

Informatics II, Spring 2023, Exercise 3

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Linear Search and Binary Search

Task 1. Consider an array A with n distinct integers that are sorted in an ascending order and an integer t .

a) The C function `linear_search` traverses the integers in A , one after another, from the beginning. If t is found in A `linear_search` returns 1, otherwise 0. Complete the C function `linear_search(int A[], int n, int t)` in `task1.c` file.

b) The C function `binary_search(int A[], int n, int t)` that employs binary search to find integer t in A . Reference the pseudocode for binary search in SL02 to implement the `binary_search` function. If t is found in A `binary_search` returns 1, otherwise 0. Complete the C function `binary_search(int A[], int n, int t)` in `task1.c` file.

c) What are the asymptotic complexity for the C functions `linear_search` and `binary_search`.

d) Compile `task1.c` file. Run your codes with the following parameters for n and t :

- $n = 1000000$, $t = 1000000$
- $n = 10000000$, $t = 10000000$
- $n = 100000000$, $t = 100000000$.

Report the run time growth for `linear_search` and `binary_search`, respectively.

Algorithmic Complexity

Task 2. Below is a pseudocode of a function named `whatDoesItDo`, which takes an array $A[1..n]$ of n integers and an integer k as inputs.

```
Algo: whatDoesItDo(A, n, k)
result = -1000
for i = 1 to n do
    current = 0
    for j = i to n by k do
        current = current + A[j]
    if current > result then
        result = current
return result
```

Note: In the above pseudocode, **for** $j = i$ **to** n **by** k means we do not increase j by 1, but each time, we increase it by k , i.e., $j = j + k$.

- a) Perform exact analysis of the running time of the algorithm.
- b) Determine the asymptotic complexity of the algorithm?

Asymptotic Complexity

Task 3. Calculate the asymptotic tight bound for the following functions and rank them by their order of growth (lowest first). Clearly work out the calculation step by step in your solution.

$$\begin{aligned}f_1(n) &= (2n + 3)! \\f_2(n) &= 2 \log(6^{\log n^2}) + \log(\pi n^2) + n^3 \\f_3(n) &= 4^{\log_2 n} \\f_4(n) &= 12\sqrt{n} + 10^{223} + \log 5^n \\f_5(n) &= 10^{\lg 20} n^4 + 8^{229} n^3 + 20^{231} n^2 + 128n \log n \\f_6(n) &= \log n^{2n+1} \\f_7(n) &= \log^2(n) + 50\sqrt{n} + \log(n) \\f_8(n) &= 14400\end{aligned}$$

Special Case and Correctness Analysis

Task 4. Consider the algorithm `algo1`. The input parameters are an array $A[1..n]$ with n distinct integers and $k \leq n$.

```
Algo: algo1(A, n, k)
sum = 0;
for i = 1 to k do
    maxi = i;
    for j = i to n do
        if A[j] > A[maxi] then
            maxi = j;
    sum = sum + A[maxi];
    swp = A[i];
    A[i] = A[maxi];
    A[maxi] = swp;
return sum
```

- Specify the pre/post conditions of the **algo1** algorithm.
- For the two **for** loops in the algorithm:
 - Determine if the loop is **up** loop or **down** loop.
 - Determine the invariants of these two loops and verify whether they hold in three stages: **initialization**, **maintenance** and **termination**.
- Identify some edge cases of the algorithm and verify if the algorithm has the correct output.
- Conduct an exact analysis of the running time of algorithm **algo1**.
- Determine the best and the worst case of the algorithm. What is the running time and asymptotic complexity in each case?

Tasks in past exams

[2021 Final Exam] Assume $f_1(n) = O(1)$, $f_2(n) = O(N^2)$, and $f_3(n) = O(N \log N)$. From these complexities it follows that $f_1(n) + f_2(n) + f_3(n) = O(N \log N)$.

Answer:

☐ True

☐ False