Assignment 7

15-359 Probability and Computing

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Section: B

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## Problem 1: Spin me right around

A. Since X and Y are independent,

$$f_{X,Y}(x,y) = f_X(x)f_Y(y) = \frac{e^{-\frac{x^2}{2}}}{\sqrt{2\pi}} \cdot \frac{e^{-\frac{y^2}{2}}}{\sqrt{2\pi}} = \boxed{\frac{1}{2\pi}e^{-\frac{x^2+y^2}{2}}}.$$

B. If r is the distance of (x, y) from the origin,  $r = x^2 + y^2$ , so that

$$f_{X,Y}(x,y) = \frac{1}{2\pi}e^{-\frac{x^2+y^2}{2}} = \frac{1}{2\pi}e^{-\frac{r^2}{2}},$$

so that  $f_{X,Y}(x,y)$  depends only on r.

### Problem 2: Unwieldy numbers

Number the tweets  $1, 2, \ldots, 340 \times 10^6$  in some order.  $\forall i \in \{1, 2, \ldots, 340 \times 10^6\}$ , let  $L_i$  be a random variable denoting the length of the  $i^{th}$  tweet. Since  $L_i \sim \text{Uniform}(10, 140)$ ,  $E[L_i] = \frac{10+140}{2} = 75$ , and  $\text{Var}[L_i] = \frac{(140-10)^2}{12} = 4225/3$ . Let L be a random variable denoting the total number of characters processed by Twitter on an average day, so that  $L = \sum_{i=1}^{340\times 10^6} L_i$ . Then, since  $340\times 10^6$  is large, and  $L_1, L_2, \ldots, L_{340\times 10^6}$  are independent and identically distributed, by the Central Limit Theorem, we can approximate the distribution of  $\frac{L-n\mu}{\sigma\sqrt{n}}$  with  $\mathcal{N}(0,1)$ , and therefore we can approximate the distribution of L as

$$L \sim \mathcal{N}(n\mu, \sigma^2 n) = \mathcal{N}(340 \times 10^6 * 75, 340 \times 10^6 * 4225/3).$$

Integrating the probability density function numerically from  $25 \times 10^9$  to  $26 \times 10^9$  gives

$$P(25 \times 10^9 \le L \le 26 \times 10^9) \approx 1.$$

### Problem 3: Exp farming

As shown in class,  $\widetilde{X}(s) = \frac{\lambda}{s+\lambda}$ . It can be shown by induction on k that,  $\forall k \in \mathbb{N}$ ,

$$\frac{d^k}{ds^k} \left( \frac{\lambda}{s+\lambda} \right) = (-1)^k \frac{k!\lambda}{(s+\lambda)^{k+1}}.$$

Thus, the  $k^{th}$  moment of X is given by

$$E[X^k] = (-1)^k \frac{d^k \widetilde{X}(s)}{d s^k} \bigg|_{s=0}$$

$$= (-1)^{2k} \frac{k! \lambda}{(s+\lambda)^{k+1}} \bigg|_{s=0}$$

$$= \frac{k! \lambda}{\lambda^{k+1}} = \left[\frac{k!}{\lambda^k}\right]$$

#### Problem 4: Random amounts of randomness

1. Since  $X_1, X_2, \ldots, X_{10}$  are independent,  $e^{-sX_1}, e^{-sX_2}, \ldots, e^{-sX_{10}}$  are independent as well, and, since  $X_1, X_2, \ldots, X_{10}$  are identically distributed,  $\widetilde{X}_1 = \widetilde{X}_2 = \ldots = \widetilde{X}_{10} = \widetilde{X}$ . Thus,

$$\widetilde{S}(s) = E\left[e^{-sS}\right] = E\left[e^{-s\left(\sum_{i=1}^{10} X_i\right)}\right] = E\left[\prod_{i=1}^{10} e^{-sX_i}\right] = \prod_{i=1}^{10} E\left[e^{-sX_i}\right] = \prod_{i=1}^{10} \widetilde{X}_i(s) = \boxed{\left(\widetilde{X}(s)\right)^{10}}.$$

2. By definition of the z-transform and the Laplace Transform,

$$\widehat{N}\left(\widetilde{X}(s)\right) = E\left[\left(\widetilde{X}(s)\right)^{N}\right] = E\left[E\left[e^{-sX}\right]^{N}\right].$$

Since  $X_1, X_2, \ldots, X_i$  are i.i.d.,  $e^{-sX_1}, \ldots, e^{-sX_n}$  are i.i.d. as well, so that

$$\widehat{N}\left(\widetilde{X}(s)\right) = E\left[E\left[e^{-sXN}\right]\right] = E\left[e^{-s\sum_{i=1}^{N}X_{i}}\right] = E\left[e^{-sS}\right] = \widetilde{S}(s) \quad \blacksquare.$$

3. By the result of part 2. above,  $\widetilde{S}(s) = \widehat{N}(\widetilde{X}(s))$ . Thus, since we know the z-transform of Geometric(p) and the Laplace Transform of  $\text{Exp}(\mu)$ ,

$$\widetilde{S}(s) = \widehat{N}\left(\widetilde{X}(s)\right) = \frac{\widetilde{X}(s)p}{1-\widetilde{X}(s)p} = \frac{\left(\frac{\mu}{\mu+s}\right)p}{1-\left(\frac{\mu}{\mu+s}\right)p} = \frac{\mu p}{\mu+z-\mu(1-p)} = \frac{\mu p}{s+\mu p}.$$

Thus, the Laplace Transform of S is that of  $\text{Exp}(\mu p)$ , so that  $S \sim \text{Exp}(\mu p)$ .

### Problem 5: Mouse in a maze

1. Let X, Y, and Z be random variables, with X denoting the time the mouse takes to return to his starting position if he goes left (so that  $X \sim \text{Exp}(1)$ ), Y denoting the time the mouse takes if he goes right and then leaves (so that  $Y \sim \text{Exp}(2)$ ), and Z denoting the time the mouse takes if he goes right and then returns to his starting position (so that  $Z \sim \text{Exp}(3)$ ). Then,

$$E[T] = \frac{1}{2} (X + E[T]) + \frac{1}{2} (\frac{1}{3}Y + \frac{2}{3} (Z + E[T])).$$

Furthermore,

$$E[T] = E[E[T]] = E\left[\frac{1}{2}\left(X + E[T]\right)\right) + \frac{1}{2}\left(\frac{1}{3}Y + \frac{2}{3}\left(Z + E[T]\right)\right)\right] = \frac{25}{36} + \frac{5}{6}E[T],$$

so that, solving for E[T] gives  $E[T] = \boxed{\frac{25}{6}}$ .

## Problem 6: Chernoff proof relents

Let  $X_p$  and  $X_q$  be random variables, counting the number of heads among n coin flips each of two coins, with probabilities p and q or getting heads, respectively (where p > q). Then, since q > p (and, consequently, 1 - q < 1 - p),  $\forall k \ge nq$ ,

$$\frac{P(X_p = k)}{P(X_q = k)} = \frac{\binom{n}{k} p^k (1 - p)^{n - k}}{\binom{n}{k} q^k (1 - q)^{n - k}} \\
= \left(\frac{p}{q}\right)^k \left(\frac{1 - p}{1 - q}\right)^{n - k} \\
= \left(\frac{p}{q}\right)^{qn} \left(\frac{1 - p}{1 - q}\right)^{n - qn} \\
\le \left(\frac{p}{q}\right)^{qn} \left(\frac{1 - p}{1 - q}\right)^{n - qn} \\
= e^{-nKL(q, p)}$$

Thus, taking the sum over all  $k \geq qn$  gives

$$P(X_n \ge qn) \le P(X_q \ge qn)e^{-nKL(q,p)} \le e^{-nKL(q,p)}$$

since  $P(X_q \ge qn)$  is a probability and therefore  $P(X_q \ge qn) \le 1$ , proving the desired result.

# Problem 7: Polling is cheap. (Tell that to a systems programmer.)

Let p be the actual fraction of people who disapprove of the president. Then, assuming it is possible for people to be polled multiple times, the probability that any particular person polled disapproves of the president is p, and  $d \sim \text{Binomial}(p,n)$ , since people's opinions are independent and each person polled has chance p of being among those who disapprove of the president. Then, E[d] = pn, so that, by Hoeffding's Inequality, for  $\epsilon = 2\% = 0.02$ ,  $19/20 = P((p-\epsilon)n \le d \le (p+\epsilon)n) \ge 2e^{-2\epsilon^2 n}$ , and thus. Solving for n gives

$$n \ge -\frac{\ln(19/40)}{2\epsilon^2} \approx \boxed{931.}$$