

Chapter 2

BCJR Algorithm

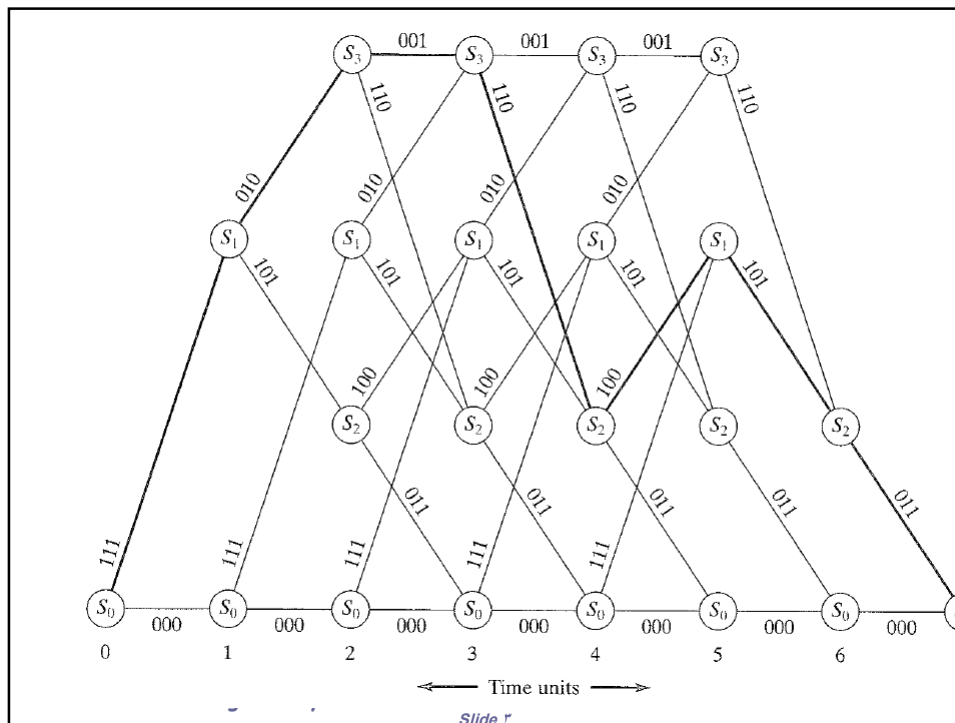
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Veterbi Algorithm (revisted)

- Consider convolutional encoder with
$$\mathbb{G}(D) = [1 + D \quad 1 + D^2 \quad 1 + D + D^2]$$
- And information sequences of length $h = 5$
- The trellis diagram has $h + m + 1$ timeslots which equals 8 in our case
- Consider received sequence as

$$\mathbb{r} = (110, 110, 110, 111, 010, 101, 101).$$



- From Fig. 12.6 on the text book, we can see that

$$\hat{\mathbf{v}} = (111, 010, 110, 011, 111, 101, 011),$$

SOVA

The Soft-Output Viterbi Algorithm (SOVA) was first introduced in 1989.

We describe SOVA for convolutional code with $R = 1/n$ on binary input, AWGN channel.

We assume that priori probabilities are not equally likely $p(u_L)$ and $L = 0, \dots, h-1$.

Log-likelihood metric

- Let us define the log-likelihood ratio or the L-value of a received symbol r at the output of channel with binary inputs $v = \pm 1$

$$L(r) = \ln \left[\frac{p(r|v = +1)}{p(r|v = -1)} \right].$$

- Similarly the L-value of an information bit u is defined as

$$L(u) = \ln \left[\frac{p(u = +1)}{p(u = -1)} \right].$$

- Using Bay's rule if v is equally likely

$$L(r) = \ln \left[\frac{p(r|v = +1)}{p(r|v = -1)} \right] = \ln \left[\frac{p(v = +1|r)}{p(v = -1|r)} \right].$$

Log-likelihood metric

- A large positive value of $L(r)$ indicates a high reliability that $v = +1$.
- A large negative value of $L(r)$ indicates a high reliability that $v = -1$.
- A close to zero value of $L(r)$ indicates a decision about the value of v based only on r is unreliable.
- The same a large positive value of $L(u)$ indicates a high reliability that $u = +1$

Log-likelihood metric

- It can be shown that the L value is equal to (left as exercise for the students)

$$L(r) = (4E_s/N_0)r,$$

- Where $L_c \equiv 4E_s/N_0$ is defined as channel reliability factor

BCJR algorithm

- Viterbi Algorithm minimizes the WER $P_w(E)$ that is $P(\hat{v} \neq v|r)$
- So it minimizes the error probability between the transmitted and received codeword.
- In BCJR algorithm, we are interested in minimizing the bit error probability.
- This is done by maximizing the posteriori probability $P(\hat{u}_l = u_l|\mathbb{r})$
- That is why BCJR decoder is also called Maximum Posteriori Probability decoder (MAP)

BCJR algorithm

- We don't assume that the information bits are equally likely.
- The algorithm calculates the a posteriori L-values
$$L(u_l) \equiv \ln \left[\frac{P(u_l = +1|\mathbb{r})}{P(u_l = -1|\mathbb{r})} \right],$$
- Called APP L-values of each information bit, the decoder output is given by

$$\hat{u}_l = \begin{cases} +1 & \text{if } L(u_l) > 0 \\ -1 & \text{if } L(u_l) < 0 \end{cases}, \quad l = 0, 1, \dots, h-1.$$

- We start our development of the BCJR algorithm by rewriting the APP value as

$$P(u_l = +1 | \mathbf{r}) = \frac{p(u_l = +1, \mathbf{r})}{P(\mathbf{r})} = \frac{\sum_{\mathbf{u} \in \mathcal{U}_l^+} p(\mathbf{r} | \mathbf{v}) P(\mathbf{u})}{\sum_{\mathbf{u}} p(\mathbf{r} | \mathbf{v}) P(\mathbf{u})},$$

- Where \mathcal{U}_l^+ is the set of all information sequences \mathbf{u} such as $u_l = 1$, \mathbf{v} is the transmitted codeword corresponding to the information sequence \mathbf{u} .

- So we can rewrite the expression of the APP L values as

$$L(u_l) = \ln \left[\frac{\sum_{\mathbf{u} \in \mathcal{U}_l^+} p(\mathbf{r} | \mathbf{v}) P(\mathbf{u})}{\sum_{\mathbf{u} \in \mathcal{U}_l^-} p(\mathbf{r} | \mathbf{v}) P(\mathbf{u})} \right],$$

- Where \mathcal{U}_l^- is the set of all information sequences \mathbf{u} such as $u_l = -1$

- The L values can be calculated using the previous formula but still it suffers from high degree of complexity.

- We can rewrite the a posteriori probability as

$$P(u_l = +1 | \mathbf{r}) = \frac{p(u_l = +1, \mathbf{r})}{P(\mathbf{r})} = \frac{\sum_{(s', s) \in \Sigma_l^+} p(s_l = s', s_{l+1} = s, \mathbf{r})}{P(\mathbf{r})},$$

- Where Σ_l^+ is the set of all state pairs $s_l = s'$ and $s_{l+1} = s$ that corresponds to the input bit $u_l = +1$.

- Reforming $P(u_l = -1 | \mathbf{r})$ in the same way and sub in the L value

$$L(u_l) = \ln \left\{ \frac{\sum_{(s', s) \in \Sigma_l^+} p(s_l = s', s_{l+1} = s, \mathbf{r})}{\sum_{(s', s) \in \Sigma_l^-} p(s_l = s', s_{l+1} = s, \mathbf{r})} \right\},$$

- Where Σ_l^- is the set of all state pairs $s_l = s'$ and $s_{l+1} = s$ that corresponds to the input bit $u_l = -1$.
- The joint pdf $p(s', s, \mathbf{r})$ can be found recursively, starting from

$$p(s', s, \mathbf{r}) = p(s', s, \mathbb{r}_{t < l}, \mathbb{r}_l, \mathbb{r}_{t > l}),$$

- Where $\mathbf{r}_{t < l}$ represents the portion of the received \mathbf{r} before the time l and Where $\mathbf{r}_{t > l}$ represents the portion of the received \mathbf{r} after the time l .
- Now application of Bay's rule

$$\begin{aligned} p(s', s, \mathbf{r}) &= p(\mathbb{r}_{t > l} | s', s, \mathbb{r}_{t < l}, \mathbb{r}_l) p(s', s, \mathbb{r}_{t < l}, \mathbb{r}_l) \\ &= p(\mathbb{r}_{t > l} | s', s, \mathbb{r}_{t < l}, \mathbb{r}_l) p(s, \mathbb{r}_l | s', \mathbb{r}_{t < l}) p(s', \mathbb{r}_{t < l}) \\ &= p(\mathbb{r}_{t > l} | s) p(s, \mathbb{r}_l | s') p(s', \mathbb{r}_{t < l}), \end{aligned}$$

- Defining

$$\begin{aligned} \alpha_l(s') &\equiv p(s', \mathbf{r}_{t < l}) \\ \gamma_l(s', s) &\equiv p(s, \mathbb{r}_l | s') \\ \beta_{l+1}(s) &\equiv p(\mathbf{r}_{t > l} | s), \end{aligned}$$

- So the joint pdf can be rewritten as

$$p(s', s, \mathbf{r}) = \beta_{l+1}(s) \gamma_l(s', s) \alpha_l(s').$$

- We can write expression for $\alpha_{l+1}(s)$ as

$$\alpha_{l+1}(s) = \sum_{s' \in \sigma_l} \gamma_l(s', s) \alpha_l(s'),$$

- Where σ_l is the set of all states at time l .

- $\beta_l(s')$ can be written as

$$\beta_l(s') = \sum_{s \in \sigma_{l+1}} \gamma_l(s', s) \beta_{l+1}(s),$$

- Where σ_{l+1} is the set of all states at time $l+1$.

- The forward recursion starts from

$$\alpha_0(s) = \begin{cases} 1, & s = \emptyset \\ 0, & s \neq \emptyset \end{cases},$$

- And the backward recursion starts from

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$$\beta_K(s) = \begin{cases} 1, & s = \emptyset \\ 0, & s \neq \emptyset \end{cases},$$

- We can write the branch metric as

$$\begin{aligned} \gamma_l(s', s) &= p(s, \mathbb{r}_l | s') = \frac{p(s', s, \mathbb{r}_l)}{P(s')} \\ &= \left[\frac{P(s', s)}{P(s')} \right] \left[\frac{p(s', s, \mathbb{r}_l)}{P(s', s)} \right] \\ &= P(s | s') p(\mathbb{r}_l | s', s) = P(u_l) p(\mathbb{r}_l | \mathbb{v}_l), \end{aligned}$$

- Which yields

$$\gamma_l(s', s) = P(u_l) p(\mathbb{r}_l | \mathbb{v}_l) = P(u_l) \left(\sqrt{\frac{E_s}{\pi N_0}} \right)^n e^{-\frac{E_s}{N_0} \|\mathbb{r}_l - \mathbb{v}_l\|^2},$$

- We can drop the constant to achieve

$$\gamma_l(s', s) = P(u_l) e^{-E_s / N_0 \|\mathbb{r}_l - \mathbb{v}_l\|^2}.$$

- The priori probability can be written as

$$\begin{aligned}
 P(u_l = \pm 1) &= \frac{[P(u_l = +1)/P(u_l = -1)]^{\pm 1}}{\{1 + [P(u_l = +1)/P(u_l = -1)]^{\pm 1}\}} \\
 &= \frac{e^{\pm L_a(u_l)}}{\{1 + e^{\pm L_a(u_l)}\}} \\
 &= \frac{e^{-L_a(u_l)/2}}{\{1 + e^{-L_a(u_l)}\}} e^{u_l L_a(u_l)/2} \\
 &= A_l e^{u_l L_a(u_l)/2},
 \end{aligned}$$

- The L value depend of the value of u, thus

$$\begin{aligned}
 \gamma_l(s', s) &= A_l e^{u_l L_a(u_l)/2} e^{-(E_s/N_0) \|\mathbf{r}_l - \mathbf{v}_l\|^2}, \\
 &= A_l e^{u_l L_a(u_l)/2} e^{(2E_s/N_0)(\mathbf{r}_l \cdot \mathbf{v}_l) - \|\mathbf{r}_l\|^2 - \|\mathbf{v}_l\|^2} \\
 &= A_l e^{-(\|\mathbf{r}_l\|^2 + n)} e^{u_l L_a(u_l)/2} e^{(L_c/2)(\mathbf{r}_l \cdot \mathbf{v}_l)} \\
 &= A_l B_l e^{u_l L_a(u_l)/2} e^{(L_c/2)(\mathbf{r}_l \cdot \mathbf{v}_l)}, \quad l = 0, 1, \dots, h-1, \\
 \gamma_l(s', s) &= P(u_l) e^{-(E_s/N_0) \|\mathbf{r}_l - \mathbf{v}_l\|^2} \\
 &= e^{-(E_s/N_0) \|\mathbf{r}_l - \mathbf{v}_l\|^2}, \\
 &= B_l e^{(L_c/2)(\mathbf{r}_l \cdot \mathbf{v}_l)}, \quad l = h, h+1, \dots, K-1,
 \end{aligned}$$

- Again if we drop the constants

$$\begin{aligned}
 \gamma_l(s', s) &= e^{u_l L_a(u_l)/2} e^{(L_c/2)(\mathbf{r}_l \cdot \mathbf{v}_l)}, \quad l = 0, 1, \dots, h-1, \\
 \gamma_l(s', s) &= e^{(L_c/2)(\mathbf{r}_l \cdot \mathbf{v}_l)}, \quad l = h, h+1, \dots, K-1,
 \end{aligned}$$

- Using the log-domain enable using

$$\max^*(x, y) \equiv \ln(e^x + e^y) = \max(x, y) + \ln(1 + e^{-|x-y|})$$

- And the log-domain metrics are

$$\gamma_l^*(s', s) \equiv \ln \gamma_l(s', s) = \begin{cases} \frac{u_l L_u(u_l)}{2} + \frac{L_c}{2} \mathbb{I}_l \cdot \mathbb{V}_l, & l = 0, 1, \dots, h-1, \\ \frac{L_c}{2} \mathbb{I}_l \cdot \mathbb{V}_l, & l = h, h+1, \dots, K-1, \end{cases}$$

$$\alpha_{l+1}^*(s) = \max_{s' \in \mathcal{S}_l} [\gamma_l^*(s', s) + \alpha_l^*(s')], \quad l = 0, 1, \dots, K-1,$$

$$\alpha_0^*(s) \equiv \ln \alpha_0(s) = \begin{cases} 0, & s = \emptyset \\ -\infty, & s \neq \emptyset, \end{cases}$$

$$\beta_l^*(s') = \max_{s \in \mathcal{S}_{l+1}} [\gamma_l^*(s', s) + \beta_{l+1}^*(s)], \quad l = K-1, K-2, \dots, 0,$$

$$\beta_K^*(s) \equiv \ln \beta_K(s) = \begin{cases} 0, & s = \emptyset \\ -\infty, & s \neq \emptyset. \end{cases}$$

- Writing the expression for the pdf $p(s', s, r)$ and the APP L-value $L(u_l)$ as:

$$p(s', s, \mathbb{F}) = e^{\beta_{l+1}^*(s) + \gamma_l^*(s', s) + \alpha_l^*(s')}$$

$$L(u_l) = \ln \left\{ \sum_{(s', s) \in \Sigma_l^+} e^{\beta_{l+1}^*(s) + \gamma_l^*(s', s) + \alpha_l^*(s')} \right\} \\ - \ln \left\{ \sum_{(s', s) \in \Sigma_l^-} e^{\beta_{l+1}^*(s) + \gamma_l^*(s', s) + \alpha_l^*(s')} \right\}$$

- Using the following math expression

$$\max^*(x, y, z) \equiv \ln(e^x + e^y + e^z) = \max^*[\max^*(x, y), z],$$

- We can formula the L value as

$$L(u_l) = \max_{(s', s) \in \Sigma_l^+}^* [\beta_{l+1}^*(s) + \gamma_l^*(s', s) + \alpha_l^*(s')] \\ - \max_{(s', s) \in \Sigma_l^-}^* [\beta_{l+1}^*(s) + \gamma_l^*(s', s) + \alpha_l^*(s')].$$

Steps of Log-Domain BCJR algorithm

- Step1: calculate the forward and backward metrics using

$$\alpha_0^*(s) \equiv \ln \alpha_0(s) = \begin{cases} 0, & s = \emptyset \\ -\infty, & s \neq \emptyset, \end{cases}$$

$$\beta_K^*(s) \equiv \ln \beta_K(s) = \begin{cases} 0, & s = \emptyset \\ -\infty, & s \neq \emptyset. \end{cases}$$

- Step 2 Compute the branch metric using

$$\gamma_l^*(s', s) \equiv \ln \gamma_l(s', s) = \begin{cases} \frac{u_l L_u(u_l)}{2} + \frac{L_c}{2} \mathbf{r}_l \cdot \mathbf{v}_l, & l = 0, 1, \dots, h-1, \\ \frac{L_c}{2} \mathbf{r}_l \cdot \mathbf{v}_l, & l = h, h+1, \dots, K-1, \end{cases}$$

Steps of Log-Domain BCJR algorithm

- Step3: calculate the forward metrics using

$$\alpha_{l+1}^*(s) = \max_{s' \in \sigma_l} [\gamma_l^*(s', s) + \alpha_l^*(s')], \quad l = 0, 1, \dots, K-1,$$

- Step 4 Compute the backward metric using

$$\beta_l^*(s') = \ln \sum_{s \in \sigma_{l+1}} e^{\gamma_l^*(s', s) + \beta_{l+1}^*(s)}$$

- Step 5 compute the APP-L values using

$$L(u_l) = \max_{(s', s) \in \Sigma_l^+} [\beta_{l+1}^*(s) + \gamma_l^*(s', s) + \alpha_l^*(s')] \\ - \max_{(s', s) \in \Sigma_l^-} [\beta_{l+1}^*(s) + \gamma_l^*(s', s) + \alpha_l^*(s')].$$

Steps of Log-Domain BCJR algorithm

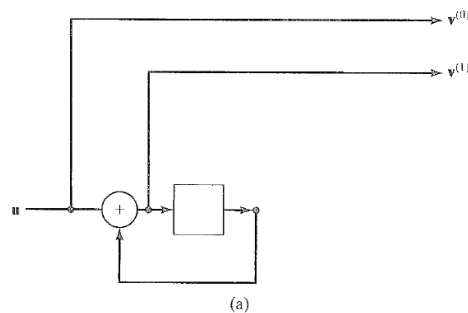
- Step6: (Optional) compute the hard decisions using

$$\hat{u}_l = \begin{cases} +1 & \text{if } L(u_l) > 0 \\ -1 & \text{if } L(u_l) < 0 \end{cases}, \quad l = 0, 1, \dots, h-1.$$

Example

- We will consider the BCJR decoding of a (2, 1, 1) systematic Recursive Convolutional code on AWGN with generator matrix

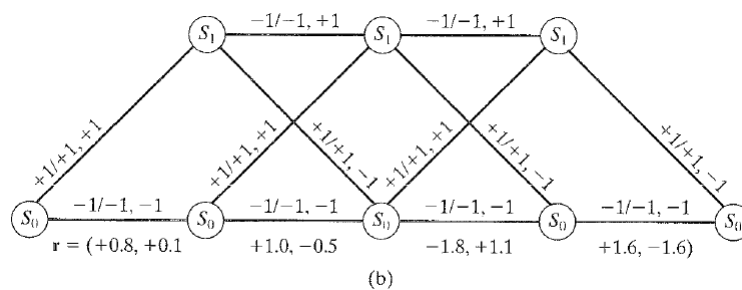
$$\mathbb{G}(D) = [1 \quad 1/(1+D)]$$



Channel Coding Theory

Slide 70

- Let $u = (u_0, u_1, u_2, u_3)$ denote the input vector of length 4 and $v = (v_0, v_1, v_2, v_3)$ denotes the codeword of length 8.
- We assume $E_s/N_0 = \frac{1}{4}(-6.02)$ dB
- The received vector $r = (+0.8, +0.1; +1.0, -0.5; -1.8, +1.1; 1.6, -1.6)$



Channel Coding Theory

Slide 71

- The rate of the terminated code is $R = h/N = 3/8$
- .
- $E_b/N_0 = E_s/RN_0 = 2/3$
- Assuming that the information bits are equally likely $L_a(u_i) = 0$
- $L_c = 4 E_b/N_0 = 1$

$$\begin{aligned}\gamma_0^*(S_0, S_0) &= \frac{-1}{2} L_a(u_0) + \frac{1}{2} r_0 \cdot v_0 \\ &= \frac{1}{2} (-0.8 - 0.1) = -0.45\end{aligned}$$

$$\begin{aligned}\gamma_0^*(S_0, S_1) &= \frac{+1}{2} L_a(u_0) + \frac{1}{2} r_0 \cdot v_0 \\ &= \frac{1}{2} (0.8 + 0.1) = 0.45\end{aligned}$$

$$\begin{aligned}\gamma_1^*(S_0, S_0) &= \frac{-1}{2} L_a(u_1) + \frac{1}{2} r_1 \cdot v_1 \\ &= \frac{1}{2} (-1.0 + 0.5) = -0.25\end{aligned}$$

$$\begin{aligned}\gamma_1^*(S_0, S_1) &= \frac{+1}{2} L_a(u_1) + \frac{1}{2} r_1 \cdot v_1 \\ &= \frac{1}{2} (1.0 - 0.5) = 0.25\end{aligned}$$

$$\begin{aligned}\gamma_1^*(S_1, S_1) &= \frac{-1}{2} L_a(u_1) + \frac{1}{2} r_1 \cdot v_1 \\ &= \frac{1}{2} (-1.0 - 0.5) = -0.75\end{aligned}$$

$$\begin{aligned}\gamma_1^*(S_1, S_0) &= \frac{+1}{2} L_a(u_1) + \frac{1}{2} r_1 \cdot v_1 \\ &= \frac{1}{2} (1.0 + 0.5) = 0.75\end{aligned}$$

$$\begin{aligned}\gamma_2^*(S_0, S_0) &= \frac{-1}{2}L_a(u_2) + \frac{1}{2}\mathbb{r}_2 \cdot \mathbb{v}_2 \\ &= \frac{1}{2}(1.8 - 1.1) = 0.35\end{aligned}$$

$$\begin{aligned}\gamma_2^*(S_0, S_1) &= \frac{+1}{2}L_a(u_2) + \frac{1}{2}\mathbb{r}_2 \cdot \mathbb{v}_2 \\ &= \frac{1}{2}(-1.8 + 1.1) = -0.35\end{aligned}$$

$$\begin{aligned}\gamma_2^*(S_1, S_1) &= \frac{-1}{2}L_a(u_2) + \frac{1}{2}\mathbb{r}_2 \cdot \mathbb{v}_2 \\ &= \frac{1}{2}(1.8 + 1.1) = 1.45\end{aligned}$$

$$\begin{aligned}\gamma_2^*(S_1, S_0) &= \frac{+1}{2}L_a(u_2) + \frac{1}{2}\mathbb{r}_2 \cdot \mathbb{v}_2 \\ &= \frac{1}{2}(-1.8 - 1.1) = -1.45\end{aligned}$$

$$\begin{aligned}\gamma_3^*(S_0, S_0) &= \frac{+1}{2}\mathbb{r}_3 \cdot \mathbb{v}_3 \\ &= \frac{1}{2}(-1.6 + 1.6) = 0\end{aligned}$$

$$\begin{aligned}\gamma_3^*(S_1, S_0) &= \frac{+1}{2}\mathbb{r}_3 \cdot \mathbb{v}_3 \\ &= \frac{1}{2}(1.6 + 1.6) = 1.60.\end{aligned}$$

Now compute the log-domain forward metrics

$$\alpha_1^*(S_0) = [\gamma_0^*(S_0, S_0) + \alpha_0^*(S_0)] = -0.45 + 0 = -0.45$$

$$\alpha_1^*(S_1) = [\gamma_0^*(S_0, S_1) + \alpha_0^*(S_0)] = 0.45 + 0 = 0.45$$

$$\begin{aligned}\alpha_2^*(S_0) &= \max^* \{ [\gamma_1^*(S_0, S_0) + \alpha_1^*(S_0)], [\gamma_1^*(S_1, S_0) + \alpha_1^*(S_1)] \} \\ &= \max^* \{ [(-0.25) + (-0.45)], [(0.75) + (0.45)] \} \\ &= \max^* (-0.70, +1.20) = 1.20 + \ln(1 + e^{-|-1.9|}) = 1.34\end{aligned}$$

$$\begin{aligned}\alpha_2^*(S_1) &= \max^* \{ [\gamma_1^*(S_0, S_1) + \alpha_1^*(S_0)], [\gamma_1^*(S_1, S_1) + \alpha_1^*(S_1)] \} \\ &= \max^* (-0.20, -0.30) = -0.20 + \ln(1 + e^{-|0.1|}) = 0.44.\end{aligned}$$

- Similarly compute the log-domain backward metrics

$$\begin{aligned}\beta_3^*(S_0) &= [\gamma_3^*(S_0, S_0) + \beta_4^*(S_0)] = 0 + 0 = 0 \\ \beta_3^*(S_1) &= [\gamma_3^*(S_1, S_0) + \beta_4^*(S_0)] = 1.60 + 0 = 1.60 \\ \beta_2^*(S_0) &= \max^* \{ [\gamma_2^*(S_0, S_0) + \beta_3^*(S_0)], [\gamma_2^*(S_0, S_1) + \beta_3^*(S_1)] \} \\ &= \max^* \{ [(0.35) + (0)], [(-0.35) + (1.60)] \} \\ &= \max^* (0.35, 1.25) = 1.25 + \ln(1 + e^{-| -0.90 |}) = 1.59 \\ \beta_2^*(S_1) &= \max^* \{ [\gamma_2^*(S_1, S_0) + \beta_3^*(S_0)], [\gamma_2^*(S_1, S_1) + \beta_3^*(S_1)] \} \\ &= \max^* (-1.45, 3.05) = 3.05 + \ln(1 + e^{-| -4.5 |}) = 3.06\end{aligned}$$

$$\begin{aligned}\beta_1^*(S_0) &= \max^* \{ [\gamma_1^*(S_0, S_0) + \beta_2^*(S_0)], [\gamma_1^*(S_0, S_1) + \beta_2^*(S_1)] \} \\ &= \max^* (1.34, 3.31) = 3.44 \\ \beta_1^*(S_1) &= \max^* \{ [\gamma_1^*(S_1, S_0) + \beta_2^*(S_0)], [\gamma_1^*(S_1, S_1) + \beta_2^*(S_1)] \} \\ &= \max^* (2.34, 2.31) = 3.02.\end{aligned}$$

- Finally we compute the app L -values

$$\begin{aligned}L(u_0) &= [\beta_1^*(S_1) + \gamma_0^*(S_0, S_1) + \alpha_0^*(S_0)] - [\beta_1^*(S_0) + \gamma_0^*(S_0, S_0) + \alpha_0^*(S_0)] \\ &= (3.47) - (2.99) = +0.48 \\ L(u_1) &= \max^* \{ [\beta_2^*(S_0) + \gamma_1^*(S_1, S_0) + \alpha_1^*(S_1)], [\beta_2^*(S_1) + \gamma_1^*(S_0, S_1) + \alpha_1^*(S_0)] \} \\ &\quad - \max^* \{ [\beta_2^*(S_0) + \gamma_1^*(S_0, S_0) + \alpha_1^*(S_0)], [\beta_2^*(S_1) + \gamma_1^*(S_1, S_1) + \alpha_1^*(S_1)] \} \\ &= \max^* [(2.79), (2.86)] - \max^* [(0.89), (2.76)] \\ &= (3.52) - (2.90) = +0.62\end{aligned}$$

$$\begin{aligned}
L(u_2) &= \max^* \{ [\beta_3^*(S_0) + \gamma_2^*(S_1, S_0) + \alpha_2^*(S_1)], [\beta_3^*(S_1) + \gamma_2^*(S_0, S_1) + \alpha_2^*(S_0)] \} \\
&\quad - \max^* \{ [\beta_3^*(S_0) + \gamma_2^*(S_0, S_0) + \alpha_2^*(S_0)], [\beta_3^*(S_1) + \gamma_2^*(S_1, S_1) + \alpha_2^*(S_1)] \} \\
&= \max^* [(-1.01), (2.59)] - \max^* [(1.69), (3.49)] \\
&= (2.62) - (3.64) = -1.02.
\end{aligned}$$

Example

In the Max-log-MAP algorithm, the branch metrics remain as computed in (12.137). The approximation $\max^*(x, y) \approx \max(x, y)$ of (12.133) affects the computation of the forward and backward metrics (see (12.138) and (12.139), respectively), as follows:

$$\begin{aligned}
\alpha_2^*(S_0) &= \max(-0.70, +1.20) = 1.20 \\
\alpha_2^*(S_1) &= \max(-0.20, -0.30) = -0.20 \\
\beta_2^*(S_0) &= \max(0.35, 1.25) = 1.25 \\
\beta_2^*(S_1) &= \max(-1.45, 3.05) = 3.05 \\
\beta_1^*(S_0) &= \max(1.34, 3.31) = 3.31 \\
\beta_1^*(S_1) &= \max(2.34, 2.31) = 2.34.
\end{aligned}$$

- Finding the L-values

$$\begin{aligned} L(u_0) &= [\beta_1^*(S_1) + \gamma_0^*(S_0, S_1) + \alpha_0^*(S_0)] - [\beta_1^*(S_0) + \gamma_0^*(S_0, S_0) + \alpha_0^*(S_0)] \\ &= (2.79) - (2.86) = -0.07 \end{aligned}$$

$$\begin{aligned} L(u_1) &= \max\{[\beta_2^*(S_0) + \gamma_1^*(S_1, S_0) + \alpha_1^*(S_1)], [\beta_2^*(S_1) + \gamma_1^*(S_0, S_1) + \alpha_1^*(S_0)]\} \\ &\quad - \max\{[\beta_2^*(S_0) + \gamma_1^*(S_0, S_0) + \alpha_1^*(S_0)], [\beta_2^*(S_1) + \gamma_1^*(S_1, S_1) + \alpha_1^*(S_1)]\} \\ &= \max[(2.79), (2.86)] - \max[(0.89), (2.76)] \\ &= (2.86) - (2.76) = +0.1000 \end{aligned}$$

$$\begin{aligned} L(u_2) &= \max\{[\beta_3^*(S_0) + \gamma_2^*(S_1, S_0) + \alpha_2^*(S_1)], [\beta_3^*(S_1) + \gamma_2^*(S_0, S_1) + \alpha_2^*(S_0)]\} \\ &\quad - \max\{[\beta_3^*(S_0) + \gamma_2^*(S_0, S_0) + \alpha_2^*(S_0)], [\beta_3^*(S_1) + \gamma_2^*(S_1, S_1) + \alpha_2^*(S_1)]\} \\ &= \max[(-1.65), (2.45)] - \max[(1.55), (2.85)] \\ &= (2.45) - (2.85) = -0.4000. \end{aligned} \tag{12.143}$$

- Then the hard decision outputs of the Max-log-MAP decoder is
- $u = (-1, +1, -1)$