

1.编写一个割线法程序，求解下面的方程组：

- (1) $x^2 - e^x = 0$
- (2) $xe^x - 1 = 0$
- (3) $\lg x + x - 2 = 0$

解：

为了避免计算导数值，使用差商来替代导数，即有：

$$x_{k+1} = x_k - \frac{f(x_k)}{f(x_k) - f(x_{k-1})} \quad (k = 1, 2, 3 \dots)$$

解题过程中，需要首先估计 x_0, x_1 ，可以用几何法或估计法大致算出根的一个边界值，然后进行迭代，知道 x_k 趋近于某个值，并且与 x_{k-1} 的差小于允许误差 eps 即迭代完成

代码如下：

```
#include <bits/stdc++.h>
using namespace std;
```

```
double eps = 1e-5; // 允许误差范围
```

```
// (1)
```

```
double f1(double x)
```

```
{
    return pow(x, 2) - exp(x);
}
```

```
// (2)
```

```
double f2(double x)
```

```
{
    return x*exp(x) - 1;
}
```

```
// (3)
```

```
double f3(double x)
```

```

{
    return log10(x) - x - 2;
}

// 割线法迭代公式

double g(double x0, double x1, double (*f)(double))
{
    return x1 - f(x1) * (x1 - x0) / (f(x1) - f(x0));
}

int main()
{
    double x0 = 0, x1 = 1;
    int cnt = 0;
    while (fabs(x1 - x0) > eps)
    {
        double x2 = g(x0, x1, f2);
        x0 = x1;
        x1 = x2;
        cnt++;

        printf("迭代%d 次的结果为: %.5lf\n", cnt, x1);
    }
    return 0;
}

```

运行结果为：

```

迭代 1 次的结果为:0.36788
迭代 2 次的结果为:0.50331
迭代 3 次的结果为:0.57862
迭代 4 次的结果为:0.56653
迭代 5 次的结果为:0.56714
迭代 6 次的结果为:0.56714

```

2. 编写用改进的平方根法解方程组 $Ax = b$ 的程序，并解下列方程组：

$$A = \begin{bmatrix} 0.5 & -0.5 & 0 & 0 & 0 & 0 \\ -0.5 & 1.5 & -0.5 & -0.25 & 0.25 & 0 \\ 0 & -0.5 & 1.5 & 0.25 & -0.25 & 0 \\ 0 & -0.25 & 0.25 & 1.5 & -0.5 & 0 \\ 0 & 0.25 & -0.25 & -0.5 & 1.5 & -0.5 \\ 0 & 0 & 0 & 0 & -0.5 & 0.5 \end{bmatrix}$$

$$b = (-1, 0, 0, 0, 0, 0)^T$$

解：

当方程组的系数是对称正定时，可以使用改进平方根法改进LU矩阵分解，加速运行效率，对于方程组 $Ax = b$ ，令 $Ly = b$ ， $Ux = y$ ，其中 U, L 分别为上三角矩阵和下三角矩阵，则有： $A = LU$ ，再用回代即可求解

对于上三角矩阵 U 矩阵有：

$$u_{ki} = a_{ki} - \sum_{q=1}^{k-1} \frac{u_{qk}u_{qi}}{u_{qq}} \quad (i = k, k+1, \dots, n)$$

对于下三角矩阵 L 矩阵有：

$$l_{ik} = a_{ki} - \sum_{q=1}^{k-1} \frac{u_{qi}u_{qk}}{u_{qq}} = \frac{u_{ki}}{u_{kk}} \quad (i = k+1, k+2, \dots, n)$$

若矩阵 A 为对称正定矩阵，则 A 一定能直接作LU分解，且 $l_{ik} = \frac{u_{ki}}{u_{kk}} \quad (k = 1, 2, \dots, n)$

代码如下：

```
#include <bits/stdc++.h>
using namespace std;

// LU decomposition using vector
void LU(vector<vector<double>>& A, vector<vector<double>>& L,
vector<vector<double>>& U) {
    int n = A.size();
    for (int i = 0; i < n; i++) {
```

```

    L[i][i] = 1;
    for (int j = i; j < n; j++) {
        double sum = 0;
        for (int k = 0; k < i; k++) {
            sum += L[i][k] * U[k][j];
        }
        U[i][j] = A[i][j] - sum;
    }
    for (int j = i + 1; j < n; j++) {
        double sum = 0;
        for (int k = 0; k < i; k++) {
            sum += L[j][k] * U[k][i];
        }
        L[j][i] = (A[j][i] - sum) / U[i][i];
    }
}

// 打印 U 矩阵

cout << "U:" << endl;
for (int i = 0; i < n; i++) {
    for (int j = 0; j < U[i].size(); j++) {
        cout << U[i][j] << " ";
    }
    cout << endl;
}

// 打印 L 矩阵

cout << "L:" << endl;
for (int i = 0; i < n; i++) {
    for (int j = 0; j < L[i].size(); j++) {
        cout << L[i][j] << " ";
    }
    cout << endl;
}
}

// solve Ax = b using LU decomposition
vector<double> solve(vector<vector<double>> &A, vector<double> &b) {
    int n = A.size();
    vector<vector<double>> L(n, vector<double>(n, 0));

```

```

vector<vector<double>> U(n, vector<double>(n, 0));
LU(A, L, U);
vector<double> y(n, 0);
vector<double> x(n, 0);
// Ly = b
for (int i = 0; i < n; i++) {
    double sum = 0;
    for (int k = 0; k < i; k++) {
        sum += L[i][k] * y[k];
    }
    y[i] = b[i] - sum;
}
// Ux = y
for (int i = n - 1; i >= 0; i--) {
    double sum = 0;
    for (int k = i + 1; k < n; k++) {
        sum += U[i][k] * x[k];
    }
    x[i] = (y[i] - sum) / U[i][i];
}
return x;
}

// main function
int main() {
    vector<vector<double>> A = {{0.5, -0.5, 0, 0, 0, 0}, {-0.5, 1.5, -0.5, -0.25, 0.25,
0},
    {0, -0.5, 1.5, 0.25, -0.25, 0}, {0, -0.25, 0.25, 1.5, -0.5, 0},
    {0, 0.25, -0.25, -0.5, 1.5, -0.5}, {0, 0, 0, 0, -0.5, 0.5}};
    vector<double> b = {-1, 0, 0, 0, 0, 0};
    vector<double> x = solve(A, b);

    cout << "x:" << endl;
    for (int i = 0; i < x.size(); i++) {
        cout << x[i] << endl;
    }
    return 0;
}

```

运行结果如下：

上三角 U 矩阵：

```
0.5 -0.5 0 0 0 0
0 1 -0.5 -0.25 0.25 0
0 0 1.25 0.125 -0.125 0
0 0 0 1.425 -0.425 0
0 0 0 0 1.29825 -0.5
0 0 0 0 0 0.307432
```

下三角 L 矩阵：

```
1 0 0 0 0 0
-1 1 0 0 0 0
0 -0.5 1 0 0 0
0 -0.25 0.1 1 0 0
0 0.25 -0.1 -0.298246 1 0
0 0 0 0 -0.385135 1
```

x：

```
-3.25275
-1.25275
-0.373626
-0.0879121
0.175824
0.175824
```

3. 编写一个用牛顿前插公式计算函数值的程序，要求先输出差分表，再计算 x 点的函数值，并应用于下面的问题：

x_i	20	21	22	23	24
y_i	1.30103	1.32222	1.34242	1.36173	1.38021

求 $x = 21.4$ 时的三次插值多项式的值

解：

先迭代求出 k 次差分表，再代入牛顿前插公式中：

$$N_n(x_0 + th) = \sum_{k=0}^n \frac{\Delta^k f_0}{k!} \prod_{j=0}^{k-1} (t - j)$$

代码如下：

```
#include <iostream>
#include <vector>

double newtonInterpolation(double x, const std::vector<double>& xi, const
std::vector<double>& yi) {
    int n = xi.size();
    std::vector<std::vector<double>> diffTable(n, std::vector<double>(n));
    for (int i = 0; i < n; i++) {
        diffTable[i][0] = yi[i];
    }

    for (int j = 1; j < n; j++) {
        for (int i = 0; i < n - j; i++) {
            diffTable[i][j] = (diffTable[i + 1][j - 1] - diffTable[i][j - 1]) / (xi[i + j] -
xi[i]);
        }
    }

    std::cout << "差分表： " << std::endl;

    for (int i = 0; i < n; i++) {
        std::cout << xi[i] << "\t";
        for (int j = 0; j <= i; j++) {
            std::cout << diffTable[i - j][j] << "\t";
        }
    }
}
```

```

        std::cout << std::endl;
    }
    std::cout << std::endl;
    double result = 0.0;
    double prod = 1.0;

    for (int i = 0; i < n; i++) {
        prod = diffTable[i][i];
        for (int j = 0; j < i; j++) {
            prod *= (x - xi[j]);
        }
        result += prod;
    }

    return result;
}

int main() {
    std::vector<double> xi = { 20, 21, 22, 23, 24 };
    std::vector<double> yi = { 1.30103, 1.32222, 1.34242, 1.36173, 1.38021 };

    double x = 21.4;
    double interpolationResult = newtonInterpolation(x, xi, yi);

    std::cout << "在 x = " << x << " 时的三次插值多项式的值为: " <<
    interpolationResult << std::endl;

    return 0;
}

```

运行结果如下:

差分表:

20	1.30103				
21	1.32222	0.02119			
22	1.34242	0.0202	-0.000495		
23	1.36173	0.01931	-0.000445	1.66667e-05	
24	1.38021	0.01848	-0.000415	1e-05	-1.66667e-06

在 x = 21.4 时的三次插值多项式的值为: 1.32908

4. 编写求超定方程组的最小二乘法解的程序，并解下列方程组：

$$\begin{cases} 2x + 4y = 10 \\ 3x - 5y = -13 \\ 10x - 12y = -26 \\ 4x + 11y = 25 \end{cases}$$

解：

对于方程组 $Ax = b$ ，可以写成 $A^T Ax = A^T b$ ，可以证明如果 A 是满秩的，则方程组 $A^T Ax = A^T b$ 存在唯一解，并把该方程组的解称为超定方程组的最小二乘法解

代码如下：

```
#include <bits/stdc++.h>
using namespace std;

// 最小二乘法解超定方程组
vector<double> leastSquare(vector<vector<double>> &A, vector<double> &b) {
    int n = A.size();
    int m = A[0].size();
    vector<vector<double>> AT(m, vector<double>(n, 0));
    vector<vector<double>> ATA(m, vector<double>(m, 0));
    vector<double> ATb(m, 0);
    for (int i = 0; i < m; i++) {
        ATb[i] = 0;
        for (int j = 0; j < n; j++) {
            AT[i][j] = A[j][i];
            ATb[i] += A[j][i] * b[j];
        }
    }
    for (int i = 0; i < m; i++) {
        for (int j = i; j < m; j++) {
            ATA[i][j] = 0;
            for (int k = 0; k < n; k++) {
                ATA[i][j] += AT[i][k] * A[k][j];
            }
            ATA[j][i] = ATA[i][j];
        }
    }
    vector<double> x(m, 0);
    for (int i = 0; i < m; i++) {
        double sum = 0;
        for (int j = 0; j < m; j++) {
            sum += ATA[i][j] * x[j];
        }
        x[i] = (ATb[i] - sum) / ATA[i][i];
    }
    return x;
}
```

```

    }
}
vector<double> x(m, 0);
// solve ATA * x = ATb
for (int i = 0; i < m; i++) {
    double sum = 0;
    for (int k = 0; k < i; k++) {
        sum += ATA[i][k] * x[k];
    }
    x[i] = (ATb[i] - sum) / ATA[i][i];
}
return x;
}

// main function
int main() {
    int n, m;
    cin >> n >> m;
    vector<vector<double>> A = {{2, 4}, {8, 4}, {2, 1}, {7, -1}, {4, 0}};
    vector<double> b = {10, -13, -26, 25};
    vector<double> x = leastSquare(A, b);
    for (int i = 0; i < x.size(); i++) {
        cout << x[i] << " ";
    }
    cout << endl;
    return 0;
}

```

运行结果如下：

```
0.284672 -2.14599
```