



Machine Learning with Graphs (MLG)

# Structure of Graphs

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# Some Tutorials

- **NetworkX**

- <https://github.com/ericmjl/Network-Analysis-Made-Simple> [Recommended!!]
- <https://github.com/CambridgeUniversityPress/FirstCourseNetworkScience>

- **PyTorch Geometric**

- A better choice, compared to DGL
- Official site: [https://github.com/rusty1s/pytorch\\_geometric](https://github.com/rusty1s/pytorch_geometric)
- <https://pytorch-geometric.readthedocs.io/en/latest/notes/introduction.html>
- <https://www.pytorchtutorial.com/pytorch-geometric-for-gnn/>
- <https://towardsdatascience.com/hands-on-graph-neural-networks-with-pytorch-pytorch-geometric-359487e221a8>

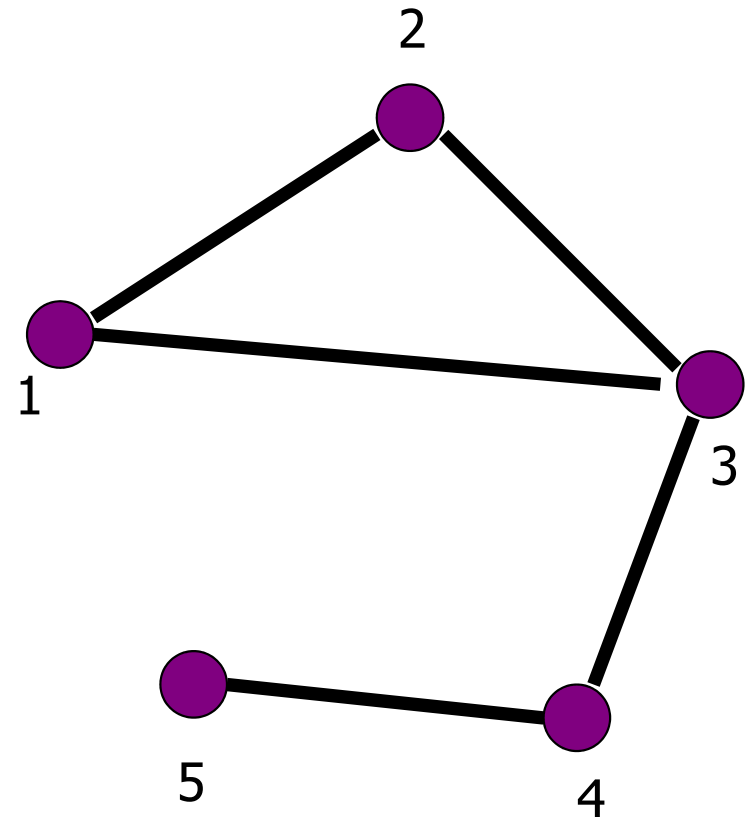
# Undirected Graph

- Graph  $G=(V,E)$ 
  - $V$  = set of vertices (nodes)
  - $E$  = set of edges

undirected graph

$V = \{1, 2, 3, 4, 5\}$

$E = \{(1,2), (1,3), (2,3), (3,4), (4,5)\}$



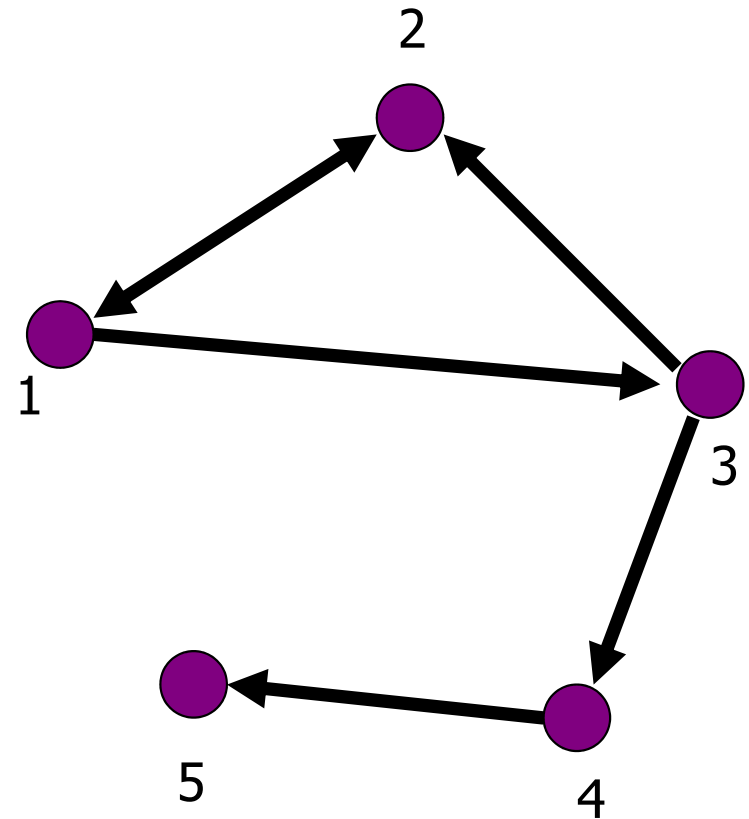
# Directed Graph

- Graph  $G=(V,E)$ 
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  - $E$  = set of edges

directed graph

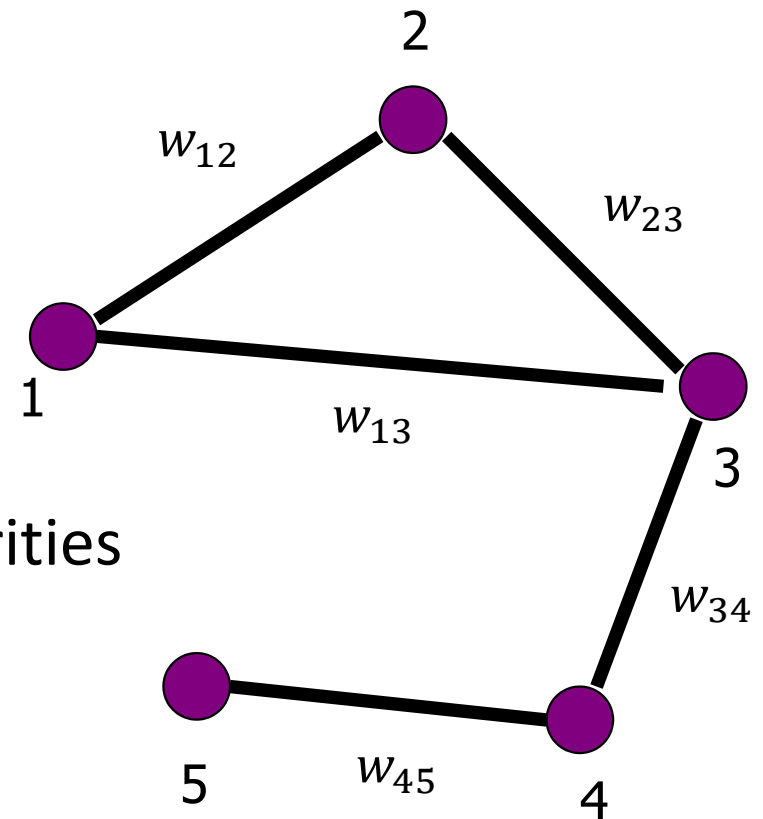
$V = \{1, 2, 3, 4, 5\}$

$E = \{\langle 1,2 \rangle, \langle 2,1 \rangle, \langle 1,3 \rangle, \langle 3,2 \rangle, \langle 3,4 \rangle, \langle 4,5 \rangle\}$



# Weighted Graph

- Graph  $G=(V,E)$ 
  - $V$  = set of vertices (nodes)
  - $E$  = set of edges and their **weights**



Weights can be either distances or similarities

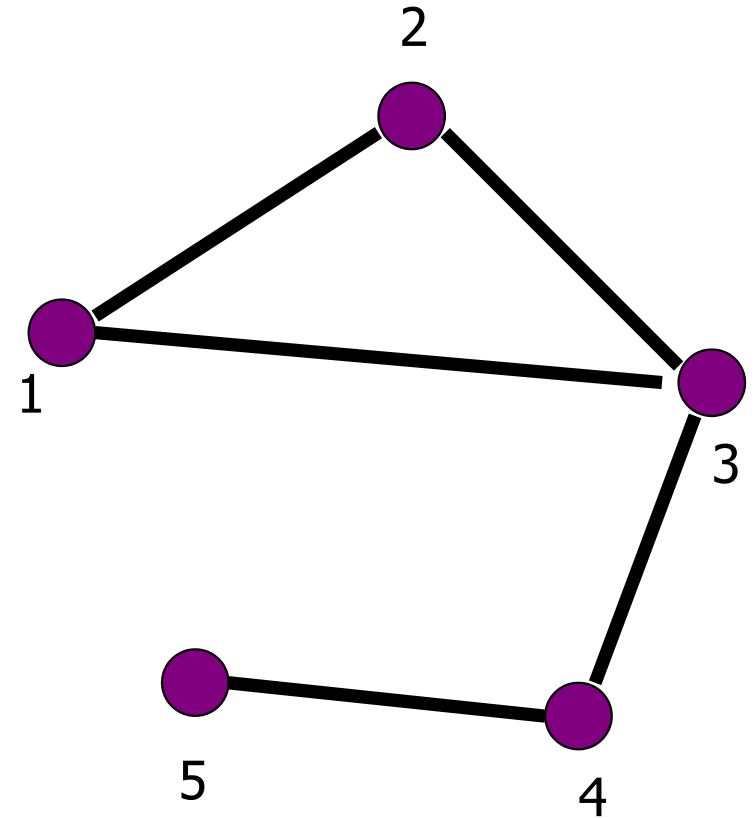
**Weighted** graph

$$V = \{1, 2, 3, 4, 5\}$$

$$E = \{(1,2,w_{12}), (1,3,w_{13}), (2,3,w_{23}), (3,4,w_{34}), (4,5,w_{45})\}$$

# Undirected graph

- **Neighborhood**  $N(i)$  of node  $i$ 
  - For any node  $i$ , in an undirected graph, the set of nodes it is connected to via an edge is called its neighborhood and is represented as  $N(i)$
- **degree**  $d(i)$  of node  $i$ 
  - Size of  $N(i)$
  - number of edges **incident** on  $i$



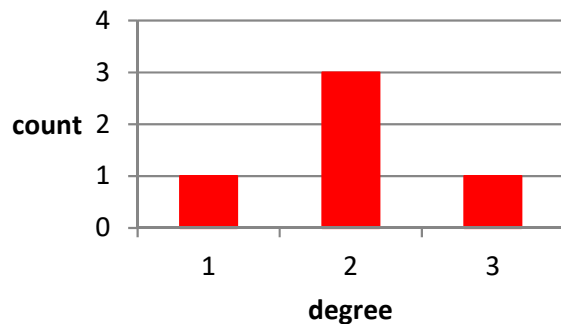
# Undirected graph

- degree sequence

- $[d(1), d(2), d(3), d(4), d(5)]$
- $[2, 2, 3, 2, 1]$

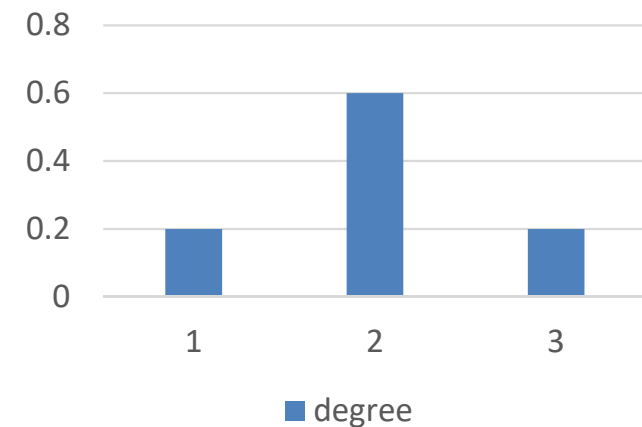
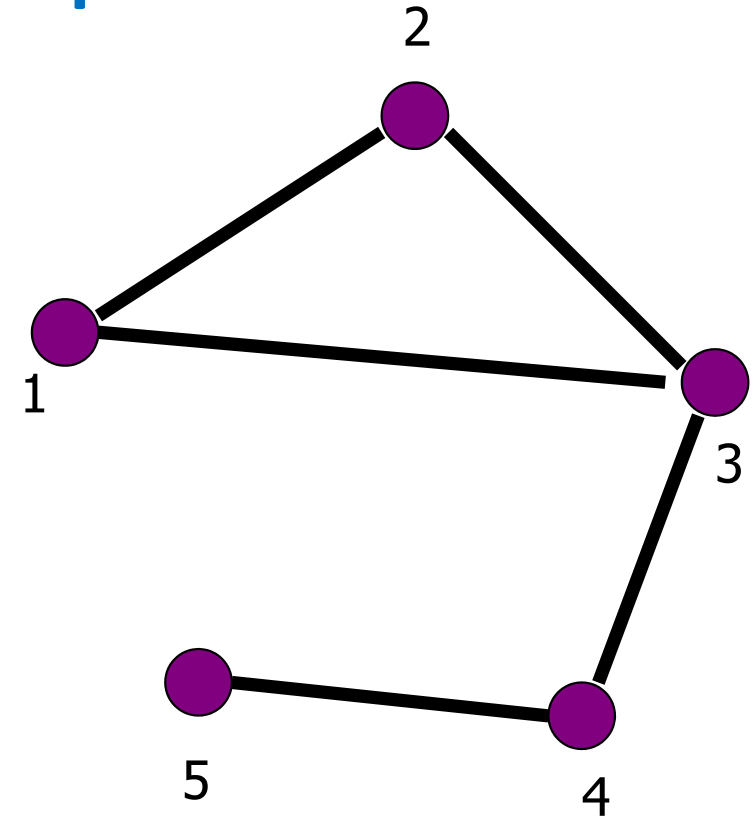
- degree histogram

- $[(1:1), (2:3), (3,1)]$



- degree distribution

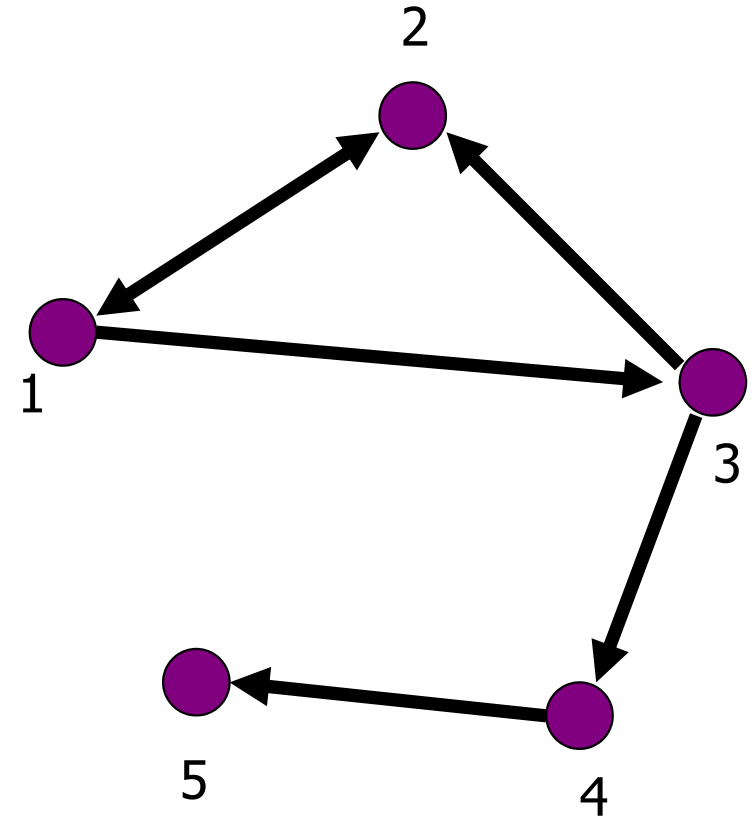
- $[(1:0.2), (2:0.6), (3,0.2)]$



# Directed Graph

In directed graphs we have incoming neighbors  $N_{in}(v)$  (nodes that connect to  $v$ ) and outgoing neighbors  $N_{out}(v)$

- **in-degree**  $d_{in}(i)$  of node  $i$ 
  - number of edges incoming to node  $i$
- **out-degree**  $d_{out}(i)$  of node  $i$ 
  - number of edges leaving node  $i$
- in-degree sequence
  - $[1, 2, 1, 1, 1]$
- out-degree sequence
  - $[2, 1, 2, 1, 0]$
- in-degree histogram
  - $[(1:4), (2:1)]$
- out-degree histogram
  - $[(0:1), (1:2), (2:2)]$





# Degree Distribution

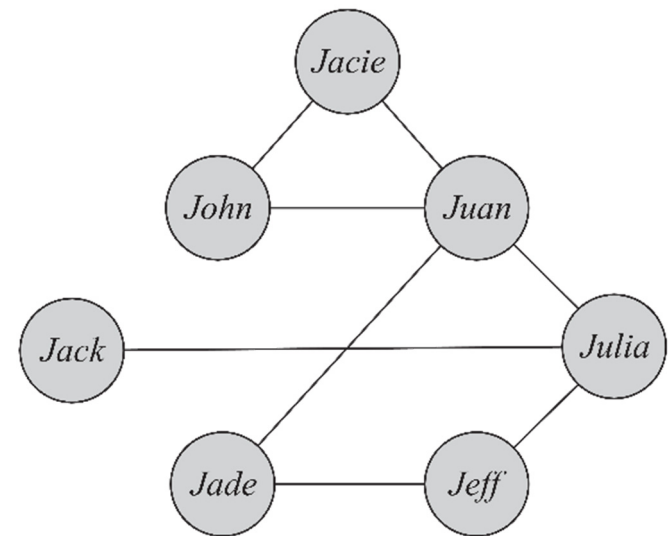
- When dealing with very large graphs, how nodes' degrees are distributed is an important concept to analyze and is called **Degree Distribution**

$$\pi(d) = \{d_1, d_2, \dots, d_n\}$$

$$p_d = \frac{n_d}{n}$$

$n_d$  is the number of nodes with degree  $d$

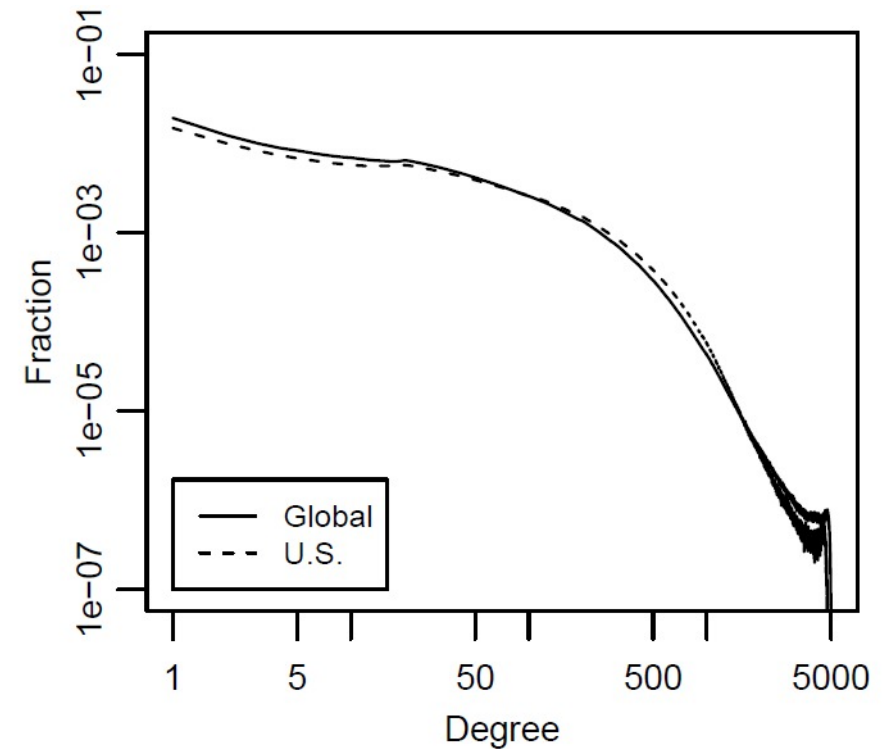
$$\sum_{d=0}^{\infty} p_d = 1$$



$$p_1 = \frac{1}{7}, p_2 = \frac{4}{7}, p_3 = \frac{1}{7}, p_4 = \frac{1}{7}$$

# Degree Distribution Plot

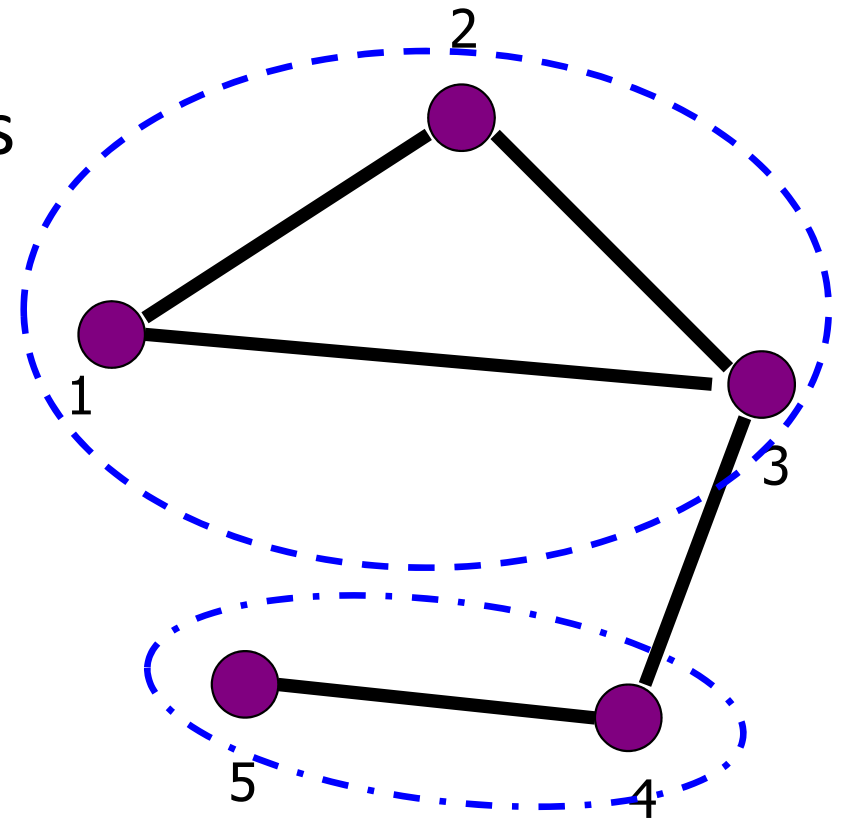
- The  $x$ -axis represents the degree and the  $y$ -axis represents the fraction of nodes having that degree
- On social networking sites
  - Many users with few connections
  - A handful of users with very large numbers of friends
- **Power-Law Degree Distribution**



Facebook Degree Distribution

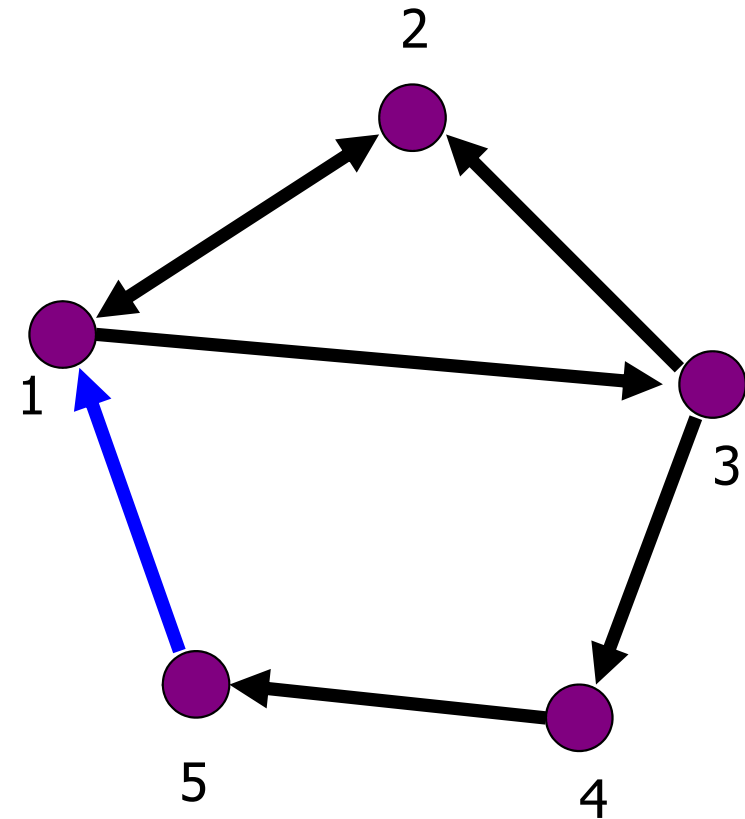
# Connectivity for Undirected graph

- **Connected** graph: a graph where there every pair of nodes is connected through a path
- **Disconnected** graph: a graph that is not connected
- **Connected Components:** subsets of vertices that are connected



# Connectivity for Directed Graph

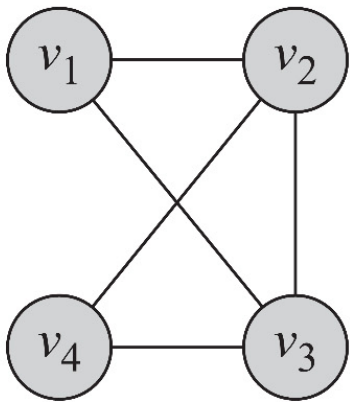
- **Strongly connected graph:** there exists a path from every  $i$  to every  $j$
- **Weakly connected graph:** If its underlying **undirected** graph is connected



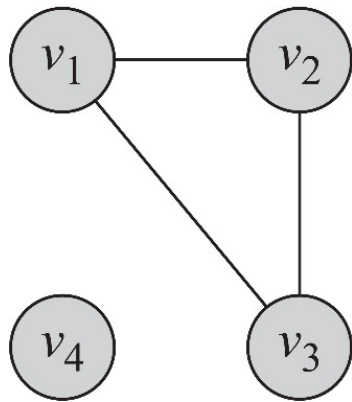
# Connectivity

- A node  $v_i$  is connected to node  $v_j$  (or reachable from  $v_j$ ) if it is adjacent to it or there exists a path from  $v_i$  to  $v_j$
- A graph is connected, if there exists a path between any pair of nodes in it
  - In a directed graph, a graph is **strongly connected** if there exists a directed path between any pair of nodes
  - In a directed graph, a graph is **weakly connected** if there exists a path between any pair of nodes, without following the edge directions
- A graph is **disconnected**, if nodes are not connected

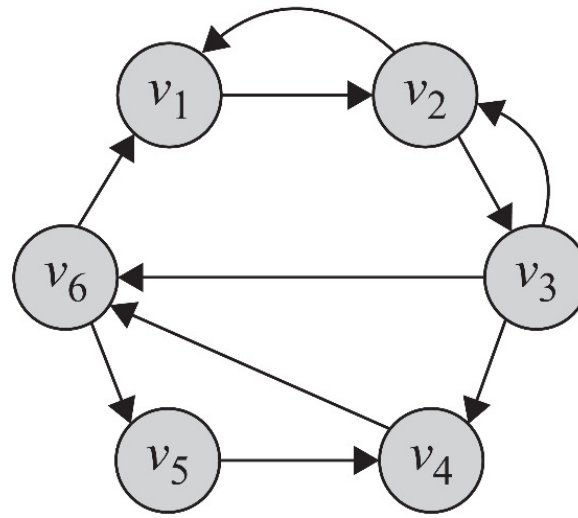
# Connectivity



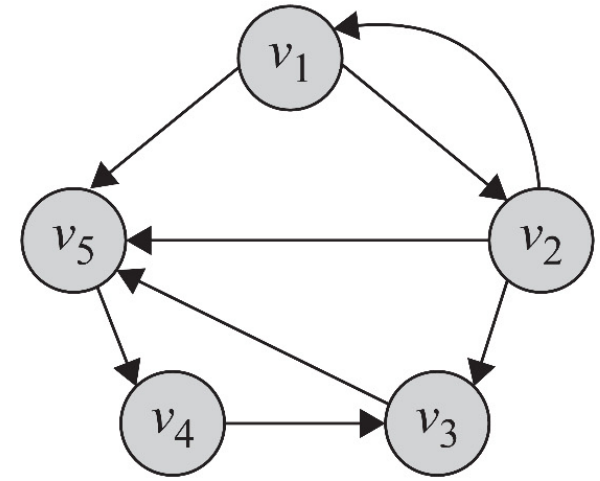
(a) Connected



(b) Disconnected



(c) Strongly connected



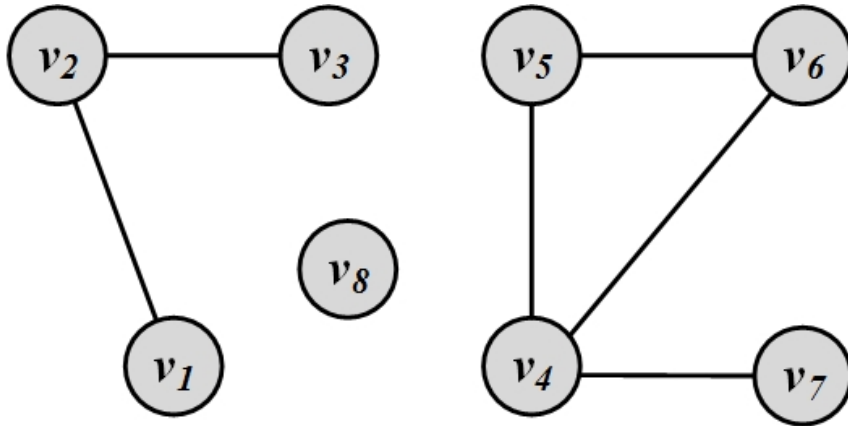
(d) Weakly connected

# Components

- A **component** in an undirected graph is a **connected subgraph**, i.e., there is a path between every pair of nodes inside the component
- In directed graphs, we have a **strongly connected component** when there is a path from  $u$  to  $v$  and one from  $v$  to  $u$  for every pair of nodes  $u$  and  $v$
- The **component is weakly connected** if replacing directed edges with undirected edges results in a connected component

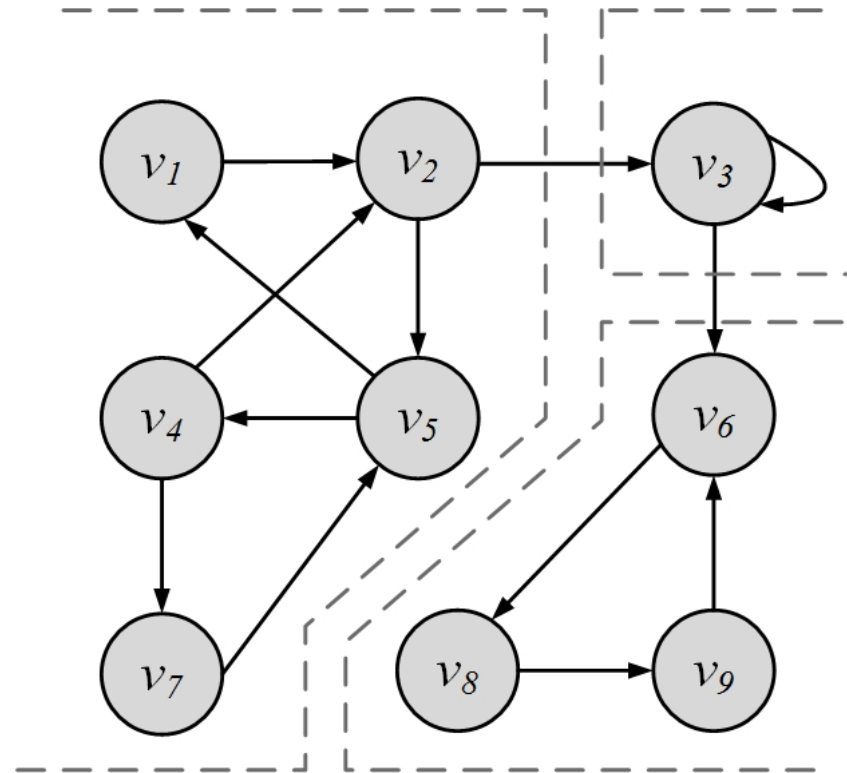
# Components

$G^1$



3 components

$G^2$

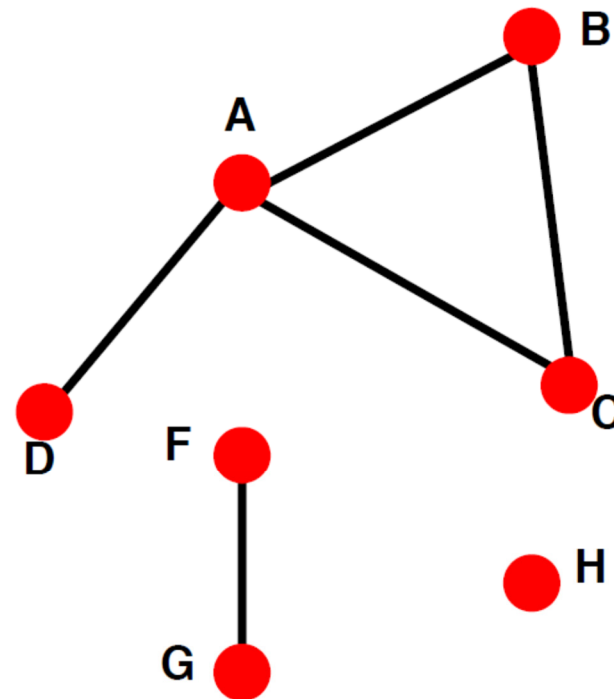
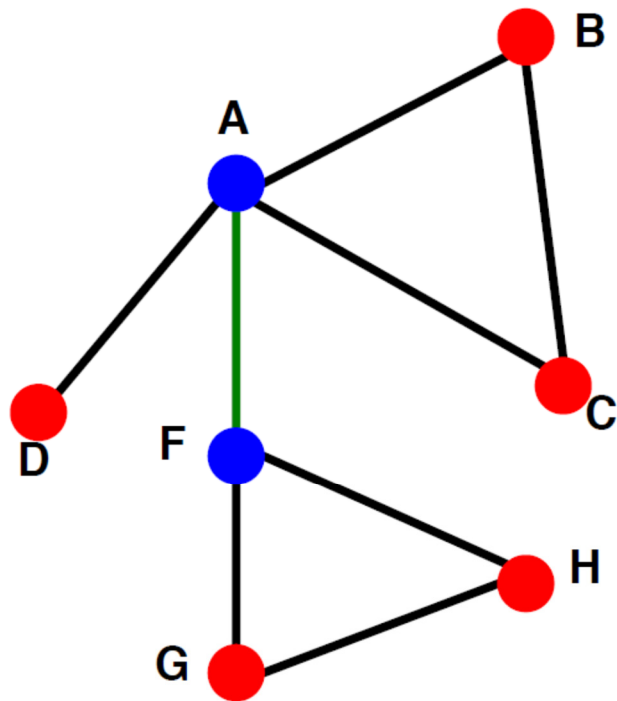


3 Strongly-connected components



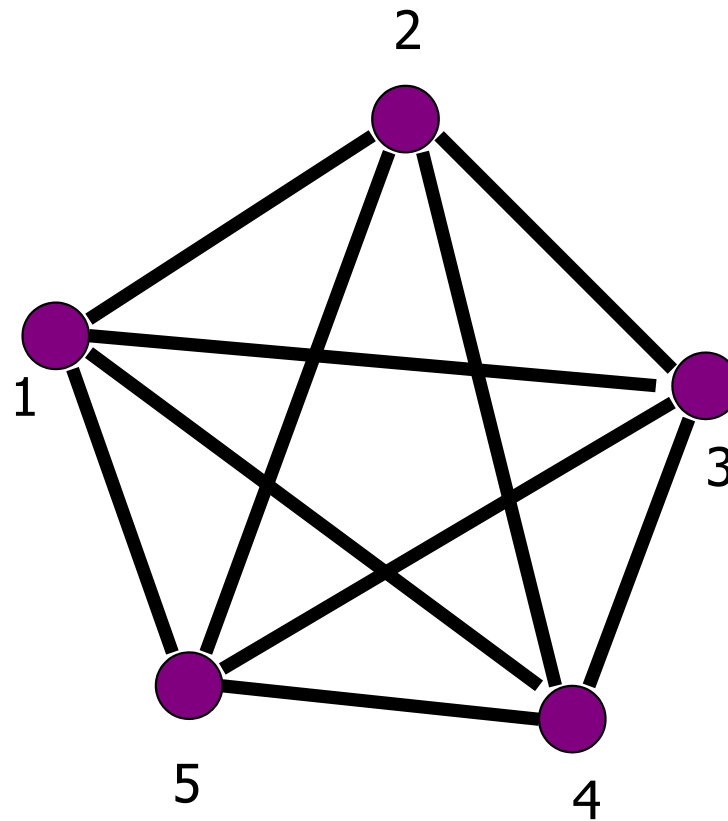
# Articulation and Bridge

- **Bridge edge**: If we erase the edge, the graph becomes disconnected
- **Articulation node**: If we erase the node, the graph becomes disconnected



# Fully Connected Graph

- Clique  $K_n$
- A graph that contains all possible  $n(n-1)/2$  edges

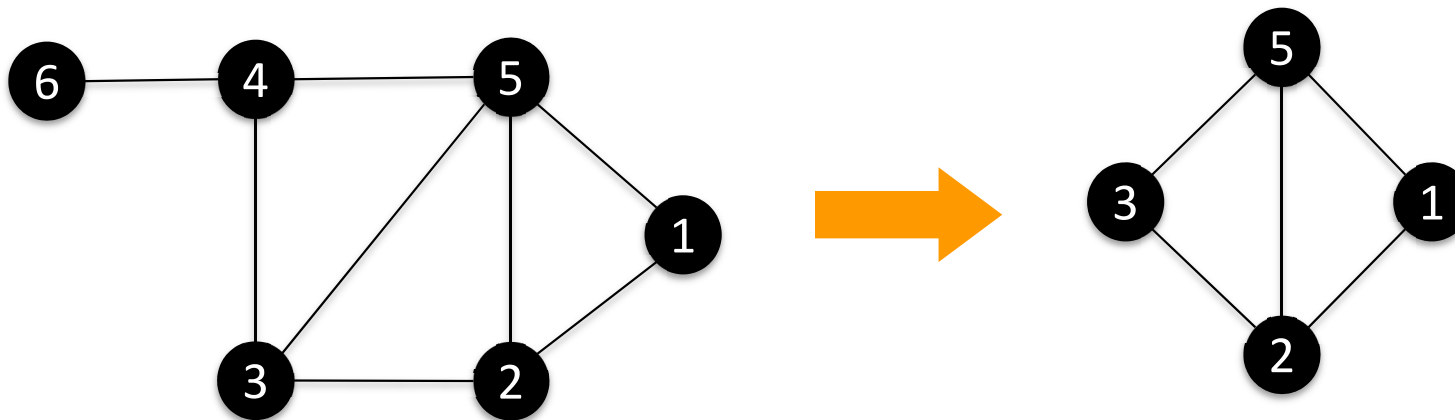


# Subgraph

- Graph  $G$  can be represented as a pair  $G(V, E)$ 
  - where  $V$  is the node set and  $E$  is the edge set
- $G'(V', E')$  is a subgraph of  $G(V, E)$

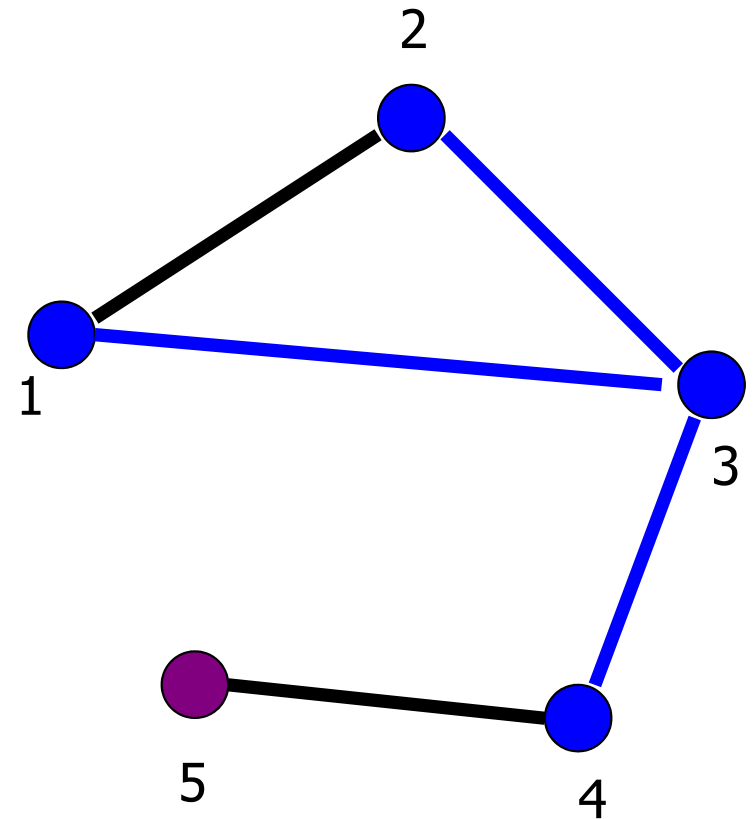
$$V' \subseteq V$$

$$E' \subseteq E$$



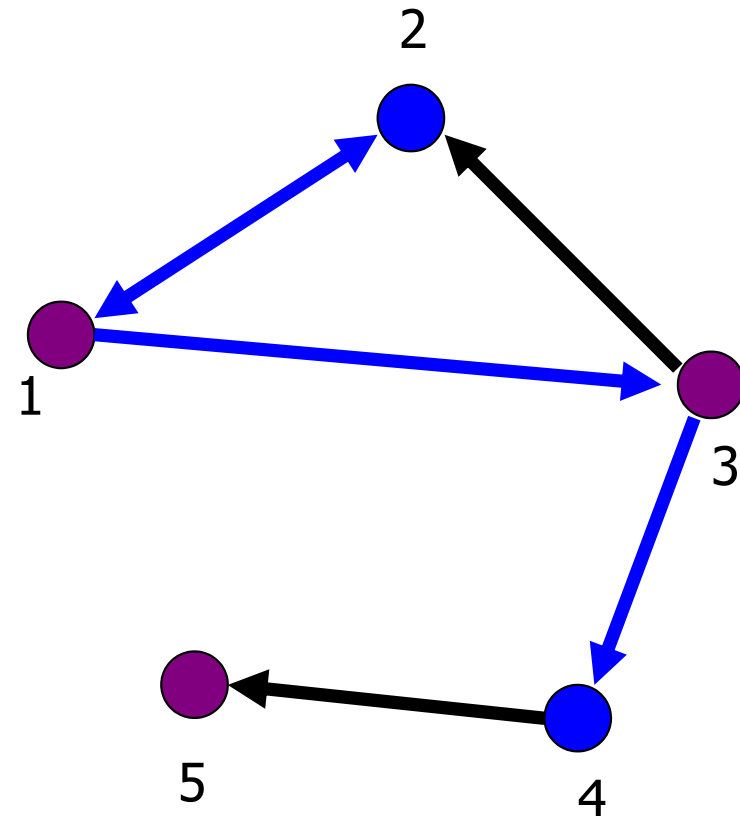
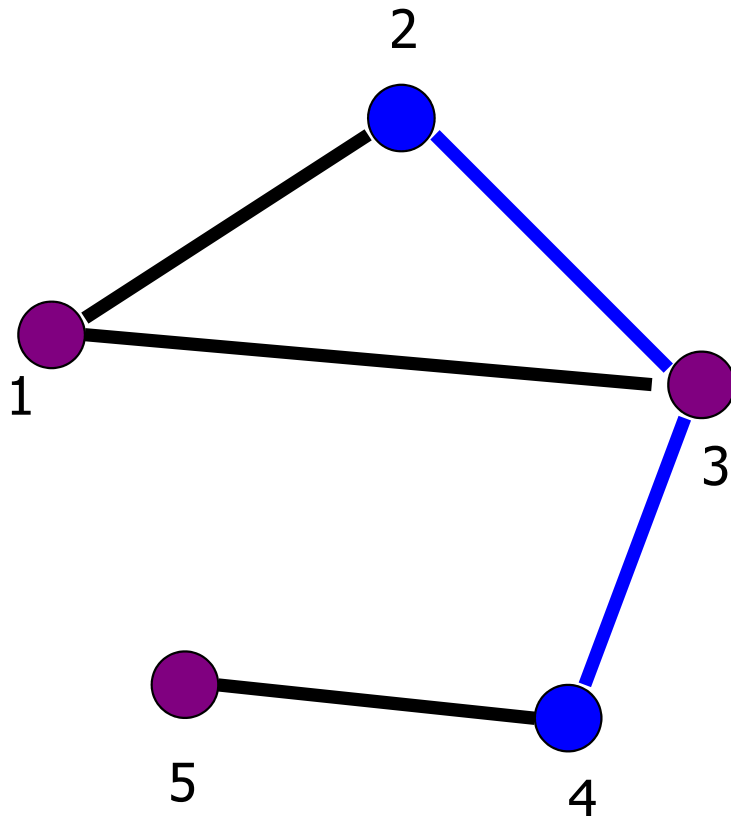
# Subgraphs

- **Subgraph:** Given  $V' \subseteq V$ , and  $E' \subseteq E$ , the graph  $G'=(V',E')$  is a subgraph of  $G$
- **Induced subgraph:** Given  $V' \subseteq V$ , let  $E' \subseteq E$  is the set of all edges between the nodes in  $V'$ . The graph  $G'=(V',E')$ , is an induced subgraph of  $G$



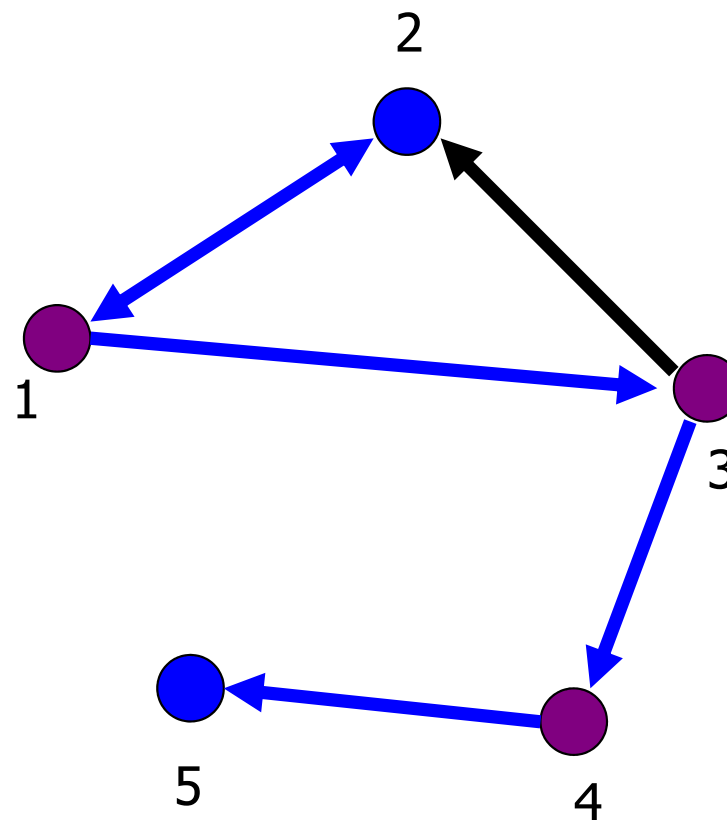
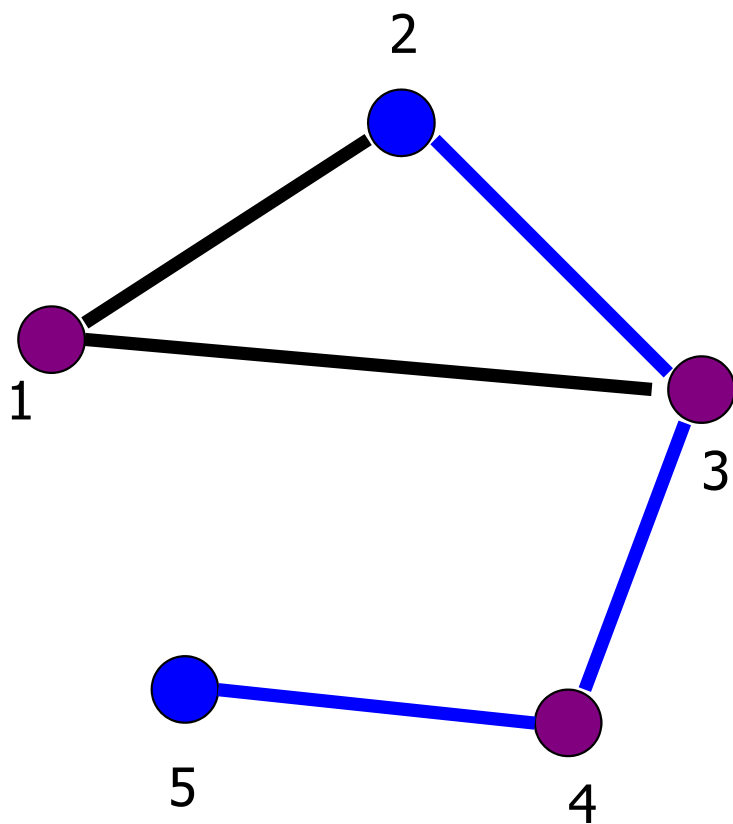
# Paths

- **Path** from node  $i$  to node  $j$ : a sequence of edges (directed or undirected from node  $i$  to node  $j$ )
  - **Path length**: number of edges on the path
  - nodes  $i$  and  $j$  are **connected** if there exists a path from  $i$  to  $j$
  - **Cycle**: a path that starts and ends at the same node



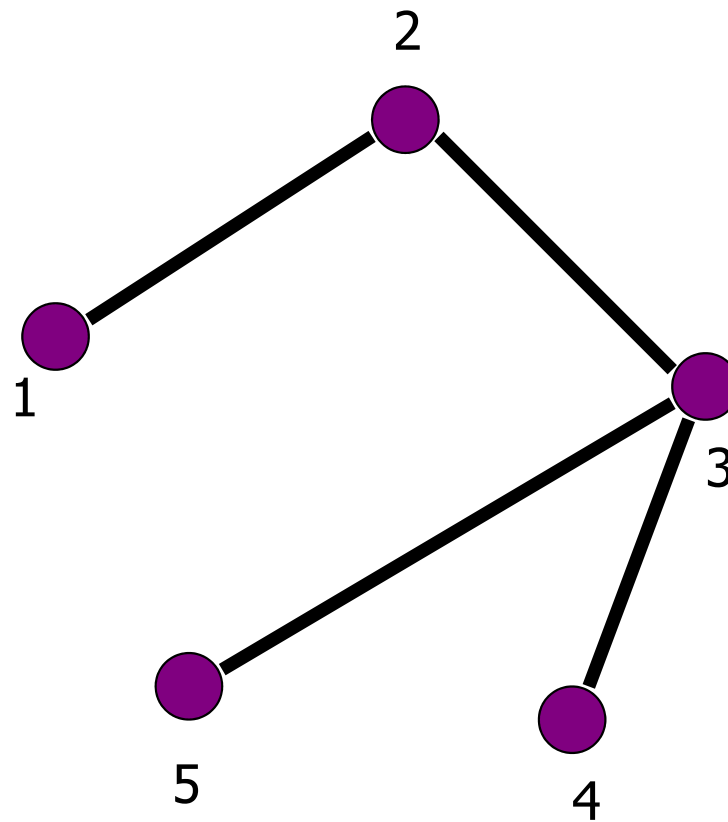
# Diameter

- The largest shortest path in the graph



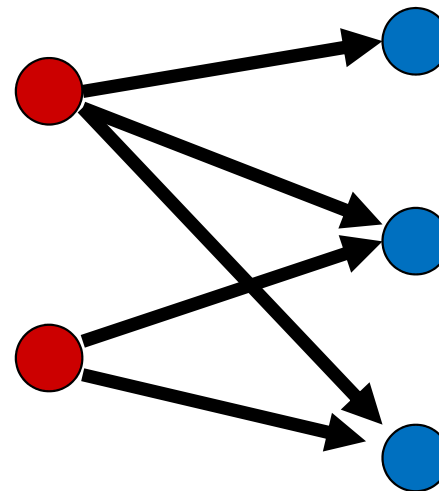
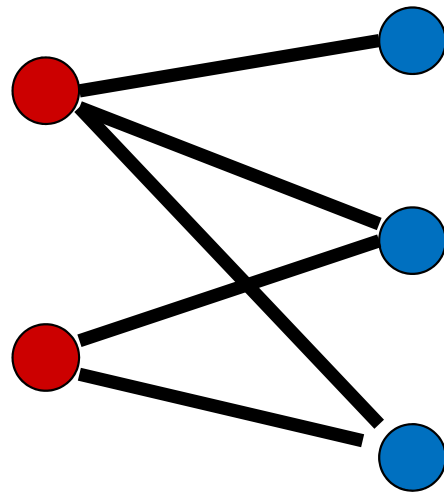
# Trees

- Connected Undirected graphs without cycles



# Bipartite Graphs

- Graphs where the set of nodes  $V$  can be partitioned into two sets  $L$  and  $R$ , such that there are edges only between nodes in  $L$  and  $R$ , and there is no edge within  $L$  or  $R$

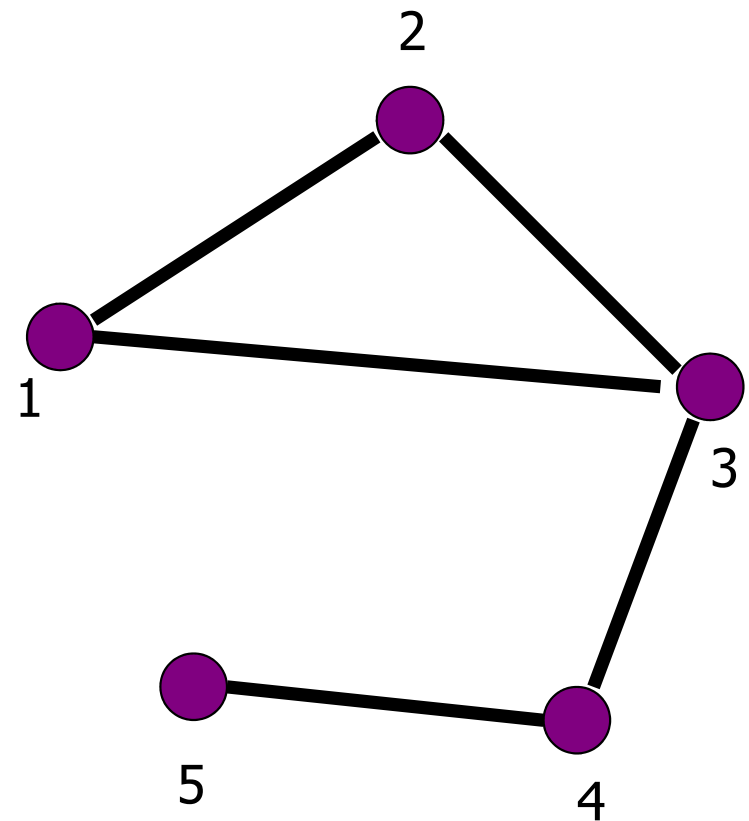




# Graph Representation

- Adjacency Matrix
  - symmetric** matrix for undirected graphs

$$A = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$



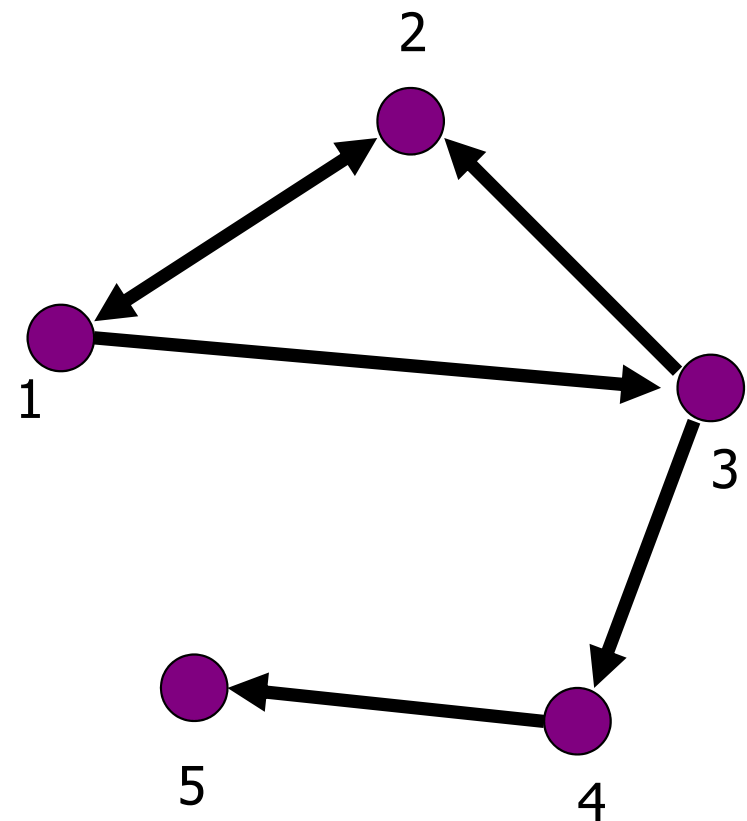
# Graph Representation

- Adjacency Matrix

However, social networks have very **sparse** Adjacency matrices

- **unsymmetric** matrix for directed graphs

$$A = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$



# Graph Representation

- Adjacency List
  - For each node in an **undirected** graph, keep a list with neighboring nodes

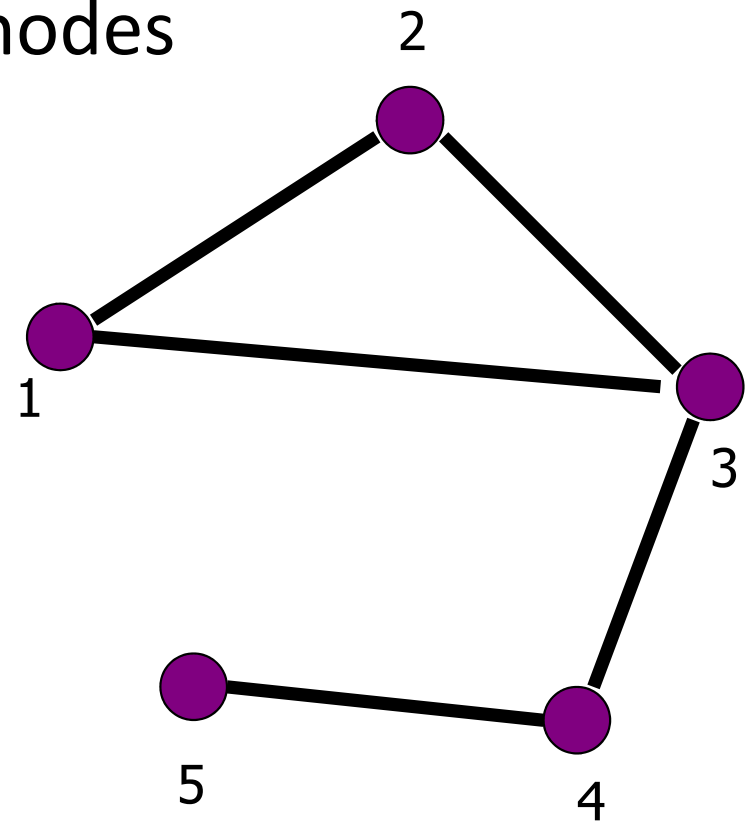
1: [2, 3]

2: [1, 3]

3: [1, 2, 4]

4: [3, 5]

5: [4]



# Graph Representation

- Adjacency List
  - For each node in a **directed** graph, keep a list of the nodes it points to

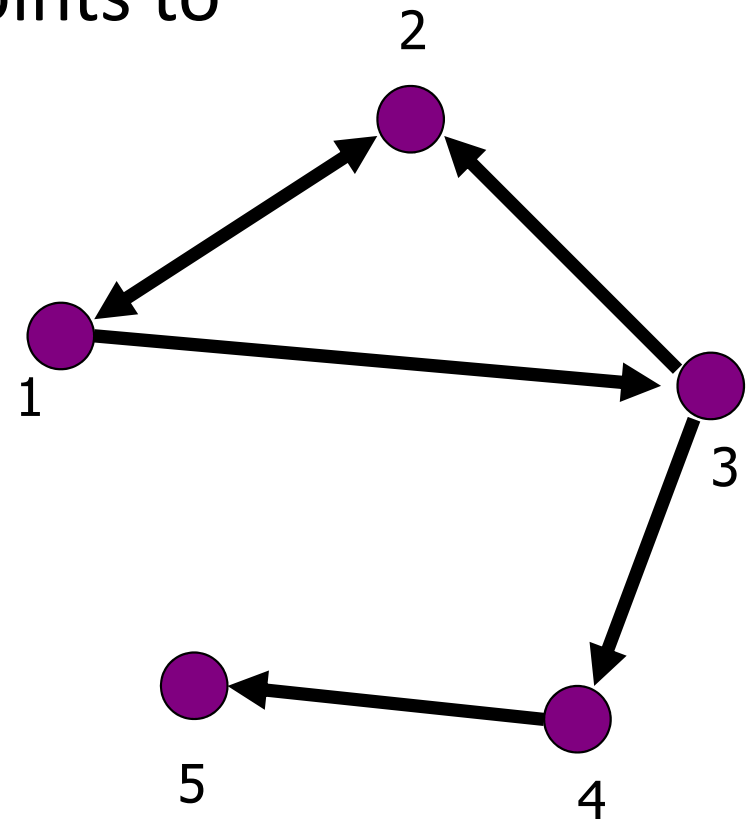
1: [2, 3]

2: [1]

3: [2, 4]

4: [5]

5: [null]



# Graph Representation

- List of edges
  - Keep a list of all the edges in the **undirected** graph

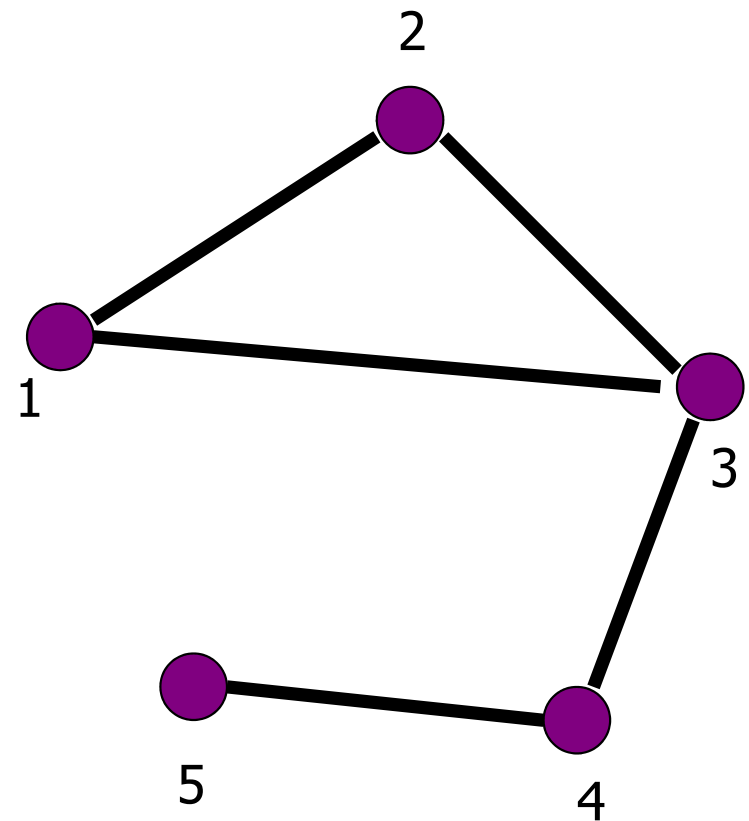
(1,2)

(2,3)

(1,3)

(3,4)

(4,5)



# Graph Representation

- List of edges
  - Keep a list of all the directed edges in the **directed** graph

(1,2)

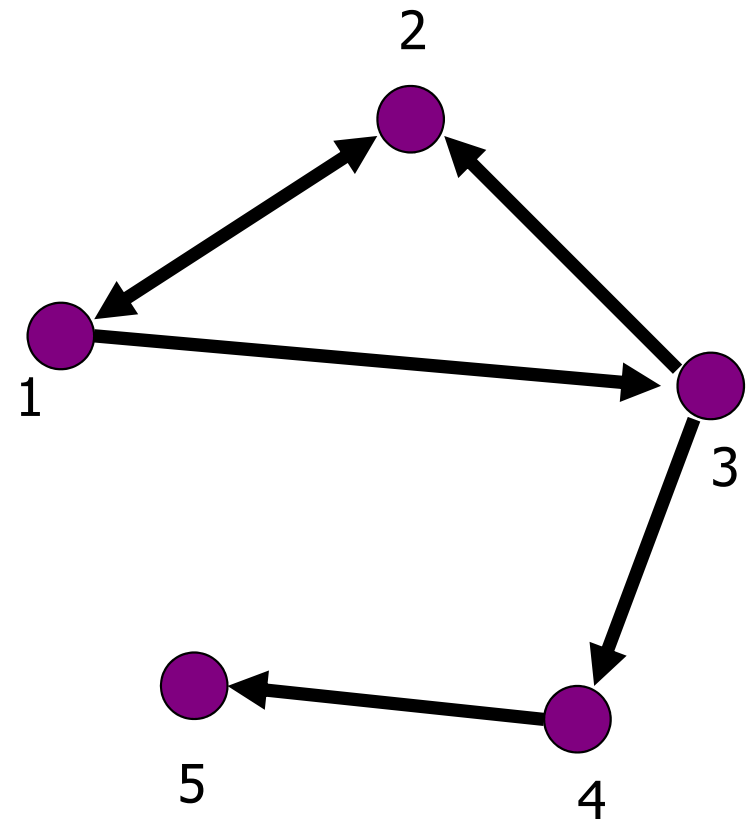
(2,1)

(1,3)

(3,2)

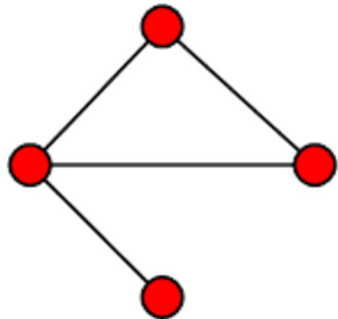
(3,4)

(4,5)

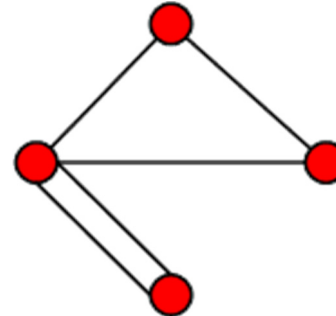


# Simple Graphs and Multigraphs

- **Simple** graphs are graphs where only a single edge can be between any pair of nodes
- **Multigraphs** are graphs where you can have multiple edges between two nodes and loops



Simple graph



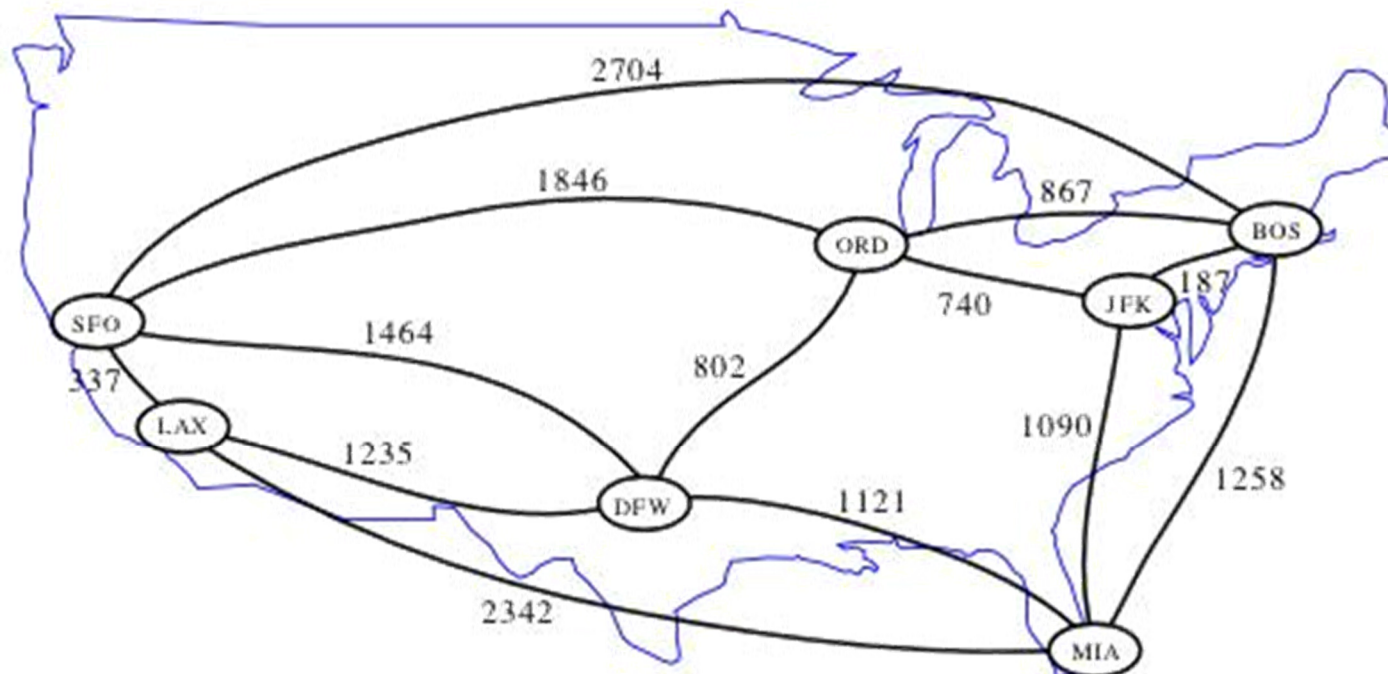
Multigraph

# Weighted Graph

- A **weighted** graph  $G(V, E, W)$  is one where edges are associated with weights

$$A_{ij} = \begin{cases} w_{ij} \text{ or } w(i, j), w \in R \\ 0, \text{ there is no edge between } v_i \text{ and } v_j \end{cases}$$

- An example on US flight network:







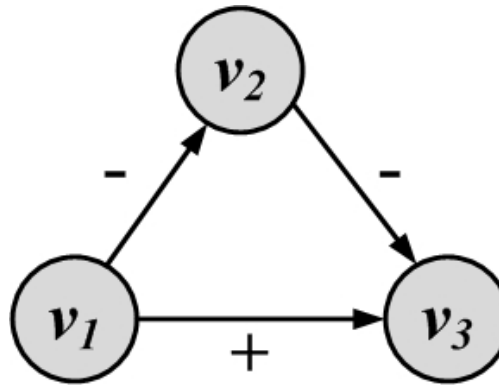
# Edge Attributes

## Possible options

- Weight (e.g. frequency of communication)
- Ranking (best friend, second best friend...)
- Type (friend, relative, co-worker)
- Sign: Friend vs. Foe, Trust vs. Distrust
- Properties depending on the structure of the rest of the graph: number of common friends

# Signed Graph

- When weights are binary (0/1, -1/1, +/-), we have a **signed** graph

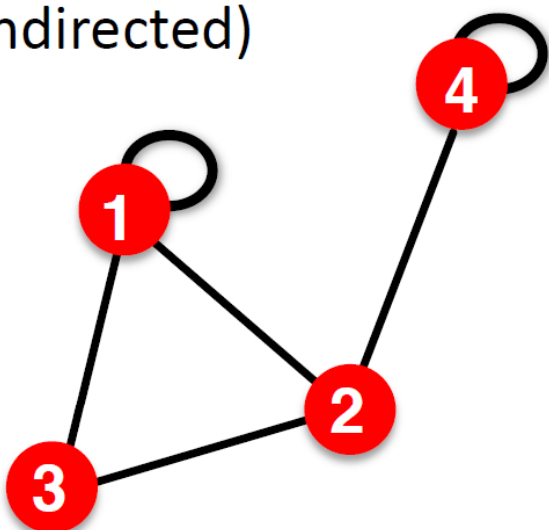


- It is used to represent **friends** or **foes**
  - Trust/Distrust (Epinions), Support/Oppose (Wikipedia)  
Like/Dislike (YouTube),
- It is also used to represent **social status**
  - $A \xrightarrow{+} B$  : B has **higher** status than A

# More Types of Graphs

## ■ Self-edges (self-loops)

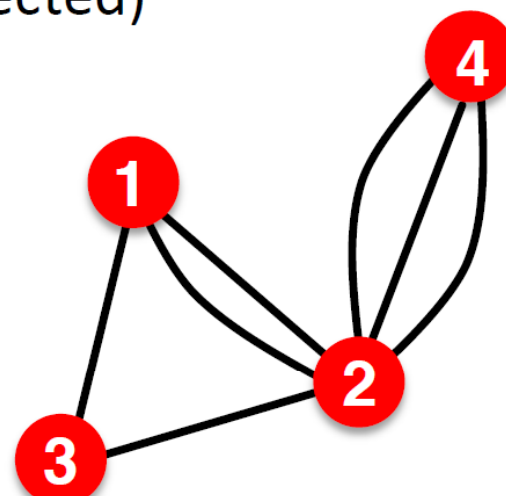
(undirected)



$$A_{ij} = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 1 \end{pmatrix}$$

## ■ Multigraph

(undirected)



$$A_{ij} = \begin{pmatrix} 0 & 2 & 1 & 0 \\ 2 & 0 & 1 & 3 \\ 1 & 1 & 0 & 0 \\ 0 & 3 & 0 & 0 \end{pmatrix}$$



# Example Graphs

Email network >> directed multigraph with self-edges

Facebook friendships >> undirected, unweighted

Citation networks >> unweighted, directed, acyclic

Collaboration networks >> undirected multigraph or weighted graph

Mobile phone calls >> directed, (weighted?) multigraph

Protein Interactions >> undirected, unweighted with self-interactions