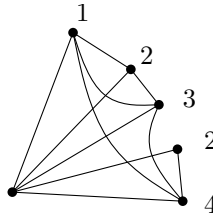


→ Kempe:

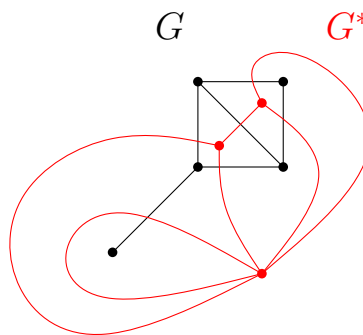


Heawood

Def: (p. 267)

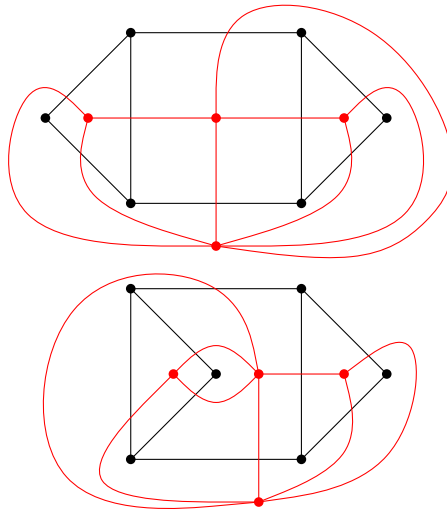
Let  $G$  be a plane graph. The dual of  $G$ , denoted  $G^*$ , is a plane multigraph. From each region of  $G$  we pick one inner point to be a vertex of  $G^*$ . For every edge of  $G$  we add a curve connecting the vertices of  $G^*$  corresponding to the regions incident with  $e$ , such that this curve intersects  $e$  once and does not intersect anything else (including itself).

Example:



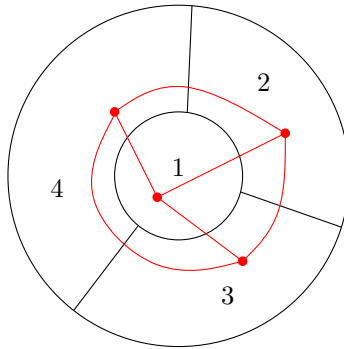
→ The name “dual” is justified.

→ We also talk about the dual of a planar graph, which is not unique.



⇒ The dual of a planar graph depends on the embedding.

Thm: “Every map can be colored in 4 colors.”



→ Taking the dual of a map gives a planar graph.

→ Every planar graph is 4-colorable

About the proof:

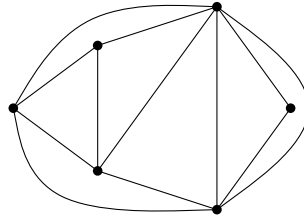
→ reducible configurations.

→ unavoidable set of reducible configurations.

5-color thm

$$\left. \begin{array}{l} 1 \text{ vertex of deg } 0 \\ 1 \text{ vertex of deg } 1 \\ \vdots \\ 1 \text{ vertex of deg } 5 \end{array} \right\} \text{unavoidable}$$

→ Triangulation:



$$2e = 3r$$

### Ramsey Theory

Def: (p. 299)

The ramsey number  $r(F_1, F_2)$  is the minimal number  $r$  such that in every red-blue coloring of the edges of  $K_r$  there is either a red copy of  $F_1$  or a blue copy of  $F_2$ .

→  $r(n, m)$  is  $r(K_n, K_m)$

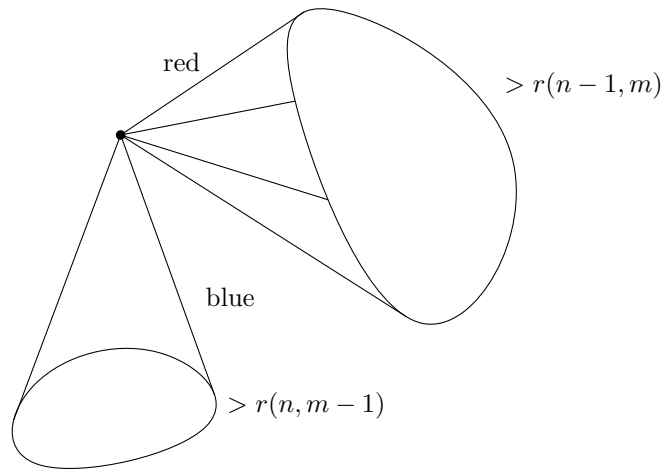
→  $r(n) = r(n, n)$

Thm (Ramsey):  $r(n, m)$  is finite.

Proof: by double induction.

$r(1, n) = r(m, 1) = 1$

→ Assume that  $r(n', m')$  is finite for all pairs  $n', m' < n, m$ .



→ Let  $r = r(n-1, m) + r(n, m-1)$

→ Fix a vertex  $v$ .

→ Let  $N_{\text{red}}$  be the set of vertices adjacent to  $v$  via a red edge.

→  $N_{\text{blue}}$ .

→ If both  $|N_{\text{red}}| < r(n-1, m)$  and  $|N_{\text{blue}}| < r(n, m-1)$  ✗

$$r = |N_{\text{red}}| + |N_{\text{blue}}| + 1 < r(n-1, m) + r(n, m-1) + 1$$

→ One of them is large enough and we can finish