

Giuseppe Peano's
Classic Mathematical Text

Arithmetices principia, nova methodo exposita

OR

“The Principles of Arithmetic, Presented by a New Method”

presented

in the original Latin

AND

in parallel English Translation

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This document is hosted at https://github.com/mdnahas/Peano_Book

About the Translation

Below is Giuseppe Peano's *Arithmetices principia* as first published¹, i.e. as "*Arithmetices principia, nova methodo exposita*"², which appeared translated to English³ in 1967 as "*The principles of arithmetic, presented by a new method*"⁴, as well as in 1973⁵. This present document⁶ is the only (to my knowledge) side-by-side Latin-English translation of the Latin original. The mathematical notation (in the English, right, column) got updated to currently canonically-used or easy-to-decrypt symbols in the international and/or English mathematical community; which is also a feature currently unseen in any reprint.

Red text is mathematical commentary.

Gray text is irrelevant for modern mathematical notation.

Dashed lines (.....) indicate pages in the original treatise.

¹H. Kennedy, *Peano. Life and Works of Giuseppe Peano*, San Francisco: Peremptory Publications, 2002, p. 41.

²G. Peano, *Arithmetices principia, nova methodo exposita*, Bocca, Torino, 1889.

³These English translations listed are the only ones (to my knowledge) and all the English translations listed in:

I. Grattan-Guinness (ed.), *Landmark Writings in Western Mathematics 1640-1940*, Amsterdam: Elsevier, 2005, p. 614.

⁴G. Peano, (1889), "The principles of arithmetic presented by a new method" in: J. van Heijenoort (ed.), *From Frege to Gödel. A source book in mathematical logic. 1879-1931*, Cambridge: Harvard University Press, 1967, p. 83-97.

⁵G. Peano, *Selected works of Giuseppe Peano*, H. Kennedy (ed.), London: George Allen & Unwin, 1973, p. 101-134.

⁶Written by Vincent Verheyen. Last updated on 17/8/2015. I encourage you to use your reason for good. If you want my support, please contact me via <http://vincentverheyen.com/contact>. It is possible to contribute to the flourishing of knowledge, even when you have an intelligence like mine. Thank you and good luck studying.

I would like to thank Mauro Allegranzo and acknowledge his support of this work and his various comments during its creation.

Arithmetices principia

Nova methodo exposita

a

Ioseph Peano

in R. Academia militari professore

Analysin infinitorum in R. Taurinensi Athenæo docente.

Labor et honor

Augustae Taurinorum

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The principles of arithmetic

Presented by a new method

by

Giuseppe Peano

professor at the Royal Military Academy

teaching Analysis of the infinite at the Royal Turin Athenaeum.

Work and honor

At Turin

Published by Libreria Bocca

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Via del Corso, 216-217.

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PRAEFATIO

Quaestiones, quae ad mathematicae fundamenta pertinent, etsi hisce temporibus a multis tractatae, satisfacienti solutione et adhuc carent. Hic difficultas maxime ex sermonis ambiguitate oritur.

Quare summi interest verba ipsa, quibus utimur attente perpendere. Hoc examen mihi proposui, atque mei studii resultatus, et arithmeticae applicationes in hoc scripto expono.

Ideas omnes quae in arithmeticae principiis occurrunt, signis indicavi, ita ut quaelibet propositio his tantum signis enuncietur.

Signa aut ad logicam pertinent, aut proprie ad arithmetica. Logicae signa quae hic occurrunt, sunt numero ad decem, quamvis non omnia necessaria. Horum signorum usus et proprietas nonnullae in priore parte communi sermone explicantur. Ipsorum theoriam fusius hic exponere nolui. Arithmeticae signa, ubi occurrunt, explicantur.

His notationibus quaelibet propositio formam assumit atque praecisionem, qua in algebra aequationes gaudent, et a propositionibus ita scriptis aliea deducuntur, idque processis qui aequationum resolutioni assimilantur. Hoc caput totius scripti.

Sique, confectis signis quibus arithmeticae propositiones scribere possim, in earum tractatione usus sum methodo, quam quia et in aliis studiis sequenda foret, breviter exponam.

Ex arithmeticae signis quae caeteris, una cum logicae signis exprimere licet, ideas

PREFACE

Questions pertaining to the foundations of mathematics, although treated by many these days, still lack a satisfactory solution. The difficulty arises principally from the ambiguity of ordinary language.

For this reason it is of the greatest concern to consider attentively the words we use. I resolved to do this, and am presenting in this paper the results of my study with applications to arithmetic.

I have indicated by signs all the idea which occur in the fundamentals of arithmetic, so that every proposition is stated with just these signs.

The signs pertain either to logic or to arithmetic. The signs of logic that occur here are about ten in number, although not all are necessary. The use of these signs and several of their properties are explained in ordinary language in the first part. I did not wish to present their theory more fully here. The signs of arithmetic are explained as they occur.

With this notation every proposition assumes the form and precision equations enjoy in algebra, and from propositions so written others may be deduced, by a process which resembles the solution of algebraic equations. That is the chief reason for writing this paper.

Having made up the signs with which I can write arithmetical propositions, in treating them I have used a method which, because it is to be followed in later studies, I shall present briefly.

Those arithmetical signs which may be expressed by using others along with signs

significant quas definire possumus. Ita omnia definivi signa, si quatuor excipias, quae in explicationibus §1 continentur. Si, ut puto, haec ulterius reduci nequeunt, ideas ipsis expressas, ideis quae prius notae supponuntur, definire non licet.

Propositiones, quae logicae operationibus a caeteris deducuntur, sunt *theoremata*; quae vero non, *axiomata* vocavi. Axiomata hic sunt novem (§1), et signorum, quae definitione carent, proprietates fundamentales exprimunt.

In §1-6 numerorum proprietates communes demonstravi; brevitatis causa, demonstrationes praecedentibus similes omisi; demonstrationum communem formam immutare oportet ut logicae signis exprimantur; haec transformatio interdum difficilior est, tamen inde demonstrationis natura clarissime patet.

In sequentibus § varia tractavi, ut huius methodi potentia magis videatur.

In §7 nonnulla theoremata, quae ad numerorum theoriam pertinent, continentur.

In §8 et 9 rationalium et irrationalium definitiones inveniuntur.

Denique, in §10, theoremata exposui nonnulla, quae nova esse puto, ad entium theoriam pertinentia, quae cl.^{mus} Cantor *Punktmenge* (*ensemble de points*) vocavit.

In hoc scripto aliorum studiis usus sum. Logicae notationes et propositiones quae in num. II, III et IV continentur, si nonnullas excipias, ad multorum opera, inter quae Boole praecipue, referenda sunt.¹

Signum ϵ , quod cum signo \supset confundere non licet, inversionis in logica applica-

of logic represent the ideas we can define. Thus I have defined every sign, if you except the four which are contained in the explanations of §1. If, as I believe, these cannot be reduced further, then the ideas expressed by them may not be defined by ideas already supposed to be known.

IV

Propositions which are deduced from others by the operations of logic are theorems; those for which this is not true I have called axioms. There are nine axioms here (§1), and they express fundamental properties of the undefined signs.

In §1-6 I have proved the ordinary properties of numbers; for the sake of brevity, I have omitted proofs which are similar to preceding ones. The ordinary form of proofs has had to be altered in order that they may be expressed with the signs of logic. This transformation is sometimes rather difficult but the nature of the proof then becomes quite clear. In the following sections I have treated various things so that the power of the method is better seen.

In the following sections I have treated various things so that the power of the method is better seen.

In §7 are several theorems pertaining to the theory of numbers. In §8 and 9 are found the definitions of rationals and irrationals.

Finally, in §10 I have given several theorems, which I believe to be new, pertaining to the theory of those entities which Professor Cantor has called *Punktmenge* (*ensemble de points*).

In this paper I have used the research of others. The notations and propositions of logic which are contained in numbers II, III, and IV, with some exceptions, represent the work of many, among them Boole especially.¹

V

The sign ϵ , which must not be confused with the sign \subset , applications of the inverse

tionem, et paucas alias institui conventiones, ut ad exprimendam quamlibet propositionem pervenirem.

In arithmeticae demonstrationibus usus sum libro: H. Grassmann, *Lehrbuch der Arithmetik*, Berlin 1861.

Utilius quoque mihi fuit recens scriptum: R. Dedekind, *Was sind und was sollen die Zahlen*; Braunschweig, 1888, in quo quaestiones, quae ad numerorum fundamenta pertinent, acute examinantur.

Hic meus libellus ut novae methodi specimen habendus est. Hisce notationibus innumeras alias propositiones, ut quae ad rationales et irrationales pertinent, enunciare et demonstrare possumus. Sed, ut aliae theoriae tractentur, nova signa, quae nova indicant entia, instituere necesse est. Puto vero his tantum logicae signis propositiones cuiuslibet scientiae exprimi posse, dummodo adiungantur signa quae entia huius scientiae representant.

in logic, and a few other conventions, I have adopted so that I could express any proposition whatever.

In the proofs of arithmetic I used the book H. Grassmann, *Lehrbuch der Arithmetik* (Berlin 1861).

Also quite useful to me was the recent work by R. Dedekind, *Was sind und was sollen die Zahlen* (Braunschweig, 1888), in which questions pertaining to the foundations of numbers are acutely examined.

My booklet should be taken as a sample of this new method. With these notations we can state and prove innumerable other propositions, such as those which pertain to rationals and irrationals. But in order to treat other theories, it is necessary to adopt new signs to indicate new entities. I believe, however, that with only these signs of logic the propositions of any science can be expressed, so long as the signs which represent the entities of the science are added.

Signorum tabula

Logicam signa

Signum	Significatio	Pag.
<i>P</i>	propositio	VII
<i>K</i>	classis	X
\cap	et	VII, X
\cup	vel	VIII, X, XI
$-$	non	VIII, X
Λ	absurdum <i>aut</i> nihil	VIII, XI
\supset	deducitur <i>aut</i> continetur	VIII, XI

Table of signs

Signs of logic

Sign	Meaning	Pag.
proposition	proposition	VII
CLASS	class	X
\wedge	and	VII, X
\vee	or	VIII, X, XI
\neg	not	VIII, X
\perp	false <i>or</i>	VIII, XI
\emptyset	nothing	
\rightarrow	one deduces <i>or</i>	VIII, XI

=	est aequalis	VIII
ϵ	est	X
[]	inversionis signum	XI
ϑ	qui <i>vel</i> [e]	XII
Th	Theorema	XVI
Hp	Hypothesis	
Ts	Thesis	
L	Logica	

Arithmeticae signa

Signum	Significatio	Pag.
Signa 1, 2, ..., =, >, <, +, -, \times vulgarem habent significationem. Divisionis signum est /.		
N	numerus integer positivus	1
R	num. rationalis positivus	12
Q	quantitas, <i>sive</i> numerus realis positivus	16
N_p	numerus primus	9
M	maximus	6
W	minimus	6
T	terminus, <i>vel</i> limes summus	15
D	dividit	9
\mathcal{C}	est multiplex	9
π	est primus cum	6

Signa composita

\subset	is contained in	
=	equals	VIII
\in	is (an element of) <i>or</i>	X
is a	is (a)	
$\{x \dots\}$	sign of the inverse	XI
\rightarrow	such that <i>or</i>	XII
isn't	isn't	
theorem	Theorem	XVI
hypothesis	Hypothesis	
thesis	Thesis	
logic	Logic	

Signs of arithmetic

Sign	Meaning	Pag.
The signs 1, 2, ..., =, >, <, +, -, * have their usual meaning. The sign of division is /.		
\mathbb{N}	positive integers	1
\mathbb{R}	postive rational numbers	12
\mathbb{Q}	quantity <i>or</i> postive real numbers	16
\mathbb{P}	prime number	9
M^{aximum}	maximum	6
m_{inimum}	minimum	6
$G^{\text{reatest bound}}$	terminus <i>or</i> greatest bound	15
divides	divides	9
is divisible	is divisible	9
is \mathbb{P} prime with	is prime with	6

Composite signs

$- <$	non est minor
$= \cup >$	est aequalis aut maior
ϑD	divisor
$M \vartheta D$	maximus divisor

\nless	is not less than
\geq	is equal to or greater than
is a divisor	is a divisor
is the greatest divisor	is the greatest divisor

VII

Logicae notationes.

I. De punctuatione.

Litteris $a, b, \dots x, y, \dots x', y', \dots$ entia indicamus indeterminata quaecumque. Entia vero determinata signis, sive litteris P, K, N, \dots indicamus.

Signa plerumque in eadem linea scribemus. Ut ordo pateat quo ea coniungere oporteat, *parenthesibus* ut in algebra, sive *punctis* $. : \cdot \vdots ::$ etc. utimur.

Ut formula punctis divisa, intelligatur, primum signa quae nullo puncto seperantur colligenda sunt, postea quae uno puncto, deinde quae duobus punctis, etc.

Ex. g. sint a, b, c, \dots signa quaecumque. Tunc $ab \cdot cd$ significat $(ab)(cd)$; et $ab \cdot cd : ef \cdot gh \cdot k$ significat $((ab)(cd))(ef)(gh))k$.

Punctuationis signa omittere licet si formulae quae diversa punctuatione existent eundem habeant sensum; vel si una tantum formula, et ipsa quam scribere volumus, sensum habeat.

Ut ambiguitatis periculum absit, aritmeticae operationum signis $. : \cdot$ nunquam utimur.

Notations of logic.

I. Punctuation.

By the letters $a, b, \dots x, y, \dots x', y', \dots$ we indicate any indeterminate entities. Determine entities are, however, indicated by the signs, or rather by the letters, P, K, N, \dots

Generally we write signs on the same line. So that it will be clear how they are to be joined, we use *parentheses*, as in algebra, or rather *points* $. : \cdot \vdots ::$ etc.

So that a formula divided by points may be understood, first the signs which are not seperated by points are taken together, then those seperated by one point, then those by two points, etc.

For example, let a, b, c, \dots be any signs. Then $(ab)(cd)$ means $(ab)(cd)$; and $\{[(ab)(cd)][(ef)(gh)]\}k$ means $\{[(ab)(cd)][(ef)(gh)]\}k$.

The signs of punctuation may be omitted if formulas having different punctuation have the same meaning, or if just one formula, that being the one we wish to write, has meaning.

To avoid the danger of ambiguity, we never use $. : \cdot$ as signs of arithmetical operations.

Parenthesum figura una est (); si in eadem formula, parentheses et puncta occurrant, primum quae parenthesibus continentur, colligantur.

II. De propositionibus.

Signo P significatur *propositio*.

Signum \cap legitur *et*. Sint a, b propositiones; tunc $a \cap b$ est simultanea affirmatio propositionum a, b . Brevitatis causa, loco $a \cap b$ vulgo scribemus $a \cdot b$.

Signum $-$ legitur *non*. Sit a quaedam P ; tunc $-a$ est negatio propositionis a .

Signo \cup legitur *vel*. Sint a, b propositiones; tunc $a \cup b$ idem est ac $- : -a \cdot -b$.

[Signo V significatur *verum*, sive *identitas*; sed hoc signo numquam utimur].

Signum Λ significat *falsum*, sive *absurdum*.

[Signum C significat *est consequentia*; ita $b \ C \ a$ legitur *b est consequentia propositionis a*. Sed hoc signo nunquam utimur].

Signum \supset significat *deducitur*; ita $a \supset b$ significat quod $b \ C \ a$. Si propositiones a, b entia indeterminata continent x, y, \dots , scilicet sunt inter ipsa entia conditiones, tunc $a \supset_{x,y,\dots} b$ significat: quaecumque sunt x, y, \dots , a propositione a deducitur b .

Si vero ambiguitatis periculum absit, loco $\supset_{x,y,\dots}$, scribemus solum \supset .

Signum $=$ significat *est aequalis*. Sint a, b propositiones; tunc $a = b$ idem significat quod $a \supset b \cdot b \supset a$; propositio $a =_{x,y,\dots} b$ idem significat quod $a \supset_{x,y,\dots} b \cdot b \supset_{x,y,\dots} a$.

III. Logicae propositiones.

Sint a, b, c, \dots propositiones. Tunc erit:

The figure of parentheses is one (); if parentheses and points occur in the same formula, whatever is contained in parentheses is to be gathered first.

II. Propositions.

The sign P means *proposition*.

The sign \wedge is read *and*. Let a, b , be propositions; then $a \wedge b$ is the simultaneous affirmation of the propositions a, b . For the sake of brevity, instead of $a \wedge b$, we ordinarily write $a \cdot b$.

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The sign \neg is read *not*. Let a be a P ; then $\neg a$ is the negation of the proposition a .

The sign \vee is read *or*. Let a, b be propositions; then $a \vee b$ is the same as $\neg[(\neg a) \wedge (\neg b)]$.

The sign \top means *true*, or *identity*, but we never use this sign.

The sign \perp means *false*, or *absurd*.

The sign \leftarrow means *is a consequence of*. Thus $b \leftarrow a$ is read *b is a consequence of the proposition a*. But we never use this sign.

The sign \rightarrow means *one deduces*; thus $a \rightarrow b$ means the same as $b \leftarrow a$. If the propositions a, b contain the indeterminate quantities x, y, \dots , that is, express conditions on these objects, then $a \xrightarrow{x,y,\dots} b$ means: whatever the x, y, \dots , from propositions a one deduces b . If indeed there is no danger of ambiguity, instead of $\xrightarrow{x,y,\dots}$, we write only \rightarrow .

The sign $=$ means *equals*. Let a, b be propositions; then $a = b$ means the same as $(a \rightarrow b) \wedge (b \rightarrow a)$; proposition $a =_{x,y,\dots} b$ means the same as $(a \xrightarrow{x,y,\dots} b) \wedge (b \xrightarrow{x,y,\dots} a)$.

III. Propositions of logic.

Let a, b, c, \dots be propositions. We have:

1. $a \supset a$
2. $a \supset b . b \supset c : \supset : a \supset c$
3. $a = b . = : a \supset b . b \supset a .$
4. $a = a$
5. $a = b . = . b = a$
6. $a = b . b \supset c : \supset . a \supset c$
7. $a \supset b . b = c : \supset . a \supset c$
8. $a = b . b = c : \supset . a = c$
9. $a = b . \supset . a \supset b$
10. $a = b . \supset . b \supset a$
11. $ab \supset a$
12. $ab = ba$
13. $a(bc) = (ab)c = abc$

-
14. $aa = a$
 15. $a = b . \supset . ac = bc$
 16. $a \supset b . \supset . ac \supset bc$
 17. $a \supset b . c \supset d : \supset . ac \supset bd$
 18. $a \supset b . a \supset c : = . a \supset bc$
 19. $a = b . c = d : \supset . ac = bd$

20. $\neg(\neg a) = a$
21. $a = b . = . \neg a = \neg b .$
22. $a \supset b . = . \neg b \supset \neg a =$

23. $a \cup b . = \therefore - : \neg a . \neg b$
24. $\neg(ab) = (\neg a) \cup (\neg b)$

1. $a \rightarrow a$
2. $[(a \rightarrow b) \wedge (b \rightarrow c)] \rightarrow (a \rightarrow c)$
3. $(a = b) = [(a \rightarrow b) \wedge (b \rightarrow a)]$
4. $a = a$
5. $(a = b) = (b = a)$
6. $[(a = b) \wedge (b \rightarrow c)] \rightarrow (a \rightarrow c)$
7. $[(a \rightarrow b) \wedge (b = c)] \rightarrow (a \rightarrow c)$
8. $[(a = b) \wedge (b = c)] \rightarrow (a = c)$
9. $(a = b) \rightarrow (a \rightarrow b)$
10. $(a = b) \rightarrow (b \rightarrow a)$
11. $(a \wedge b) \rightarrow a$
12. $(a \wedge b) = (b \wedge a)$
13. $(a \wedge (b \wedge c)) = ((a \wedge b) \wedge c) = (a \wedge b \wedge c)$

-
14. $(a \wedge a) = a$
 15. $(a = b) \rightarrow [(a \wedge c) = (b \wedge c)]$
 16. $(a \rightarrow b) \rightarrow [(a \wedge c) \rightarrow (b \wedge c)]$
 17. $[(a \rightarrow b) \wedge (c \rightarrow d)] \rightarrow [(a \wedge c) \rightarrow (b \wedge d)]$
 18. $[(a \rightarrow b) \wedge (a \rightarrow c)] = [(a \rightarrow (b \wedge c))]$
 19. $[(a = b) \wedge (c = d)] \rightarrow [(a \wedge c) = (b \wedge d)]$

20. $\neg(\neg a) = a$
21. $(a = b) = [(\neg a) = (\neg b)]$
22. $(a \rightarrow b) = [(\neg b) \rightarrow (\neg a)]$

23. $(a \vee b) = \{[(\neg a) \wedge (\neg b)]\}$
24. $[\neg(a \wedge b)] = [(\neg a) \vee (\neg b)]$

IX

25. $\neg(a \cup b) = (\neg a)(\neg b)$
26. $a \supset . a \cup b$
27. $a \cup b = b \cup a$
28. $a \cup (b \cup c) = (a \cup b) \cup c = a \cup b \cup c$
29. $a \cup a = a$
30. $a(b \cup c) = ab \cup ac$
31. $a = b . \supset . a \cup c = b \cup c$
32. $a \supset b . \supset . a \cup c \supset b \cup c$
33. $a \supset b . c \supset d : \supset : a \cup c . \supset . b \cup d$
34. $b \supset a . c \supset a : = . b \cup c \supset a$

35. $a - a = \Lambda$
36. $a \Lambda = \Lambda$
37. $a \cup \Lambda = a$
38. $a \supset \Lambda . = . a = \Lambda$
39. $a \supset b . = . a - b = \Lambda$
40. $\Lambda \supset a$
41. $a \cup b = \Lambda . : : a = \Lambda . b = \Lambda$

42. $a \supset . b \supset c : = : ab \supset c$
43. $a \supset . b = c : = . ab = ac$

.....

Sit α quoddam relationis signum (ex. gr. $=$, \supset), ita ut $a \alpha b$ sit quaedam propositio. Tunc loco $\neg . a \alpha b$ scribemus $a - \alpha b$; scilicet:

$$a - = b . = : - . a = b$$

$$a - \supset b . = : - . a \supset b$$

Ita signum $- =$ significat *non est aequalis*. Si propositio a indeterminatum continet x , $a - =_x \Lambda$ significat: sunt x quae conditioni a satisfaciunt. Signum $- \supset$

25. $[\neg(a \vee b)] = [(\neg a) \vee (\neg b)]$
26. $a \rightarrow (a \vee b)$
27. $(a \vee b) = (b \vee a)$
28. $[a \vee (b \vee c)] = [(a \vee b) \vee c] = (a \vee b \vee c)$
29. $(a \vee a) = a$
30. $[a \wedge (b \vee c)] = [(a \wedge b) \vee (a \wedge c)]$
31. $(a = b) \rightarrow [(a \vee c) = (b \vee c)]$
32. $(a \rightarrow b) \rightarrow [(a \vee c) \rightarrow (b \vee c)]$
33. $[(a \rightarrow b) \wedge (c \rightarrow d)] \rightarrow [(a \vee c) \rightarrow (b \vee d)]$
34. $[(b \rightarrow a) \wedge (c \rightarrow a)] = [(b \vee c) \rightarrow a]$

35. $[a \wedge (\neg a)] = \perp$
36. $(a \wedge \perp) = \perp$
37. $(a \vee \perp) = a$
38. $(a \rightarrow \perp) = (a = \perp)$
39. $(a \rightarrow b) = [(a \wedge \neg b) = \perp]$
40. $\perp \rightarrow a$
41. $[(a \vee b) = \perp] = [(a = \perp) \wedge (b = \perp)]$

42. $[a \rightarrow (b \rightarrow c)] = [(a \wedge b) \rightarrow c]$
43. $[a \rightarrow (b = c)] = [(a \wedge b) = (a \wedge c)]$

X

.....

Let α be the sign of some relation (eg., $=$, \rightarrow) so that $a \alpha b$ is a proposition. Then instead of $\neg(a \alpha b)$, we write $a \not\alpha b$. Thus:

$$(a \neq b) = \neg(a = b)$$

$$(a \not\rightarrow b) = \neg(a \rightarrow b)$$

Thus the sign \neq means *is not equal to*. If the proposition a contains the indeterminate x , $a \neq_x \perp$ means: there is an x which satisfies condition a . The sign $\not\rightarrow$ means

significat *non deducitur*.

Similiter, si α et β sunt relationis signa, loco $a \alpha b$, et $a \alpha b \cup a \beta b$ scribere possumus $a \cdot \alpha \beta \cdot b$ et $a \cdot \alpha \cup \beta \cdot b$. Ita, si a et b sunt propositiones, formula $a \cdot \supset - = \cdot b$ dicit: ab a *deducitur* b , sed non vice versa.

$$a \cdot \supset - = \cdot b : = : a \supset b \cdot b -$$

\supset

a

Formulae:

$$a \supset b \cdot b \supset c \cdot a - \supset c : = \Lambda$$

$$a = b \cdot b = c \cdot a - = c : = \Lambda$$

$$a \supset b \cdot b \supset - = c : \supset \cdot a \supset - = c$$

$$a \supset - = b \cdot b \supset c : \supset \cdot a \supset - = c$$

Sed his notationibus raro utimur.

IV. De classibus.

Signo \mathbf{K} significatur *classis*, sive entium aggregatio.

Signum \in significat *est*. Ita $a \in b$ legitur *a est quoddam b*; $a \in \mathbf{K}$ significat *a est quaedam classis*; $a \in \mathbf{P}$ significat *a est quaedam propositio*.

Loco $\neg(a \in b)$ scribemus $a - \in b$; signum $- \in$ significat *non est*; scilicet:

$$44. \quad a - \in b \cdot = : - \cdot a \in b$$

Signum $a, b, c \in m$ significat: a, b et c sunt m ; scilicet:

$$45. \quad a, b, c \in m \cdot = : a \in m \cdot b \in m \cdot c \in m$$

Sit a classis; tunc $\neg a$ significatur classis indiviuis constituta quae non sunt a .

$$46. \quad a \in \mathbf{K} \cdot \supset : x \in - a \cdot = \cdot x - \in a$$

one does not deduce.

Similarly, if α and β are signs of relations, instead of $(a \alpha b) \wedge (a \beta b)$, and $(a \alpha b) \vee (a \beta b)$ we may write $a (\alpha \wedge \beta) b$ and $a (\alpha \vee \beta) b$. Thus, if a and b are propositions, the formula $a (\rightarrow \wedge \neq) b$ says: from a *one deduces* b , but not vice versa.

$$[a (\rightarrow \wedge \neq) b] = [(a \rightarrow b) \wedge (b \neq a)]$$

Formulas:

$$[(a \rightarrow b) \wedge (b \rightarrow c) \wedge (a \neq c)] = \perp$$

$$[(a = b) \wedge (b = c) \wedge (a \neq c)] = \perp$$

$$\{(a \rightarrow b) \wedge [b (\rightarrow \wedge \neq) c]\} \rightarrow [a (\rightarrow \wedge \neq) c]$$

$$[a (\rightarrow \wedge \neq) b] \wedge (b \rightarrow c) \rightarrow [a (\rightarrow \wedge \neq) c]$$

But we shall rarely use these notations.

IV. Classes.

COMMENTARY: Peano's "Class" is similar to modern "sets".

The sign \mathbf{CLASS} means *class*, or aggregate of entities.

The sign \in means *is an element of*. Thus $a \in b$ is read *a is an element of b*; A is a \mathbf{CLASS} means *A is a class*; a is a proposition means *a is a proposition*.

Instead of $\neg(a \in b)$ we shall write $a \not\in b$. The sign $\not\in$ means *is not*; thus:

$$44. \quad (a \not\in b) = [\neg(a \in b)]$$

The sign $(a, b, c) \in m$ means: a, b , and c are in m ; thus:

$$45. \quad [(a, b, c) \in m] = [(a \in m) \wedge (b \in m) \wedge (c \in m)]$$

Let A be a class. Then \overline{A} means that class made up of individuals that are not in A .

$$46. \quad A \text{ is a CLASS} \rightarrow [(x \in \overline{A}) = (x \not\in A)]$$

Sint a, b classes; $a \cap b$, sive $a \cap b$, est classis individuis constituta

quae eodem tempore sunt a et b ; $a \cup b$ est classis individuis constituta qui sunt a vel b .

$$47. \quad a, b \in K. \cap \therefore a \cap b : = : x \in a . x \in b$$

$$48. \quad a, b \in K. \cup \therefore a \cup b : = : x \in a . \cup . x \in b$$

Signum Λ indicat classem quae nullum continet individuum. Ita:

$$49. \quad a \in K. \cap \therefore a = \Lambda : = : x \in a . =_x \Lambda$$

[Signo \forall , quod classem ex omnibus individuis constitutam, de quibus quaestio est, indieat, non utimur].

Signum \supset significat *continetur*. Ita $a \supset b$ significat *classis a continetur in classi b*.

$$50. \quad a, b \in K. \supset \therefore a \supset b : = : x \in a . \supset_x . x \in b$$

[Formula $b \subset a$ significare potest *classis b continet classem a*; at signo \subset non utimur].

Hic signa Λ et \supset significationem habent quae paullo a praecedenti differt; sed nulla orietur ambiguitas. Nam si de propositionibus agatur, haec signa legantur *absurdum* et *deducitur*; si vero de classibus, *nihil* et *continetur*.

Formula $a = b$, si a et b sint classes, significat $a \supset b . b \supset a$. Itaque

$$51. \quad a, b \in K. = \therefore a = b : = : x \in a . =_x . x \in b$$

Propositiones 1...41 quoque subsistunt, si $a, b...$ classes indicant; praeterea est:

$$52. \quad a \in b . \cap . b \in K$$

$$53. \quad a \in b . \cap . b = = \Lambda$$

$$54. \quad a \in b . b = c : \cap . a \in c$$

$$55. \quad a \in b . b \supset c : \cap . a \in c$$

Sit s classis, et k classis quae in s contineatur; tunc dicimus k esse individuum

Let A, B be classes. Then $A \cap B$, or AB , is the class composed of individuals

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which are at the same time in A and B ; $A \cup B$ is the class composed of individuals which are in A or B .

$$47. \quad A, B \text{ are CLASSES} \rightarrow \{[(x \in (A \cap B)) = [(x \in A) \wedge (x \in B)]]\}$$

$$48. \quad A, B \text{ are CLASSES} \rightarrow \{[(x \in (A \cup B)) = [(x \in A) \vee (x \in B)]]\}$$

The sign \emptyset indicates the class which contains no individuals. Thus:

$$49. \quad A \text{ is a CLASS} \rightarrow \{(A = \emptyset) = [(x \in A) =_{\forall x} \perp]\}$$

[We shall not use the sign true , which indicates the class composed of all individuals being considered].

The sign \subset means *is contained*. Thus $a \subset b$ means *the class a is contained in the class b*.

$$50. \quad A, B \text{ are CLASSES} \rightarrow \{(A \subset B) = [(x \in A) \rightarrow_{\forall x} (x \in B)]\}$$

[The formula $B \supset A$ could mean *the class B contains the class A*, but we shall not use the sign \supset].

The signs (\perp/\emptyset) and (\rightarrow/\subset) have meanings here which are slightly different from the preceding, but no ambiguity will arise, for if propositions are being considered, the signs are read *absurd* (\perp) and *one deduces* (\rightarrow), but if classes are being considered, they are read *empty* (\emptyset) and *is contained* (\subset).

The formula $A = B$, if A and B are classes, means $(A \subset B) \wedge (B \subset A)$. Thus

$$51. \quad A, B \text{ are CLASSES} \rightarrow \{(A = B) = [(x \in A) =_{\forall x} (x \in B)]\}$$

Propositions 1-41 also hold if $a, b, ...$ indicate classes. In addition, we have:

$$52. \quad (a \in B) \rightarrow B \text{ is a CLASS}$$

$$53. \quad (a \in B) \rightarrow (B \neq \emptyset)$$

$$54. \quad [(a \in B) \wedge (B = C)] \rightarrow (a \in C)$$

$$55. \quad [(a \in B) \wedge (B \subset C)] \rightarrow (a \in C)$$

Let C be a class, and C_{subset} be a class which is contained in C ; then we say

classis s , si k ex uno tantum constat individuo. Itaque:

$$56. \quad s \in K . k \supset s : \supset :: k \in s . = \therefore k = \Lambda : x, y \in k . \supset x, y . x = y$$

V. De inversione.

Inversionis signum est $\{ \}$, eiusque usum in sequenti numero explicabimus. Hic tantum casus particulares exponimus.

1. Sit a propositio, indeterminatum continens x ; tunc scriptura $[x] \in a$, quae legitur *ea x quibus a , sive solutiones, vel radices conditionis a , classem significat individuis constitutam, quae conditioni a satisfaciunt*. Itaque:

$$57. \quad a \in P . \supset : [x \in] a . \in K$$

$$58. \quad a \in K . \supset \therefore [x \in] . x \in a : = a$$

$$59. \quad a \in P . \supset \therefore x \in . [x \in] a : = a$$

Sint α, β propositiones indeterminatum continentes x ; erit:

$$60. \quad [x \in] (\alpha \beta) = ([x \in] \alpha) ([x \in] \beta)$$

$$61. \quad [x \in] - \alpha = - [x \in] \alpha$$

$$62. \quad [x \in] (\alpha \cup \beta) = [x \in] \alpha \cup [x \in] \beta$$

$$63. \quad \alpha \supset_x \beta . = . [x \in] \alpha \supset [x \in] \beta$$

$$64. \quad \alpha =_x \beta . = . [x \in] \alpha = [x \in] \beta$$

2. Sint x, y entia quaecumque; system ex ente x et ex ente y compositum ut novum ens consideramus, et signo (x, y) indicamus; similiterque si entium numerus maior fit. Sit α propositio indeterminata continens x, y ; tunc $[(x, y) \in] \alpha$ significat classem entibus (x, y) constitutam, quae conditioni α satisfaciunt. Erit:

that C_{subset} is an individual of the class C , if C_{subset} consists of only one individual. That is:

$$56. \quad [C \text{ is a CLASS} \wedge$$

$$(C_{subset} \subset C) \rightarrow \left\{ (C_{subset} \in C) = \{ (C_{subset} \neq \emptyset) \wedge \{ [(x, y) \in C_{subset}] \xrightarrow{\forall x, y} (x = y) \} \} \right\}$$

V. The inverse.

COMMENTARY: By inversion, Peano means going backwards from a proposition to a class. It is not a mathematical inverse; it is “set-builder notation”.

The sign of the inverse is $\{x|\dots\}$, and we shall explain its use in the following section. Here we give some particular examples.

1. Let a be a proposition containing the indeterminate x ; then the expression $\{x|a\}$, which is read *those x such that a* , or *solutions*, or *roots* of the condition a , indicates the class consisting of individuals which satisfy the condition a . That is:

$$57. \quad a \text{ is a proposition} \rightarrow \{x|a\} \text{ is a CLASS}$$

$$58. \quad A \text{ is a CLASS} \rightarrow \{x|x \in A\} = A$$

$$59. \quad a \text{ is a proposition} \rightarrow (x \in \{x|a\}) = a$$

Let α, β be propositions containing the indeterminate x . We will have:

$$60. \quad \{x|\alpha \wedge \beta\} = \{x|\alpha\} \cap \{x|\beta\}$$

$$61. \quad \{x|\neg \alpha\} = \overline{\{x|\alpha\}}$$

$$62. \quad \{x|\alpha \vee \beta\} = \{x|\alpha\} \cup \{x|\beta\}$$

$$63. \quad \alpha \xrightarrow{\forall x} \beta = \{x|\alpha\} \subset \{x|\beta\}$$

$$64. \quad (\alpha =_{\forall x} \beta) = (\{x|\alpha\} = \{x|\beta\})$$

2. Let x, y be any entities. We shall consider the system composed of the entity x and the entity y as a new entity, and indicate it by the sign (x, y) ; and similarly if the number of entities becomes larger. Let α be a proposition containing the indeterminate x, y ; then $\{(x, y)|\alpha\}$ indicates the class of entities (x, y) which satisfy the condition α . We have:

$$65. \quad \alpha \supset_{x,y} \beta . = . [(x,y) \in] \alpha \supset [(x,y) \in] \beta$$

$$66. \quad [(x,y) \in] \alpha - = \Lambda . = \therefore [x \in] . [y \in] \alpha - = \Lambda : - = \Lambda$$

3. Sit $x \alpha y$ relatio inter indeterminata x et y (ex. g. in logica relationes $x = y$, $x - = y$, $x \supset y$; in arithmetica $x < y$, $x > y$, etc). Tunc signo $[\epsilon \alpha]$ y ea x indicamus, quae relationi $x \alpha y$ satisfaciunt. Commoditatis causa, loco $[\epsilon]$, signo \ni utimur. Ita $\ni \alpha y . = : [x \in] . x \alpha y$, et signum \ni legitur *qui*, vel *quae*. Ex. gr. sit y numerus; tunc $\ni < y$ classem indicat numeris x compositam qui conditioni $x < y$ satisfaciunt, scilicet, *qui sunt minores y*, vel simpliciter *minores y*. Similiter, quum signum D significet *dividit*, vel *est divisor*, formula $\ni D$ significat *qui dividunt* vel divisores. Deducitur $x \in \ni \alpha y = x \alpha y$.

4. Sit α formula indeterminate continens x . Tunc scriptura $x'[x]\alpha$, quae legitur *x' loco x in α substituto*, formulam indicat quae obtinetur si in α , loco x , x' legimus. Deducitur $x[x]\alpha = \alpha$.

5. Sit α formula, quae indeterminata x, y, \dots continet. Tunc

$$(x', y', \dots) [x, y, \dots] \alpha,$$

quae legitur *$x' y', \dots$ loco x, y, \dots in α substitutis*, formulam indicat quae obtinetur si in α loco x, y, \dots , litterae $x' y', \dots$ scribantur. Deducitur $(x, y) [x, y] \alpha = \alpha$.

VI. De functionibus.

Logicae notationes quae praecedunt exprimendae cuilibet arithmeticae proposi-

$$65. \quad \alpha \xrightarrow{\forall x, y} \beta = \{(x, y) | \alpha\} \subset \{(x, y) | \beta\}$$

$$66. \quad (\{(x, y) | \alpha\} \neq \emptyset) = (\{x | \{y | \alpha\} \neq \emptyset\} \neq \emptyset)$$

3. Let $x \alpha y$ be a relation between the indeterminates x and y (eg. in logic, the relations $x = y$, $x \neq y$, $x \rightarrow y$; in arithmetic, $x < y$, $x > y$, ...). Then the sign $[\epsilon \alpha]$ y denotes the x that satisfy the relation $x \alpha y$. For the sake of convenience, we use the sign \ni instead of the sign $[\epsilon]$. Thus, $\ni \alpha y = \{x | x \alpha y\}$, and the sign \ni is read *the objects that*. For example, let y be a number; then $\ni < y$ denotes the class formed by the numbers x that satisfy the condition $x < y$, that is, *the objects that are smaller than y*, or simply *the objects smaller than y*. Similarly, if the sign D means *divides* or *is a divisor of*, the formula $\ni D$ means *the objects that divide* or *the divisors*. It follows that $x \in (\ni \alpha y) = x \alpha y$.

4. Let α be a formula containing the indeterminate x . Then the expression $\alpha[x := x']$, which is read *x' being substituted for x in α* , denotes the formula obtained if, in α , we read x' for x . It follows that $\alpha[x := x] = \alpha$.

5. Let α be a formula that contains the indeterminates x, y, \dots . Then

$$\alpha[x := x', y := y', \dots],$$

which is read *x', y', \dots being substituted for x, y, \dots in α* , denotes the formula obtained if, in α , the letters x', y', \dots are written for x, y, \dots . It follows that $\alpha[x := x, y := y] = \alpha$.

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VI. Functions.

COMMENTARY: Peano has a very structural or syntactical version of a function. There is no parameter. The function is text like “2 + “ or “+ 3” where if you append a number to the start or end, you have a valid expression.

The symbols of logic introduced above suffice to express any proposition of arith-

tioni sufficiunt, iisdemque tantum utimur. Hic notationes alias nonnullas breviter explicamus, quae utiles fieri possunt.

Sit s quaedam classis; supponimus aequalitatem inter entia systematis s definitam, quae conditionibus satisfaciat:

$$a = a$$

$$a = b \text{ . } = \text{ . } b = a$$

$$a = b \text{ . } b = c : \supset \text{ . } a = c$$

Sit φ signum, sive signorum aggregatus, ita ut si x est ens classis s , scriptura φ x novum indicet ens; supponimus quoque aequalitatem inter entia φ x definitam; et si x et y sunt entia classis s , et est $x = y$, supponimus deduci posse φ $x = \varphi$ y . Tunc signum φ dicitur esse *functionis praesignum in classi* s , et scribemus $\varphi \ni F' s$.

$$s \in K \text{ . } \supset :: \varphi F' s \text{ . } = \text{ . } \therefore x, y \in s \text{ . } x = y : \supset_{x,y} \text{ . } \varphi x = \varphi y$$

Verum si, cum sit x quodlibet ens classis s , scriptura $x\varphi$ novum indicet ens, et, ex, $x = y$ deducitur $x\varphi = y\varphi$, tunc dicimus φ esse *functionis postsignum in classi* s et scribemus $\varphi \in s'F$.

$$s \in K \text{ . } \supset :: \varphi s'F \text{ . } = \text{ . } \therefore x, y \in s \text{ . } x = y : \supset_{x,y} \text{ . } x\varphi = y\varphi$$

Exempla. Sit a numerus; tunc $a +$ est functionis praesignum in numerorum classe, et $+ a$ est functionis postsignum; quicumque enim est numerus x , formulae $a + x$ et $x + a$ novos indicant numeros, et ex $x = y$ deducitur $a + x = a + y$, et $x + a = y + a$. Itaque

$$a \in N \text{ . } \supset : a + \text{ . } \in \text{ . } F' N$$

$$a \in N \text{ . } \supset : + a \text{ . } \in \text{ . } N' F$$

Sit φ functionis praesignum in classe s . Tunc $[\varphi]y$ classem significat iis x constitutam, quae conditioni $\varphi x = y$ satisfaciunt; scilicet:

$$\text{Def. } s \in K \text{ . } \varphi \in F' s : \supset : [\varphi]y \text{ . } = \text{ . } [x \in] (\varphi x = y)$$

metic, and we shall only these. We explain here briefly some other symbols that may be useful.

Let S be a class; we assume that equality is defined between the objects of the system s so as to satisfy the conditions:

$$a = a.$$

$$(a = b) = (b = a)$$

$$[(a = b) \wedge (b = c)] \rightarrow a = c$$

Let φ be a sign or an aggregate of signs such that, if x is an object of the class S , the expression φ x denotes a new object; we assume also that equality is defined between the objects φ x ; further, if x and y are objects of the class S and if $x = y$, we assume it is possible to deduce φ $x = \varphi$ y . Then the sign φ is said to be a *function presign in the class* S , and we write $\varphi \ni F' s$.

...

If, x being any object of the class S , the expression $x\varphi$ denotes a new object and $x\varphi = y\varphi$ follows from $x = y$, then we say that φ is a *function postsign in the class* S , and we write $\varphi \in s'F$.

...

Examples. Let a be a number; then $a +$ is a function presign in the class of numbers, and $+ a$ is a function postsign; for any number x , formulas $a + x$ and $x + a$ denote new numbers; $a + x = a + y$ and $x + a = y + a$ follow from $x = y$. Thus

...

Let φ be a function presign in the class S . Then $[\varphi]y$ denotes the class composed of x that satisfy the condition $\varphi x = y$; that is,

Def.

Classis $[\varphi]y$ vel unum vel plura, vel etiam nullum individuum continere potest.

Erit:

$$s \in K. \varphi \in F' : \mathcal{D} : y = \varphi x. = . x \in [\varphi]y$$

Si vero φy uno tantum constat individuo, erit $y = \varphi x. = . x = [\varphi]y$

Sit φ functions postsignum; similiter ponimus:

$$s \in K. \varphi \in s'F : \mathcal{D} : y \mid \varphi \mid = |x| (x\varphi = y).$$

Signum $[]$ dicitur *inversionis signum*, eiusque usus nonnullos in logica iam exposuimus. Nam si α est propositio indeterminatum continens x , atque A est classis individuis x composita quae conditioni α satisfaciunt, erit $x \in \alpha. = \alpha$, tunc $\alpha = [x \in] \alpha$, ut in V, i.

Sit α formula indeterminata continens x , sitque φ functionis praesignum, quod litterae x praepositum, formulam α gignat; scilicet sit $\alpha = \varphi x$; tunc erit $\varphi = \alpha[x]$, et si x' est novum ens, erit $\varphi x' = \alpha[x]x'$, scilicet, si α est formula indeterminatum continens x , tunc $\alpha[x]x'$ significat id quod obtinetur si in α , loco x, x' ponatur.

Similiter, sit α formula indeterminata continens x , sitque φ functionis postsignum, ut $x\varphi = \alpha$; deducitur $\varphi = [x]\alpha$; tunc, si x' est novum ens, erit $x'\varphi = x'[x]\alpha$, scilicet $x'[x]\alpha$ rursum indicat id quod obtinetur si in α , loco x, x' legatur, ut in V, 4.

Alios quoque usus in logica signum $[]$ habere potest, quos breviter esponimus, quum ipsis non utamur. Sint a et b duae classes; tunc $[a \cap]b$ sive $b[\cap a]$ classes indicat x , quae conditioni $b = a \cap x$, sive $b = x \cap a$ satisfaciunt. Si b in a non continetur, nulla classis huic conditioni satisfacit; si b in a continetur, signum $b[\cap a]$ omnes indicat classes quae b continent atque in $b \cup -a$ continentur.

In Arithmetica, sint a, b numeri; tunc $[b + a]$ sive $[a +]b$ numerum indicat x , qui conditioni $b = x + a$, sive $b = a + x$ satisfacit, nempe $b - a$. Similiter erit $b[\times a] =$

The class $[\varphi]y$ may contain one or several individuals, or none at all. We have:

...

But if φy consists of just one individual, we have $y = \varphi x. = . x = [\varphi]y$

Let φ be a function postsign; we write similarly:

...

The sign $\{x|\dots\}$ is called *inversion sign*, and we have already presented some of its uses in logic. If α is a proposition containing the indeterminate x and A is a class composed of the individuals x that satisfy the condition α , we have $x \in A = \alpha$, and then $A = \{x|\alpha\}$, as in V, i.

Let α be a formula containing the indeterminate x and let φ be a function presign that yields the formula α when written before the letter x ; that is, let $\alpha = \varphi x$; then we have $\varphi = \alpha[x := ?]$, and if x' is a new object, we have $\varphi x' = \alpha[x := x']$, that is, if α is a formula containing the indeterminate x , then $\alpha[x := x']$ means what is obtained when, in α , we put x' for x .

Similarly, let α be a formula containing the indeterminate x and let φ be a function postsign, such that $x\varphi = \alpha$; it follows that $\varphi = \alpha[x := ?]$. Then, if x' is a new object, have $x'\varphi = \alpha[x := x']$; that is, $x'[x]\alpha$ again denotes what is obtained, when, in α , read x' for x , as in V, 4..

The sign $[]$ can have other uses in logic, which we present only briefly, since we shall not use it in these ways. Let A and B be two classes; then $[A \cap]B$ or $B[\cap A]$ denotes the classes X that satisfy the condition $B = A \cap X$, or $B = X \cap A$. If B is not contained in A , no class satisfies this condition; if B is contained in A , the sign $B[\cap A]$ denotes all classes that contain B and are contained in $B \cup \bar{A}$.

In arithmetic, let a and b be numbers; then $[b + a]$ or $[a +]b$ denotes the number x that satisfies the condition $b = x + a$, or $b = a + x$, that is $b - a$. Similarly we have

$[a \times]b = b/a$. Et in analysi hoc signum usuvenire potest; itaque

$$y = \sin x \text{ . } = \text{ . } x \in [\sin] y \quad (\text{loco } x = \arcsin y).$$

$$dF(x) = f(x)dx \text{ . } = \text{ . } F(x) \in [d] \quad (\text{loco } F(x) = \int f(x)dx).$$

$f(x)dx$

Sit rursum φ functionis praesignum in classi s , sitque k classis

in s contenta; tunc φk classem indicat omnibus φx compositam, ubi x sunt entia classis k ; scilicet

$$\text{Def. } s \in K \text{ . } k \in K \text{ . } k \supset s \text{ . } \varphi \in F' s : \supset \text{ . } \varphi k = [y \in] (k \text{ . } [\varphi]y : - = \Lambda)$$

$$\text{Sive } s \in K \text{ . } k \in K \text{ . } k \supset s \text{ . } \varphi \in F' s : \supset \text{ . } \varphi k = [y \in] ([x \in] : x \in k \text{ . } [\varphi]x = y \therefore - = \Lambda)$$

$$\text{Def. } s \in K \text{ . } k \in K \text{ . } k \supset s \text{ . } \varphi \in s' F : \supset \text{ . } k \varphi = [y \in] (k \text{ . } y [\varphi] : - = \Lambda)$$

Itaque, si $\varphi \in F' s$, tunc φs classem indicat omnibus φx constitutam, ubi x sint entia classis s . Erit:

$$s \in K \text{ . } \varphi \in F' s \text{ . } y \in \varphi s : \supset : \varphi[\varphi]y = y$$

$$s \in K \text{ . } a, b \in K \text{ . } a \supset s \text{ . } b \supset s \text{ . } \varphi \in F' s : \supset \text{ . } \varphi(a \cup b) = (\varphi a) \cup (\varphi b)$$

$$s \in K \text{ . } \varphi \in F' s : \supset \text{ . } \varphi \Lambda = \Lambda$$

$$s \in K \text{ . } a, b \in K \text{ . } b \supset s \text{ . } a \supset b \text{ . } \varphi \in F' s : \supset \text{ . } \varphi a \supset \varphi b$$

$$s \in K \text{ . } a, b \in K \text{ . } a \supset s \text{ . } b \supset s \text{ . } \varphi \in F' s : \supset \text{ . } \varphi(ab) \supset (\varphi a)(\varphi b)$$

Sit a quaedam classis; tunc $a \cap K$, sive $K \cap a$, sive Ka , classes omnes indicat formae $a \cap x$, sive $x \cap a$, xa , ubi x est classis quacumque; scilicet Ka indicat classes quae in a continentur. Formula $x \in Ka$ idem significat quod $x \in K \text{ . } x \supset a$. Hac conventionem quandoque utimur; ita KN isgnificat *numerorum classem*.

Similiter, si a est classis, $K \cup a$ indicat classes quae a continent. Sit a numerus; tunc $a + N$, sive $N + a$, *numeros* indicat *numero a maiores*; $a \times N$, sive $N \times a$, sive Na indicat *multiplices numeri a*; a^N indicat *potestas numeri a*; N^2, N^3, \dots indicat *numeros quadratos, vel numeros cubos,...*

$b[\times a] = [a \times]b = b/a$. This sign can even find a use in analysis; thus

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Let φ again be a function presign in a class S and let C be a class

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contained in S ; then φC denotes the class consisting of all φx , where the x are the objects of class C ; that is

Def.

Or

Def.

Thus, if $\varphi \in F' S$, then φs denotes the class composed of all φx , where the x are objects of the class S . We have:

...

...

...

...

...

Let a be a class; then $a \cap K$, or $K \cap a$, or Ka , denotes all classes of the form $a \cap x$, or $x \cap a$, xa , where x is any class; that is Ka denotes the classes that are contained in a . The formula $x \in Ka$ means the same as $x \in K \text{ . } x \supset a$. We shall sometimes use this convention; thus KN means *a class of numbers*.

Similarly, if A is a class, $K \cup A$ indicates the classes that contain A . Let a be a number; then $a + N$, or $N + a$, denotes *the numbers greater than the number a*; $a \times N$, or $N \times a$, or Na denotes *the multiples of the number a*; a^N denotes *the powers of the number a*; N^2, N^3, \dots denote *the squares, the cubes,...*

Functional signorum aequalitatem, productum, potestas, ita definire licet:

Def. $s \in K . \varphi, \psi \in F's : \supset \therefore \varphi = \psi : = : x \in s . \supset . \varphi x = \psi x$

Def. $s \in K . \varphi \in F's . \psi \in F'\varphi s . x \in s : \supset . \psi \varphi x = \psi(\varphi x)$

Itaque, in definitionis hypothesi, erit $\psi\varphi$ novum functionis praesignum; idque *productum signorum ψ et φ* vocatur.

Similiterque, si φ, ψ sunt functionis postsigna.

Haec valet propositio:

$$s \in K . \varphi \in F's . \varphi s \supset s : \supset : \varphi \varphi s \supset s . \varphi \varphi \varphi s \supset s . \text{ etc.}$$

Functiones $\varphi\varphi, \varphi\varphi\varphi, \dots$ *iteratae* vocantur, et communiter signis $\varphi^2, \varphi^3, \dots$ indicantur, ut operationis φ potestates.

Si vero φ est functionis postsignum, haec faciliori notatione, absque ambiguitate, uti licet:

Def. $s \in K . \varphi \in s'F . s\varphi \supset s : \supset : \varphi 1 = \varphi . \varphi 2 = \varphi\varphi . \varphi 3 = \varphi\varphi\varphi . \text{ etc.}$

In definitionis hypothesi, si $m, n \in N$, erit $\varphi(m+n) = (\varphi m)(\varphi n)$; $(\varphi m)n = \varphi(mn)$

Si hac definitione in Arithmetica utimur, haec invenimus. Numerum qui sequitur numerum a signo faciliori $a+$ indicare possumus; tunc $a+1, a+2, \dots$ et, si b est numerus, $a+b$, sensum habent $a+, a++ , \dots$ quod a definitione in §1 patet. Propositionem 6 in §1 scribere possumus $N+ \supset N$. Si a, b, c sunt numeri, tunc $a : +b . c$ significat $a+bc$, et $a : \times b . c$ significat ab^c .

Multi aliis proprietatibus gaudent functionem signa, praesertim si conditioni satisfaciunt: $\varphi x = \varphi y . \supset . x = y$. Functionis signum quod huic conditioni satisfacit vocatur a clarrissimo Dedekind *simile* (ähnliche Abbildung).

Sed his exponendis locus deest.

Equality, product and powers can be defined thus for function signs:

Def.

Def.

Thus, if we assume this definition, we have the new function presign $\psi\varphi$; it is called the *product of the signs ψ and φ* .

Similarly if φ, ψ are function postsigns.

The following proposition holds:

...

The functions $\varphi\varphi, \varphi\varphi\varphi, \dots$ are said to be *iterated* and are generally denoted by the signs $\varphi^2, \varphi^3, \dots$ as powers of the operation φ .

XVI

But if φ is a function postsign, we can use the following more convenient notation without ambiguity:

Def.

Assuming this definition, if $m, n \in N$, we have $\varphi(m+n) = (\varphi m)(\varphi n)$; $(\varphi m)n = \varphi(mn)$

If we use this definition in arithmetice, we obtain the following. We can denote the number that follows the number a by the more convenient sign $a+$; then $a+1, a+2, \dots$, and, if b is a number, $a+b$, have the meaning of $a+, a++ , \dots$, which is clear from the definition in §1 below. Proposition 6 in §1 can be written $N+ \supset N$.

If a, b, c are numbers, then $a : +b . c$ means $a+bc$, and $a : \times b . c$ means ab^c .

Function signs possess many other properties, especially if they satisfy the condition: $\varphi x = \varphi y . \supset . x = y$. A function sign that satisfies this condition is called *equivalent*² by Dedekind (ähnliche Abbildung).

But we lack the space to present these properties.

Declarationes.

Defenitio, vel breviter *Def.* est propositio formam habens $x = a$, sive $\alpha \supset . x = a$, ubi α est signorum aggregatus sensum habens notum; x est signum, vel signorum aggregatus significatione adhuc carnes; a vero est conditio sub qua definitio datur.

Theorema, (Theor. vel Th) est propositio quae demonstratur. Si theorema formam habet $\alpha \supset \beta$, ubi α et β sunt propositiones, tunc α dicitur *Hypothesis* (Hyp. vel breviter Hp.), β vero *Thesis* (Thes. vel Ts.). Hyp. ac Ts. a Theorematis forma pendent; nam si loco $\alpha \supset \beta$ scribemus $-\beta \supset -\alpha$, erit $-\beta$ Hp, et $-\alpha$ Ts.; si vero scribemus $\alpha - \beta = \Lambda$, Hp. ac Ts. absunt.

In quolibet § signum P quod quidam numerus sequatur, propositionem indicat eiusdem § hoc numero signatam. Logicae propositiones indicantur signo L et propositiones numero.

Formulae quae in una linea non continentur, in altera linea, nullo interposito signo, sequuntur.

Remarks.

A *Definition*, or *Def.* for short, is a proposition of the form $x = a$ or $\alpha \supset . x = a$, where a is an aggregate of signs having a known meaning, x is a sign or an aggregate of signs, hitherto without meaning, and α is the condition under which the definition is given.

A *theorem* (Theor. or Th.) is a proposition that is *proved*. If a theorem has the form $\alpha \supset \beta$, where α and β are propositions, then α is called the *hypothesis* (Hyp., or even shorter, Hp.) and β the *thesis* (Thes. or Ts.). Hyp. and Ts. depend on the form of the theorem; in fact, if we write $\alpha \supset \beta$ instead of $-\beta \supset -\alpha$, then $-\beta$ is the Hp, and $-\alpha$ the Ts.; if we write $\alpha - \beta = \Lambda$, Hp. and Ts. do not exist.

In any § below, the sign P followed by a number denotes the proposition indicated by that number in the same §. Propositions of logic are indicated by the sign L and the number of the proposition.

Formulas that do not fit on one line are continued on the next line without any intervening sign.

ARITHMETICES PRINCIPIA.

§1. De numeris et de additione.

Explicationes.

Signo	N	significatur	numerus (integer positivus).
»	1	»	unitas.
»	$a + 1$	»	sequens a , sive a plus 1.
»	=	»	est aequalis. Hoc ut novum signum considerandum est, etsi logicae signi figuram habeat.

Axiomata.

THE PRINCIPLES OF ARITHMETIC.

§1. Numbers and addition.

Explanations.

COMMENTARY: Peano starts his natural numbers at 1; most modern versions start at 0. Also, while Peano calls it “successor” (“sequens”), he does not use the modern convention of using “S” for it.

The sign	N	means	number (positive integer).
”	1	”	unity.
”	$a + 1$	”	the successor of a , or a plus 1.
”	=	”	is equal to. This must be considered as a new sign, although it has the appearance of a sign of logic.

Axioms.

1. $1 \in N$...
2. $a \in N . \supset . a = a$...
3. $a, b, c \in N . \supset : a = b . = . b = a$...
4. $a, b \in N . \supset : a = b . b = c : \supset . a = c$...
5. $a = b . b \in N : \supset . a \in N$...
6. $a \in N . \supset . a + 1 \in N$...
7. $a, b \in N . \supset : a = b . = . a + 1 = b + 1$...
8. $a \in N . \supset . a + 1 \dashv = 1$...
9. $k \in K : 1 \in k : x \in N . x \in k : \supset_x . x + 1 \in k :: \supset . N \supset k$...

Definitiones.

10. $2 = 1 + 1; 3 = 2 + 1; 4 = 3 + 1; \text{ etc.}$...

2

Theoremata.

11. $2 \in N.$...

Demonstratio:

- | | | |
|--|-----------------------------------|----------|
| $P1 . \supset :$ | $1 \in N$ | (1) |
| $1[a](P6) . \supset :$ | $1 \in N . \supset . 1 + 1 \in N$ | (2) |
| $(1)(2) . \supset :$ | $1 + 1 \in N$ | (3) |
| $P10 . \supset :$ | $2 = 1 + 1$ | (4) |
| $(4).(3).(2, 1 + 1)[a, b](P5) : \supset :$ | $2 \in N$ | (Theor.) |

Nota. - Huius facillimae demonstrationis gradus omnes ecplícite scripsimus. Brevitatis causa ipsam ita scribemus:

$P1 . 1[a](P6) : \supset : 1 + 1 \in N . P10 . (2, 1 + 1)[a, b](P5) : \supset : Th$

Definitions.

Proof:

- | | |
|------|-----|
| | (1) |
| ... | (2) |
| ... | (3) |
| ... | (4) |

Note. - We have explicitly written every step of this very easy proof. For the sake of brevity, we shall write it as follows:

...

vel		...
	$P1.P6:\supset:1+1\in N.P10.P5:\supset:Th$...
12.	$3,4,\dots\in N$...
13.	$a,b,c,d\in N.a=b.b=c.c=d:\supset:a=d$...
Dem.	$Hyp.P4:\supset:a,c,d\in N.a=c.c=d.P4:\supset:Thes.$...
14.	$a,b,c\in N.a=b.b=c.a\neq c:=\Lambda$...
Dem.	$P4.L39:\supset.Theor.$...
15.	$a,b,c\in N.a=b.b\neq c:\supset.a\neq c$...
16.	$a,b\in N.a+1=b+1:\supset.a=b$...
17.	$a,b\in N.\supset:a\neq b.=.a+1\neq b+1$...
Dem.	$P7.L21:\supset.Theor$...

Definitio.

18.	$a,b\in N.\supset.a+(b+1)=(a+b)+1$...
-----	------------------------------------	-----

Nota. - Hanc definitionem ita legere oportet: si a et b sunt numeri, et $(a+b)+1$ sensum habet (scilicet si $a+b$ est numerus), sed $a+(b+1)$ nondum definitus est, tunc $a+(b+1)$ significat numerum qui $a+b$ sequitur.

Ab hac definitione, et a praecedentibus deducitur:

$$a\in N.\supset:a+2=a+(1+1)=(a+1)+1$$

$$a\in N.\supset:a+3=a+(2+1)=(a+2)+1, \text{ etc.}$$

Definition.

Note. - This definition should be read: if a and b are numbers, and $(a+b)+1$ has meaning (that is, if $a+b$ is a number), but $a+(b+1)$ has not yet been defined, then $a+(b+1)$ indicates the number that follows $a+b$.

From this definition, and the preceding, we deduce that:

...

Theoremata.

19.	$a,b\in N.\supset.a+b\in N.$...
Dem.	$a\in N.P6:\supset:a+1\in N:\supset:1\in[b\in]Ts.$	(1) ...
	$a\in N.\supset:b\in N.b\in[b\in]Ts:\supset:a+b\in N.P6$	(2) ...

Theorems.

...
...
...

	$\vdash (a + b) + 1 \in \mathbb{N}$. P18				
	$\vdash a + (b + 1) \in \mathbb{N} : \vdash (b + 1) \in [b \in] \text{Ts}$.				
	$a \in \mathbb{N} . (1) . (2) . \vdash 1 \in [b \in] \text{Ts} \therefore b \in \mathbb{N} . b \in [b \in] \text{Ts}$	(3)
	$\vdash b + 1 \in [b \in] \text{Ts} \therefore ([b \in] \text{Ts})[k] \text{P9}$				
	$\vdash \vdash \mathbb{N} \supset [b \in] \text{Ts} . (\text{L50}) \vdash \vdash b \in \mathbb{N} . C \text{Ts}$.				
	(3) . (L42) : $C : a, b \in \mathbb{N} . \vdash$. Thesis	(Th.)
20. Def.	$a + b + c = (a + b) + c$	
21.	$a, b, c \in \mathbb{N} . \vdash a + b + c \in \mathbb{N}$	
22.	$a, b, c \in \mathbb{N} . \vdash a = b . = . a + c = b + c$	
Dem.	$a, b \in \mathbb{N} . \text{P7} : \vdash 1 \in [c \in] \text{Ts}$.	(1)
	$a, b \in \mathbb{N} . \vdash \vdash c \in \mathbb{N} . c \in [c \in] \text{Ts}$	(2)
	$\therefore \vdash \therefore a = b . = . a + c = b + c : a + c, b + c \in \mathbb{N} :$				
	$a + c = b + c . = . a + c + 1 = b + c + 1 : \vdash \therefore a = b . =$				
	$. a + (c + 1) = b + (c + 1) : \vdash \therefore (c + 1) \in [C \in] \text{Ts}$.				
	$a, b \in \mathbb{N} . (1) . (2) : \vdash 1 \in [c \in] \text{Ts} \therefore c \in [c \in] \text{Ts}$	(3)
	$\vdash \vdash (c + 1) \in [c \in] \text{Ts} \vdash \vdash c \in \mathbb{N} . \vdash \text{Ts}$.				
23.	$a, b, c \in \mathbb{N} . \vdash a + (b + c) = a + b + c$	
Dem.	$a, b \in \mathbb{N} . \text{P18} . \text{P20} : \vdash 1 \in [c \in] \text{Ts}$.	(1)
	$a, b \in \mathbb{N} . \vdash \therefore c \in \mathbb{N} . c \in [c \in] \text{Ts}$	(2)
	$\vdash \vdash a + (b + c) = a + b + c . \text{P7}$				
	$\vdash \vdash a + (b + c) + 1 = a + b + c + 1 . \text{P18}$				
	$\vdash \vdash a + (b + (c + 1)) = a + b + (c + 1) : \vdash c + 1 \in [c \in]$				
	Ts .				
	(1)(2) (P9) . \vdash . Theor.	
24.	$a \in \mathbb{N} . \vdash 1 + a = a + 1$	
Dem.	$\text{P2} . \vdash 1 \in [a \in] \text{Ts}$.	(1)
	$a \in \mathbb{N} . a \in [a \in] \text{Ts} : \vdash 1 + 1 = a + 1 : \vdash$	(2)
	$1 + (a + 1) = (a + 1) + 1 : \vdash (a + 1) \in [a \in] \text{Ts}$.				

	(1)(2). \odot . Theor.	
24'.	$a, b \in \mathbf{N} . \odot . 1 + a + b = a + 1 + b$	
Dem.	Hyp. P24 : $\odot \odot 1 + a = a + 1$. P22 : \odot . Thesis.	
..... 4					
25.	$a, b \checkmark \in \mathbf{N} . \odot . a + b = b + a .$	
Dem.	$a \in \mathbf{N} .$ P24 : $\odot : 1 \in [b \in]$ Ts.	(1)
	$a \in \mathbf{N} . \odot . b \in \mathbf{N} . b \in [b \in]$ Ts : $\odot : a + b = b + a .$ P7 :	(2)
	$\odot : (a + b) + 1 = (b + a) + 1 . (a + b) + 1 = (a + (b + 1)) . (b +$ $a) + 1 = 1 + (b + a) . 1 + (b + a) = (1 + b) + a . (1 + b) + a =$ $(b + 1) + a : \odot : a + (b + 1) = (b + 1) + a : \odot : (b + 1) \in [b \in]$ Ts.				
	(1)(2). \odot . Theor.	
26.	$a, b, c \in \mathbf{N} . \odot : a = b . = . c + a = c + b .$	
27.	$a, b, c \in \mathbf{N} . \odot : a + b + c = a + c + b .$	
28.	$a, b, c, d \in \mathbf{N} . a = b . c = d : \odot . a + c = b + d .$	

§2. De subtractione.

§2. Substraction.

Explicationes.

Signo	—	legitur	<i>minus.</i>
»	<	»	<i>est minor.</i>
»	>	»	<i>est maior.</i>

Explanations.

The sign	—	is read	<i>minus.</i>
”	<	”	<i>is less than.</i>
”	>	”	<i>is greater than.</i>

Definitiones.

1. $a, b \in \mathbf{N} . \odot : b - a = \mathbf{N}[x \in](x + a = b) .$

Definitions.

2. $a, b \in \mathbb{N} . \supset : a < b . = . b - a \in \mathbb{N} .$

3. $a, b \in \mathbb{N} . \supset : b > a . = . a < b .$

$$a + b - c = (a + b) - c ; a - b + c =$$

$$(a - b) + c ; a - b - c = (a - b) - c .$$

Theoremata.

4. $a, b, a', b' \in \mathbb{N} . a = a' . b = b' : \supset : b - a = b' - a' .$

Dem. Hyp $\supset : x + a = b . = . x + a' = b' : \supset$. Thesis.

5. $a, b \in \mathbb{N} . \supset : a < b . = . b - a \in \mathbb{N} .$

Dem. $a, b \in \mathbb{N} : \supset \therefore x, y \in b - a . \supset_{x,y} : x, y \in \mathbb{N} . x + a =$ (1)

$$b . y + a = b . \S 1 \text{ P22 } : \supset : x = y .$$

$$a, b \in \mathbb{N} . a < b . \text{ P2 . (1) } : \supset \therefore b - a \in \mathbb{N} : x, y \in$$
 (2)

$$b - a . \supset . x = y : (N, b - a)[s, k] \text{ (L56)}$$

$$\therefore \supset \therefore b - a \in \mathbb{N} .$$

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$$a, b \in \mathbb{N} . b - a \in \mathbb{N} . \text{ (L56)} \quad (3)$$

$$: \supset : b - a \in \mathbb{N} : \supset : a < b .$$

$$(2)(3) . \supset . \text{ Theor.}$$

6. $a, b \in \mathbb{N} . a < b : \supset . b - a + a = b .$

Dem. Hyp . P5 . P1

$$: \supset : b - a \in \mathbb{N} . (b - a) \in [x \in](x + a = b) : \supset : \text{Thes.}$$

7. $a, b, c \in \mathbb{N} . \supset : c = b - a . = . c + a = b .$

Dem. Hyp . §1 P22 . P6

$$: \supset : c = b - a . = . c + a = b - a + a . = . c + a = b .$$

8. $a, b \in \mathbb{N} . \supset . a + b - a = b .$

Dem. $(a + b, b)[b, c] \text{ P7 } . \supset . \text{ Theor.}$

9. $a, b, c \in \mathbb{N} . a < b : \supset : c + (b - a) = c + b - a .$

Dem.

Hyp . P6

$$:\supset : (b - a) + a = b : \supset : c + (b - a) + a = c + b . \text{ P7}$$

$:\supset :$ Thesis.

10.

$$a, b, c \in \mathbf{N} . a > b + c : \supset . a - (b + c) = a - b - c .$$

11.

$$a, b, c \in \mathbf{N} . b > c . a > b - c : \supset . a - (b - c) = a + c - b .$$

12.

$$a, b, a', b' \in \mathbf{N} . a = a' . b = b' : \supset : a < b . = . a' < b' .$$

Dem.

Hyp

$$. \supset . b - a = b' - a' . \supset . b - a \in \mathbf{N} = b' - a' \in \mathbf{N} . \supset .$$

Thes.

13.

$$a, b \in \mathbf{N} . \supset . a < a + b .$$

Dem.

$$\text{Hyp . P8 : } \supset : a + b - a = b : \supset . a + b - a \in \mathbf{N} . \text{ P5}$$

$:\supset :$ Thesis.

14.

$$a, b, c \in \mathbf{N} . a < b . b < c : \supset . a < c .$$

Dem.

$$\text{Hyp . } \supset : b - a \in \mathbf{N} . c - b \in \mathbf{N} : \supset : (b - a) + (c - b) \in \mathbf{N} :$$

$$\supset : c - a \in \mathbf{N} : \supset . \text{ Thesis.}$$

15.

$$a, b, c \in \mathbf{N} . \supset : a < b . = . a + c < b + c .$$

Dem.

$$\text{Hyp . } \supset : a < b . = . b - a \in \mathbf{N} . = . (b + c) - (a + c) \in$$

$$\mathbf{N} . = . a + c < b + c .$$

16.

$$a, b, a', b' \in \mathbf{N} . a < b . a' < b' : \supset . a + a' < b + b' .$$

Dem.

$$\text{Hyp . } \supset : a + a' < b + a' . b + a' < b + b' : \supset . \text{ Thesis.}$$

17.

$$a, b, c \in \mathbf{N} . a < b < c : \supset . c - a > c - b .$$

Dem.

Hyp

$$. \supset b - a \in \mathbf{N} . c - b \in \mathbf{N} . (c - b) + (b - a) = c - a : \supset .$$

Thesis.

18.

$$a \in \mathbf{N} . \supset : a = 1 . \cup . a > 1 .$$

Dem.

$$1 \in [a \in] \text{ Thesis.}$$

$$a \in \mathbf{N} . \text{ P13 : } \supset : a + 1 > 1 : \supset : a + 1 \in [a \in] \text{ Thesis.}$$

$$(1)(2) . \supset . \text{ Theor.}$$

19. $a, b \in \mathbb{N} . \odot . a + b = b . \checkmark$
- Dem.* $a \in \mathbb{N} . \S 1 \text{ P8} : \odot : a + 1 = 1 : \odot : 1 \in [b \in] \text{ Thesis.} \quad (1)$
- $a \in \mathbb{N} . b \in \mathbb{N} . b \in [b \in] \text{ Ts} : \odot : a + b = b . \S 1 \text{ P17} \quad (2)$
- $: \odot : a + (b + 1) = b + 1 : \odot : b + 1 \in [b \in] \text{ Ts.}$
- $(1)(2) . \odot . \text{Theor.}$
20. $a, b \in \mathbb{N} . a < b . a = b := \Lambda .$
- Dem.* $\text{Hyp} : \odot : b - a \in \mathbb{N} . (b - a) + a = a . \text{P19} : \odot : \Lambda .$
21. $a, b \in \mathbb{N} . a > b . a = b := \Lambda .$
22. $a, b \in \mathbb{N} . a > b . a < b := \Lambda .$
23. $a, b \in \mathbb{N} : \odot : a < b . \cup . a = b . \cup . a > b .$
- Dem* $a \in \mathbb{N} . \text{P18} : \odot . 1 \in [b \in] \text{ Ts.} \quad (1)$
- $a, b \in \mathbb{N} . a < b : \odot . a < b + 1 . \quad (2)$
- $a, b \in \mathbb{N} . a = b : \odot . a < b + 1 . \quad (3)$
- $a, b \in \mathbb{N} . a > b : \odot : a - b \in \mathbb{N} . \text{P18} \quad (4)$
- $: \odot : a - b = 1 . \cup . a - b > 1 .$
- $a, b \in \mathbb{N} . a - b = 1 : \odot . a = b + 1 . \quad (5)$
- $a, b \in \mathbb{N} . a - b > 1 : \odot . a > b + 1 . \quad (6)$
- $a, b \in \mathbb{N} . a > b . (4)(5)(6) : \odot : a = b + 1 . \cup . a > b + 1 . \quad (7)$
- $a, b \in \mathbb{N} : a > b . \cup . a = b . \cup . a > b : \quad (8)$
- $(2)(3)(7) . \therefore \odot : a < b + 1 . \cup . a = b + 1 . \cup . a > b + 1 .$
- $a, b \in \mathbb{N} . b \in [b \in] \text{ Ts} . (8) : \odot : b + 1 \in [b \in] \text{ Ts.} \quad (9)$
- $(1)(9) . \odot . \text{Theor.}$

§3. De maximis et minimis.

§3. Maxima and minima.

Explicationes.

Sit $a \in KN$, hoc est sit a quaedam numerorum classis; tunc Ma legatur *maximus inter* a , et Wa legatur *minimus inter* a .

Definitiones.

1. $a \in KN . \supset : Ma = [x \in](x \in a . \therefore a . \in > x := \Lambda . \checkmark$
2. $a \in KN . \supset : Wa = [x \in](x \in a . \therefore a . \in < x := \Lambda . \checkmark$

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Theoremata.

3. $n \in N . a \in KN . a \text{ - } = \Lambda . a \supset > n = \Lambda : \supset . Ma \in N . \checkmark$
- Dem.* $a \in KN . a \text{ - } = \Lambda . a \supset > 1 = \Lambda : \supset : a = 1 : \supset . Ma = 1 : \quad (1)$
- $\supset . Ma \in N . \checkmark$
- $(1) \supset : 1 \in [n \in] \text{ (Hp } \supset \text{ Ts).} \quad (2)$
- $n \in N . a \in KN . a \supset > n + 1 = \Lambda . n + 1 \in a : \supset : \quad (3)$
- $n + 1 = Ma : \supset : Ma \in N .$
- $n \in N . a \in KN . a \supset > n + 1 = \Lambda . n + 1 \text{ - } \in a : \supset : \quad (4)$
- $a \supset > n = \Lambda .$
- $n \in [n \in] \text{ (Hp } \supset \text{ Ts)} \quad (5)$
- $. a \in KN . a \supset > n + 1 = \Lambda . n + 1 \text{ - } \in a : \supset : Ma \in N .$
- $n \in [n \in] \text{ (Hp } \supset \text{ Ts)} . (6) : \supset . (n + 1) \in [n \in] \text{ (Hp } \supset \text{ Ts).} \quad (7)$
- $(2)(7) . \S 1 \text{ P9} : \supset : n \in N . \supset . \text{ Hp } \supset \text{ Ts.} \quad (\text{Th.})$
4. $a \in KN . a \text{ - } = \Lambda : \supset . Wa \in N .$
5. $a \in KN . \supset : Wa = M[x \in](a \supset < x = \Lambda) .$

Explanations.

Let $a \in K N$, that is, let a be a class of numbers; then Ma is read *greatest among* a , and Wa is read *least among* a .

Definitions.

§4. De multiplicatione.

Definitiones.

1. $a \in \mathbb{N} . \supset . a \times 1 = a . \checkmark$
2. $a, b \in \mathbb{N} . \supset . a \times (b + 1) = a \times b + a .$
 $ab = a \times b ; ab + c = (ab) + c ; abc = (ab)c .$

...

Theoremata.

3. $a, b \in \mathbb{N} . \supset . ab \in \mathbb{N} . \checkmark$
- Dem.* $a, b \in \mathbb{N} . \text{P1} : \supset : a \times 1 \in \mathbb{N} : \supset . 1 \in [b \in] \text{Ts} .$
 $a, b \in \mathbb{N} . b \in [b \in] \text{Ts} : \supset : a \times b \in \mathbb{N} . \S 1 \text{ P19}$ (2)
 $: \supset : ab + a \in \mathbb{N} . \text{P1}$
 $: \supset : a(b + 1) \in \mathbb{N} : \supset : b + 1 \in [b \in] \text{Ts} .$
(1)(2). $\supset . \text{Theor} .$

8

§4. Multiplication.

Definitions.

4. $a, b, c \in \mathbb{N} . \supset . (a + b)c = ac + bc . \checkmark$
- Nota.* Haec est prop. 5^a Euclidis elem.
libri VII.

- Dem.* $a, b \in \mathbb{N} . \text{P1} : \supset : 1 \in [c \in] \text{Ts} .$ (1)
 $a, b, c \in \mathbb{N} . c \in [c \in] \text{Ts}$ (2)
 $: \supset : (a + b)c = ac + bc . \S 1 \text{ P22}$
 $: \supset : (a + b)c + a + b = ac + bc + a + b .$
 $\text{P2} : \supset : (a + b)(c + 1) =$
 $a(c + 1) + b(c + 1) : \supset : c + 1 \in [c \in] \text{Ts} .$
(1)(2). $\supset . \text{Theor} .$

5. $a \in \mathbb{N} . \supset . 1 \times a = a .$

Note.

This is prop. 5^a of Euclid's elem.
book VII.

$$Dem. \quad 1 \in [a \in] \text{Ts.} \quad (1)$$

$$a \in [a \in] \text{Ts} . \circ . 1 \times a = a . \circ . 1 \times a + 1 = \quad (2)$$

$$a+1 . \times . 1 \times (a+1) = a+1 . \circ . a+1 \in [a \in]$$

Ts.

(1)(2). \circ . Theor.

$$6. \quad a, b \in \mathbb{N} . \circ . ba + a = (b + 1)a.$$

$$7. \quad a, b \in \mathbb{N} . \circ . ab = ba. \quad (\text{Eucl. VII,}$$

16)

$$Dem. \quad a \in \mathbb{N} . P5 . P1 \quad (1)$$

$$: \circ . a \times 1 = a = 1 \times a : \circ : 1 \in [b \in] \text{Ts.}$$

$$a, b \in \mathbb{N} . b \in [b \in] \text{Ts} \quad (2)$$

$$: \circ : ab = ba : \circ : ab + a = ba + a . P1$$

$$.P6: \circ : a(b + 1) = (b + 1)a : \circ : b + 1 \in$$

$[b \in] \text{Ts.}$

(1)(2). \circ . Theor.

$$8. \quad a, b, c \in \mathbb{N} . \circ . a(b + c) = ab + ac.$$

$$Dem. \quad P4 . P7 : \circ . \text{Theor.}$$

$$9. \quad a, b, c \in \mathbb{N} . a = b : \circ : ac = bc.$$

$$Dem. \quad a, b \in \mathbb{N} . a = b :: \circ :: 1 \in [c \in] \text{Ts}$$

$$\therefore c \in [c \in] \text{Ts}$$

$$. \circ : ac = bc . a = b : \circ : ac + a = bc + b :$$

$$\circ : a(c + 1) = b(c + 1) : \circ : c + 1 \in [c \in]$$

$$\text{Ts} :: \circ : c \in \mathbb{N} . \circ . \text{Ts.}$$

$$10. \quad a, b, c \in \mathbb{N} : a < b : \circ . (b - a)c = bc - ac. \quad (\text{Eucl. VII,}$$

7)

$$Dem. \quad \text{Hyp} . \circ : b - a \in \mathbb{N} . (b - a) + a = b : \circ :$$

$$(b - a)c + ac = bc : \circ : (b - a)c = bc - ac.$$

11. $a, b, c \in \mathbb{N} . a < b : \supset : ac < bc .$
- Dem.* Hyp $\supset : b - a \in \mathbb{N} .$ P3
 $\supset : (b - a)c \in \mathbb{N} .$ P10
 $\supset : bc - ac \in \mathbb{N} : \supset$ Thesis.
12. $a, b, c \in \mathbb{N} . \supset \therefore a < b . = . ac < bc : a =$
 $b . = . ac = bc : a > b . = . ac > bc .$
13. $a, b, a', b' \in \mathbb{N} . a < a' . b < b' : \supset :$
 $ab < a'b' .$
14. $a, b \in \mathbb{N} : \supset . ab . > \cup = . a .$
15. $a, b, c \in \mathbb{N} . \supset . a(bc) = abc .$

- Dem.* $a, b \in \mathbb{N} .$ P1 $\supset : 1 \in [c \in] \text{Ts} .$ (1)
- $a, b, c \in \mathbb{N} . c \in [c \in] \text{Ts}$ (2)
- $\supset : a(bc) = abc : \supset : a(bc) + ab =$
 $abc + ab : \supset : a(bc + b) = ab(c + 1) : \supset :$
 $a(b(c + 1)) = ab(c + 1) : \supset : c + 1 \in [c \in]$
Ts.
(1)(2). \supset . Theor.

§5. De potestatibus.

Definitiones.

1. $a \in \mathbb{N} . \supset . a^1 = a .$
2. $a, b \in \mathbb{N} . \supset . a^{b+1} = a^b a .$

§5. Powers.

Definitions.

Theoremata.

3. $a, b \in \mathbb{N} . \supset . a^b \in \mathbb{N} .$
- Dem.* $a \in \mathbb{N} . \text{P1} : \supset . 1 \in [b \in] \text{Ts} . \quad (1)$
- $a, b \in \mathbb{N} . b \in [b \in] \text{Ts} \quad (2)$
- $: \supset : a^b \in \mathbb{N} . \S 4 \text{P3} : \supset : a^b a \in \mathbb{N} .$
- $\text{P1} : \supset : a^{b+1} \in \mathbb{N} : \supset : b+1 \in [b \in] \text{Ts} .$
- (1)(2). $\supset . \text{Theor} .$
4. $a \in \mathbb{N} . \supset . 1^a = 1 .$
5. $a, b, c \in \mathbb{N} . \supset . a^{b+c} = a^b a^c .$
6. $a, b, c \in \mathbb{N} . \supset . (ab)^c = a^c b^c .$
7. $a, b, c \in \mathbb{N} . \supset . (a^b)^c = a^{bc} .$
8. $a, b, c \in \mathbb{N} . \supset . a < b . = . a^c < b^c : a =$
 $b . = . a^c = b^c : a > b . = . a^c > b^c .$
9. $a, b, c \in \mathbb{N} . a > 1 . \supset . b < c . =$
 $. a^b < a^c : b = c . = . a^b = a^c : b > c . =$
 $. a^b > a^c .$

§6. De divisione.

Explicationes.

Signum	/	legatur	<i>divisus per.</i>
»	D	»	<i>dividit, sive est divisor.</i>
»	¶	»	<i>est multiplex.</i>
»	Np	»	<i>numerus primus.</i>
»	π	»	<i>est primus cum.</i>

Theorems.

§6. Division.

Explications.

The sign	/	is	<i>divided by.</i>
		read	
”	D	”	<i>divides, or is a divisor of.</i>
”	¶	”	<i>is a multiple of.</i>
”	Np	”	<i>prime number.</i>
”	π	”	<i>is prime with.</i>

Definitiones.

1. $a, b \in N . \supset : b/a = N[x \in](xa = b).$
2. $a, b \in N . \supset : aDb . = . b/a = \Lambda.$
3. $a, b \in N . \supset : b \sqsubset a . = . aDb.$
4. $Np = N[x \in](\exists Dx . \exists > 1 . \exists < x := \Lambda).$
5. $a, b \in N . \supset :: a\pi b . \therefore =$
 $\therefore \exists Da . \exists Db . \exists > 1 := \Lambda.$
6. $a, b \in N . \supset \therefore \exists D(a, b) :=: \exists Da . \cap . \exists Db.$
7. $a, b \in N . \supset \therefore \exists \sqsubset(a, b) :=:$
 $\exists \sqsubset a . \cap . \exists \sqsubset b.$
 $ab/c = (ab)/c; a/b/c = (a/b)/c; a/b \times c =$
 $(a/b)c.$

Theoremata.

Nota. Haec theoremata ut in subtractione demonstrantur.

8. $a, b, a', b' \in N . a = a' . b = b' : \supset . a/b =$
 $a'/b'.$
9. $a, b, a', b' \in N . a = a' . b = b' : \supset :$
 $aDb . = . a'Db'.$
10. $a, b, c \in N . \supset : ac = b . = . c = b/a.$
11. $a, b \in N . \supset : aDb . = . b/a \in N.$
12. $a \in N . \supset . a/1 = a.$
13. $a \in N . \supset . a/a = 1.$
14. $a \in N . \supset . 1Da.$
15. $a \in N . \supset . aDa.$

*Definitions.**Theorems.*

Note.

These theorems are proved as for subtraction.

16. $a, b \in \mathbb{N} . ab/b = a.$
17. $a, b \in \mathbb{N} . aDb : \supset . a(b/a) = b.$
18. $a, b, c \in \mathbb{N} . cDb : \supset . a(b/c) = ab/c.$
19. $a, b, c \in \mathbb{N} . a \sqsubset bc : \supset . a/(bc) = a/b/c.$
20. $a, b, c \in \mathbb{N} . a \sqsubset b . b \sqsubset c : \supset . a/(b/c) =$
 $a/b \times c.$
21. $a, m, n \in \mathbb{N} . m > n : \supset . a^m/a^n = a^{m-n}.$
22. $a, b \in \mathbb{N} . \supset . aDab.$
23. $a, b, c \in \mathbb{N} . aDb . bDc : \supset . aDc.$
24. $a, b, c \in \mathbb{N} . aDbDc : \supset . c/a \sqsubset c/b.$
25. $a, b, c \in \mathbb{N} . cDa . cDb : \supset . (a + b)/c =$
 $a/c + b/c.$
26. $a, b, c \in \mathbb{N} . cDa . cDb . a > b : \supset :$
 $(a - b)/c = a/c - b/c.$
27. $a, b, c \in \mathbb{N} . cDa . cDb : \supset . cDa + b.$
28. $a, b, c \in \mathbb{N} . cDa . cDb . a > b :$
 $\supset . cDa - b.$

29. $a, b, c, m, n \in \mathbb{N} . cDa . cDb :$
 $\supset . cDma + nb.$
30. $a, b, c, m, n \in \mathbb{N} . cDa . cDb . ma > nb :$
 $\supset . cDma - nb.$
31. $a, b \in \mathbb{N} . aDb : \supset . a . < \cup = . b.$
- Dem.* Hyp . P11 . P17 . §4 P14 : $\supset : b/a \in$
 $\mathbb{N} . a(b/a) = b . a < \cup = a(b/a) : \supset .$
 Thesis.
32. $a, b \in \mathbb{N} . aDb . bDa : \supset . a = b.$

$$33. \quad a \in N . \supset . M \exists D a = a.$$

$$34. \quad a, b \in N . a > b : \supset . \exists D(a, b) = \\ \exists D(b, a - b).$$

$$Dem. \quad Hyp. P28 \quad (1)$$

$$: \supset \therefore xDa . xDb : \supset : xDb . xD(a - b)$$

$$Hyp. P27 : \supset \therefore xDb . xD(a - b) : \supset : \quad (2)$$

$$xD b . xD(b + (a - b)) : \supset : xDb . xDa .$$

$$(1)(2) \supset : Hyp. \quad (Th.)$$

$$\supset \therefore xDa . xDb := : xDb . xD(a - b).$$

$$35. \quad a, b \in N . \supset : M \exists D(a, b) \in N.$$

$$Dem. \quad 1Da . 1Db : \supset : \exists D(a, b) = \Lambda. \quad (1)$$

$$\exists D(a, b) . \exists > a := \Lambda. \quad (2)$$

$$(1)(2) . \S 3 P3 : \supset . Th.$$

$$36. \quad a, b \in N . \supset . \exists D(a, b) = \exists DM \exists D(a, b). \quad (Eucl. VII, \\ 2)$$

$$Dem. \quad k = N[c \in] (Hp. a < c . b < c : \supset . Ts.). \quad (1)$$

$$a \in N . b \in N . a < 1 . b < 1 := \Lambda. \quad (2)$$

$$(1)(2) . \supset . 1 \in K. \quad (3)$$

$$a, b \in N . a < c + 1 . b < c + 1 : \quad (4)$$

$$\supset \therefore a < c . b < c : \cup : a = c . b < c : \cup :$$

$$a < c . b = c : \cup : a = c . b = c.$$

$$c \in k . a, b \in N . a < c . b < c : \supset : Ts. \quad (5)$$

$$c \in k . a, b \in N . a = c . b < c : \supset : c \in \quad (6)$$

$$k . b < c / pa - b < c . \exists D(a, b) =$$

$$\exists D(b, a - b) : \supset : \exists D(b, a - b) =$$

$$\exists Dm \exists D(b, a - b) : \supset : \exists D(a, b) =$$

$$\exists DM \exists D(a, b) : \supset : Ts.$$

$$(a, b)[b, a](6) \supset . c \in k . a, b \in \quad (7)$$

$$\mathbf{N}.a < c.b = c : \mathfrak{O} : \mathbf{T}s.$$

$$c \in k.a, b \in \mathbf{N}.a = c.b = c : \mathfrak{O} : \quad (8)$$

$$\mathfrak{a}D(a, b) = \mathfrak{a}Dc = \mathfrak{a}DM \mathfrak{a}Dc =$$

$$\mathfrak{a}DM \mathfrak{a}D(a, b) : \mathfrak{O} : \mathbf{T}s.$$

$$(4)(5)(6)(7)(8).\mathfrak{O}.c \in k.a, b \in \quad (9)$$

$$\mathbf{N}.a < c+1.b < c+1 : \mathfrak{O} : \mathbf{T}s.$$

$$(9).\mathfrak{O}.c \in k.\mathfrak{O}.(c+1) \in k. \quad (10)$$

$$(1)(10).\mathfrak{O}.\therefore c \in \mathbf{N}. \mathbf{Hp}. a < c.b < c : \mathfrak{O} : \quad (11)$$

$$\mathbf{T}s.$$

$$(a+b)[c](11).\mathfrak{O} : \mathbf{Hp}. \mathfrak{O}. \mathbf{T}s. \quad (\mathbf{Th}.)$$

37.

$$a, b, m \in \mathbf{N}.\mathfrak{O}.M \mathfrak{a}D(am, bm) =$$

$$m \times M \mathfrak{a}D(a, b).$$

§7. Theoremata varia.

$$1. \quad a, b \in \mathbf{N}.a^2 + b^2 \mathfrak{D}7 : \mathfrak{O} : a \mathfrak{D}7.b \mathfrak{D}7.$$

$$2. \quad x \in \mathbf{N}.\mathfrak{O}.x(x+1) \mathfrak{D}2.$$

$$3. \quad x \in \mathbf{N}.\mathfrak{O}.x(x+1)(x+2) \mathfrak{D}6.$$

$$4. \quad x \in \mathbf{N}.\mathfrak{O}.x(x+1)(2x+1) \mathfrak{D}6.$$

$$5. \quad x \in \mathbf{N}.\mathfrak{O} : x.\pi.x+1.$$

$$6. \quad x \in \mathbf{N}.\mathfrak{O} : 2x-1.\pi.2x+1.$$

$$7. \quad x \in \mathbf{N}.\mathfrak{O}.(2x+1)^2 - 1 \mathfrak{D}8.$$

$$8. \quad a \in \mathbf{N}.a > 1 : \mathfrak{O} : \mathbf{Np}. \mathfrak{a} > 1. \mathfrak{a}Da : - = \quad (\mathbf{Eucl. VII},$$

$$\Lambda. \quad 31)$$

$$9. \quad a, b \in \mathbf{N}.\therefore b^2 > a.\therefore \mathfrak{a}Da.\mathfrak{a} > 1.\mathfrak{a} < b :=$$

$$\Lambda :: \mathfrak{O}.a \in \mathbf{Np}.$$

§7. Various theorems.

10. $a, b \in \mathbb{N} . a \in \mathbb{Np} . a \text{ - } Db : \supset : a \pi b .$ (Eucl. VII,
29)
11. $a, b, c \in \mathbb{N} . aDbc . a \pi b : \supset . aDc .$
12. $a, b \in \mathbb{N} . m = M \exists D(a, b) : \supset :$
 $a/m \pi . b/m .$
13. $a \in \mathbb{Np} . b, c \in \mathbb{N} . aDbc : \supset :$ (Eucl. VII,
 $aDb . \cup . aDc .$ 30)
14. $a \in \mathbb{Np} . b, n \in \mathbb{N} : \supset : aDb^n . = . aDb .$ (Eucl. IX,
12)
15. $a, b, c \in \mathbb{N} . a \pi b . cDa : \supset : c \pi b .$ (Eucl. VII,
23)
16. $a, b, c \in \mathbb{N} . \supset : . a \pi b . a \pi c := . a \pi bc .$ (Eucl. VII,
24)
17. $a, b, c \in \mathbb{N} . b \pi c . bDa . cDa : \supset . bcDa .$
18. $a, b, c \in \mathbb{N} \text{ } p a \pi b : \supset : \exists D(ac, b) =$
 $\exists D(c, b) .$
19. $a, b \in \mathbb{N} . \supset . \mathbb{W} \exists \mathbb{D}(a, b) \in N .$
20. $a, b \in \mathbb{N} . \supset . \mathbb{W} \exists \mathbb{D}(a, b) =$ (Eucl. VII,
 $ab/M \exists D(a, b) .$ 34)
21. $a, b, c \in \mathbb{N} . c \mathbb{D} a . c \mathbb{D} b : \supset :$ (Eucl. VII,
 $c \mathbb{D} \mathbb{W} \exists D(a, b) .$ 35)
22. $x \in \mathbb{N} . x < 41 : \supset . 41 - x + x^2 \in \mathbb{Np} .$
23. $M . \mathbb{Np} := \Lambda .$ (Eucl. IX,
20)
23. $n \in \mathbb{Np} . a \in \mathbb{N} . a \text{ - } \mathbb{D} n : \supset . a^{n-1} - 1 \mathbb{D} n .$ (Fermat)

§8. Numerorum rationes.

Explicationes.

Si $p, q \in \mathbb{N}$, tunc $\frac{p}{q}$ legitur *ratio numeri p numero q*.

Signum R legitur *duorum numerorum ratio*, et indicat numeros rationales positivos.

Definitiones.

1. $m, p, q \in \mathbb{N} . \odot . m \frac{p}{q} = mp/q.$
2. $p, q, p', q' \in \mathbb{N} . \odot :: \frac{p}{q} = \frac{p'}{q'} . = \therefore x \in$
 $\mathbb{N} . x \frac{p}{q}, x \frac{p'}{q'} \in \mathbb{N} : \odot_x . x \frac{p}{q} = x \frac{p'}{q'}.$
3. $R = :: [x \in] \therefore p, q \in \mathbb{N} . \frac{p}{q} = x : - = \Lambda.$
4. $p \in \mathbb{N} . \odot . \frac{p}{1} = p.$

Theoremata.

5. $p, q, p', q' \in \mathbb{N} . \odot :: \frac{p}{q} = \frac{p'}{q'} . = . pq' =$ (Eucl. VII,
 $p'q.$ 19)
- Dem.* Hp. $\frac{p}{q} = \frac{p'}{q'} : \odot \therefore qq', qq' \frac{p}{q}, qq' \frac{p'}{q'} \in \mathbb{N} .$ (1)
P2 $\therefore \odot :: qq' \frac{p}{q} = qq' \frac{p'}{q'} . qq' \frac{p}{q} =$
 $pq' . qq' \frac{p'}{q'} = p'q \therefore \odot \therefore pq' = p'q.$
Hp. $pq' = p'q \therefore \odot \therefore x \in \mathbb{N} . x \frac{p}{q}, x \frac{p'}{q'} \in$ (2)
 $\mathbb{N} : \odot_x : xp'q' = xp'q : \odot : (x \frac{p}{q})qq' =$
 $(x \frac{p'}{q'})qq' : \odot : x \frac{p}{q} = x \frac{p'}{q'}.$
(1)(2). $\odot .$ Th.

§8. Rational numbers.

Explications.

If $p, q \in \mathbb{N}$, then $\frac{p}{q}$ is read *the ratio of the number p to the number q*.

The sign R is read *ratio of two numbers*, and indicates the positive rational numbers.

Definitions.

Theorems.

6. $m, p, q \in \mathbb{N} . \odot . \frac{p}{q} = \frac{mp}{mq} .$ (Eucl. VII,
17)
7. $p, q \in \mathbb{N} . m \in \mathbb{N} . mDp . mDq : \odot . \frac{p}{q} =$
 $\frac{p/m}{q/m} .$
8. $p, q, p', q' \in \mathbb{N} . p\pi q . p'\pi q' . \frac{p}{q} = \frac{p'}{q'} : \odot :$
 $p = p' . q = q' .$
9. $p, q, p', q' \in \mathbb{N} . p'\pi q' . \frac{p}{q} = \frac{p'}{q'} : \odot :$
 $p'/p = q'/q = M \exists D(p, q) .$
10. $p, q, p', q' \in \mathbb{N} . \frac{p}{q} = \frac{p'}{q'} . p\pi q . q' < q :=$ (Eucl. VII,
 $\Lambda .$ 21)
11. $p, q, p', q' \in \mathbb{N} : \odot : \frac{p}{q} = \frac{p'}{q'} . = . \frac{p}{p'} =$ (Eucl. VII,
 $\frac{q}{q'} . = . \frac{q}{p} = \frac{q'}{p'} .$ 13)
12. $p, q \in \mathbb{N} . \odot :: [m \in \mathbb{N}] : m \in \mathbb{N} . m \frac{p}{q} \in$
 $\mathbb{N} \therefore \text{-} = \Lambda .$
- 12'. $a\check{\alpha} \in \mathbb{R} . \odot :: [m \in \mathbb{N}] : m \in \mathbb{N} . ma \in$
 $\mathbb{N} \therefore \text{-} = \Lambda .$

13. $p, q, p', q' \in \mathbb{N} . \odot :: [(r, s, l) \in \mathbb{N}] : r, s, t \in$
 $\mathbb{N} . \frac{p}{q} = \frac{r}{t} . \frac{p'}{q'} = \frac{s}{t} \therefore \text{-} = \Lambda .$
- 13'. $a, b \in \mathbb{R} . \odot :: [(r, s, t) \in \mathbb{N}] : r, s, t \in \mathbb{N} . a =$
 $\frac{r}{t} . b = \frac{s}{t} \therefore \text{-} = \Lambda .$
14. $a, b, c \in \mathbb{R} . \odot :: [(m, n, p, q) \in \mathbb{N}] :$
 $m, n, p, q \in \mathbb{N} . a = \frac{m}{q} . b = \frac{n}{q} . c =$
 $\frac{p}{q} \therefore \text{-} = \Lambda .$
15. $p, q, r \in \mathbb{N} . a = \frac{p}{r} . b = \frac{q}{r} : \odot : a = b . =$
 $. p = q .$
16. $m \in \mathbb{N} . a, b \in \mathbb{R} . a = b . ma \in \mathbb{N} :$

$$\mathfrak{O}.mb \in N.$$

$$17. \quad a, b, c \in \mathbf{R} . \mathfrak{O} \therefore a = a.$$

$$\mathfrak{O} \therefore a = b . = . b = a.$$

$$\mathfrak{O} \therefore a = b . b = c : \mathfrak{O} . a = c.$$

$$18. \quad N \supset \mathbf{R}.$$

Definitiones.

Definitions.

$$19. \quad a, b \in \mathbf{R} . \mathfrak{O} :: a < b . = \therefore x \in \mathbf{N} . xa, xb \in$$

$$\mathbf{N} : \mathfrak{O} . xa < xb.$$

$$20. \quad a, b \in \mathbf{R} . \mathfrak{O} : b > a . = . a < b.$$

Theoremata.

Theorems.

$$21. \quad p, q, r \in \mathbf{N} . a = \frac{p}{r} . b = \frac{q}{r} : \mathfrak{O} : a < b . =$$

$$. p < q.$$

$$22. \quad p, q, p', q' \in \mathbf{N} . \mathfrak{O} : \frac{p}{q} < \frac{p'}{q'} . =$$

$$. pq' < p'q.$$

$$23. \quad p, q, r \in \mathbf{N} . a = \frac{r}{p} . b = \frac{r}{q} : \mathfrak{O} : a < b . =$$

$$. p > q.$$

$$24. \quad p, q, p', q' \in \mathbf{N} . \frac{p}{q} < \frac{p'}{q'} :$$

$$\mathfrak{O} . \frac{p}{q} < \frac{p+p'}{q+q'} < \frac{p'}{q'}.$$

$$25. \quad a \in \mathbf{R} . \mathfrak{O} \therefore ? . \mathfrak{O} > a : \text{---} = \Lambda.$$

$$26. \quad a \in \mathbf{R} . \mathfrak{O} \therefore \mathbf{R} . \mathfrak{O} < a : \text{---} = \Lambda.$$

$$27. \quad a, b \in \mathbf{R} . a < b : \mathfrak{O} \therefore \mathbf{R} . \mathfrak{O} > a . \mathfrak{O} < b :$$

$$\text{---} = \Lambda.$$

$$28. \quad a, b \in \mathbf{R} : \mathfrak{O} \therefore a < b . a = b := \Lambda.$$

$$\mathfrak{O} \therefore a > b . a = b := \Lambda.$$

$$\mathfrak{O} \therefore a < b . a > b := \Lambda.$$

$$\supset \therefore a < b . a =$$

$$b . a > b := \Lambda.$$

$$29. \quad a, b, c \in \mathbb{R} : \supset \therefore a < \cup = b . b < c : \supset :$$

$$a < c.$$

$$\supset \therefore a < b . b < \cup = c : \supset :$$

$$a < c.$$

Definitiones.

$$30. \quad a, b \in \mathbb{R} . \supset . a + b = [c \in](c \in \mathbb{R} \therefore x \in$$

$$\mathbb{N} . xa, xb, xc \in \mathbb{N} : \supset_x . xa + xb = xc).$$

$$31. \quad a, b \in \mathbb{R} . \supset :: b - a = \therefore [x \in](x \in$$

$$\mathbb{R} . a + x = b).$$

$$32. \quad a, b \in \mathbb{R} . \supset . ab = [c \in](c \in \mathbb{R} \therefore x \in$$

$$\mathbb{N} . xa, (xa)b, xc \in \mathbb{N} : \supset_x . (xa)b = xc).$$

$$33. \quad a, b \in \mathbb{R} . \supset . b/a = [x \in](x \in \mathbb{R} . ax = b).$$

Theoremata.

$$34. \quad p, q, r \in \mathbb{N} . \supset \frac{p}{r} + \frac{q}{r} = \frac{p+q}{r}.$$

$$35. \quad a, b \in \mathbb{R} . \supset . a + b \in \mathbb{R}.$$

$$36. \quad p, q, r \in \mathbb{N} . p < q : \supset . \frac{q}{r} - \frac{p}{r} = \frac{q-p}{r}.$$

$$37. \quad a, b \in \mathbb{R} . a < b : \supset . b - a \in \mathbb{R}.$$

$$38. \quad p, q, p', q' \in \mathbb{N} . \supset . \frac{p}{q} \frac{p'}{q'} = \frac{pp'}{qq'}.$$

$$39. \quad a, b \in \mathbb{R} . \supset . ab \in \mathbb{R}.$$

$$40. \quad p, q, p', q' \in \mathbb{N} . \supset . \frac{p}{q} / \frac{p'}{q'} = \frac{pq'}{p'q}.$$

$$41. \quad a, b \in \mathbb{R} . \supset . b/a \in \mathbb{R}.$$

$$42. \quad p, q \in \mathbb{N} . \supset . \frac{p}{q} = \frac{p}{q}.$$

Definitions.

Theorems.

§9. Rationalum systemata. Irrationales.

Explicatio.

Si $a \in \mathbf{K R}$, signum $T a$ legitur *terminus summus*, vel *limes summus classis a*.

Supra hoc novum ens relationes ac operationes tantum definimus.

Definitiones.

1. $a \in \mathbf{K R} . x \in \mathbf{R} : \mathcal{O} :: x < T a . =$
 $\therefore a . \mathcal{O} > x : - = \Lambda .$
2. $a \in \mathbf{K R} . x \in \mathbf{R} : \mathcal{O} :: x = T a . = : :$
 $a . \mathcal{O} > x := \Lambda :: u \in \mathbf{R} . u < x :$
 $\mathcal{O}_x \therefore a . \mathcal{O} > u : - = \Lambda .$
3. $a \in \mathbf{K R} . x \in \mathbf{R} : \mathcal{O} : x > T a . = :$
 $x - - < T a . x - = T a .$

Theorema.

4. $x \in \mathbf{R} . \mathcal{O} :: x = \therefore T : \mathbf{R} . \mathcal{O} < x .$

Explicatio.

Signum Q legitur *quantitas*, numerosque indicat reales positivos, rationales aut irrationales, 0 et ∞ exceptis.

Definitiones.

§9. The system of rationals. Irrationals.

Explanation.

If $a \in \mathbf{K R}$, the sign $T a$ is read *upper boundary*, or *upper limit of the class a*. We shall define only a few relations and operations on this new entity.

Definitions.

Theorem.

Explanation.

The sign Q is read *quantity*, and indicates the positive real numbers, rational or irrational, with the exception of 0 and ∞ .

Definitions.

5. $Q = [x \in](a \in K R : a - = \Lambda : R \in$
 $> T a . \min = \Lambda : T a = x . \therefore - = \Lambda).$
6. $a, b \in Q . \supset :: a = b . = \therefore R . \exists < a :=:$
 $R . \exists < b .$
7. $a, b \in Q . \supset :: a < b . =$
 $\therefore R . \exists > a . \exists < b := - = \Lambda .$
8. $a, b \in Q . \supset : b > a . = . a < b .$

Theoremata.

9. $a \in Q . \supset \therefore R . \exists < a : - = \Lambda .$
10. $a \in Q . \supset \therefore R . \exists > a : - = \Lambda .$
11. $R \supset Q .$
- Subsistunt quoque propositiones
 quae a P17, 28, 29 in §8 obtinentur,
 si loco R legatur Q.

Definitiones.

12. $a, b \in Q . \supset . a + b = T[z \in]([x, y) \in] :$
 $x, y \in R . x < a . y < b . x + y = z . \therefore - =$
 $\Lambda).$
13. $a, b \in Q . \supset . ab = T[z \in]([x, y) \in] : x, y \in$
 $R . x < a . y < b . xy = z . \therefore - = \Lambda).$
- Ut valeant hae definitiones,
 demonstrandum est subsistere
 propositiones 12 et 13, si $a, b \in R .$

Substractionem et divisionem ut

Theorems.

The propositions obtained from P17,
 28, 29 in §8 also hold, by reading Q
 for R.

Definitions.

In order for these definitions to have
 meaning, it must be proved that
 propositions 12 and 13 hold, if
 $a, b \in R .$

Subtraction and division could be

operationes inversas additiones et
multiplicationis definire licet,
illarumque proprietates demonstrare.

defined as the inverse operations to
addition and multiplication, and
their properties could be proved.

§10. Quantitatum systemata.

§10. System of quantities.

Explicationes.

Si $a \in K Q$, signa $I a, E a, L a$ leguntur: *interior, exterior, limes classis a*.

Explanations.

If $a \in K Q$, the signs $I a, E a, L a$ are read: *interior, exterior, limit of the class a*.

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Definitiones.

1. $a \in K Q . \supset I a = Q[x \in]([u, v) \in] ::$
 $u, v \in Q . \therefore u < x < v . \therefore \exists > u . \exists < v : \supset :$
 $a : . : - = \Lambda).$
2. $a \in K Q . \supset . E a = I(-a).$
3. $a \in K Q . \supset . L a = (-I a)(-E a).$

Definitions.

4. $a \in K Q . x, u, v \in$
 $Q . u < x < v . (\exists > u . \exists < v : \supset a) : \supset . x \in$
 $I a.$
 5. $a \in K Q . x \in I a : \supset : [(u, v) \in](u, v \in$
 $Q . \therefore u < x < v . \therefore \exists > u . \exists < v : \supset : a) - =$
 $\Lambda.$
- Dem.* P1 = (P4)(P5).

Theorems.

6. $a \in K Q . u, v \in Q . (\exists > u . \exists < v : \neg a) \therefore \neg \therefore \exists > u . \exists < v : \neg I a .$
- Dem.* P6 = P4.
7. $a \in K Q . \neg . I a \neg a .$
8. $a \in K Q . \neg . II a = I a .$
- Dem.* Hp. (Ia)[a] P7 : $\neg . II a \neg I a$ (1)
- Hp. $x, u, v \in$ (2)
- $Q . u < x < v . (\exists > u . \exists < v : \neg a) .$ P6
- $: \neg : u, v \in Q . u < x < v . (\exists > u . \exists < v : \neg I a)$
- Hp. $x \in I a . (2) : \neg : x \in II a$ (3)
- Hp. (3) : $\neg : I a II a$ (4)
- Hp. (1).(4) : $\neg : Ts .$ (Theor.)
9. $a, b \in K Q . a \neg b : \neg . I a \neg I b$
- Dem.* Hp. (1)
- $x, u, v \in Q . u < x < v . (\exists > u . \exists < v : \neg a) : \neg \therefore \exists > u . \exists < v : \neg b$
- Hp. $x \in I a : \neg : x \in I b$ (Theor.)
10. $a, b \in K Q : \neg : I(ab) \neg I a$
- Dem.* $(ab, a)[a, b]$ P9 . = . P10
11. $a, b \in K Q . \neg . I(ab) \neg (I a)(I b)$
- Dem.* P11 =: P10 . $\cap . (b, a)[a, b]$ P10
12. $a, b \in K Q . \neg . I a \neg I(a \cup b)$
13. $a, b \in K Q . \neg . I a \cup I b \neg I(a \cup b)$
14. $a, b \in K Q . \neg . I(ab) = (I a)(I b)$
- Dem.* Hp. P11 : $\neg . I(ab) \neg (I a)(I b)$ (1)

$$\text{Hp. } x \in \mathbf{Q} . u, v \in \quad (2)$$

$$\mathbf{Q} . u < x < v . (\exists > u . \exists < v :$$

$$\supset a) . u', v' \in$$

$$\mathbf{Q} . u' < x < v' . (\exists > u' . \exists < v' :$$

$$\supset b) . u'' = M(u \cup u') . v'' = W(v, v') : \supset :$$

$$u'', v'' \in$$

$$Q . u'' < x < v'' . (\exists > u'' . \exists > v'' : \supset : ab)$$

$$\text{Hp. } x \in \mathbf{I} a . x \in \mathbf{I} b . (2) : \supset . x \in \mathbf{I}(ab) \quad (3)$$

$$\text{Hp. } (3) : \supset : (\mathbf{I} a)(\mathbf{I} b) \supset \mathbf{I}(ab) \quad (4)$$

$$\text{Hp. } (1) \text{ \& } (4) : \supset . \text{Ts.}$$

$$15. \quad a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{E} a \supset \neg a$$

$$\text{Dem.} \quad \text{P15} = (\neg a)[a] \text{ P7}$$

$$16. \quad a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} a . \mathbf{E} a := \Lambda$$

$$\text{Dem.} \quad \text{Hp. P7 . P15}$$

$$: \supset . \mathbf{I} a . \mathbf{E} a : \supset : a \neg a := \Lambda$$

$$17. \quad a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} \mathbf{E} a = \mathbf{E} a$$

$$\text{Dem.} \quad \text{P17} = (\neg a)[a] \text{ P8}$$

$$18. \quad a, b \in \mathbf{K} \mathbf{Q} . b \supset a : \supset . \mathbf{E} a \supset \mathbf{E} b$$

$$\text{Dem.} \quad \text{P18} = (\neg a, \neg b)[a, b] \text{ P9}$$

$$19. \quad a, b \in \mathbf{K} \mathbf{Q} . \supset : \mathbf{E} a \cup \mathbf{E} b . \supset \mathbf{E}(ab)$$

$$20. \quad a, b \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{E}(a \cup b) = (\mathbf{E} a)(\mathbf{E} b)$$

$$\text{Dem.} \quad \text{P20} = (\neg a, \neg b)[a, b] \text{ P14}$$

$$21. \quad a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{L}(\neg a) = \mathbf{L} a$$

$$22. \quad a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} a . \mathbf{L} a := \Lambda$$

$$\supset . \mathbf{E} a . \mathbf{L} a := \Lambda$$

$$\supset . \neg \mathbf{I} a . \neg \mathbf{E} a . \neg \mathbf{L} a := \Lambda$$

$$\text{Dem.} \quad \text{P22} = \text{P3}$$

$$23. \quad a \in \mathbf{K} \mathbf{Q} . \supset : a \supset . \mathbf{I} a \cup \mathbf{L} a$$

24. $a \in K Q . \supset . I(a L a) = \Lambda$

Dem. Hp. P14 . P7 . P22 : $\supset : I(a L a) . =$
 $. I a I L a . \supset . I a L a . = . \Lambda$

25. $a, b \in K Q . a \supset b : \supset : L a . \supset . I b \cup L b$

Dem. Hp. P18 : $\supset : E b \supset E a : \supset :$
 $I a \cup L a . \supset . I b \cup L b : \supset . Ts.$

26. $a, b \in K Q . \supset :$

$L(ab) \supset . I a L b \cup I b L a \cup L a L b$

Dem. Hp. $\supset : ab \supset a . ab \supset b . P25$
 $: \supset : L(ab) \supset I a \cup L a . L(ab) \supset I b \cup$
 $L b : \supset : L(ab) \supset (I a \cup L a)(I b \cup$
 $L b) . L(ab)(I a)(I b) = L(ab) I(ab) =$
 $\Lambda : \supset : Ts.$

26' $a, b \in K Q . \supset . L(ab) \supset L a \cup L b$

27. $a, b \in K Q . \supset : L(a \cup b) =$
 $L a E b \cup L b E a \cup L a L b$

Dem. P27 = (-a, -b)[a, b] P26

27'. $a, b \in K Q . \supset : L(a \cup b) \supset L a \cup L b$

28. $a \in K Q . \supset . L I a \supset L a$

Dem. Hp. P7 : $\supset : I a \supset a . P25$ (1)
 $: \supset : L I a \supset I a \cup L a$
 Hp. P8 . P22 (2)
 $: \supset . L I a I a = L I a I I a = \Lambda$
 (1)(2). $\supset .$ Theor.

28'. $a \in K Q . \supset L E a \supset L a$

29. $a \in K Q . \supset . L L a \supset L I a \cup L E a$

<i>Dem.</i>	Hp. $\alpha C : \mathbf{L} \mathbf{L} a = \mathbf{L}(\mathbf{I} a \cup \mathbf{E} a)$. P27'		
	$: \supset$. Ts.		
29'.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{L} \mathbf{L} a \supset \mathbf{L} a$		
<i>Dem.</i>	P29 . P28 . P28' $: \supset$. Theor.		
30.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{L} a = \mathbf{I} \mathbf{L} a \cup \mathbf{L} \mathbf{L} a$		
<i>Dem.</i>	Hp. P23 $: \supset . \mathbf{L} a \supset \mathbf{I} \mathbf{L} a \cup \mathbf{L} \mathbf{L} a$	(1)	
	Hp. P7 $: \supset . \mathbf{I} \mathbf{L} a \supset \mathbf{L} a$	(2)	
	Hp. P29' $: \supset . \mathbf{L} \mathbf{L} a \supset \mathbf{L} a$	(3)	
	(1)(2)(3). \supset . Theor.		
31.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{L} \mathbf{I} \mathbf{L} a \supset \mathbf{L} \mathbf{L} a \checkmark$		
<i>Dem.</i>	P31 = $(\mathbf{L} a)[a]$ P28		
32.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} \mathbf{L} \mathbf{L} a = \Lambda$		
<i>Dem.</i>	Hp. P29'		
	$: \supset : \mathbf{L} \mathbf{L} a = \mathbf{L} a \mathbf{L} \mathbf{L} a . (\mathbf{L} a)[a]$ P24 $: \supset$		
	Ts.		
33.	$a \in \mathbf{K} \mathbf{Q} . \supset : \mathbf{I} \mathbf{L} \mathbf{I} \mathbf{L} a = \Lambda$		
<i>Dem.</i>	P31 . P32 $: \supset . P33$		
34.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{L} \mathbf{L} \mathbf{L} a = \mathbf{L} \mathbf{L} a$		
<i>Dem.</i>	$(\mathbf{L} a)[a]$ P30 /p P32 $: \supset$. Theor.		
35.	$a, b \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} a \mathbf{L} b \supset \mathbf{L}(ab)$		
<i>Dem.</i>	Hp. P14	(1)	(1)
	$: \supset . \mathbf{I} a \mathbf{L} b \mathbf{I}(ab) = \mathbf{I} a \mathbf{I} b \mathbf{L} b = \Lambda$		
	Hp. P2 . P14	(2)	(2)
	$: \supset . \mathbf{I} a \mathbf{L} b \mathbf{E}(ab) = \mathbf{I} a \mathbf{L} b \mathbf{I}(\neg a \cup$		
	$\neg b) = \mathbf{I}(a \neg b) \mathbf{L} b = \mathbf{I} a \mathbf{E} b \mathbf{L} b = \Lambda$		
	(1)(2) \supset Theor.		
36.	$a, b \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} a \mathbf{L} b \cup \mathbf{I} b \mathbf{L} a \supset \mathbf{L} ab$.	(Vide P26)	(Cf. p. 26)
<i>Dem.</i>	P36 =: P35 . $(b, a)[a, b]$ P35		

37.	$a, b \in$	(Vide P27)	(Cf. p. 27)
	$\mathbf{K} \mathbf{Q} . \supset . \mathbf{E} a \mathbf{L} b \cup \mathbf{E} b \mathbf{L} a \cup \mathbf{L}(a \cup b).$		
<i>Dem.</i>	$\text{P37} = (-a, -b)[a, b] \text{ P36}$		
38.	$a, b \in$	(Vide P13)	(Cf. p. 13)
	$\mathbf{K} \mathbf{Q} . \supset . \mathbf{I}(a \cup b) \supset \mathbf{I} a \cup \mathbf{I} b \cup \mathbf{L} a \mathbf{L} b$		
<i>Dem.</i>	$\text{Hp. } \supset . \mathbf{I}(a \cup b) \supset (\mathbf{I} a \cup \mathbf{L} a \cup$	(1)	(1)
	$\mathbf{E} a)(\mathbf{I} b \cup \mathbf{L} b \cup \mathbf{E} b)$		
	$\text{Hp. P20 . P16 : } \supset . \mathbf{I}(a \cup b) \mathbf{E} a \mathbf{E} b =$	(2)	(2)
	$\mathbf{I}(a \cup b) \mathbf{E}(a \cup b) = \Lambda$		
<hr/>			
		20	
	$\text{Hp. P37: } \supset : \mathbf{I}(a \cup b) \checkmark (\mathbf{E} a \mathbf{L} b \cup$	(3)	(3)
	$\mathbf{E} b \mathbf{L} a) . \supset . \mathbf{I}(a \cup b) \mathbf{L}(a \cup b) . = \Lambda$		
	$(1)(2)(3) . \supset . \text{Theor.}$		
38'.	$a, b \in$	(Vide P19)	(Cf. p. 19)
	$\mathbf{K} \mathbf{Q} . \supset . \mathbf{E}(ab) \supset \mathbf{E} a \cup \mathbf{E} b \cup \mathbf{L} a \mathbf{L} b$		
39.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} \mathbf{L} a \mathbf{L} \mathbf{I} a = \Lambda$		
<i>Dem.</i>	Hp. P36		
	$: \supset : \mathbf{I} \mathbf{L} a \mathbf{L} \mathbf{I} a \supset \mathbf{L}(\mathbf{L} a \mathbf{I} a) = \Lambda$		
40.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{L} \mathbf{I} a \supset \mathbf{L} \mathbf{L} a$		
<i>Dem.</i>	$\text{Hp. P28 . P30 . P39 : } \supset \text{Theor.}$		
40'.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{L} \mathbf{E} a \supset \mathbf{L} \mathbf{L} a$		
41.	$a \in \mathbf{K} \mathbf{Q} . \supset \mathbf{L} \mathbf{L} a = \mathbf{L} \mathbf{I} a \cup \mathbf{L} \mathbf{E} a$		
<i>Dem.</i>	$\text{P29 . P40 . P40' : } \supset . \text{Theor.}$		
42.	$a \in \mathbf{K} \mathbf{Q} . \supset . \mathbf{I} \mathbf{L} \mathbf{I} a = \Lambda$		
	$\supset . \mathbf{I} \mathbf{L} \mathbf{E} a = \Lambda$		
	$\supset . \mathbf{L} \mathbf{L} \mathbf{I} a = \mathbf{L} \mathbf{I} a$		
	$\supset . \mathbf{L} \mathbf{L} \mathbf{E} a = \mathbf{L} \mathbf{E} a$		

43.	$a, b \in K \mathcal{Q} . \supset . \mathbf{I}(\mathbf{I} a \cup \mathbf{I} b) = \mathbf{I} a \cup \mathbf{I} b$	
<i>Dem.</i>	Hp. P7 : $\supset . \mathbf{I}(\mathbf{I} a \cup \mathbf{I} b) \supset \mathbf{I} a \cup \mathbf{I} b$	(1)
	Hp. P8 . P13 : $\supset : \mathbf{I} a \cup \mathbf{I} b . =$	(2)
	$. \mathbf{I} \mathbf{I} a \cup \mathbf{I} \mathbf{I} b . \supset . \mathbf{I}(\mathbf{I} a \cup \mathbf{I} b)$	
	(1)(2) \supset Theor.	
44.	$a, b \in K \mathcal{Q} . \supset . \mathbf{I}(\mathbf{L} \mathbf{L} a \cup \mathbf{L} \mathbf{L} b) = \Lambda$	
<i>Dem.</i>	Hp. P38 . P32 . P34 :	(1)
	$\supset . \mathbf{I}(\mathbf{L} \mathbf{L} a \cup \mathbf{L} \mathbf{L} b) \supset \mathbf{L} \mathbf{L} a \mathbf{L} \mathbf{L} b \supset \mathbf{L} \mathbf{L} a$	
	Hp. (1) . P8	
	: $\supset . \mathbf{I}(\mathbf{L} \mathbf{L} a \cup \mathbf{L} \mathbf{L} b) \cup \mathbf{I} \mathbf{L} \mathbf{L} a = \Lambda$	
45.	$a \in K \mathcal{Q} . \supset . \mathbf{I}(\mathbf{I} a \cup \mathbf{E} a) = \mathbf{I} a \cup \mathbf{E} a$	
<i>Dem.</i>	P8 . P17 . (-a)[b] P43 : \supset . Theor.	
45'.	$a \in K \mathcal{Q} . \supset . \mathbf{E} \mathbf{L} a = \mathbf{I} a \cup \mathbf{E} a$	
46.	$a \in K \mathcal{Q} . \supset . \mathbf{E} \mathbf{I} a = \neg(\mathbf{I} a \cup \mathbf{L} \mathbf{I} a)$	
46'.	$a \in K \mathcal{Q} . \supset . \mathbf{E} \mathbf{E} a = \neg(\mathbf{E} a \cup \mathbf{L} \mathbf{E} a)$	46'.

END

Endnotes

1 Giuseppe Peano's footnote (original):

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Vide:

C. S. Pierce, *On the Algebra of logic*; American Journal, III, 15; VII, 180.

Jevons, *The principles of science*, London, 1883.

Mc.Coll., *The calculus of equivalent statements*, Proceedings of the London Math. Society, 1878, Vol. IX, 9. Vol X, 16.

1 Giuseppe Peano's footnote (translated):

Boole:

The mathematical analysis of logic ... (Cambridge, 1847.)

'The calculus of logic,' *Camb. and Dublin Math. J.*, 3 (1848), 193-98.

An investigation of the laws of thought ... (London, 1854).

E. Schröder:

Der Operationskreis des Logikkalkulus (Leipzig, 1877).

He had already treated several matters pertaining to logic in a preceding work.

Lehrbuch der Arithmetik und Algebra ... (Leipzig, 1873).

I gave a very brief presentation of the theories of Boole and Schröder in my book *Calcolo geometric* etc. (Torino, 1888).

Cf:

C. S. Pierce, 'On the Algebra of logic,' *American J. Math.*, 3 (1880), 15-57; 7 (1885), 180-202.

Jevons, *The principles of science* (London, 1883).

Mc.Coll., 'The calculus of equivalent statements,' *Proc. London Math. Soc.*, 9 (1878), 9-20; 10 (1878), 16-28.

2 The 2 other translation, mentioned at the beginning of this current document, translated "ähnlich" literally to "similar", instead of "equivalent". However, additional information can be found in a footnote of the first translation:

"Today "similar" has another meaning and instead we would say "equivalent"."

G. Peano, (1889), "The principles of arithmetic presented by a new method" in: J. van Heijenoort (ed.), *From Frege to Gödel. A source book in mathematical logic. 1879-1931*, Cambridge: Harvard University Press, 1967, p. 93.