

Review

- Search and sort are essential functions to process data
- Linear search
- Binary search
- Selection sort
- Insertion sort

Recursion

Lecture 11-1

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Recursion

- Function that calls itself during execution -

Recursion

- Let's implement factorial function ($n! = 1 \times 2 \times 3 \times \dots \times (n-1) \times n$)
 - `>>> def facto(n: int) -> int:`
 - `... ans = 1`
 - `... for i in range(1,n+1):`
 - `... ans = ans * i`
 - `... return ans`
- How about this?
 - `>>> def facto(n: int) -> int:`
 - `... if n == 0:`
 - `... return 1`
 - `... else:`
 - `... return n*facto(n-1)`

Recursion

- Recursion can happen when solving a problem includes solving **subproblems** having the **same structure**
 - Easier to implement (if you can think of this way ever)
 - Results of subproblems can be reused (called dynamic programming, out of scope)
- Structure
 - `>>> def facto(n: int) -> int:`
 - `... if n == 0:` *#Conditional statements check for base cases*
 - `... return 1` *#Base case (evaluated without recursive calls)*
 - `... else:`
 - `... return n*facto(n-1)` *#Recursive case (evaluated with recursive calls)*

Example – Fibonacci Sequence

- Implement **Fibonacci sequence**, starting from $n=1$
 - 1,2,3,5,8,13,21,34,55,89 ...
- What are
 - (1) the conditional statement,
 - (2) the base case, and
 - (3) the recursive case?

Example – Fibonacci Sequence

- Another example: $\text{Fibonacci}(n) = \text{Fibonacci}(n-1) + \text{Fibonacci}(n-2)$
 - 1, 2, 3, 5, 8, 13, 21, 34 ...
 - `>>> def fibonacci(n: int) -> int:`
 - ... `if n == 1 or n == 2:` *#Conditional statements*
 - ... `return n` *#Base case*
 - ... `else:`
 - ... `return fibonacci(n-1) + fibonacci(n-2)` *#Recursive case*

Example – Fibonacci Sequence

- Another example: $\text{Fibonacci}(n) = \text{Fibonacci}(n-1) + \text{Fibonacci}(n-2)$
 - 1, 2, 3, 5, 8, 13, 21, 34 ...
 - `>>> def fibonacci(n: int) -> int:`
 - ... `if n == 1 or n == 2:` *#Conditional statements*
 - ... `return n` *#Base case*
 - ... `else:`
 - ... `return fibonacci(n-1) + fibonacci(n-2)` *#Recursive case*
- Is Fibonacci implemented correctly?
 - Verify the base case
 - Assuming that `fibonacci(n-1)` and `fibonacci(n-2)` are correct, verify if `fibonacci(n)` is correct

Merge Sort

Lecture 11-2

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Motivation

- Insertion sort and selection sort work but too slow – proportional to n^2
 - Does not matter when handling small data, but we want to handle **big data**!
- Recall linear search vs. binary search – Divide the whole task into **two parts**
 - Is there a way something similar?

Merge sort!

index	0	1	2	3	4	5	6	7
values	5	-2	0	100	-6	7	4	9

Merge Sort – Idea

- Step 1: **Divide** the whole list into two sub-lists

	<i>Sublist1</i>				<i>Sublist2</i>			
index	0	1	2	3	4	5	6	7
values	5	-2	0	100	-6	7	4	9

Merge Sort – Idea

- Step 1: **Divide** the whole list into two sub-lists
- Step 2: **Sort** the left sublist and the right sublist separately
 - Smells like **binary** something...

	<i>Sublist1 – sorted!</i>				<i>Sublist2 – sorted!</i>			
index	0	1	2	3	4	5	6	7
values	-2	0	5	100	-6	4	7	9

Merge Sort – Idea

- Step 1: **Divide** the whole list into two sub-lists
- Step 2: **Sort** the left sublist and the right sublist separately
 - Smells like **binary** something...
- Step 3: **Merge** the two sorted sublist in a sorted way

Merge sublist1 and sublist2!

index	0	1	2	3	4	5	6	7
values	-6	-2	0	4	5	7	9	100

How to sort sublists?

Merge Sort – Idea

- Step 1: **Divide** the whole list into two sub-lists
- Step 2: **Sort** the left sublist and the right sublist separately, by using merge sort
 - Smells like **binary** something...
- Step 3: **Merge** the two sorted sublists in a sorted way

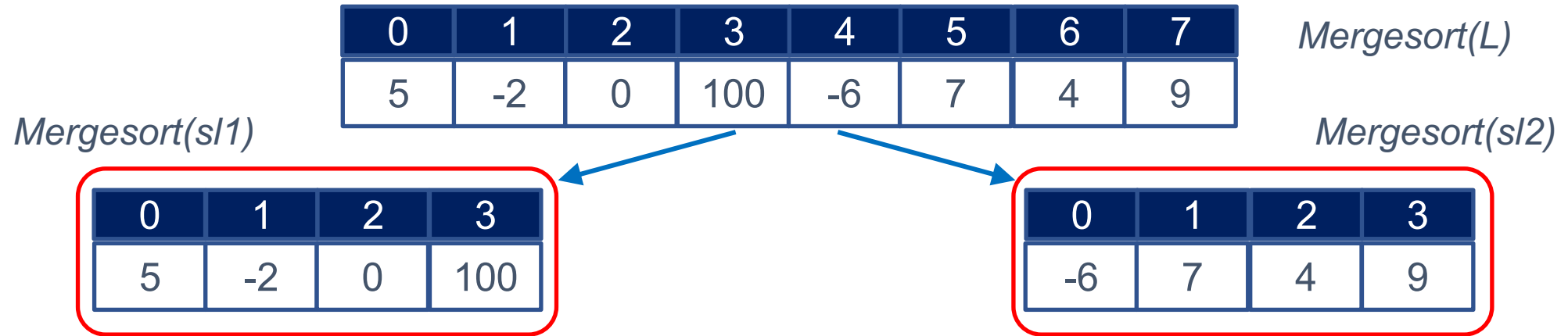
	<i>Sublist1 – mergesort!</i>				<i>Sublist2 – mergesort!</i>			
index	0	1	2	3	4	5	6	7
values	5	-2	0	100	-6	7	4	9

Merge Sort – Operation (Breakdown)

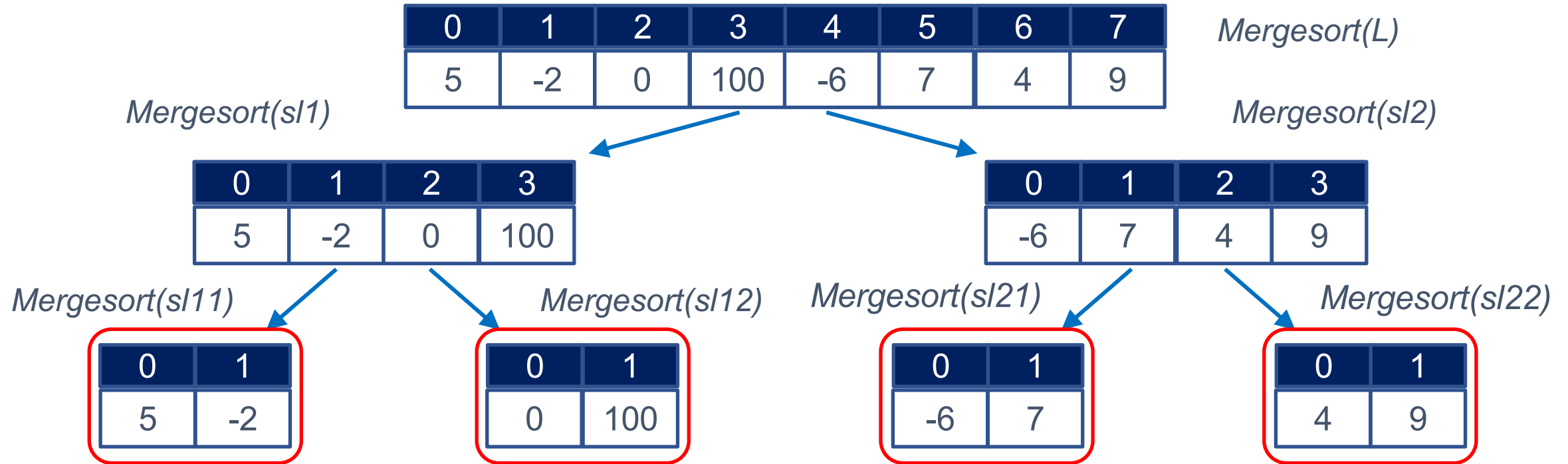
0	1	2	3	4	5	6	7
5	-2	0	100	-6	7	4	9

Mergesort(L)

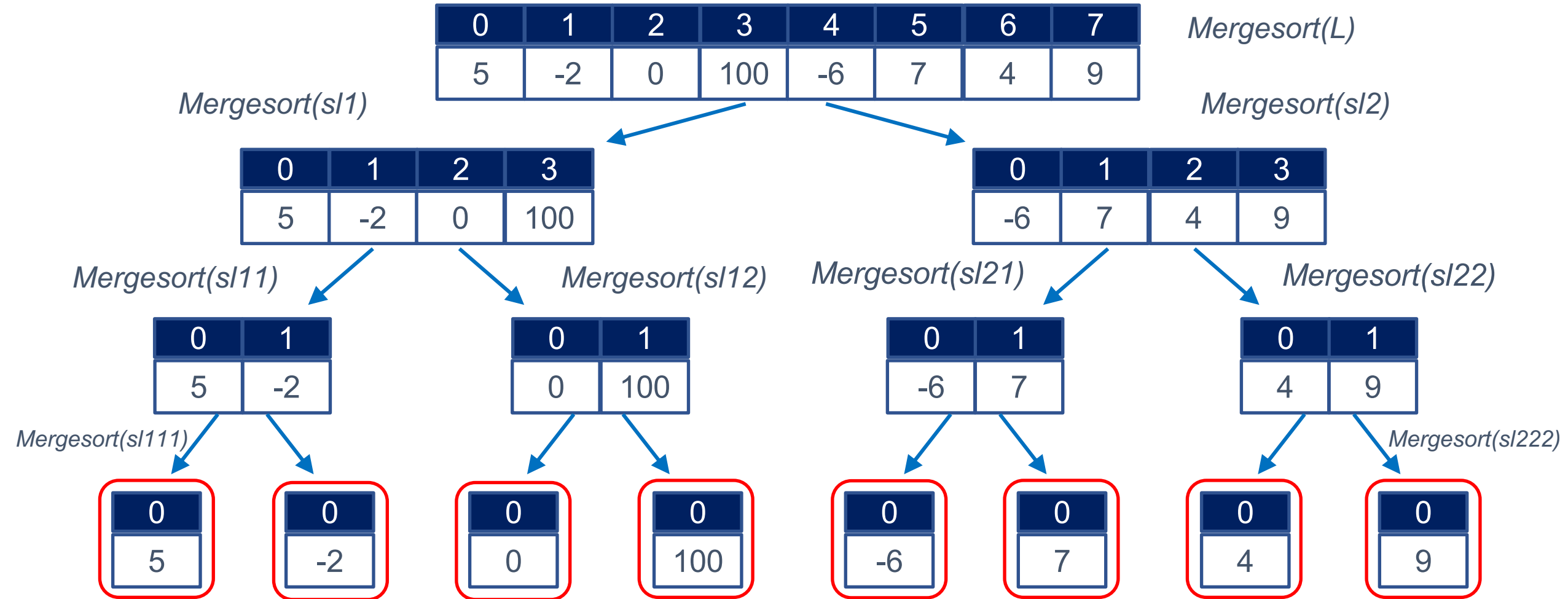
Merge Sort – Operation (Breakdown)



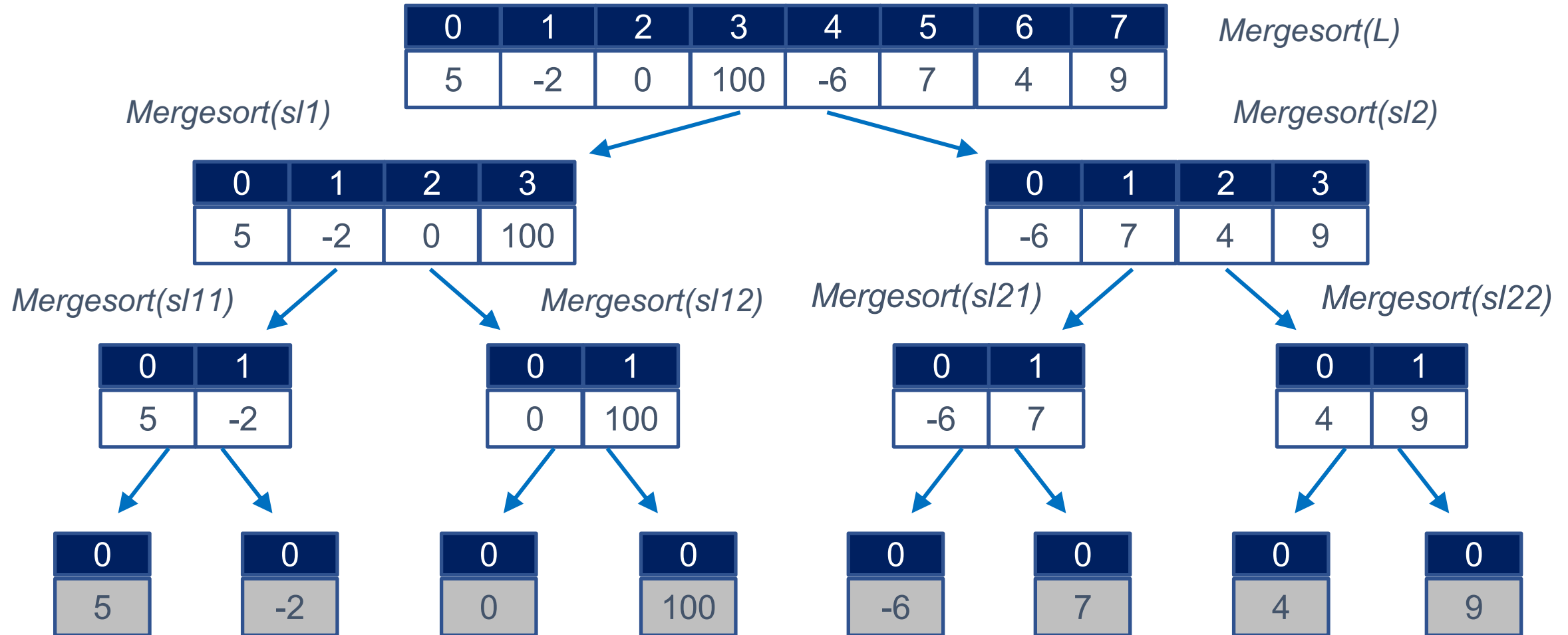
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Merge Sort – Operation (Breakdown)

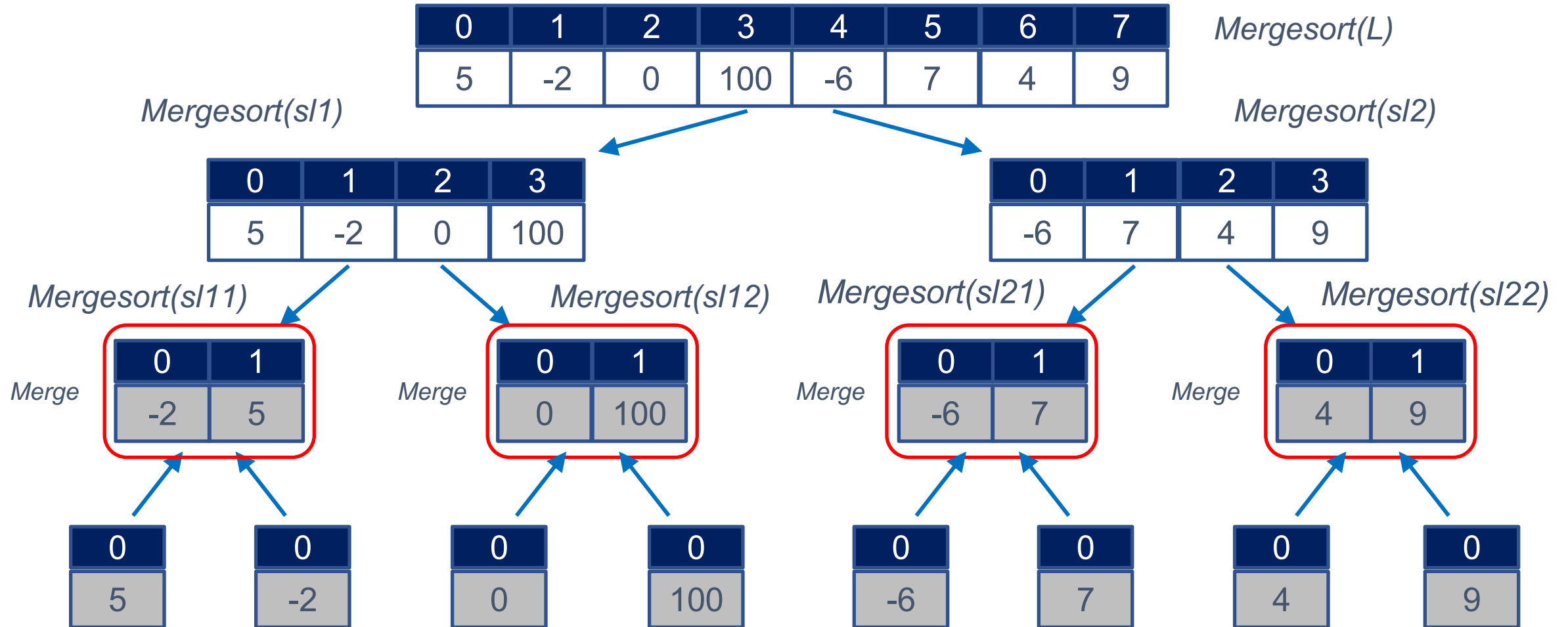


Merge Sort – Operation (Breakdown)

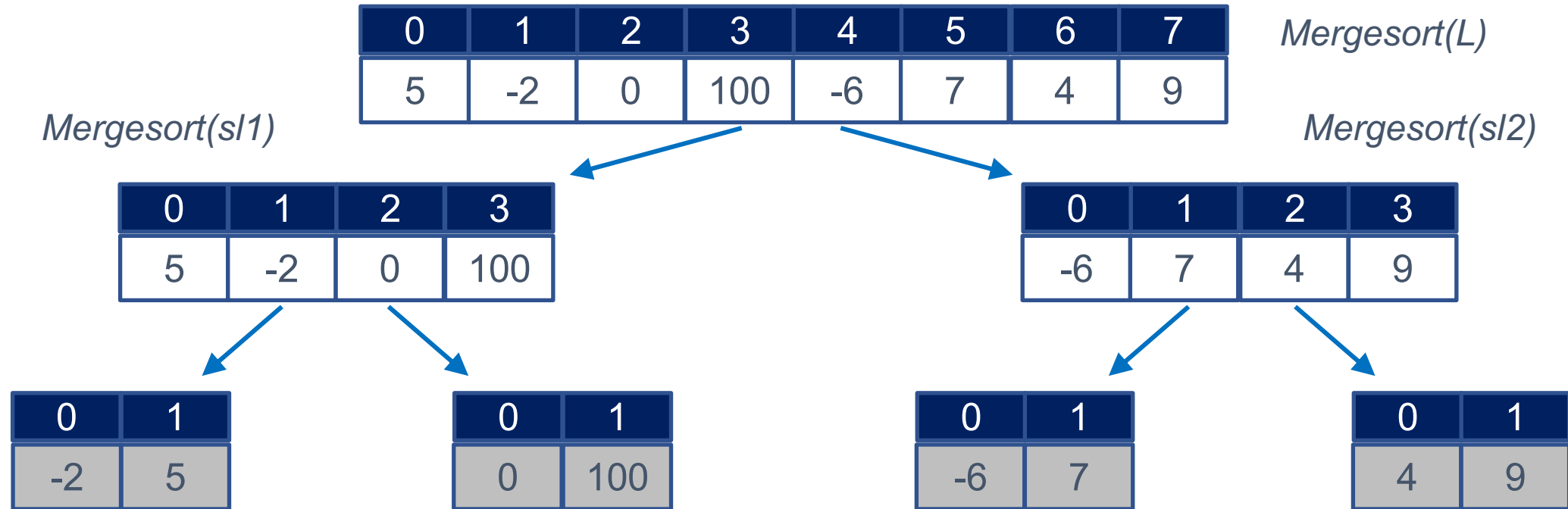


Time to go upwards!

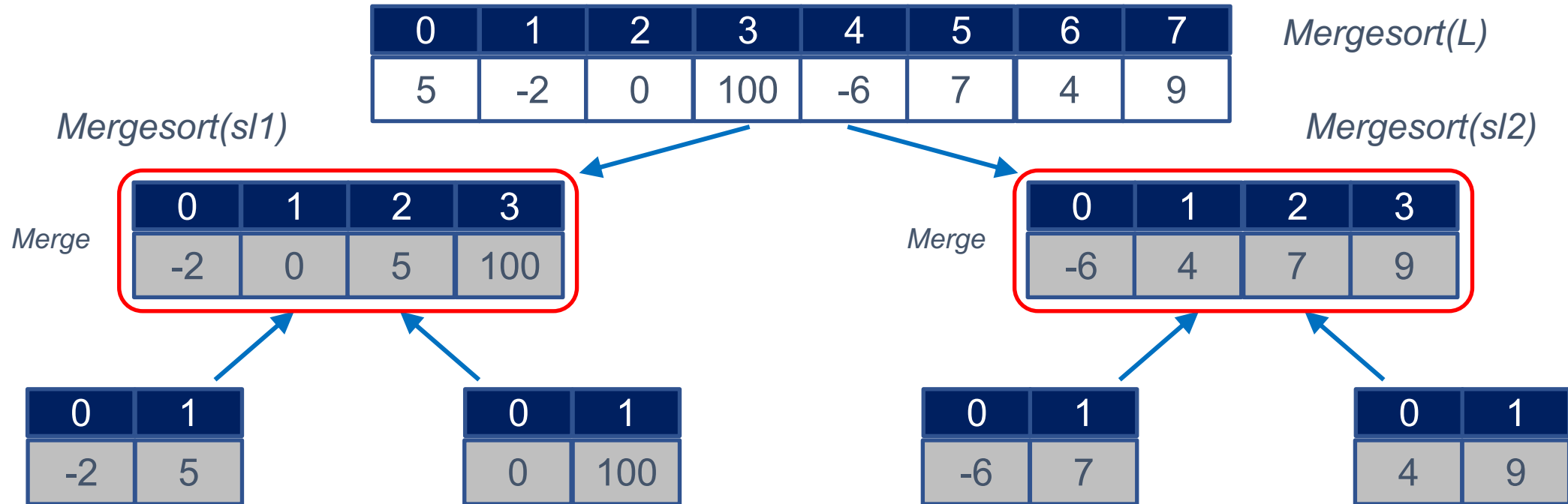
Merge Sort – Operation (Merge)



Merge Sort – Operation (Merge)

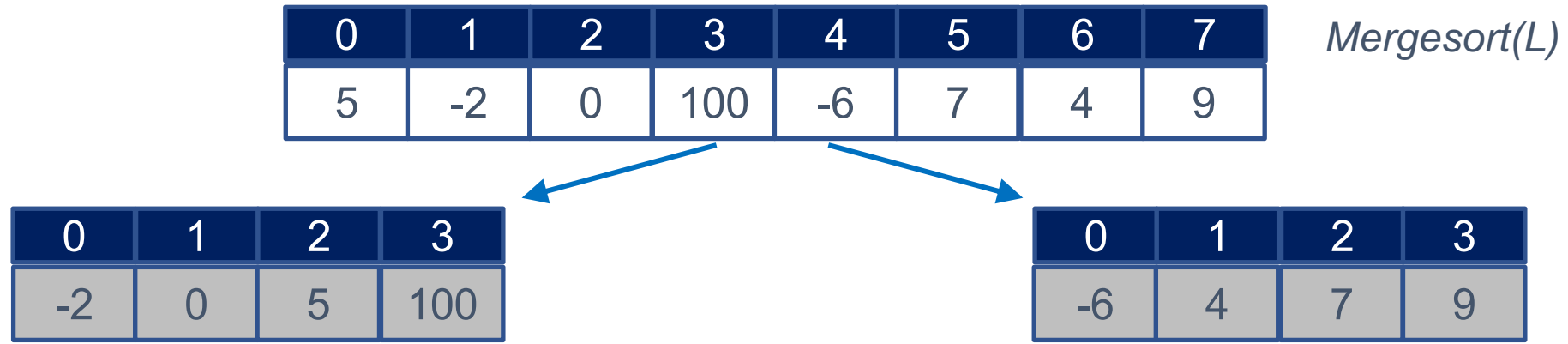


Merge Sort – Operation (Merge)

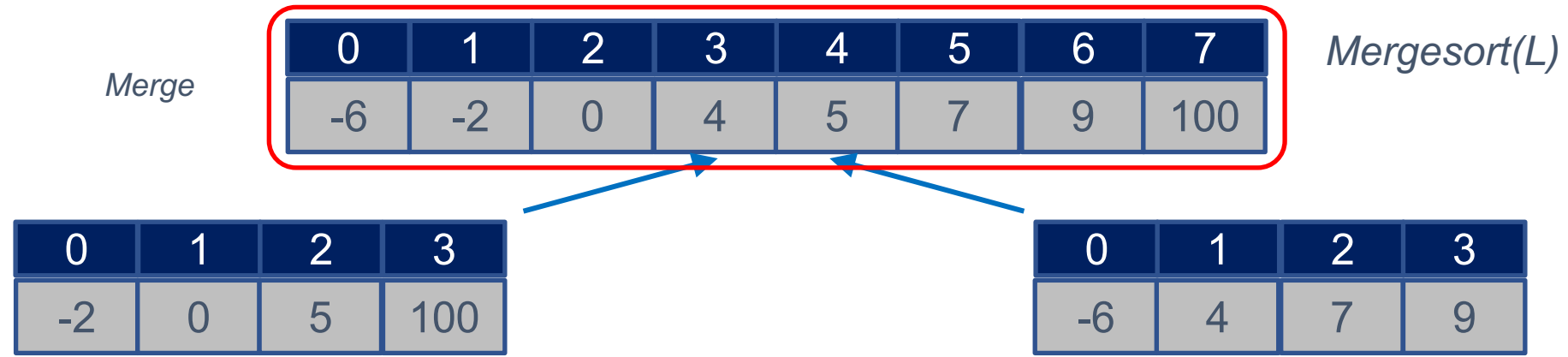


Compare the **leftmost items** of the two sublists, given that the two lists are **already sorted**!

Merge Sort – Operation (Merge)



Merge Sort – Operation (Merge)



Again, compare the **leftmost items** of the two sublists,
given that the two lists are **already sorted**!

Merge Sort – Operation (Merge)

0	1	2	3	4	5	6	7
-6	-2	0	4	5	7	9	100



Summary

- Mergesort – a sorting algorithm using recursion

Merge Sort Implementation

Lecture 11-3

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Merge Sort – Recursive Call

- `def mergeSort(L: list) -> None:`
- `mergeSortHelp(L, 0, len(L) - 1)`

Merge Sort – Recursive Call

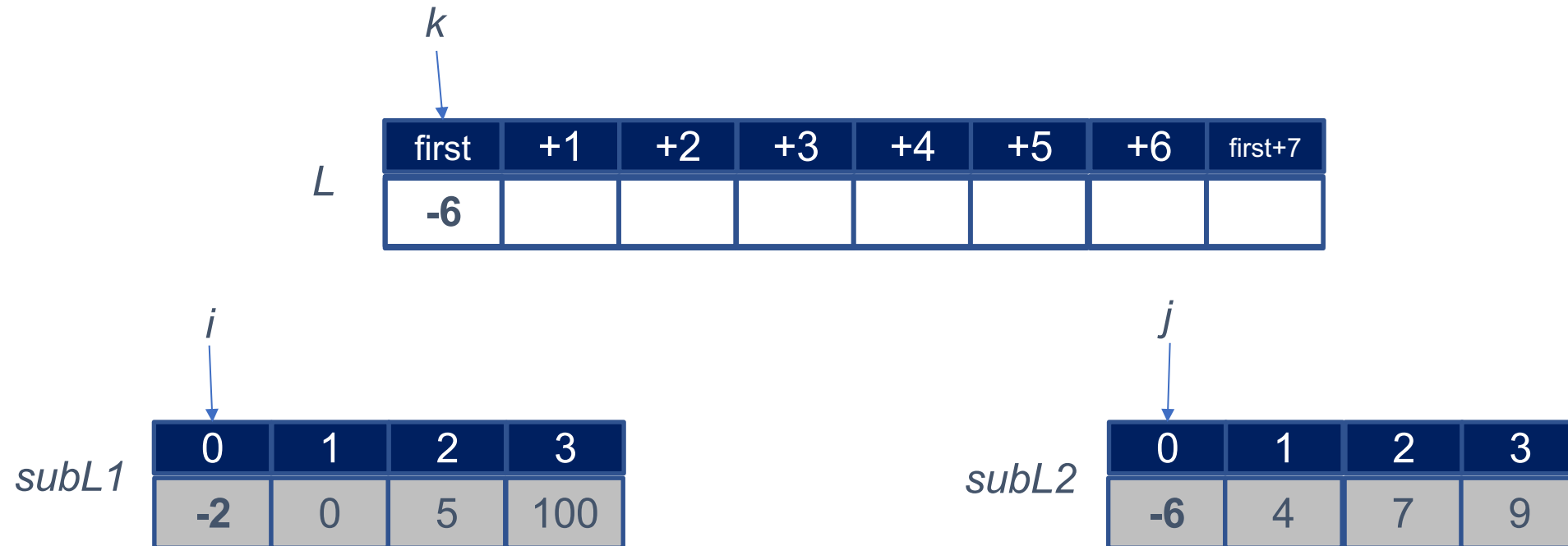
- `def mergeSortHelp(L: list, first: int, last: int) -> None:`
- `if first == last:` `# Conditional statements`
- `return` `# Base case`
- `else:`
- `mid = first + (last - first) // 2`
- `mergeSortHelp(L, first, mid)` `# Recursive call for sublist1`
- `mergeSortHelp(L, mid+1, last)` `# Recursive call for sublist2`
- `merge(L, first, mid, last)` `# Merge the two (sorted) sublists`

*Parameters to indicate where are
two sublists and the whole list*



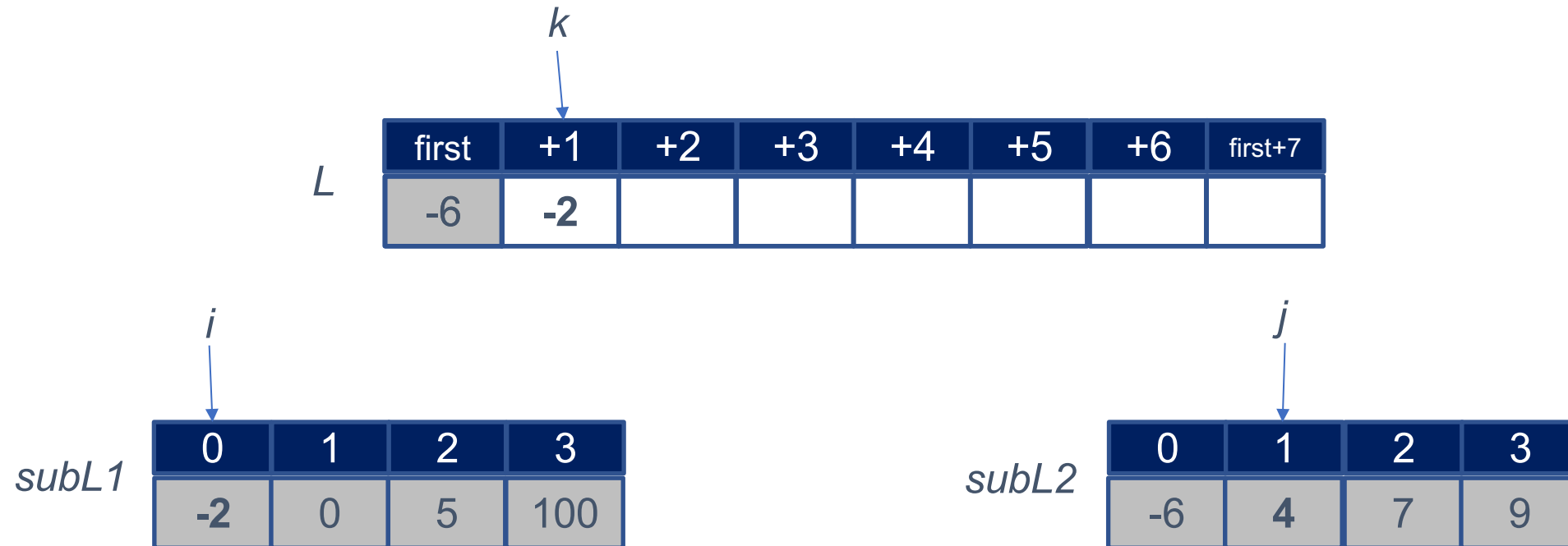
Merge Sort – Merge Algorithm

- Merge(L, first, mid, last)
 - Memory complexity: $O(\text{len}(L))$ for $\text{subL1} = L[\text{first}:\text{mid}+1]$ and $\text{subL2} = L[\text{mid}+1:\text{last}+1]$



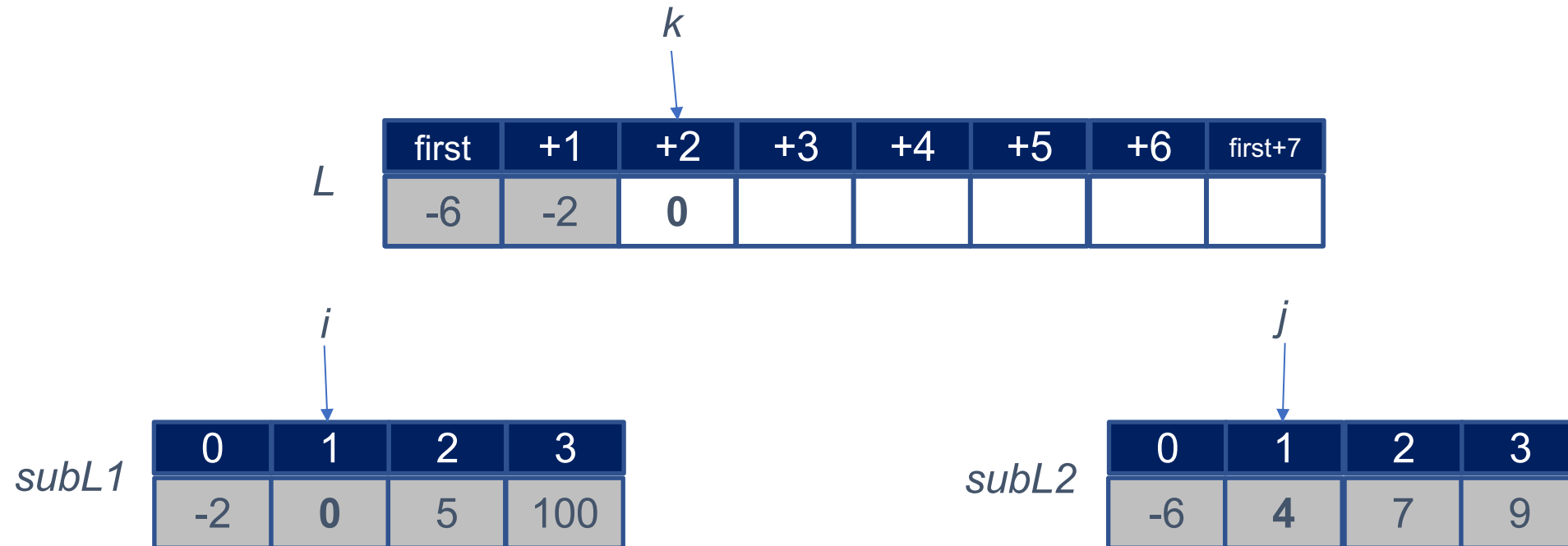
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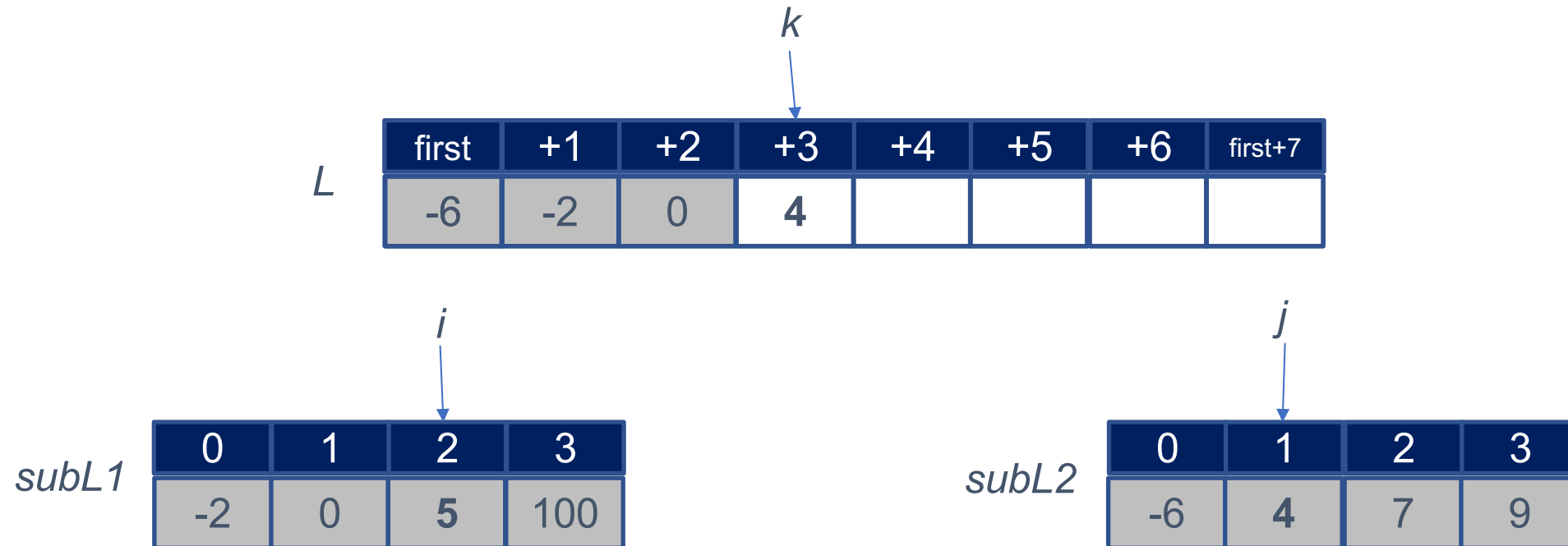
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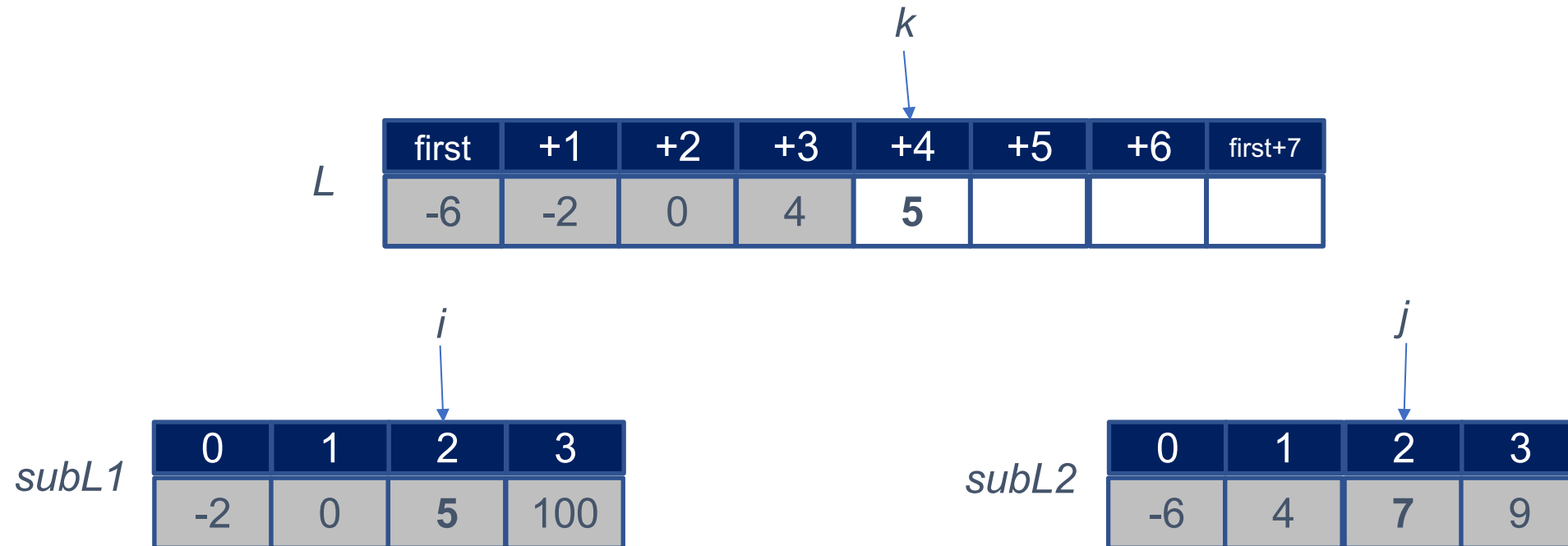
Merge Sort – Merge Algorithm

- Merge(L , first, mid, last)
 - Memory complexity: $O(\text{len}(L))$ for $\text{subL1}=L[\text{first}:\text{mid}+1]$ and $\text{subL2}=L[\text{mid}+1:\text{last}+1]$



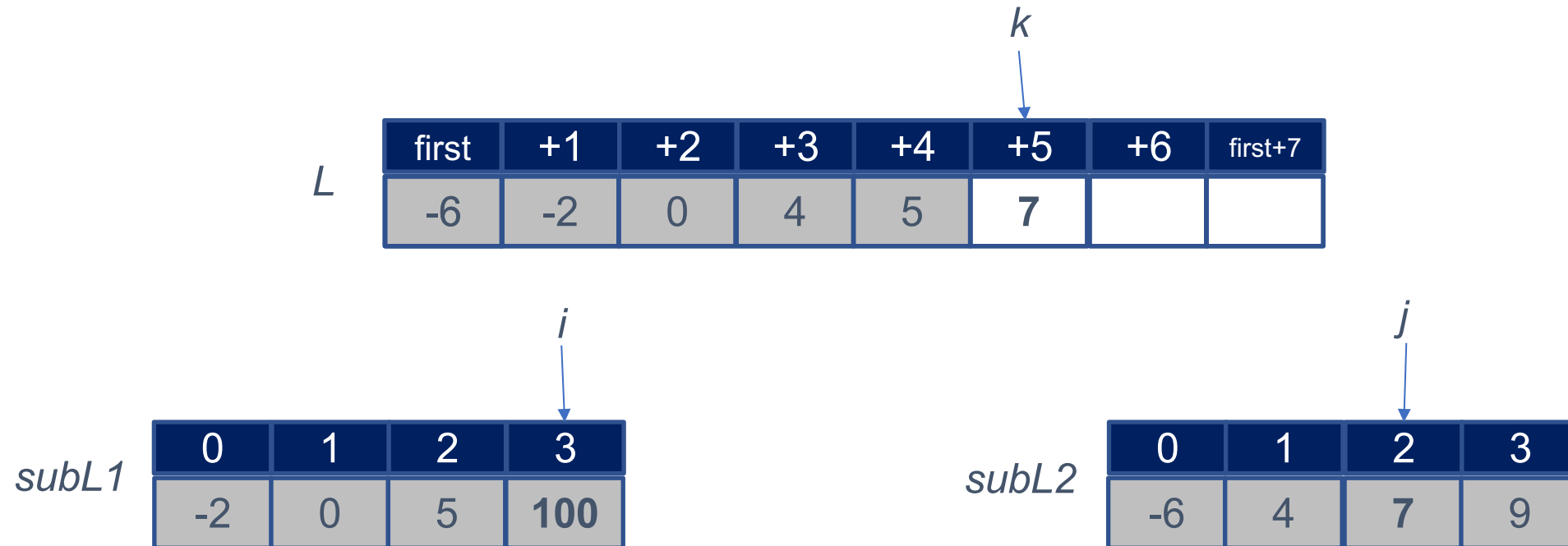
Merge Sort – Merge Algorithm

- Merge(L, first, mid, last)
 - Memory complexity: $O(\text{len}(L))$ for $\text{subL1}=L[\text{first}:\text{mid}+1]$ and $\text{subL2}=L[\text{mid}+1:\text{last}+1]$



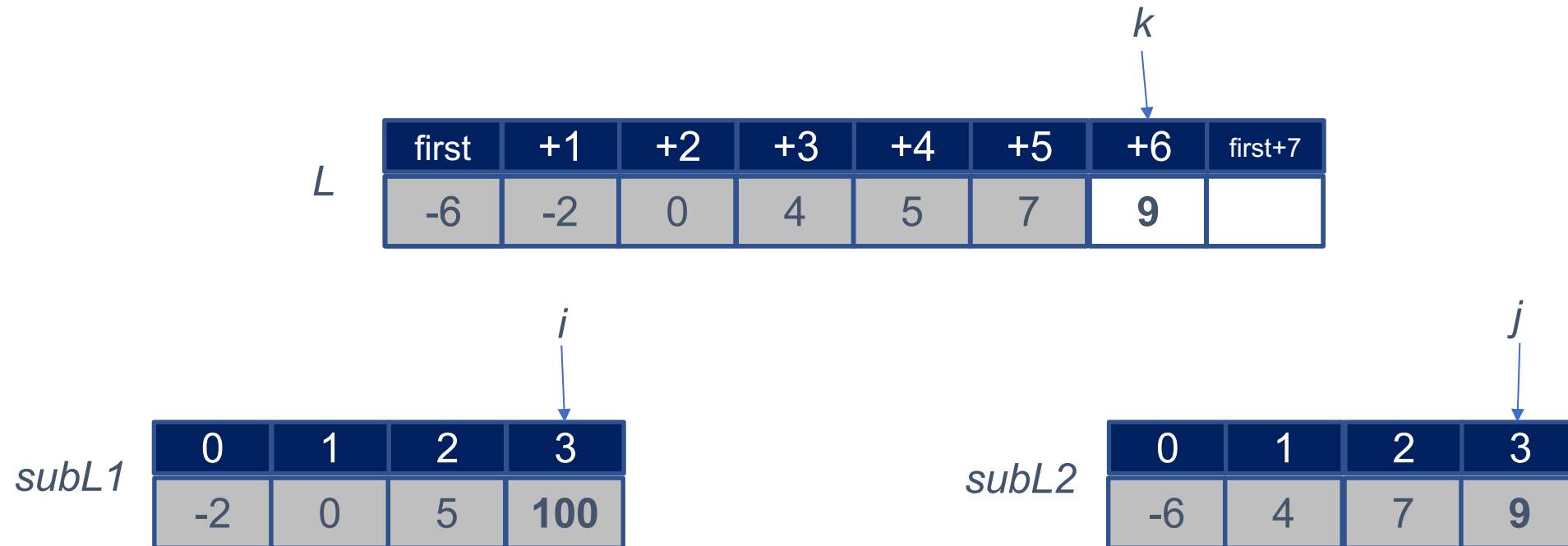
Merge Sort – Merge Algorithm

- Merge(L , first, mid, last)
 - Memory complexity: $O(\text{len}(L))$ for $\text{subL1}=L[\text{first}:\text{mid}+1]$ and $\text{subL2}=L[\text{mid}+1:\text{last}+1]$



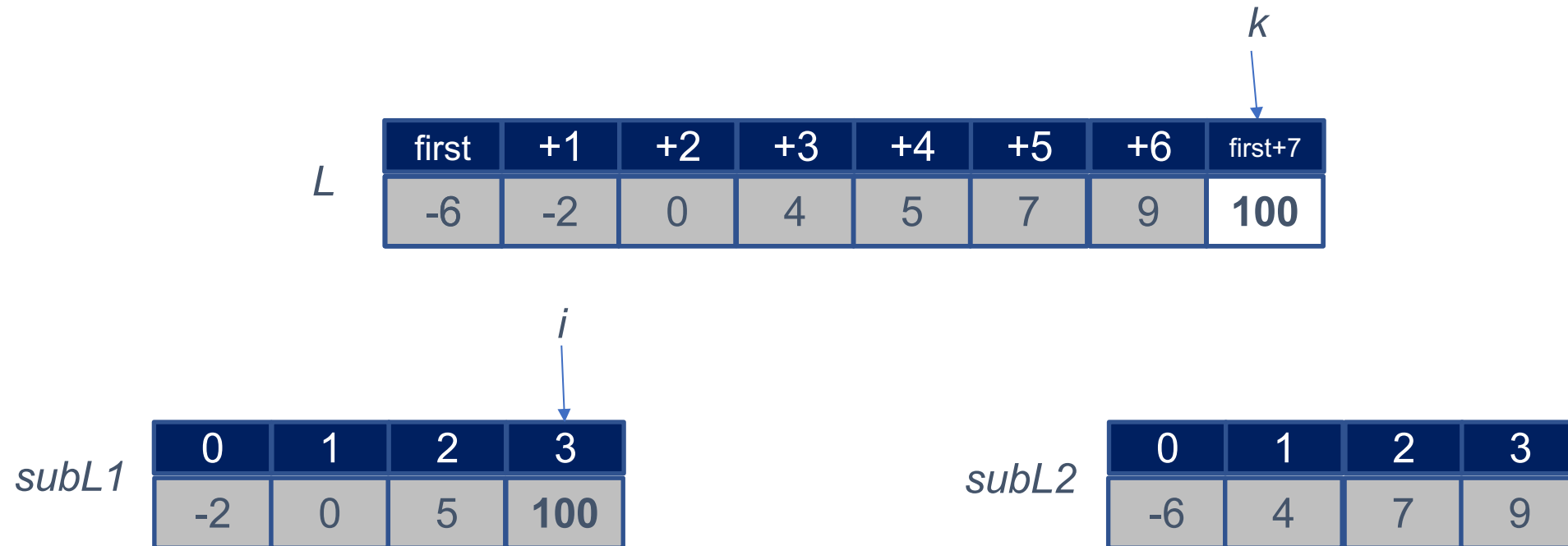
Merge Sort – Merge Algorithm

- Merge(L , first, mid, last)
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Merge Sort – Merge Algorithm

- Merge(L , first, mid, last)
 - Memory complexity: $O(\text{len}(L))$ for $\text{subL1}=L[\text{first}:\text{mid}+1]$ and $\text{subL2}=L[\text{mid}+1:\text{last}+1]$



Merge Sort – Merge Algorithm

- Merge(L, first, mid, last)
 - Memory complexity: $O(\text{len}(L))$ for $\text{subL1}=L[\text{first}:\text{mid}+1]$ and $\text{subL2}=L[\text{mid}+1:\text{last}+1]$
 - Time complexity of $O(\text{len}(L))$, instead of $O(\text{len}(L)^2)$

L

first	+1	+2	+3	+4	+5	+6	first+7
-6	-2	0	4	5	7	9	100

subL1

0	1	2	3
-2	0	5	100

subL2

0	1	2	3
-6	4	7	9

Merge Sort – Merge Code

```
• >>> def merge(L: list, first: int, mid: int, last: int) -> None:
•     ...     k = first
•     ...     sub1 = L[first:mid+1]
•     ...     sub2 = L[mid+1:last+1]
•     ...     i = j = 0
•     ...     while i < len(sub1) and j < len(sub2):
•         ...         if sub1[i] <= sub2[j]:
•             ...             L[k] = sub1[i]
•             ...             i = i+1
•         ...         else:
•             ...             L[k] = sub2[j]
•             ...             j = j+1
•         ...         k = k+1
•     ...     # Checking if any element is left
•     ...     if i < len(sub1):
•         ...         L[k:last+1] = sub1[i:]
•     ...     elif j < len(sub2):
•         ...         L[k:last+1] = sub2[j:]
```

Merge Sort – Time Complexity

$$\sim N \log_2 N$$

$$8 \times 1 = 8$$

0	1	2	3	4	5	6	7

$$8 \times \log_2^8 = 24$$

$$4 \times 2 = 8$$

0	1	2	3

0	1	2	3

$$2 \times 4 = 8$$

0	1

0	1

0	1

0	1

0

0

0

0

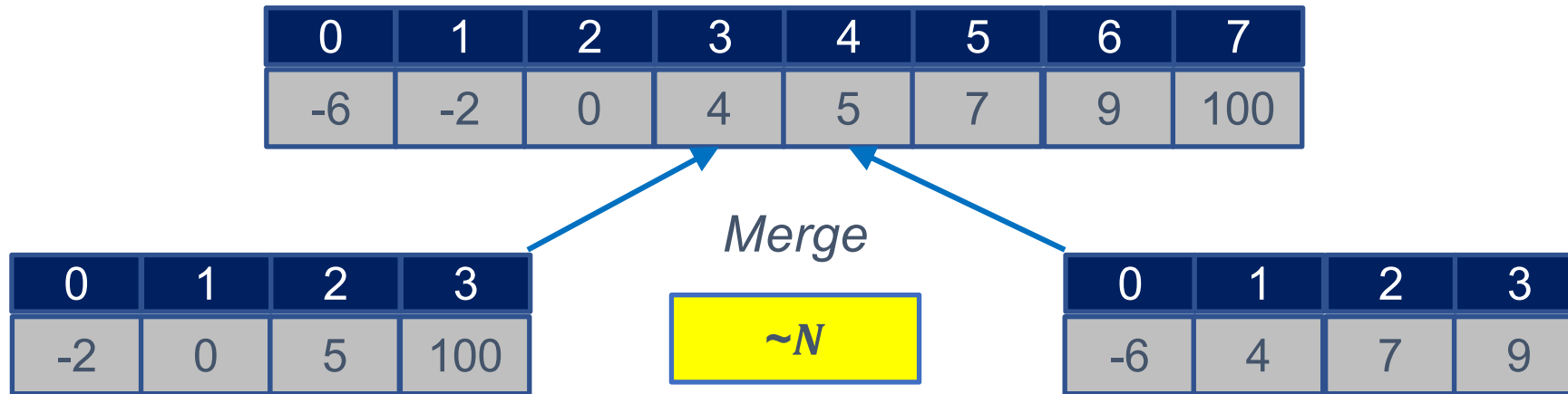
0

0

0

0

Merge Sort – Memory Complexity



Performance Comparison

- Despite messy implementation and somewhat complex logic, Merge Sort is much faster than selection/insertion sort ($N \log_2 N$ vs. N^2)
- Built-in sorting function is still faster, but its complexity **grows similar** to Merge Sort

List Length	Selection sort	Merge Sort	list.sort
1000	148	7	0.3
2000	583	15	0.6
3000	1317	23	0.9
4000	2337	32	1.3
5000	3699	41	1.6
10000	14574	88	3.5

So... What Sort Algorithm is Used for Python?

- **Tim Sort** in 2002 – a hybrid sorting algorithm (merge sort + insertion sort)
 - Divide and conquer like merge sort
 - When a sublist becomes smaller than a threshold, sort the sublist by using insertion sort
 - Insertion sort is faster than merge sort for a small list
- Visualization
 - <https://www.youtube.com/watch?v=NVIjHj-lrT4>

Summary

- Mergesort
 - Divide and recursive calls
 - Merge sorted sublists
- Time complexity $O(N \log N)$
 - $O(\log N)$ levels for recursion
 - $O(N)$ for merging at each level
- Memory complexity $O(N)$
 - For copying sublists

Thanks!