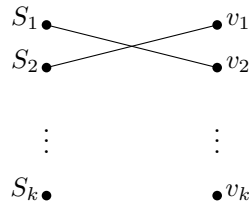


Tutte's Thm:

$k_o(G - S)$, G contains a perfect matching iff $k_o(G - S) \leq |S| \forall S \subseteq V(G)$. \odot

Proof:

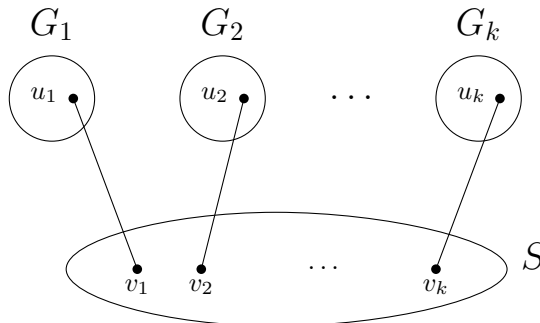
- Induction on number of vertices in G .
 - Assume that G has \odot .
- Let S be a maximal set of vertices having $k_o(G - S) = |S|$.
- S is not empty.
 - Let G_1, \dots, G_k be the components of $G - S$, then $|G_i|$ is odd $\forall 1 \leq i \leq k$.
 - Let S_i be the set of neighbors of G_i in S .
 - S_i is not empty. (G_i is odd and all the connected components of G are even).



- $\odot\odot$ For every $1 \leq t \leq k$ and every t of the S_i 's, the union of these S_i 's is of size at least t .
- Otherwise, let T be the union of the S_i 's.

$$k_o(G - T) \geq t > |T| \nrightarrow \odot$$

- Consider the bipartite graph with sides $\{S_1, \dots, S_k\}$ and $S = \{v_1, \dots, v_k\}$. There is an edge in H between S_j and v_i iff $\forall v_i \in S_j$
- by $\odot\odot$ +Hall's Theorem, there is a perfect matching in H .
- Let u_i be the neighbor of v_i in G_i (assuming without loss of generality that the perfect matching matched v_i to S_i).



→ Need to show that $\forall 1 \leq i \leq k$ and $\forall W \subseteq V(G_i - u_i)$

$$k_o(G_i - u_i - W) \leq |W|$$

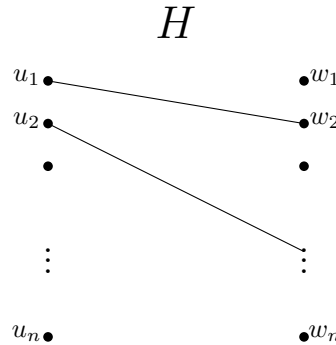
→ Assume otherwise: $|W| < k_o(G_i - u_i - W)$ for some i and W .

→ Since $G_i - u_i$ is even, the parity of $|W|$ and $k_o(G_i - u_i - W)$ is the same.

→ Consider $S' = S \cup W \cup \{u_i\}$

$$|S'| \geq k_o(G - S') = k_o(G - S) + k_o(G_i - u_i - W) - 1 \geq |S| + |W| + 2 - 1 = |S'| \nrightarrow \text{maximality of } S.$$

by \odot



→ By the induction hypothesis, there is a perfect matching M_i in G_i .

$$M = \left(\bigcup_{i=1}^k M_i \right) \cup \{v_i u_i\} \text{ is a perfect matching in } G$$

Theorem (Tutte-Berge formula): For every graph G , the size of a maximum matching is

$$\min_{S \subseteq V(G)} \frac{(|S| - k_o(G - S) + |V|)}{2}$$

Def: Let k be a positive integer. A k -factor in a graph G is a spanning subgraph which is k -regular.

Example: A perfect matching is a 1-factor.

Theorem (Petersen): A graph G can be decomposed into 2-factors F_1, \dots, F_k if and only if G is

$$2k\text{-regular.} \quad \left(G = \bigcup_{i=1}^k F_i \right)$$

Proof's idea: one direction is easy (decomposition $\implies 2k$ -regular).

→ Assume that G is $2k$ -regular. By Euler's Theorem, there is an Eulerian circuit C .

→ (Def. H)

→ H is k -regular

→ By Hall's Theorem and counting every regular bipartite graph contains a perfect matching.

→ Every perfect matching in H corresponds to a 2-factor of G .

→ repeat. ■