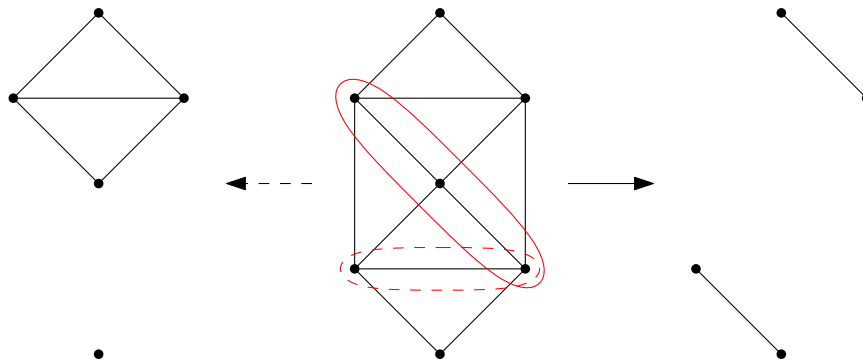


(Havel-Hakimi)

Def: (p. 115-116):

- A vertex-cut in a graph G is a set U such that $G - U$ is disconnected.
- A vertex-cut of minimal cardinality is called a minimal vertex cut.
- A vertex-cut U such that no proper subset of U is a vertex-cut is called a minimal vertex-cut.



- Every minimum vertex-cut is minimal.
- A graph contains a vertex-cut iff it is not complete.
- The vertex-connectivity of a graph G , denoted by $\kappa(G)$, is the size of a smallest set U such that $G - U$ is disconnected or trivial.
- The number of connected components in G will be denoted $k(G)$.
- A graph is said to be k -connected if $\kappa(G) \geq k$.

Def (p. 116-117):

- An edge-cut $X \subseteq E(G)$ is a set of edges such that $G - X$ is disconnected.
- An edge-cut with minimal size is called a minimum edge-cut.
- An edge-cut for which no proper subset is an edge-cut is called a minimal edge-cut.
- $\lambda(G)$, the edge-connectivity of G , is the size of a minimal $X \subseteq E(G)$ such that $G - X$ is disconnected or trivial.
- G is k -edge-connected if $\lambda(G) \geq k$.

Property (theorem 5.11, Whitney): For all G , $\kappa(G) \leq \lambda(G) \leq \delta(G)$

Proof:

→ If G is disconnected or trivial, $\kappa(G) = \lambda(G) = 0$ ✓

→ If G is complete graph then $\kappa(K_n) = n - 1 = \lambda(K_n)$

→ removing all $n - 1$ edges incident with one vertex disconnects the graph.

→ Let X be an edge-cut and assume that $G - X$ has two components of size j and $n - j$.

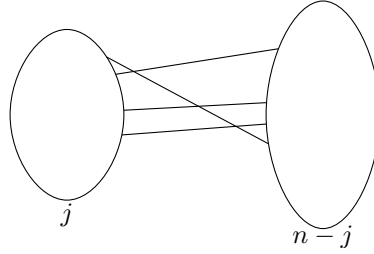
→ $|X| = j(n - j)$

→ both components are not empty, so $j \geq 1, n - j \geq 1$

→ $0 \leq (j - 1)(n - j - 1) = j(n - j) - j - n + j + 1 = j(n - j) - n + 1$

⇒ $|X| = j(n - j) \geq n - 1$

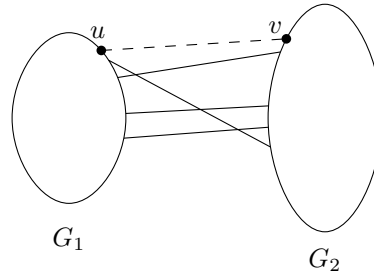
⇒ $\lambda(K_n) = n - 1 = \delta(K_n)$.



→ Assume that G is non-trivial, connected, not complete and with at least three vertices.

→ $\lambda(G) \leq \delta(G)$ ✓ (removing all edges incident with a vertex of minimum degree disconnects the graph).

→ Let X be a minimum edge-cut, and let G_1 and G_2 be the two components of $G - X$.



→ If all of the edges between G_1 and G_2 are in G , then $|X| = j(n - j)$ where j is the number of vertices in G_1 . Then $j \geq 1$ and $n - j \geq 1 \Rightarrow j(n - j) \geq n - 1$ contradicting the facts that $\delta(G \not\cong K_n) < n - 1$ and $\lambda(G) \leq \delta(G)$.

→ Since not all the edges between G_1 and G_2 are in G , then we have $u \in G_1$ and $v \in G_2$ such that $uv \notin E(G)$.

→ Define U as follows. For every $e \in X$, pick a vertex incident with e as follows.

→ if $u \in e$, pick $e \cap V(G_2)$

→ otherwise, pick $e \cap V(G_1)$

- $|U| \leq |X|$
- $X \cap E(G - U) = \emptyset$

