

Tutorial on Orthogonal Frequency Division Multiplexing (OFDM)



Discrete Fourier Transform (DFT)

- In practice, there is a huge demand on processing **finite duration signals**
- Given a finite duration signal $\{x[0], x[1], \dots, x[N-1]\}$, its Fourier transform is

$$X(e^{j\omega}) = \sum_{n=0}^{N-1} x[n]e^{-j\omega n}$$

- **Drawback:** The spectrum of DTFT is continuous \Rightarrow Cannot be handle by computer.
- **Discrete Fourier Transfrom (DFT)** is developed for digital processing of finite duration signals

$$\tilde{X}[k] = \frac{1}{N} \sum_{n=0}^{N-1} x[n]e^{-jk(2\pi/N)n} \quad k = < N >$$

- **DFT: frequency sampling of DTFT**

$$\tilde{X}[k] = \frac{1}{N} X(e^{j2k\pi/N}) \quad k = < N >$$

Inverse DFT

- Equation system of DFT:

$$\begin{pmatrix} \tilde{X}[0] \\ \tilde{X}[1] \\ \vdots \\ \tilde{X}[N-1] \end{pmatrix} = \underbrace{\frac{1}{N} \begin{pmatrix} e^{-j0} & e^{-j0} & \dots & e^{-j0} \\ e^{-j0} & e^{-j2\pi/N} & \dots & e^{-j2(N-1)\pi/N} \\ \vdots & \vdots & \ddots & \vdots \\ e^{-j0} & e^{-j2(N-1)\pi/N} & \dots & e^{-j2(N-1)(N-1)\pi/N} \end{pmatrix}}_{\mathbf{F}} \begin{pmatrix} x[0] \\ x[1] \\ \vdots \\ x[N-1] \end{pmatrix}$$

- Transform matrix \mathbf{F} is full rank.
- Observation:** $\{\tilde{X}[k] | k = 0, 1, \dots, N-1\}$ maintains all the information of $\{x[0], x[1], \dots, x[N-1]\}$
- Inverse DFT is feasible:

$$x[n] = \sum_{k=0}^{N-1} \tilde{X}[k] e^{jk(2\pi/N)n} \quad n = 0, 1, \dots, N-1$$

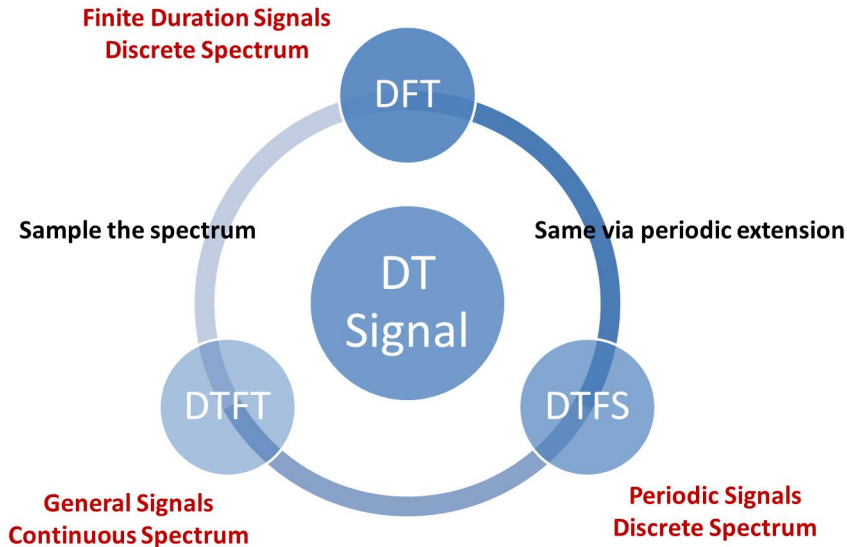
DFT and DTFS

- DFT is for finite duration signals; DTFS is for periodic signals
- Define $\tilde{x}[n]$ as the periodic extension of $x[n]$: Repeat $x[n]$ with period N
- Fourier series of $\tilde{x}[n]$:

$$\frac{1}{N} \sum_{n=0}^{N-1} \tilde{x}[n] e^{-jk(2\pi/N)n} = \frac{1}{N} \sum_{n=0}^{N-1} x[n] e^{-jk(2\pi/N)n} \Rightarrow \text{DFT of } x[n]$$

- DFT of a finite-duration signal = DTFS of its periodic extension
- Reference on DFT:
 - ▶ Textbook: Problem 5.53, 5.54
 - ▶ http://en.wikipedia.org/wiki/Discrete_Fourier_transform

Comparison



Periodic Convolution of Finite Duration Signals

Periodic Convolution

Let x and y be two finite duration signals with duration N , \tilde{x} and \tilde{y} be the associated periodic extension, then the periodic convolution of finite duration signals is defined as

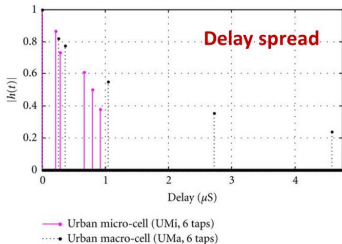
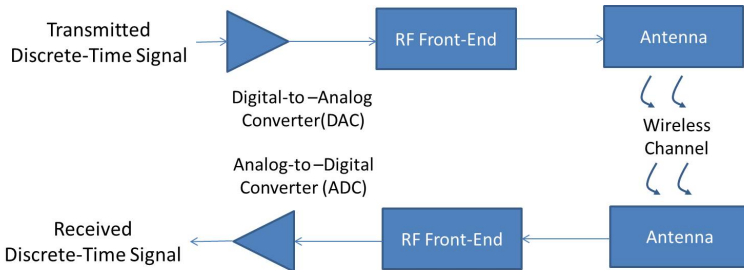
$$x[n] \circledast y[n] := \tilde{x}[n] \circledast \tilde{y}[n] = \sum_{k=\langle N \rangle} \tilde{x}[k] \tilde{y}[n-k] \quad n = 0, 1, 2, \dots, N-1$$

Convolution Property of DFT

Time domain periodic convolution is equivalent to frequency domain multiplication, thus,

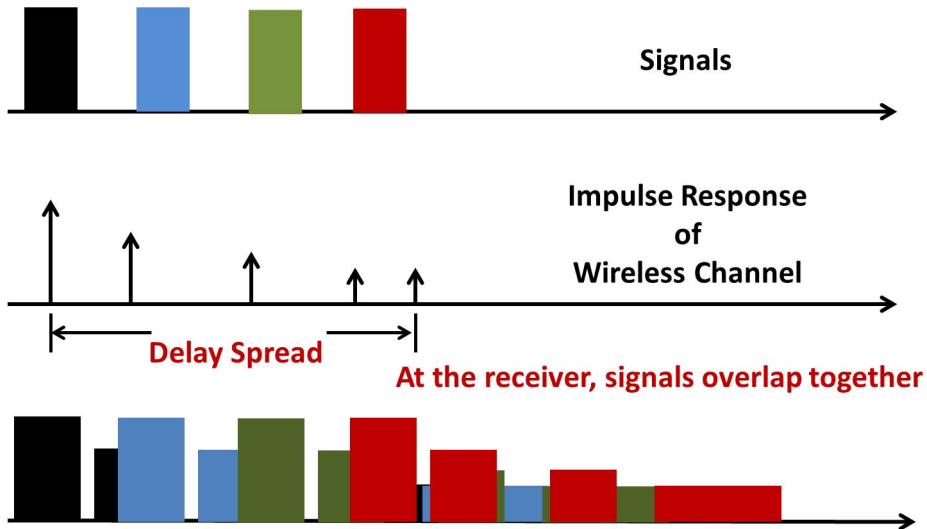
- $x[n] \circledast y[n] \longleftrightarrow N \tilde{X}[k] \tilde{Y}[k]$

Wireless Systems



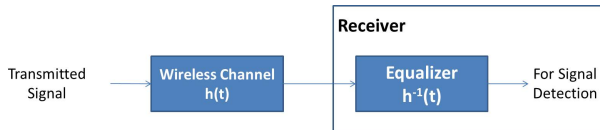
Delay spread causes inter-symbol interference

Inter-Symbol Interference (1/2)



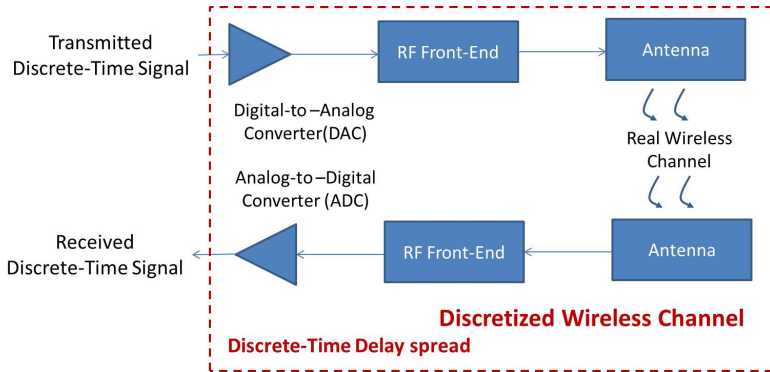
Inter-Symbol Interference (2/2)

- How to deliver signals without inter-symbol interference?
 - ▶ Suppose the duration of delay spread is ΔH seconds
- Approach 1: Send data on every ΔH seconds
- Approach 2: Channel equalizer



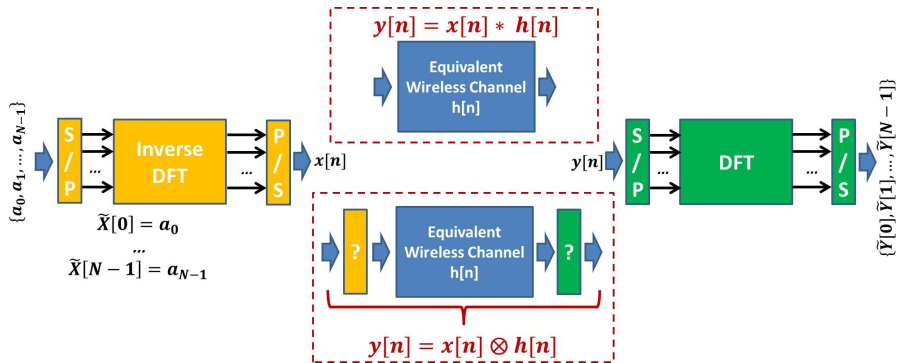
- Approach 3: Orthogonal Frequency Division Multiplexing
 - ▶ Pre-processing at the transmitter + post-processing at the receiver

Discretetize Wireless Channel



- Tx front-end + wireless channel + Rx front-end: approximately discrete-time LTI system
- Denote the impulse response as $h[n]$

OFDM at First Glance



- Signals are loaded in frequency domain
- Some mechanism is necessary to generate the effect of periodic convolution

OFDM Analysis

- Let $\{a_0, a_1, \dots, a_{N-1}\}$ be the sequence of bits to be delivered from the transmitter to the receiver.
- At the transmitter

- ▶ we let

$$\tilde{X}[0] = a_0, \tilde{X}[1] = a_1, \dots, \tilde{X}[N-1] = a_{N-1}$$

- ▶ Take inverse DFT:

$$\{x[0], x[1], \dots, x[N-1]\} = DFT^{-1}\{\tilde{X}[0], \tilde{X}[1], \dots, \tilde{X}[N-1]\}$$

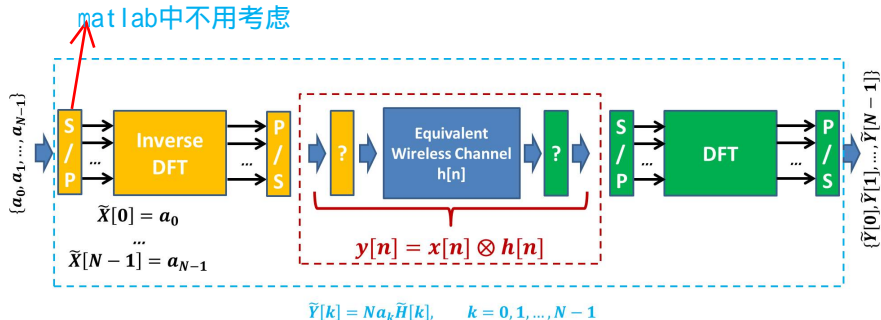
- With certain mechanism, the received signal becomes

$$y[n] = x[n] \circledast h[n]$$

- At the receiver, take DFT on $y[n]$:

$$\tilde{Y}[k] = N\tilde{X}[k]\tilde{H}[k] = Na_k\tilde{H}[k] \quad k = 0, 1, \dots, N-1$$

So, ...



- How to detect $\{a_k | \forall k\}$ from $\{\tilde{Y}[k] | \forall k\}$?
- How to design the blocks "?" ?

Channel Estimation & Signal Detection

- **Channel estimation:** Estimate the channel gain $\tilde{H}[k]$ ($k = 0, 1, \dots, N - 1$)
 - ▶ The transmitter sends the common information to the receiver $\{a_0^c, a_1^c, \dots, a_{N-1}^c\}$:

$$\tilde{Y}^c[k] = Na_k^c \tilde{H}[k] \quad k = 0, 1, \dots, N - 1$$

- ▶ Evaluate channel via

$$\tilde{H}[k] = \tilde{Y}^c[k] / (Na_k^c) \quad k = 0, 1, \dots, N - 1$$

- **Signal detection:** Detect the transmitter's signal via the knowledge of $\tilde{H}[k]$ ($k = 0, 1, \dots, N - 1$):
 - ▶ The transmitter sends the information to the receiver $\{a_0, a_1, \dots, a_{N-1}\}$:

$$\tilde{Y}[k] = Na_k \tilde{H}[k] \quad k = 0, 1, \dots, N - 1$$

- ▶ Detect signal via

$$a_k = \tilde{Y}[k] / (N\tilde{H}[k]) \quad k = 0, 1, \dots, N - 1$$

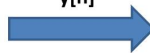
Convolution vs. Periodic Convolution

Transmission Signal
 $x[n]$

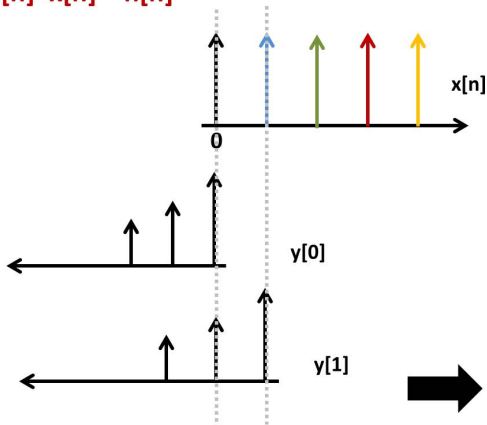
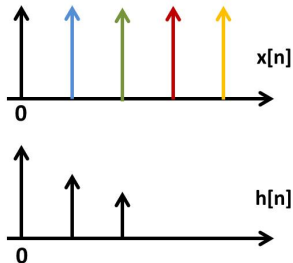


Equivalent Wireless Channel
DT LTI System $h[n]$

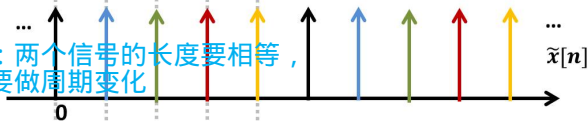
Received Signal
 $y[n]$



$$y[n] = x[n] * h[n]$$

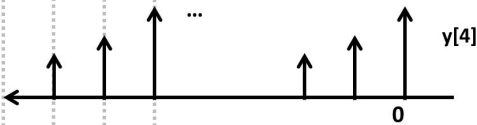
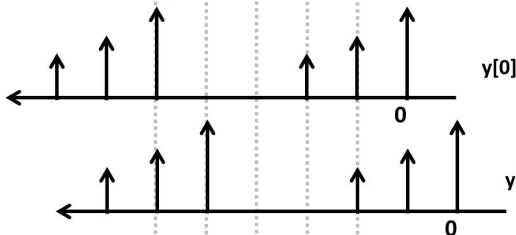
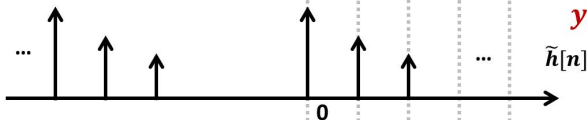


周期卷积的前提：两个信号的长度要相等，
同时两个信号还要做周期变化



$$y[n] = x[n] \otimes h[n]$$

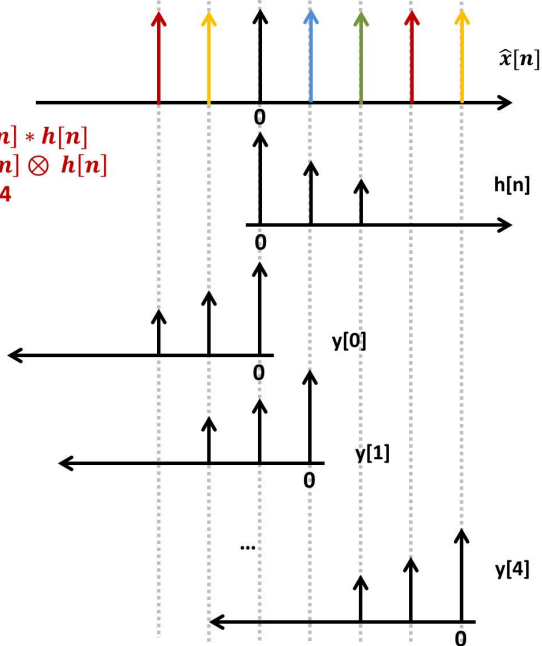
周期卷积的结果也是周期



$$y[n] = \hat{x}[n] * h[n]$$

$$= x[n] \otimes h[n]$$

$$n=0,1,2,3,4$$



Observations

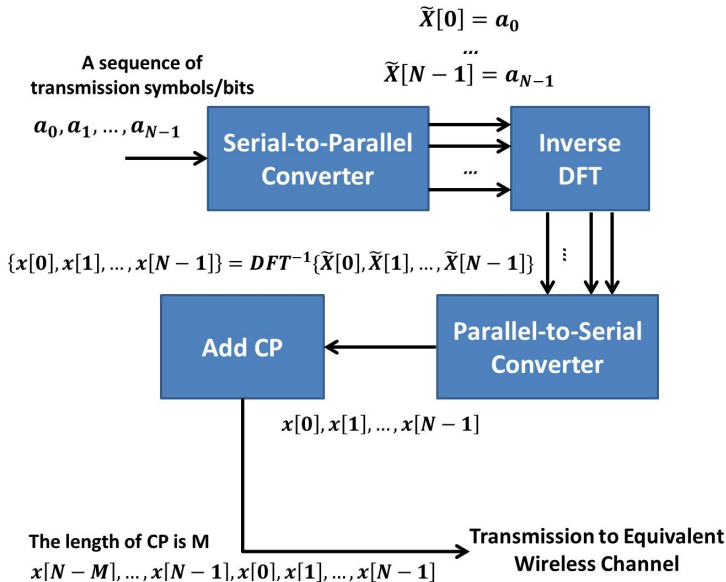
- Periodic convolution can be different from convolution
- In order to guarantee that

$$y[n] = x[n] \circledast h[n], \text{ for } n = 0, 1, 2, \dots, N - 1,$$

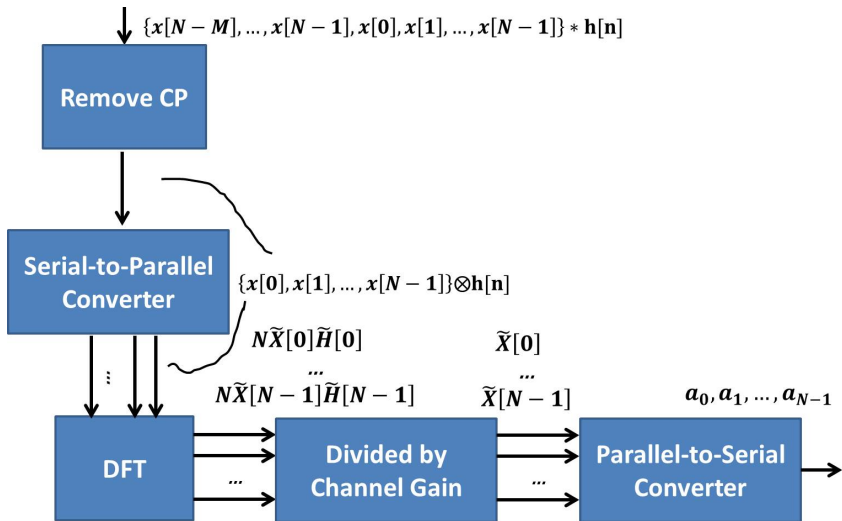
we should copy the last few samples of $x[n]$ to the beginning, which is named as **cyclic prefix**

- ▶ How many samples should we copy?

Transmitter Structure of OFDM



Receiver Structure of OFDM



Reference

- Reference

- ▶ www.gaussianwaves.com/2011/05/introduction-to-ofdm-orthogonal-frequency-division-multiplexing-2/
- ▶ www.wirelesscommunication.nl/reference/chaptr05/ofdm/ofdmmath.htm

