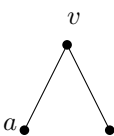


→ Brooks Theorem: G is connected, not complete, not odd cycle $\chi(G) \leq \Delta(G)$.

→ Assume that G is 2-connected.

→ Assume that $\Delta(G) \geq 3$

→ Assume that G is Δ -regular



→ Find a, b and $G - a - b$ is connected.

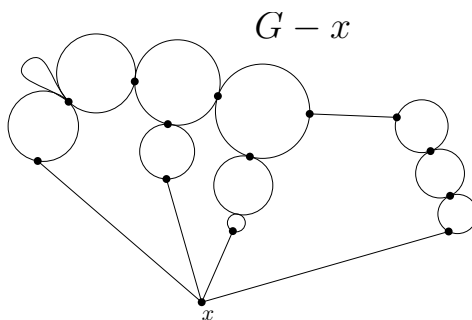
→ Once we have 3 vertices a, b, v such that $av, bv \in E(G)$, $ab \notin E(G)$ and $G - a - b$ is connected, color a and b by color 1. Find a spanning tree for $G - a - b$; root the tree at v . Color vertices according to their place in the tree from leaves towards the root. This can be done because every vertex has at most $\Delta - 1$ colored neighbors (its parent is not colored).

→ When we try to color v , it has at most $\Delta - 2$ colored neighbors besides a, b . But a and b are both colored 1.

→ Consider a vertex x that is not adjacent to all other vertices.

→ If $G - x$ is still 2-connected, find a vertex of distance 2 from x (call it y). Let v be a common neighbor of x and y . Letting x and y have the roles of a and b works. Indeed, $G - x - y$ is connected because $G - x$ is 2-connected.

→ Assume that $G - x$ is not 2-connected. Consider the block decomposition of $G - x$.



→ We have a tree of blocks, there are at least two end blocks (because every tree has at least two leaves).

→ An end block B_i has a vertex j_i such that every other block B_k is either disjoint from B_i , or they have j_i as their only common vertex.

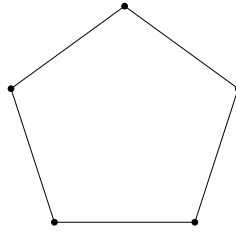
→ Let B_1 and B_2 be two end blocks. There are two vertices $b_1 \in B_1, b_2 \in B_2$ such that $b_1 \neq j_1, b_2 \neq j_2$, and $xb_1 \in E(G)$ and $xb_2 \in E(G)$.

→ otherwise, if such b_1 does not exist, then j_1 is a cut vertex of $G \nmid G$ is 2-connected.

→ x, b_1, b_2 can have the roles of v, a, b .

→ b_1 and b_2 are neighbors of x , by the above. They are not adjacent because they are in different blocks in G .

→ $G - b_1 - b_2 - x$ is connected, because neither b_1 nor b_2 was a joint and $d_G(x) \geq 3$. ■



$$\chi(C_5) = 3$$

$$\omega(C_5) = 2$$

Theorem 10.10: \forall integer k there is a triangle free graph with chromatic number k .

\rightarrow For every forest F , $\chi(F) \leq 2$.

Theorem (Erdős): For all integers k, ℓ , $\exists G$ such that $\text{girth}(G) > \ell$ and $\chi(G) > k$.

Proof: \rightarrow set $0 < \theta < \frac{1}{\ell}$ constant

\rightarrow define: $p = n^{-1+\theta}$

\rightarrow Consider a graph on n vertices such that every possible edge is in G with probability P , independently of all other edges.

Let X be the number of short ($\leq \ell$) cycles in G . X is a random variable.

$$\mathbb{E}[x] = \sum_{i=3}^{\ell} (\# \text{ of } i\text{-cycles in } K_n) \cdot p^i = \sum_{i=3}^{\ell} \frac{n(n-1)\cdots(n-i+1)}{2 \cdot \underset{\substack{\uparrow \\ \text{starting point} \\ + \text{direction}}}{i}} \leq \sum_{i=3}^{\ell} n^i p^i \leq 2 \cdot (np)^\ell = 2n^{\theta\ell} \underset{\substack{n \text{ is large} \\ \text{enough}}}{<} \frac{n}{\log n}$$

$$\Pr[X \geq n/2] \leq \frac{\mathbb{E}[X]}{n/2} \leq \frac{2n}{\log n \cdot n} = \frac{2}{\log n} \xrightarrow{n \rightarrow \infty} 0$$

Markov's inequality: if X is a non-negative random variable with expectation then for any positive real a

$$\Pr[X > a] \leq \frac{\mathbb{E}[X]}{a}$$