

## Week 12.1 — Sorting & Recursion

## Recursion — reminder

- ▶ Functions call themselves (possibly indirectly)
- ▶ Make sure the recursion will stop eventually (base case/s).
- ▶ Recursive calls should be to smaller / simpler instances.
- ▶ Beware stack overflow.

# Why recursion?

“It's easier”

- ▶ ... to read?
- ▶ ... to write?
- ▶ ... for whom?

Is the person saying it:

- ▶ a theoretical computer scientist?
- ▶ a fan of “functional programming” languages<sup>1</sup>?
  - ▶ Languages without loops = have to use recursion

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<sup>1</sup>Lisp, Haskell, Erlang, ...

# Why not iteration?

Can't we use iteration<sup>2</sup> instead?

Sort of.

There is a theoretical result that yes, recursive algorithms can be converted into algorithms which use loops (and a stack).

But, there are some algorithms which really are more clearly expressed in a recursive form.

$$F(1) = 1$$

$$F(2) = 1$$

$$F(n) = F(n-1) + F(n-2)$$

Got an iterative Formulation?

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<sup>2</sup>loops

# Sorting

We'll look at sorting algorithms as an example of a task which recursion can make easier.

# Sorting a List in Java

```
List li;  
...  
li.sort(null);  // use default ordering
```

See `Sort.java`

Doesn't give us much insight into how it works though.

# Merge Sort

Task: Sort `Arr` — an array of  $N$  ints.

- ▶ Suppose we know how to sort  $N/2$  ints.
- ▶ Split `Arr` in half.
- ▶ Sort each half
- ▶ Merge the two sorted parts into one sorted whole.

See `MSort.java`, Merge Sort video.

Works because a single element array is automatically sorted.  
Sorts “bottom up”.

# Quick Sort

Suppose we have a routine  $\text{Part}^3(\text{Arr})$ :

- ▶ returns the index of one element in  $\text{Arr}$  which is guaranteed to be in the correct place.
- ▶ everything to the left of that element is  $\leq$  than it.
- ▶ everything to the right of that element is  $\geq$  than it.

See `QSort.java`, Quick Sort video

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<sup>3</sup>partition



# Which is Better?

It depends:

<b>Criterion</b>	<b>Better</b>
Simplicity	Mergesort - easier to understand
Worst case performance	Mergesort - $O(n \log n)$
Expected performance	Quicksort - $O(n^2)$
Memory requirements	Quicksort - only one array

Why? See an algorithms course.