reg <= a_next;
    r_reg <= a_next;
end if;
end process;</pre>
```



```
-- next-state, logic/output logic and data path routing
process (start, state_reg, a_reg, n_reg, r_reg, a_in, b_in, adder_out, n_next) is
begin
     -- default value
                                                                   idle
     a next <= a reg;
                                                                       ready <= 1
     n next <= n reg;
     r next <= r reg;
                                                                         start=
     ready <= '0';
                                                                             yes
     case state_reg is
         when idle =>
                                                                yes
                                                                                    no
                                                                        a in = 0
                                                                        or \overline{b} in = 0
            if start = '1' then
                                                       ab0
                if (a in = "00000000" or
                                                           n < -b_{in}
                                                                                       n < -b_{in}
                   b in = "00000000") then
                                                           a < -a in
                                                                                       a < -a in
                                                            r < -0
                                                                                        r < -\overline{0}
                   state next <= ab0;
                else
                   state next <= load;
                                                                                       r < -r+a
                end if;
            else
                                                                                       n < -n-1
                state next <= idle;
            end if;
            ready <= '1';
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                                                                                                247
```

```
when ab0 =>
                                                                idle
       a next <= a in;
                                                                    ready <= 1
       n next <= b in;
       r next \leq x"0000";
                                                                      start=
       state next <= idle;
                                                                         yes
   when load =>
                                                              yes
                                                                                no
                                                                     a in = 0
       a next <= a in;
                                                                    or \overline{b}_i =
       n_next <= b_in;
                                                     ab0
       r next \leq x"0000";
                                                         n \leftarrow b_i
                                                                                   n \leftarrow b_in
       state next <= op1;
                                                         a \le -a in
                                                                                   a \le -a in
                                                          r < -0
                                                                                    r < -0
   when op1 =>
       r next <= adder out;
       state next <= op2;
                                                                                   r < -r + a
   when op2 =>
       n next = adder out (WIDTH -1 downto 0);
                                                                                   n < -n-1
       if (n next = "00000000") then
          state_next <= idle;
       else
          state next <= op1;
       end if:
   end case;
                     两段式结构不够用了,需要再
                    加一个第三段
end process;
```

```
-- datapath input routing and functional units
process (state_reg, r_reg, a_reg, n_reg) is
begin
    if (state_reg = op1) then
        adder_src1 <= r_reg;
        adder_src2 <= "00000000" & a_reg;
    else -- for op2 state
        adder_src1 <= "00000000" & n_reg;
        adder_src2 <= x"FFFF";
    end if;
end process;
adder_out <= adder_src1 + adder_src2;
-- output
r <= r_reg;
end architecture sharing arch;</pre>
```

addder_out更新后 ,上面的寄存器 process被重新激活 ,再op1,2state的 时候更新这些值

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Chapter 8: Parameterized Design

- Goal: Design reuse
 - Ideally, we want to design some common modules that can be shared by many applications.
 - Since every application is different, it is desirable that a module can be customized to some degree to meet the specific need of an application.
 - Customization is normally specified by explicit or implicit parameters

Types of Parameters

Width Parameters

- The widths of data signals normally can be modified to meet different requirement.
- The width parameters of a parameterized design specify the sizes (i.e., number of bits) of the relevant data signal.

Feature Parameters

- Specify the structure or organization of a design.
- Defined on an ad hoc basis. 根据实际功能设计参数
- To include or exclude certain functionalities (i.e., features) from implementation or to select one particular version of the implementation

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8.1 Generics

- The generic construct of VHDL is a mechanism to pass information into an entity and a component.
 - They are first declared in entity and component declaration and later assigned a value during component instantiation

```
entity para_binary_counter is

generic (WIDTH: natural);

port (
    clk, reset: in std_logic;
    q: out std_logic_vector(WIDTH-1 downto 0)
    );
end entity para_binary_counter;
```

- After the declaration, the generic can be used in the associated architecture bodies.
- A generic cannot be modified inside the architecture body and thus functions like a constant
 - It is sometimes referred to as a generic constant.

```
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```

```
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```

```
architecture arch of para_binary_counter is
    signal reg, reg_next : std_logic_vector (WIDTH-1
        downto 0);
begin
    process (clk, reset) is
    begin
        if reset = '1' then
            reg <= (others =>'0');
        elsif clk'event and clk='1' then
            reg <= reg next;
        end if;
    end process;
-- next-state logic
    reg next <= reg + 1;
    q <= std logic vector(reg);</pre>
end architecture arch;
```

用在多层次系统中

- To use the parameterized free-running binary counter in a hierarchical design, a similar component declaration should be included in the architecture declaration.
- The generic can then be assigned a value in the generic mapping section when a component instance is instantiated.
- Example of the use of generics

```
这个电路只能被调用和实例化,没法直接单独使用,需要参数才行
```

```
library ieee;

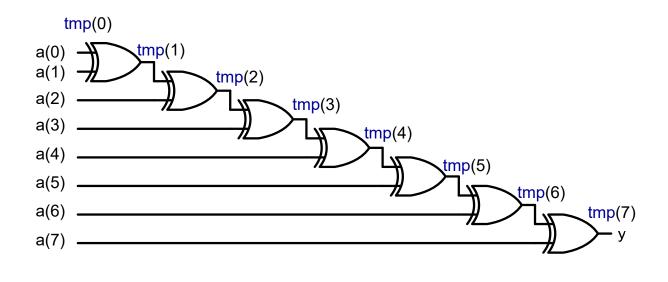
use ieee.std_logic_1164.all;
entity generic_demo is

port(
    clk, reset: in std_logic;
    q_4: out std_logic_vector(3 downto 0);
    q_12: out std_logic_vector(11 downto 0)
    );
end entity generic_demo;
```

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```
architecture arch of generic demo is
    component para binary counter is
       generic (WIDTH: natural);
       port (
           clk, reset: in std logic;
           q: out std logic vector(WIDTH-1 downto 0)
       );
    end component para binary counter;
begin
   four bit: para binary counter
       generic map (WIDTH => 4)
       port map (clk => clk, reset => reset, q => q 4);
   twelve bit: para binary count
       generic map (WIDTH => 12)
       port map (clk => clk, reset => reset, q => q 12);
end architecture arch:
```

Example: Reduced-xor circuit



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不是组合逻辑

 $tmp(i) \le tmp(i-1) xor a(i)$

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```
library ieee; use ieee.std logic 1164.all;
  Parameterized
                       entity reduced xor is
  reduced-xor
                           generic (WIDTH: natural); -- generic declaration
 circuit using a
                           port(
 generic
                               a: in std_logic_vector(WIDTH-1 downto 0);
                               y: out std logic
                       end entity reduced xor;
                       architecture loop linear arch of reduced xor is
                           signal tmp: std logic vector(WIDTH-1 downto 0);
                                                              第一次进入process将
                       begin
                                                              tmpO的值赋值,在delta
tmp(0)
                           process (a, tmp) is
                                                              后它得到值,但是因为
a(0) ک
     tmp(1)
a(1)
                           begin
                                                              tmp(i -1)i 等于1的时候
         <u>tmp(2)</u>
a(2)
             tmp(3)
                                                              tmp(0)的值还没被更新
                               tmp(0) \le a(0); -- boundary bit
a(3)
                                                               后面Loop算的一直是
a(4)
                               for i in 1 to (WIDTH-1) loop
                                                              错的,只有tmpO正确
a(5)
                                                             第二次进入process
a(6)
                                   tmp(i) \le a(i) xor tmp(i-1);
a(7)
                                                              tmp(1)的值被更新正确
                               end loop;
                                                                其他的temp仍旧错
                           end process;
                                                                   -轮轮更新直到正
                                                                ,因为是组合逻辑
                           y \le tmp(WIDTH-1);
                                                              这样实际被综合后还
     Slide 293
                       end architecture loop linear arch;
                                                                当于一步步xor,
                                                              我们参数化设计需要
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                                                              样Ioop写,不能直接写成xor因为不知道长度,
```

Array attribute更简便,不需要 8.2 Array attribute generic,但是更容易出错,不建议使用

- A VHDL attribute provides information about a named item, such as a data type or a signal.
- We have used the 'event attribute, as in clk'event, express the changing edge of the clk signal.
- A set of attributes is associated with an object of an array data type. Let s be a signal with an array data type.
 - s'left, s'right: the left and right bounds of the index range of s.
 - s'low, s'high: the lower and upper bounds of the index range of s.
 - s'length: the length of the index range of s.
 - s'range: the index range of s.
 - s'reverse_range: the reversed index range of s.

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 The attributes can be applied to the signal defined with std_logic_vector, unsigned and signed:

```
signal s1: std_logic_vector (31 downto 0);
signal s2: std_logic_vector (8 to 15);
```

The attributes of s1 are

```
s1'left = 31; s1'right = 0;
s1'low = 0; s1'high =31;
s1'length = 32;
s1'range = 31 downto 0
s1'reverse_range = 0 to 31
```

The attributes of s2 are

```
reduced-xor
circuit using
an attribute
```

```
Parameterized architecture loop_linear_arch of reduced xor is
                        signal tmp: std logic vector(a'length-1 downto 0);
                    begin
                        process (a, tmp) is
                        begin
                            tmp(0) \le a(0);
                            for i in 1 to (a'length-1) loop
                                 tmp(i) \le a(i) xor tmp(i-1);
                            end loop;
                        end process;
                        y <= tmp<mark>(a'length</mark>-1);
                    end architecture loop linear arch;
```

The range of the for loop can also be expressed as:

```
for i in a'low+1 to a'high loop
for i in a'right+1 to a'left loop
```

可能会出问 题这样+1

The last signal assignment

```
- v <= tmp (tmp'left);</p>
```

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8.3 Unconstrained Array

- Unconstrained array is defined as an array type with specified data type of the index value, but without specified exact bounds of the index value.
- Example:

```
type std logic vector is array (natural range <>) of
std logic
```

- Similarly, we have unsigned and signed data types.
- If an object is declared with an unconstrained array data type, we must specify its index range when the data type is used, as 15 downto 0 in

```
signal x: std logic vector(15 downto 0);
```

- A special case: the unconstrained array can be declared without specifying the range in port declaration.
- Example:

```
library ieee; use ieee.std logic 1164.all;
entity unconstrain dff is
    port( clk: std logic;
         d: in std logic vector;
                                   -- the actual range is inferred
         q: out std_logic vector -- when an instance of
                                     -- unconstrain dff is instantiated.
    );
end entity unconstrain dff;
architecture arch of unconstrain dff is
begin
    process (clk) is
    begin
       if (clk'event and clk='1') then q <= d; end if;
    end process:
end architecture arch;
```

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The ranges of the actual signals become the ranges of d and g signals. Example: the dff16 instance is instantiated as a 16-bit register

```
signal din, gout: std logic vector(15 downto 0);
signal clk: std logic;
dff16: unconstrain dff
    port map( clk => clk, d => din, q => gout);
```

 Since no range is specified for d and q, the boundaries of the two signal will not be check in the analysis stage.

```
signal din: std logic vector(15 downto 0);
signal gout: std logic(7 downto 0); -- syntactically correct.
               -- error may be detected during the synthesis
dff16: unconstrain dff
    port map( clk => clk, d => din, q => qout);
```

```
Parameterized library ieee; use ieee.std logic 1164.all;
                   entity unconstrain reduced xor is
reduced-xor
                      port(
circuit using
                          a: in std logic vector;
an
                          y: out std logic
unconstrained
                   end entity unconstrain reduced xor;
array
                   architecture arch of unconstrain reduced xor is
                      constant WIDTH: natural := a'length; 这块只是一个普通常数
The code
                      signal tmp: std logic vector(WIDTH-1 downto 0);
appears to
                   begin
                      process (a, tmp) is
be correct at
                      begin
first glance
                          tmp(0) \le a(0);
                          for i in 1 to (WIDTH-1) loop
                              tmp(i) \le a(i) xor tmp(i-1);
                          end loop;
                      end process;
                      y \le tmp(WIDTH-1);
                   end architecture arch:
```

 If we map the a signal to an actual signal with the type of std_logic_vector of 8 bits during component instantiation, we may have a to be:

```
std_logic_vector(7 downto 0);
std_logic_vector(0 to 7);
std_logic_vector(15 downto 8);
std_logic_vector(8 to 15);
```

The code does not work properly for the last two formats.

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Improved parameterized reduced-xor circuit using an unconstrained array

```
architecture better_arch of unconstrain_reduced_xor is
    constant WIDTH: natural := a'length;
    signal tmp: std_logic_vector(WIDTH-1 downto 0);
    signal aa: std_logic_vector(WIDTH-1 downto 0);

begin
    aa <= a;
    process (aa, tmp) is
    begin
        tmp(0) <= aa(0);
        for i in 1 to (WIDTH-1) loop
            tmp(i) <= aa(i) xor tmp(i-1);
        end loop;
    end process;
    y <= tmp(WIDTH-1);
end architecture better_arch;</pre>
```

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8.3 Comparison

- The unconstrained array mechanism uses attributes to infer the relevant information from the actual signal.
 - More general and flexible than the generic mechanism, but also
 更灵活更易错
 - More opportunities for errors.
 - Requires comprehensive error-checking code
- Generic mechanism is preferred, unless a module is extremely general and widely used.
 - More rigid
 - It clearly specifies the range, direction and width of each signal and avoids many subtle erroneous conditions.

8.4 Generate Statement

- The generate statements are concurrent statements with embedded internal concurrent statement, which can be interpreted as a circuit part.
- Two types of generated statements:
 - for generate statement: used to create a circuit by replicating the hardware part
 - conditional or if generate statement: used to specify whether or not to create an optional hardware part.

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8.4.1 For Generate Statement

- Many digital circuits can be implemented as a repetitive composition of basic building blocks, exhibiting a regular structure, such as a one-dimensional cascading chain, a tree-shaped connection or a two-dimensional mesh.
- For generate statement syntax

```
gen_label: -- mandatory to identify to this -- particular generate statement
```

for loop_index in loop_range generate
 concurrent statements;

-- describe a stage of the iterative circuit

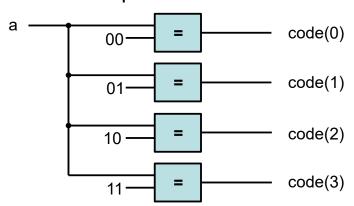
end generate;

 The loop_range has to be static. It is normally specified by the width parameters.

相当于把很多普通的并行语句,因为基本相似,直接通过I oop写 $_{FE35}$ 在一起了,本质还是并行语句

Example: Binary decoder

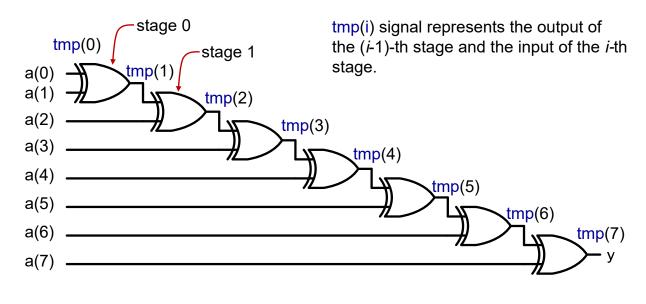
- A binary n-to- 2^n decoder is circuit that asserts one of the 2^n possible output signal according to an n bit input signal.
- One way to view the binary decoder is to treat each bit of the decoded output as the result of a constant comparator.



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```
Parameterized library ieee;
binary
                  use ieee.std logic 1164.all;
decoder using use ieee.numeric_std.all;
a for generate entity bin decoder is
                      generic(WIDTH: natural);
statement
                      port(
                         a: in std logic vector(WIDTH-1 downto 0);
                         code: out std logic vector(2**WIDTH-1 downto 0)
                  end bin decoder;
                  architecture gen arch of bin decoder is
                  begin
                      comp gen: 标签
                      for i in 0 to (2**WIDTH-1) generate
                         code(i) <= '1' when i = to integer(unsigned(a)) else
                                   '0';
                      end generate;
  Slide 293
                  end architecture gen arch;
```

Example: Reduced-xor circuit



 $tmp(i+1) \le tmp(i) xor a(i+1)$

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```
Parameterized
                     architecture gen linear arch of reduced xor is
 reduced-xor
                         signal tmp: std logic vector(WIDTH-1 downto 0);
 circuit using a
                     begin
for generate
                         tmp(0) \le a(0);
                         xor gen:
 statement
                             for i in 1 to (WIDTH-1) generate
                                  tmp(i) \le a(i) xor tmp(i-1);
                             end generate;
                          y \le tmp(WIDTH-1);
                     end architecture gen linear arch;
        -stage 0
  tmp(0)
             stage 1
a(0)
a(1)
             tmp(2)
a(2)
                  tmp(3)
a(3)
                       tmp(4)
a(4)
a(5)
                                                          Slide 293
a(6)
a(7)
```

 In an iterative structure, the boundary stages interface to the external input and output signals, and sometimes their connections are different from the regular blocks.

8.4.2 Conditional Generate Statement

- The conditional generate statement is used to specify an optional circuit that can be included or excluded in the final implementation.
- Conditional generate statement syntax

```
gen_label: -- mandatory
if boolean_exp generate -- boolean_exp must be static
    concurrent statements;
end generate;
```

- There is no else branch in conditional generate statement.
- If we want to include one of the two possible circuits in an implementation, we must use two separate if generate statements.

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Reduced-xor circuit revisited

- One common use of the conditional generate statement is to describe the "irregular" stages in a for generate statement.
- For example, two statements

```
tmp(0) <= a(0);
y <= tmp(WIDTH-1);
```

are used to rename the input and output signals in the for generate statement examples.

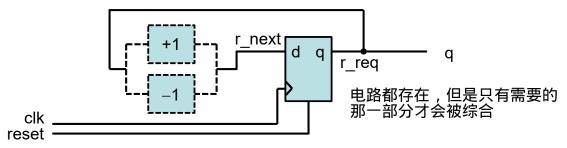
 To eliminate these statements, we can use conditional generate statements inside the for generate statement.

```
Parameterized
                   architecture gen if arch of reduced xor is
 reduced-xor
                       signal tmp: std logic_vector(WIDTH-2 downto 1);
 circuit with a
                   begin
 conditional
                       xor gen:
                       for i in 1 to (WIDTH-1) generate
 generate
                           -- leftmost stage
 statement
                          left gen: if i = 1 generate
                                                      就不需要特殊处理了
                               tmp(i) \le a(i) xor a(0);
                           end generate;
将所有情况都包含在for里面
                           -- middle stage
                          middle gen: if (i>1) and (i<(WIDTH -1)) generate
                               tmp(i) \le a(i) xor tmp(i-1);
                           end generate;
                           -- rightmost stage
                          right gen: if i = (WIDTH - 1) generate
                               y \le a(i) xor tmp(i-1);
                           end generate;
                       end generate;
                   end architecture gen if arch;
```

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Example: Up-or-down free-running binary counter

- An up-or-down binary counter is a counter that can be instantiated in a specific mode.
- Note that the "or" here means that only one mode of operation, either counting up or counting down but not both, can be implemented in the final circuit.



• We use the UP generic as the feature parameter to specify the desired mode.

```
Up-or-down
free-running
binary
counter
```

```
library ieee;
                 use ieee.std logic 1164.all;
                 use ieee.numeric std.all;
                 entity up or down counter is
                     generic(WIDTH: natural; UP: natural);
                     port(clk, reset: in std logic;
                              : out std logic vector(WIDTH-1 downto 0)
                 end up or down counter;
                 architecture arch of up or down counter is
                     signal r reg, r next: unsigned(WIDTH-1 downto 0);
                 begin
                     -- register
                     process (clk, reset)
                     begin
                         if (reset = '1') then
                              r reg <= (others => '0')
                         elsif (clk'event and clk='1') then
                              r reg <= r next;
                         end if;
EE332 Digital System Design, end, processo
                                                                          279
```

```
-- next-state logic
    inc gen: -- incrementor
    if UP = 1 generate
        r next <= r reg + 1;
    end generate;
    dec gen: -- decrementor
    if UP /= 1 generate
        r \text{ next} \leq r \text{ req} - 1;
    end generate;
    q <= std logic vector(r reg); -- output logic
end architecture arch;
```

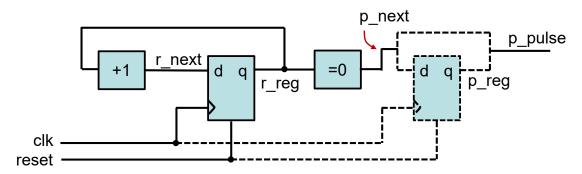
```
Up-and-down
                 library ieee;
free-running
                 use ieee.std logic_1164.all;
                 use ieee.numeric_std.all;
binary
                 entity up and down counter is
counter
                     generic map (WIDTH: natural)
                     port map(clk, reset: in std_logic; mode: in std_logic;
                             : out std logic vector(2**WIDTH-1 downto 0)
                 end up and down counterr;
```

```
architecture arch of up and down counter is
    signal r reg, r next: unsigned(WIDTH-1 downto 0);
begin
    -- register
    process (clk, reset)
    begin
        if (reset = '1') then
            r reg <= (others => '0')
        elsif (clk'event and clk='1') then
             r reg <= r next;
        end if;
    end process;
    -- next-state logic
    r next <= r reg + 1 when mode ='0' else
              r reg -1;
    -- output logic
    q <= std logic vector(r reg);</pre>
end architecture arch;
```

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Counter with an optional output buffer

- An output buffer can remove glitches from the signal.
- Since the buffer is only needed for certain application, it will be convenient to include the buffer as an optional part of the circuit.



Counter with an optional output buffer

```
library ieee;
use ieee.std_logic_1164.all;
use ieee.numeric std.all;
entity op buf counter is
    generic(WIDTH: natural; BUFF: natural);
    port(clk, reset: in std logic;
        pulse: out std logic);
end op buf counterr;
architecture arch of op buf counter is
    signal r reg, r next: unsigned(WIDTH-1 downto 0);
    signal p reg, p next: std logic;
begin
    -- register
    process (clk, reset)
    begin
        if (reset = '1') then r reg <= (others => '0')
        elsif (clk'event and clk='1') then r reg <= r next;
        end if:
    end process;
```

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```
-- next-state logic
    r next <= r reg + 1;
    -- output logic
    p next <= '1' when r reg = 0 else '0';
    buf gen:
                              gernerate也是并行语
    if BUFF = 1 generate
                               句 , 所以内部可以直
        process (clk, reset)
                              接写process
        begin
            if (reset = '1') then p reg <= '0'
            elsif (clk'event and clk = '1') then p reg <= p next;
            end if:
       end process;
        pulse <= p reg;
    end generate:
    no buf gen: -- without buffer,
                                                         p next
    if BUFF /= 1 generate
                                                                      p pulse
        pulse <= p next;
                                            d
                                             q
                                                r reg
                                                                  p reg
    end generate;
end architecture arch;
                            clk
                          reset
```

8.5 Comparison

- To create a circuit with a selectable feature:
 - use conditional generate statement
 - a full-featured circuit with some input control signal connected to constant values to permanently enable the desired feature
 - use the configuration construct structure结构化设计,做两个模块,一个up,一个down
- Assume we need a 16-bit up counter in a design.

Difference

- The up-or-down counter instance
 - creates a circuit with only the needed features.
 - The selected portion of code is passed to the synthesis stage, i.e., the synthesis software only needs to synthesize the selected portion.
- The up-and-down counter instance
 - creates a circuit that consists of all features and uses an external control signal to selectively enable a portion of the circuit.
 - The entire VHDL code is passed to synthesis stage. The synthesis software eliminates the unused portion through logic optimization.
- In general, use of the feature parameters and conditional generate statement is better than the full-featured approach.

- The selected hardware creation can also be achieved by configuration where multiple architecture bodies are constructed, each containing a specific feature, e.g., architectures up_arch and down_arch of the same entity updown_counter, for counting up and counting down, respectively.
- And the following instantiation can be used to select the counting up circuit.

```
count16up: work.updown_counter(up_arch)
generic map(WIDTH =>16)
port map (clk => clk, reset => reset, q => q);
```

```
Up-or-down counter with two architecture bodies
```

```
library ieee;
use ieee.std logic 1164.all;
use ieee.numeric std.all;
entity updown counter is
    generic(WIDTH: natural);
    port(clk, reset: in std logic;
        q: out std logic vector(WIDTH-1 downto 0)
end updown counter;
architecture up arch of updown counter is
    signal r reg, r next: unsigned(WIDTH-1 downto 0);
begin
    -- register
    process (clk, reset)
    begin
        if (reset = '1') then r reg <= (others => '0')
        elsif (clk'event and clk='1') then r reg <= r next;</pre>
        end if:
    end process;
```

```
-- next-state logic
                 r next <= r reg + 1;
                 -- output logic
                 q <= std logic vector(r reg);</pre>
            end architecture up arch;
            architecture down arch of updown counter is
                 signal r reg, r next: unsigned(WIDTH-1 downto 0);
            begin
                 -- register
                 process (clk, reset)
                 begin
                     if (reset = '1') then r reg <= (others => '0')
                     elsif (clk'event and clk='1') then r reg <= r next;</pre>
                     end if:
                 end process;
                 -- next-state logic
                 r \text{ next} \leq r \text{ req} - 1;
                 -- output logic
                 q <= std logic vector(r reg);</pre>
            end architecture down arch;
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```

- Conversely, we can merge the logic from several architecture bodies into a single body and use a feature generic and conditional generate conditions to select the desired portion.
- There is no rule about when to use a feature parameter and when to use a configuration construct. In general,
 - code with a feature parameter is more difficult to develop and comprehend, but on the other hand, if we use a separate architecture body for each distinctive feature, the number of architecture bodies will grow exponentially and becomes difficult to manage.
 - when a feature parameter leads to significant modification or addition of the no-feature codes and starts to make the code incomprehensible, it is probably a good idea to use separate architecture bodies and the configuration construct.

8.6 For Loop Statement

• The for loop statement is a sequential statement and is the only sequential loop construct that can be synthesized.

must be static

for index in loop_range loop
sequential statements;
for loop, while loop
, loop只有for loop能
综合

end loop;

 The basic way to synthesize a for loop statement is to unroll or flatten the loop. Unrolling a loop means to replace the loop structure by explicitly listing all iterations.

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Example: Binary decoder

The code is very similar to the for generate version Slide 273

```
architecture loop_arch of bin_decoder is
begin
    process (a)
    begin
        for i in 0 to (2**WIDTH-1) loop
            if i = to_integer(unsigned(a)) then code(i) <= '1';
            else code(i) <= '0';
            end if;
        end loop;
        end process;
end architecture gen_arch;</pre>
```

Example: Reduced-xor circuit

• For loop version: Slide 259 需要process

For generate version: Slide 275 不需要 process

Example: Priority Encoder

- Recall that a signal can be assigned with multiple times inside process and only the last assignment takes effect.
- A priority encoder is a circuit that returns the binary code of the highest-priority request.
- Assume that the input is an array of r(WIDTH-1 downto 0), and r(WIDTH-1) has the highest priority.

architecture loop_linear_arch of prio_encoder is

```
constant B: natural := log2c(WIDTH);
    signal tmp: std logic vector(WIDTH-1 downto 0);
begin
    process (r) --binary code
    begin
       bcode <= (others => '0');
                                    向高次循环,更高位的编
       for i in 0 to (WIDTH-1) loop 码会覆盖以前的
            if r(i) = '1' then
                bcode <= std logic vector(to unsigned(i, B));</pre>
            end if;
                                     从高位往下时要用到exi t语句
       end loop:
                                     ,这个语句无法综合
    end process;
    process(r, tmp) -- reduced - or circuit
    begin
       tmp(0) \le r(0);
       for i in 1 to (WIDTH – 1) loop
                                      全部做或
            tmp(i) \le r(i) \text{ or } tmp(i-1);
       end loop;
   end process:
   valid <= tmp (WIDTH-1);
end architecture loop linear arch;
```

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8.7 Comparison

- Both the for generate and for loop statements are used to describe replicated structures.
- For generate statement:
 - can only use concurrent statements.
 - start a design with a conceptual diagram of a few stages; the diagram is used to identify the basic building block and connection pattern, and then the code of the loop body is derived.
- For loop statement:
 - can only use sequential statements.
 - the body of the loop statement can be more general and versatile.
 - may lead to unnecessarily complex implementation or even an unsynthesizable description.

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Chapter 9: Pipelined Design 流水线设计

- Pipeline is an important technique to increase the performance of a system.
- The basic idea is to overlap the processing of several tasks so that more tasks can be completed in the same amount of time.

同样的时间里可以 完成更多的任务

9.1 Delay and throughput

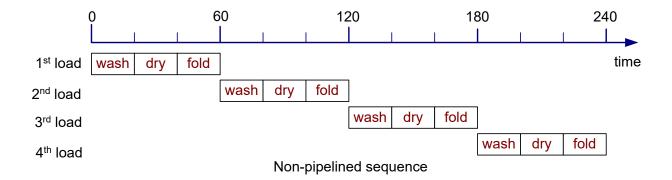
- Delay and throughput are the two criteria used to examine the performance of a system
 - Delay: the time required to complete one task.
 - Throughput: the number of tasks that can be 吞吐量 completed per unit time.
- Adding pipeline to a combinational circuit can increase the throughput but not reduce the delay.

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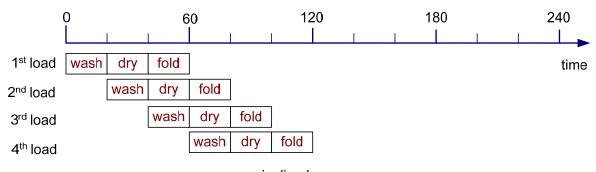
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9.2 Overview on pipelined design

- The pipelining technique can be applied to a task that is processed in stages.
 - An example: Pipelined laundry.
 - A complete laundry includes the stages of washing, drying and folding, for example.



- For non-pipelined process, a new load cannot start until the previous load is completed.
 - It takes 240 minutes to complete the four loads.
 - The delay of processing one load is 60 minutes.
 - The throughput is 1/60 load per minute.



pipelined sequence

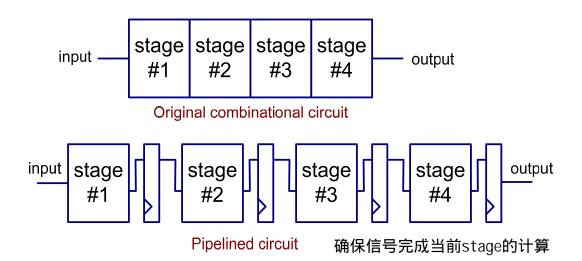
- For pipelined process,
 - It takes 120 minutes to complete the four loads.
 - The delay in processing one load **remains** 60 minutes.
 - The throughput increases to 2/60 load per minute.
 - To process k loads, it will take 40+20k minutes.
 - The throughput becomes k/(40+20k) load per minute -> 1/20 load per minute when k is large.

Pipelined combinational circuit

- A combinational circuit can be divided into stages so that the processing of different tasks can be overlapped.
- To ensure that the signals in each stage flow in the correct order and prevent any potential race, registers must be added between successive stages.
- The registers ensures that the signals can be passed to the next stage only at predetermined points.

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Assume: propagation delay for each stage: T1, T2, T3 and T4.

$$T_{max} = max(T1, T2, T3, T4);$$

Thus, the minimum clock period has to accommodate the longest delay plus the overhead introduced by the buffer register in each stage:

$$T_c = T_{max} + T_r$$
;

The effectiveness of the circuit

- Propagation delay:
 - non-pipelined circuit: $T_{comb} = T1+T2+T3+T4$
 - pipelined circuit: $T_{pipe} = 4T_c = 4T_{max} + 4T_r$
- Throughput:

延时增大

- non-pipelined circuit: 1/T_{comb}
- pipelined circuit: $k/(3T_c+kT_c) \rightarrow 1/T_c$.

k/((N-1)Tc+kTc) N:级数

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Ideally, for an N-stage circuit

- The propagation delay of each stage is identical (i.e., $T_{\text{max}} = T_{\text{comb}}/N$
- The register overhead (T_r) is comparably small

但是实际情况一般每

- I_{pipe} - $IVI_c = IVI_{max} = T_{comb}$ - Throughput: $1/T_c = 1/T_{max} = N/T_{comb}$

个stage delay都不一样且Tr不能忽略

- Ideally, it is desirable to have more stages in the pipeline. However, when N becomes large,
 - the propagation delay of each stage becomes smaller, but Tr remains the same; its effect cannot be ignored.
 - In reality, it is difficult to keep dividing the original combinational circuit into smaller and smaller stages.

9.3 Adding pipeline to a combinational circuit

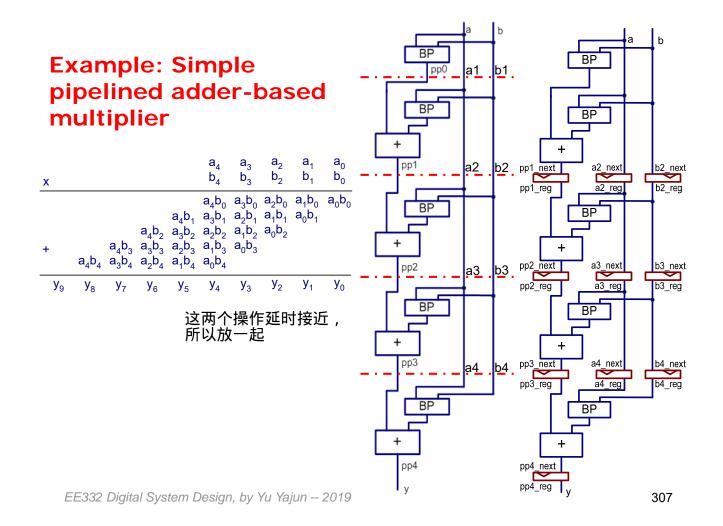
- The candidate circuits for effective pipeline design should include the following characteristics:
 - There is enough input data to feed the pipeline circuit.
 - The throughput is the main performance criterion. 主要希望改善的指标
 - The combinational circuit can be divided into stages with similar propagation delay.
 - The propagation delay of a stage is much longer than the delay incurred due to the register.

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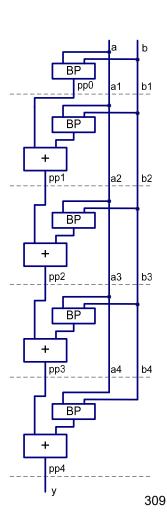
The procedure to convert a combinational circuit to a pipelined design

- Derive the block diagram of the original combinational circuit and arrange it as a cascading chain.
- Identify the major components and estimate the relative propagation delays of these components.
- Divide the chain into stages of similar propagation delays.
- identify the signals that cross the boundary of the chain.
- Insert registers for these signals in the boundary.

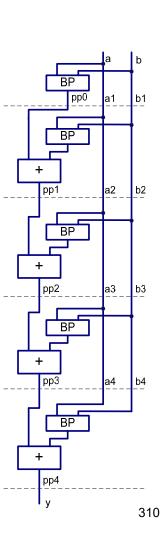


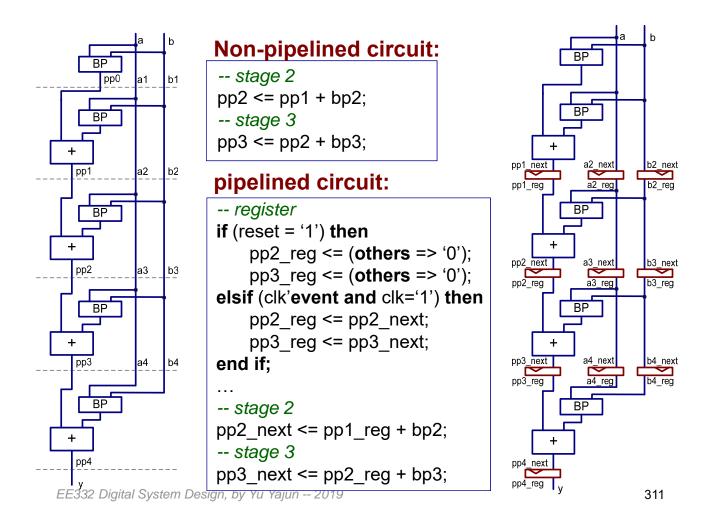
```
Non-pipelined multiplier in cascading stages
   library IEEE;
                                                                    RP
   use ieee.std logic 1164.all;
   use ieee.std logic arith.all;
                                                                           b2
   use ieee.std logic unsigned.all;
   entity mult5 is
   port (clk, reset : in std logic;
                                                                   pp2
         a, b: in std logic vector(4 downto 0);
                                                                    BP
         y: out std logic vector(9 downto 0));
   end entity mult5;
   architecture comb arch of mult5 is
   constant WIDTH: integer := 5;
   signal a1, a2, a3, a4 : std logic vector (WIDTH-1 downto 0);
   signal b1, b2, b3, b4 : std logic vector (WIDTH-1 downto 0);
   signal bv0, bv1, bv2, bv3, bv4: std logic vector (WIDTH-1 downto 0);
   signal bp0, bp1, bp2, bp3, bp4: std logic vector (2*WIDTH-1 downto 0);
   signal pp0, pp1, pp2, pp3, pp4: std logic vector (2*WIDTH-1 downto 0);
```

```
begin
    -- stage 0
    bv0 \le (others => b(0));
    bp0 <= "00000" & (bv0 and a);
    pp0 \le bp0;
    a1 <= a;
    b1 \le b;
    -- stage 1
    bv1 \le (others => b1(1));
    bp1 <= "0000" & (bv1 and a1) & "0";
    pp1 <= pp0 + bp1; 往左移一位
    a2 <= a1;
    b2 <= b1;
    -- stage 2
    bv2 \le (others => b2(2));
    bp2 <= "000" & (bv2 and a2) & "00";
```



```
pp2 \le pp1 + bp2;
    a3 \le a2:
    b3 <= b2;
    -- stage 3
   bv3 \le (others => b3(3));
   bp3 <= "00" & (bv3 and a3) & "000";
    pp3 \le pp2 + bp3;
    a4 <= a3:
    b4 <= b3:
    -- stage 4
    bv4 \le (others => b4(4));
    bp4 <= "0" & (bv4 and a4) & "0000";
    pp4 \le pp3 + bp4;
    -- output
   y \le pp4;
end architecture comb arch;
```





Pipelined multiplier BP architecture pipe arch of mult5 is **constant** WIDTH: integer := 5; BP signal a2 reg, a3 reg, a4 reg, b2 reg, b3 reg, b4 reg; b2_next b2_reg pp1_next pp1_reg a2_next a2_reg std logic vector (WIDTH-1 downto 0); signal a1, a2 next, a3 next, a4 next: BP std logic vector (WIDTH-1 downto 0); signal b1, b2 next, b3 next, b4 next: a3_next a3_reg b3_next pp2_next std logic vector (WIDTH-1 downto 0); pp2_reg b3_reg **signal** bv0, bv1, bv2, bv3, bv4: BP std logic vector (WIDTH-1 downto 0); signal bp0, bp1, bp2, bp3, bp4: std logic vector (2*WIDTH-1 downto 0); pp3_next pp3_reg a4_next b4_next a4_reg signal pp1 reg, pp2 reg, pp3 reg, pp4 reg: BF std logic vector (2*WIDTH-1 downto 0); signal pp0, pp1 next, pp2 next, pp3 next, pp4 next: std logic vector (2*WIDTH-1 downto 0); pp4 next pp4_reg EE332 Digital System Design, by Yu Yajun -- 2019 312

```
a2 reg \leq a2 next;
begin
                                                                           BP
                                            a3 reg <= a3 next;</pre>
 -- pipeline registers
                                            a4 reg <= a4 next;
 process (clk, reset)
                                            b2 reg \le b2 next;
 begin
                                                                           BP
                                            b3 reg \le b3 next;
    if (reset = '1') then
                                            b4 reg \le b4 next;
       pp1 req <= (others => '0');
                                                                             a2_next
a2_reg
                                                                                      b2_next
                                                                   pp1_next
                                          end if;
       pp2 reg <= (others => '0');
                                                                                      b2_reg
                                                                   pp1_reg
                                        end process;
       pp3 reg <= (others => '0');
                                                                           BP
       pp4 reg <= (others => '0');
       a2 reg <= (others => '0');
                                                                   pp2_next
pp2_reg
                                                                             a3_next
       a3 reg <= (others => '0');
                                                                                      b3_next
                                                                             a3_reg
                                                                                      b3 reg
       a4 reg <= (others => '0');
       b2 reg <= (others => '0');
                                                                           ΒP
       b3 reg <= (others => '0');
       b4 reg <= (others => '0');
                                                                                      b4_next
                                                                   pp3_next
                                                                             a4 next
    elsif (clk'event and clk = '1') then
                                                                   pp3_reg
                                                                             a4_reg
                                                                                      b4_reg
       pp1 reg <= pp1 next;
                                                                           BF
       pp2 reg <= pp2 next;
       pp3 reg <= pp3 next;
       pp4 reg <= pp4 next;
                                                                   pp4_reg
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                                                                                     313
```

```
-- merged stage 0 & 1 for pipeline

bv0 <= (others => b(0));

bp0 <= "00000" & (bv0 and a);

pp0 <= bp0;

a1 <= a;

b1 <= b;

--

bv1 <= (others => b1(1));

bp1 <= "0000" & (bv1 and a1) & "0";

pp1_next <= pp0 + bp1;

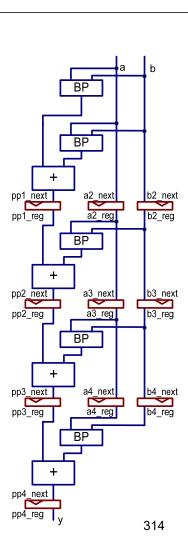
a2_next <= a1;

b2_next <= b1;

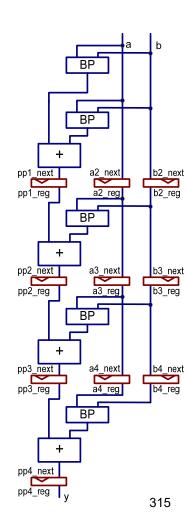
-- stage 2

bv2 <= (others => b2_reg(2));

bp2 <= "000" & (bv2 and a2_reg) & "00";
```



```
pp2 next \le pp1 reg + bp2;
    a3 next <= a2 reg;
    b3 \text{ next} \leq b2 \text{ reg};
    -- stage 3
    bv3 \le (others => b3 reg(3));
    bp3 <= "00" & (bv3 and a3 reg) & "000";
    pp3 next \le pp2 reg + bp3;
    a4 next <= a3 reg;
    b4 \text{ next} \le b3 \text{ reg};
    -- stage 4
    bv4 \le (others => b4 reg(4));
    bp4 <= "0" & (bv4 and a4 reg) & "0000";
    pp4 next \le pp3 reg + bp4;
    -- output
    y \le pp4 reg;
end architecture pipe arch;
```



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More efficient Pipelined multiplier

- Use a smaller (n+1)-bit adder to replace the 2n-bit adder in an n-bit multiplier.
- Reduce the size of the partial-product register
- Reduce the size of the registers that hold the b signal.

```
architecture efficient pipe arch of mult5 is
signal a2 reg, a3 reg, a4 reg; std logic vector(WIDTH-1 downto 0);
signal a1, a2 next, a3 next, a4 next:std logic vector(WIDTH-1 downto 0);
signal b1: std logic vector(4 downto 1);
                                                                  BP
signal b2 next, b2 reg: std logic vector (4 downto 2);
signal b3 next, b3 reg: std logic vector (4 downto 3);
signal b4 next, b4 reg: std logic vector (4 downto 4);
signal bv0, bv1, bv2, bv3, bv4:
    std logic vector (4 downto 0);
                                                                          b2_reg
signal bp0, bp1, bp2, bp3, bp4:
   std logic vector (5 downto 0);
signal pp0: std logic vector (5 downto 0);
signal pp1 next, pp1 reg:std logic vector (6 downto 0);
signal pp2 next, pp2 reg: std logic vector (7 downto 0);
signal pp3 next, pp3 reg: std logic vector (8 downto 0);
signal pp4 next, pp4 reg: std logic vector (9 downto 0);
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                                                                        317
```

```
begin
    -- pipeline registers
    process (clk, reset)
    begin
       if (reset = '1') then
            pp1 reg <= (others => '0');
                                                     pp1 reg <= pp1 next;
            pp2 reg <= (others => '0');
                                                     pp2 reg <= pp2 next;
            pp3 reg <= (others => '0');
                                                     pp3 reg <= pp3 next;
            pp4 reg <= (others => '0');
                                                     pp4 reg <= pp4 next;
            a2 reg <= (others => '0');
                                                     a2 reg \leq a2 next;
            a3 req <= (others => '0');
                                                     a3 reg \le a3 next;
            a4 reg <= (others => '0');
                                                     a4 reg <= a4 next;
            b2 reg <= (others => '0');
                                                     b2 req <= b2 next;
            b3 reg <= (others => '0'):
                                                     b3 req \le b3 next;
            b4 reg <= (others => '0');
                                                     b4 reg <= b4 next;
       elsif (clk'event and clk = '1') then
                                               end if:
                                             end process;
```

```
-- merged stage 0 & 1 for pipeline
bv0 \le (others => b(0));
bp0 \le "0" \& (bv0 and a);
pp0 \le bp0;
a1 <= a;
b1 <= b (4 downto 1);
bv1 \le (others => b1(1));
bp1 \le "0" \& (bv1 and a1);
pp1 next(6 downto 1) \leq ("0" & pp0(5 downto 1)) + bp1;
pp1 next(0) \le pp0(0);
a2 next <= a1;
b2 next <= b1(4 downto 2);
-- stage 2
bv2 \le (others => b2 reg(2));
bp2 \le "0" \& (bv2 and a2 reg);
pp2 next(7 downto 2) \leq ("0" & pp1 reg(6 downto 2)) + bp2;
pp2 next(1 downto 0) \leq pp1 reg(1 downto 0);
                                                               pp4_next
a3 next <= a2 reg;
                                                              pp4_reg \
b3 next \leq b2 reg(4 downto 3);
```

```
-- stage 3
   bv3 <= (others => b3_reg(3));
   bp3 <= "0" & (bv3 and a3 reg);
   pp3 next(8 downto 3) <=
            ("0" & pp2 reg(7 downto 3)) + bp3;
   pp3 next(2 downto 0) \le pp2 reg(2 downto 0);
   a4 next <= a3 reg;
   b4 next \leq b3 reg(4);
    -- stage 4
   bv4 \le (others => b4 reg(4));
   bp4 <= "0" & (bv4 and a4 reg);
   pp4 next(9 downto 4) <=
                                                             BF
           ("0" & pp3 reg(8 downto 4)) + bp4;
   pp4 next(3 downto 0) \leq pp3 reg(3 downto 0);
    -- output
   y \le pp4 reg;
                                                             BP
end architecture efficient pipe arch;
```

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Chapter 10. Subprograms, packages and libraries

- For complex design, VHDL provides mechanics for structuring programs, reusing modules
 - Subprograms such as functions and procedures for encapsulating commonly used operations
 - Packages and libraries for sharing large bodies of code.

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10.1 Functions

- Functions are used to compute a value based on the values of the input parameters.
- Function declaration:

function func_name (parameter_list) return data_type;

formal parameter return value type

Parameter values in functions are used, but not changed within the function.

Structure of a function

```
function rising_edge (clock: std_logic) return boolean is
--
-- declarative region: declare variables local to the function
--
begin
-- body
-- sequential statement;
-- sequential statement;
-- return (value);
end function rising edge;
```

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Calling a function in a VHDL module:

```
rising_edge (enable); -- positional association
```

actual parameter, the type must match to the formal parameter, i.e., enable must be a **signal** of type std logic.

rising_edge (clock => enable); -- name association

- Functions execute in zero simulation time.
 - Wait statement are not permitted in functions.

```
entity dff is
                 D in, CLK: in std logic;
        port (
                 D out: out std logic);
        end entity dff;
        architecture behavior of dff is
            function rising edge(signal clock: in std logic) return boolean is
                 variable edge : boolean :=FALSE;
function
             begin
definition
                 edge := (clock='1' and clock'event);
                 return (edge);
             end function rising edge;
        begin
            process (CLK) is
            begin
                 if (rising edge(CLK)) then
                          D out \leq D in;
                 end if:
            end process;
        end architecture behavior;
```

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A majority function

- It returns '1' if two or more of the 3 input parameters, a, b and c are '1'.
- It can be treated as a shorthand for the a·b+a·c+b·c expression

```
architecture arch of ...
    -- declaration
function maj(a, b, c: std_logic) return std_logic is
    variable result: std_logic;
begin
    result := (a and b) or (a and c) or (b and c);
    return result;
end function maj;
signal i1, i2, i3, i4, x, y: std_logic;
begin
    ...
    x <= maj(i1, i2, i3) or i4;
    y <= i1 when maj(i2, i3, i4)='1' else ...
end arch;</pre>
```

Type conversion functions

 To make assignment from an object of one type to an object of another type.

```
- for example: bit_vector and std_logic_vector.
function to_bitvector (svalue: std_logic_vector) return bit_vector is
  variable outvalue: bit_vector(svalue'length - 1 downto 0);
begin
  for i in svalue' range loop -- scan all elements of the array
      case svalue(i) is
      when '0' => outvalue (i) := '0';
      when '1' => outvalue (i) := '1';
      when others => outvalue (i) := '0';
      end case;
  end loop;
  return outvalue;
end function to bitvector;
```

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A Function performing [log₂n]

```
function log2c (n: integer) return integer is variable m, p: integer;
begin m := 0;
p:= 1;
while p < n loop
m:= m+1;
p:= p*2;
end loop;
return m;
end function log2c;
```

Summary

- Unlike entity and architecture, functions (and procedures) are not design units and thus cannot be processed independently.
- In synthesis, functions should not be used to specify the design hierarchy, but should be treated as a shorthand for simple, repeatedly used operations.
- A function can be thought of as an extension of the expression and can be "called" whenever an expression is used.

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9.2 Package

- The primary purpose of a package is to collect elements that can be shared (globally) among two or more design units. A package is represented by:
 - a package declaration and, optionally,
 - a package body.

contains a set of declarations that may possibly be shared by many design units. It defines items that can be made visible to other design units, for example, a function declaration.

Package declaration and package body are design units of VHDL.

An example of a package declaration

```
package SYNTH_PACK is
   constant LOW2HIGH : TIME := 20ns;
   type ALU_OP is (ADD, SUB, MUL, DIV, EQL);
   type MVL is ('U', '0', '1', 'Z');
   component NAND2
        port (A, B : in MVL; C : out MVL);
   end component;
   -- subprogram, type, constant, signal, variable, component ...,
   -- and use clause can be declared in package declaration
end package SYNTH_PACK;
```

 If the declarations include items such as functions or procedure declarations, the behavior of the function and procedure are specified in a separate design unit called the package body.

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```
use WORK.SYNTH_PACK.all;

package PROGRAM_PACK is
    constant PROP_DELAY : TIME;
    function ISZERO(A: MVL) return boolean;
end package PROGRAM PACK;
```

- In this case, a package body is required.
- A package body primarily contains the behavior of the subprograms declared in a package declaration. It may also contain other declarations.

An example of a package body

```
package body PROGRAM_PACK is

constant PROP_DELAY : TIME := 15ns;

function ISZERO(A: MVL) return boolean is

begin

if (A='0') return TRUE;

else return FALSE;

end if;
end function ISZERO;
end package body PROGRAM PACK;
```

package body name: must be the same as of its corresponding package declaration.

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Note:

 An item declared inside a package body has its scope restricted to be within the package body, and this item cannot be made visible in other design unit. (This is in contrast to items declared in a package declaration).

10.4 Design libraries

- Each design unit entity, architecture, configuration, package declaration, package body is analyzed (complied) and placed in a design library for subsequent use.
- To use a design library, the library must be declared by its logical name.

library logical-library-name1, logical-library-name2,...;

In VHDL, the libraries STD and WORK are implicitly declared.

standard packages provided with the

VHDL distributions

working directory

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 Once a library is declared, all of the functions, procedures, and type declaration of a package in that library can be made accessible to a VHDL model through the use clause.

library IEEE use IEEE.std logic 1164.all;

 These clauses apply only to the immediate entityarchitecture pair! Visibility must be established for other design units separately!