

Informatics II, Spring 2023, Exercise 4

Publication of exercise: March 13, 2023

Publication of solution: March 19, 2023

Exercise classes: March 20 - March 24, 2023

Note, in this exercise $\log n$ refers to log base 2, $(\log_2 n)$, if not stated otherwise.

Task 1: Divide and Conquer

The maximum-subarray algorithm finds the contiguous subarray that has the largest sum within an **unsorted** array A with n integers. For example, for array $A = [-2, -5, 6, -2, -3, 1, 5]$, the maximum subarray is $[6, -2, -3, 1, 5]$.

The algorithm works as follows:

Firstly, it divides the input array into two equal partitions: **I** ($A[0] \dots A[mid]$) and **II** ($A[mid+1] \dots A[n-1]$). Afterwards, it calls itself recursively on both partitions to find the maximum subarray of each partition. The combination step decides the maximum-subarray by comparing three arrays: the maximum-subarray from the left part, the maximum-subarray from the right part, and the maximum-subarray that overlaps the middle. The maximum-subarray that overlaps the middle is determined by considering all elements to the left and all elements to the right of the middle.

1.1 Based on the above algorithm description, draw a tree that illustrates the process of determining the maximum subarray in array $A = [-1, 2, -4, 1, 9, -6, 7, -3, 5]$.

1.2 Provide a C code for maximum-subarray algorithm.

1.3 What is the recurrence relation of the algorithm and its asymptotic tight bound.

Recurrences

Task 2: Substitution Method

Consider the recurrence T_a for the following questions:

$$T_a(n) = \begin{cases} 1, & \text{if } n = 1 \\ 2T(n-1) + c_1, & \text{otherwise} \end{cases} \quad (1)$$

- 2.1** In order to find an upper bound for the recurrence $T_a(n)$ with the substitution method, you need to make an educated guess what a good upper bound could be. To make it easier for you to make a good guess, write a short C program that computes the values for $T(n)$ and that prints out the results after each recursive step. What function of n could result in the values that you observe? Vary n and c_1 to get a better intuition.
- 2.2** Now use the substitution method to find an upper bound for $T_a(n)$. Use the guessed upper bound from above.

Task 3: Repeated Backward Substitution

Consider the recurrence T_b for the following questions, defined as:

$$T_b(n) = \begin{cases} 1, & \text{if } n = 1 \\ 2T(n/2) + n, & \text{otherwise} \end{cases}$$

Use the repeated backward substitution method to find an upper bound for $T_b(n)$.

Task 4: Recursion Tree Method

Consider the following Algorithm for the subsequent questions.

<div style="border-bottom: 1px solid black; margin-bottom: 5px;"> Algo: RecursiveAlgo(n) </div> <div style="margin-bottom: 5px;"> 1 if $n \leq 1$ then </div> <div style="margin-bottom: 5px;"> 2 return </div> <div style="margin-bottom: 5px;"> 3 int t = 0; </div> <div style="margin-bottom: 5px;"> 4 for $i=0$ to $\max(n,1)$ do </div> <div style="margin-bottom: 5px;"> 5 t = t + i ; </div> <div style="margin-bottom: 5px;"> 6 RecursiveAlgo(n/10) ; </div> <div style="margin-bottom: 5px;"> 7 RecursiveAlgo(9n/10) ; </div>

- 4.1** Analyse the algorithm above and find the recurrence relation $T(n)$ that determines its runtime. What is the base case?
- 4.2** Use the recursion tree method to find an upper bound for the runtime of the algorithm.

Task 5: Master Method

Use the Master Method to calculate the asymptotic tight bound for the following recurrences. Write down which case applies, as well as a , b and $f(n)$.

5.1 $T(n) = 2T(\frac{n}{4}) + \sqrt{n} + 5$

5.2 $T(n) = 12T(\frac{n}{8}) + n^3$