

# Solution to Homework#0

## 1 Probability and Statistics

### 1.1 (combinatorics)

**Proof:** To prove the equation  $C_n^k = \frac{n!}{k!(n-k)!}$ , **mathematical induction can be applied there.**

S1 when  $n = 0$ ,  $k$  can only value as 0 or 1.

In those cases,  $C_n^k = 1 = \frac{n!}{k!(n-k)!}$  holds

S2 Assume that it holds when  $n = t$ , as  $C_n^k = \frac{n!}{k!(n-k)!}$ .

Then for  $n = t + 1$

$$C_{t+1}^k = C_t^k + C_t^{k-1} = \frac{t!}{k!(t-k)!} + \frac{t!}{(k-1)!(t-k+1)!} = \frac{t!}{k!(t-k+1)}(t-k+1+k) = \frac{(t+1)!}{k!(t+1-k)!}$$

hence the equation holds when  $n = t + 1$

So

$$C_n^k = \frac{n!}{k!(n-k)!}$$

for  $n \geq 1$  and  $0 \leq k \leq n$

□

### 1.2 (counting)

Let random variable  $X$  represent the number of the heads, and  $X \sim B(10, \frac{1}{2})$ , then

$$P(X = 4) = C_{10}^4 \left(\frac{1}{2}\right)^4 \left(\frac{1}{2}\right)^6 = \frac{105}{512}$$

Let event A be the event of getting a full house(XXXYY) from a deck of 52 cards, then

$$P(A) = \frac{A_4^2 \cdot C_{13}^3 \cdot C_{13}^2}{C_{52}^5} = \frac{429}{4165}$$

### 1.3 (conditional probability)

Let event A represent the event of all three tosses resulting in head, while event B represent the event of at least one of the tosses resulting in head. Then

$$P(A|B) = \frac{P(AB)}{P(B)} = \frac{\left(\frac{1}{2}\right)^3}{1 - \left(\frac{1}{2}\right)^3} = \frac{1}{7}$$

### 1.4 (Bayes theorem)

$$\begin{aligned} P(X < 0) &= \frac{P(bit=1, X=-1)}{P(bit=0, X=1) + P(bit=1, X=-1)} \\ &= \frac{\frac{1}{2} \times \frac{1}{4}}{\frac{1}{2} \times \frac{1}{8} + \frac{1}{2} \times \frac{1}{4}} = \frac{2}{3} \end{aligned}$$

## 1.5 (union/intersection)

As applications of the principle of Inclusion-exclusion,

$$P(A \cap B) = P(A) + P(B) - P(A \cup B)$$

and

$$P(A \cup B) = P(A) + P(B) - P(A \cap B)$$

Hence

$$P(A \cap B)_{\max} = \min[P(A), P(B)] = P(A) = 0.3$$

$$P(A \cap B)_{\min} = \max[0, P(A) + P(B) - 1] = 0$$

$$P(A \cup B)_{\max} = \min[P(A) + P(B), 1] = 0.7$$

$$P(A \cup B)_{\min} = \max[P(A), P(B)] = P(B) = 0.4$$

## 2 Linear Algebra

### 2.1 (rank)

$$r \begin{pmatrix} 1 & 2 & 1 \\ 1 & 0 & 3 \\ 1 & 1 & 2 \end{pmatrix} = 2$$

### 2.2 (inverse)

$$\begin{pmatrix} 0 & 2 & 4 \\ 2 & 4 & 2 \\ 3 & 3 & 1 \end{pmatrix}^{-1} = \begin{pmatrix} 0 & 2 & 3 \\ 2 & 4 & 3 \\ 4 & 2 & 1 \end{pmatrix}$$

### 2.3 (eigenvalue/eigenvectors)

Table 1: Eigenvalues and their respective eigenvectors

$\alpha_1=4$	$p_1 = \begin{bmatrix} -1 & -2 & 1 \end{bmatrix}^T$
$\alpha_2=2$	$p_2 = \begin{bmatrix} -1 & 0 & 1 \end{bmatrix}^T$
$\alpha_3=2$	$p_3 = \begin{bmatrix} -1 & 1 & 0 \end{bmatrix}^T$

### 2.4 (singular value decomposition)

(a) For  $\Sigma$  is a rectangular diagonal matrix, so is  $\Sigma^\dagger$

$$\Sigma \Sigma^\dagger_{[i][j]} = \sum_{k=1}^m \Sigma_{[i][k]} \cdot \Sigma^\dagger_{[k][j]}$$

Apparently,  $\Sigma \Sigma^\dagger$  will be an  $n \times n$  identity matrix. Since  $U$  and  $V$  are both unitary matrix,

$$\begin{cases} UU^T = U^T U = I \\ VV^T = V^T V = I \end{cases}$$

$$\begin{aligned}
MM^\dagger M &= (U \Sigma V^T)(V \Sigma^\dagger U^T)(U \Sigma V^T) \\
&= U \Sigma \Sigma^\dagger \Sigma V^T \\
&= U \Sigma V^T \\
&= M
\end{aligned}$$

(b)

$$MM^\dagger = (U\Sigma V^T)(V\Sigma^\dagger U^T) = U\Sigma\Sigma^\dagger U^T = UU^T = I$$

Analogously,

$$M^\dagger M = I$$

So,

$$M^{-1} = M^\dagger$$

## 2.5 (PD/PSD)

(a)

$$\mathbf{x}^T Z Z^T \mathbf{x} = (\mathbf{x}Z)(\mathbf{x}Z)^T \geq 0$$

So  $Z Z^T$  is a postive semi-defined matrix.

(b) Since  $A$  is a real symmetric matrix, the eigenvectors of  $A$  (presented as  $v_i$ , and their respectful eigenvalues are defined as  $\lambda_i$ ) can be orthogonal and unitary.

Let  $\mathbf{x} = \sum k_i v_i$  and  $\mathbf{x}^T = \sum k_i v_i^T$

$$\begin{aligned}\mathbf{x}^T A \mathbf{x} &= (\sum k_i v_i^T) A (\sum k_i v_i) \\ &= (\sum k_i v_i^T) (\sum k_i \lambda_i v_i) \\ &= \sum \lambda_i k_i^2 > 0\end{aligned}$$

Hence,  $A$  is postive-defined.

## 2.6 (inner product)

When  $\mathbf{x}$  and  $\mathbf{u}$  have the same direction,  $\mathbf{x}^T \mathbf{u}$  reaches its maximum  $\|\mathbf{x}\|$ .

When  $\mathbf{x}$  and  $\mathbf{u}$  have the opposite direction,  $\mathbf{x}^T \mathbf{u}$  reaches its minimum  $\|\mathbf{x}\|$ .

When  $\mathbf{x}$  and  $\mathbf{u}$  are orthogonal,  $|\mathbf{x}^T \mathbf{u}|$  reaches its minimum 0.

## 3 Calculus

### 3.1 (differential and partial differential)

$$\frac{df(x)}{dx} = \frac{-2}{e^{2x} + 1}$$

$$\frac{\partial g(x, y)}{\partial x} = 2e^{2y} + 6ye^{3xy^2}$$

### 3.2 (chain rule)

$$\frac{\partial f}{\partial v} = -y \sin v - x \cos v$$

### 3.3 (gradient and Hessian)

When  $u = 1$  and  $v = 1$ , the gradient of  $E$

$$\begin{aligned}\nabla E &= -2(e^v + 2ve^{-u})(ue^v - 2ve^{-u})\mathbf{i} + 2(ue^v - 2e^{-u})(ue^v - 2ve^{-u})\mathbf{j} \\ &= 2(4e^{-2} - e^2)\mathbf{i} + 2(2e^{-1} - e)^2\mathbf{j}\end{aligned}$$

Meanwhile, its Hessian matrix  $\nabla^2 E =$

$$\begin{pmatrix} 2e^2 - 12 + 16e^{-2} & 0 \\ 0 & 4e^2 - 12 + 8e^{-2} \end{pmatrix}$$

### 3.4 (Taylor's expansion)

$$\begin{aligned} E &= 2(4e^{-2} - e^2)(u - 1) + 2(2e^{-1} - e)^2(v - 1) \\ &\quad + 2e^2 - 6 + 8e^{-2})(u - 1)^2 + 4(e^2 - 3 + 4e^{-2})(v - 1)^2 \\ &\quad + o \end{aligned}$$

### 3.5 (optimization)

From the inequality of arithmetic and geometric means

$$\begin{aligned} Ae^\alpha + Be^{-2\alpha} &= \frac{1}{2}Ae^\alpha + \frac{1}{2}Ae^\alpha + Be^{-2\alpha} \geq 3\sqrt[3]{\frac{A}{2}e^\alpha \cdot \frac{A}{2}e^\alpha \cdot Be^{-2\alpha}} \\ &\geq 3\sqrt[3]{\frac{A^2B}{4}} \end{aligned}$$

with equality if and only if  $\frac{A}{2}e^\alpha = Be^{-2\alpha}$ , at that time,  $\alpha = \frac{1}{3}(\ln 2B - \ln A)$ .

### 3.6 (vector calculus)

$$\frac{dE}{d\mathbf{w}} = \frac{1}{2}(A\mathbf{w} + A\mathbf{w}) + \mathbf{b} = A\mathbf{w} + \mathbf{b}$$

Then,

$$\nabla E(\mathbf{w}) = A\mathbf{w} + \mathbf{b}$$

$$\frac{d^2E}{d\mathbf{w}^2} = A$$

So,

$$\nabla^2 E(\mathbf{w}) = A$$