
CSE 150A/250A. Assignment 3

3.2 Variable Elimination Algorithm

(a) Computation using variable elimination

We start with the following factors:

- $f_0(T)$ for $P(T)$
- $f_1(F)$ for $P(F)$
- $f_2(T, F, A)$ for $P(A|T, F)$
- $f_3(F, S)$ for $P(S|F)$
- $f_4(A, L)$ for $P(L|A)$
- $f_5(L, R)$ for $P(R|L)$

Eliminating the evidence variables S and R yields factors $f'_3(F) = f_3(F, S=1)$ and $f'_5(L) = f_5(L, R=1)$:

F	$f'_3(F)$	L	$f'_5(L)$
0	0.01	0	0.01
1	0.9	1	0.75

To eliminate L , we can define $f_6(A) = \sum_L f_4(A, L) \cdot f'_5(L)$. This involves 4 multiplications and 2 additions:

A	$f_6(A)$
0	$f_4(A=0, L=0)f'_5(L=0) + f_4(A=0, L=1)f'_5(L=1) = \mathbf{0.01074}$
1	$f_4(A=1, L=0)f'_5(L=0) + f_4(A=1, L=1)f'_5(L=1) = \mathbf{0.6612}$

To eliminate A , we can define $f_7(T, F) = \sum_A f_6(A) \cdot f_2(T, F, A)$. This involves 8 multiplications and 4 additions:

T	F	$f_7(T, F)$
0	0	$f_6(A=0)f_2(T=0, F=0, A=0) + f_6(A=1)f_2(T=0, F=0, A=1) = \mathbf{0.01080505}$
0	1	$f_6(A=0)f_2(T=0, F=1, A=0) + f_6(A=1)f_2(T=0, F=1, A=1) = \mathbf{0.6546954}$
1	0	$f_6(A=0)f_2(T=1, F=0, A=0) + f_6(A=1)f_2(T=1, F=0, A=1) = \mathbf{0.563631}$
1	1	$f_6(A=0)f_2(T=1, F=1, A=0) + f_6(A=1)f_2(T=1, F=1, A=1) = \mathbf{0.33597}$

To eliminate F , we can define $f_8(T) = \sum_F f_7(T, F) \cdot f_1(F) \cdot f'_3(F)$. This involves 8 multiplications and 2 additions:

T	$f_8(T)$
0	$f_7(T=0, F=0)f_1(F=0)f'_3(F=0) + f_7(T=0, F=1)f_1(F=1)f'_3(F=1) = \mathbf{0.00599923}$
1	$f_7(T=1, F=0)f_1(F=0)f'_3(F=0) + f_7(T=1, F=1)f_1(F=1)f'_3(F=1) = \mathbf{0.00860368}$

An alternate calculation of $f_8(T)$ uses 6 multiplications and 2 additions by computing

- $f'_8(F) = f_1(F) \cdot f'_3(F)$, via 2 multiplications, followed by
- $f_8(T) = \sum_F f'_8(F) \cdot f_7(T, F)$, via 4 multiplications, and 2 additions.

To combine the remaining T factors, we can define $f_9(T) = f_0(T) \cdot f_8(T)$. This involves 2 multiplications:

T	$f_9(T)$
0	$f_0(T=0)f_8(T=0) = \mathbf{0.00587924}$
1	$f_0(T=1)f_8(T=1) = \mathbf{0.00017207}$

Lastly, we normalize the factor f_9 to compute a distribution over T :

$$P(T=0|R=1, S=1) = \frac{f_9(T=0)}{\sum_T f_9(T)} = \frac{0.00587924}{0.00587924 + 0.00017207} = \mathbf{0.97156429}$$

$$P(T=1|R=1, S=1) = \frac{f_9(T=1)}{\sum_T f_9(T)} = \frac{0.00017207}{0.00587924 + 0.00017207} = \mathbf{0.02843571}$$

(b) Counting calculations used by the variable elimination algorithm

Phase of algorithm	# multiplications	# additions	# divisions
Eliminate S (evidence)	0	0	0
Eliminate R (evidence)	0	0	0
Eliminate L	4	2	0
Eliminate A	8	4	0
Eliminate F	6 or 8	2	0
Combine T factors	2	0	0
Normalize distribution over T	0	1	2
Total	20 or 22	9	2

(c) Counting calculations used by the enumeration algorithm

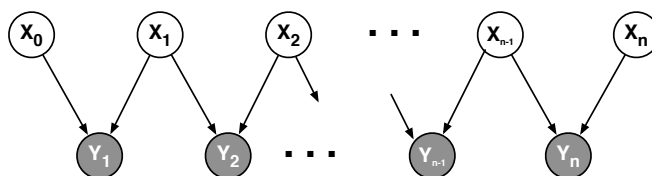
After marginalizing over F , A , and L , applying the product rule, and using the network's conditional independence relationships, we have

$$P(T = 0, S = 1, R = 1) = \sum_f \sum_a \sum_l P(T = 0)P(F = f)P(A = a|T = 0, F = f) \\ \cdot P(S = 1|F = f)P(L = l|A = a)P(R = 1|L = l)$$

Here, there are $2^3 = 8$ terms, each with 6 numbers multiplied together. So each term requires 5 multiplications, and they are summed together using 7 additions. The same process is repeated for $T = 1$, and we end by normalizing to get a distribution over T .

Phase of algorithm	# multiplications	# additions	# divisions
Compute $P(T = 0, S = 1, R = 1)$	$5 \times 2^3 = 40$	$2^3 - 1 = 7$	0
Compute $P(T = 1, S = 1, R = 1)$	$5 \times 2^3 = 40$	$2^3 - 1 = 7$	0
Normalize distribution over T	0	1	2
Total	80	15	2

3.3 More inference in a chain (10 pts)



(a) First term in numerator (2 pts)

$$\begin{aligned} P(Y_1|X_1) &= \sum_x P(Y_1, X_0=x|X_1) \quad \text{(marginalization)} \\ &= \sum_x P(X_0=x|X_1) P(Y_1|X_0=x, X_1) \quad \text{(product rule)} \\ &= \sum_x P(X_0=x) P(Y_1|X_0=x, X_1) \quad \text{(conditional independence)} \end{aligned}$$

(b) **Denominator (2 pts)**

$$\begin{aligned}
 P(Y_1) &= \sum_{x'} P(Y_1, X_1 = x') \quad \text{(marginalization)} \\
 &= \sum_{x'} P(X_1 = x') P(Y_1 | X_1 = x') \quad \text{(product rule)} \\
 &= \sum_{x, x'} P(X_0 = x) P(X_1 = x') P(Y_1 | X_0 = x, X_1 = x') \quad \text{(substitution from (a))}
 \end{aligned}$$

(c) **Second term in numerator (1 pt)**

$$P(X_n | Y_1, Y_2, \dots, Y_{n-1}) = P(X_n) \quad \text{(marginal independence)}$$

(d) **First term in numerator (3 pts)**

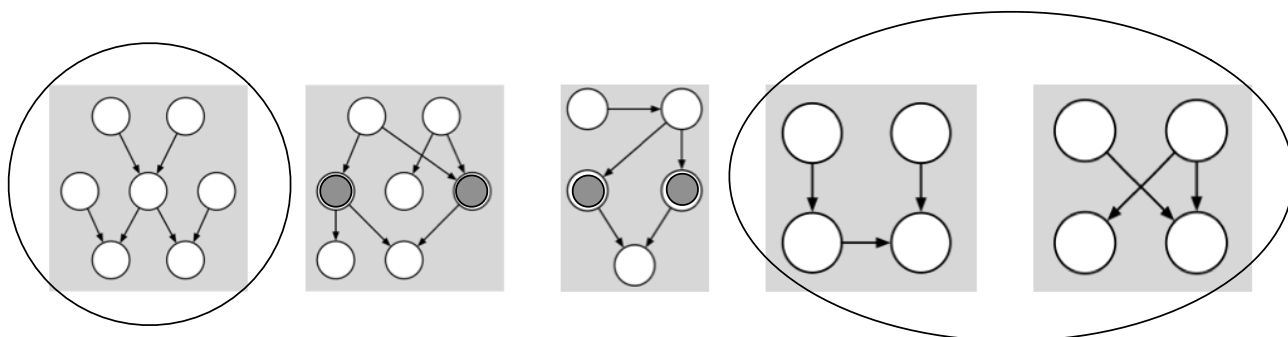
$$\begin{aligned}
 P(Y_n | X_n, Y_1, Y_2, \dots, Y_{n-1}) &= \sum_x P(Y_n, X_{n-1} = x | X_n, Y_1, Y_2, \dots, Y_{n-1}) \quad \text{(marginalization)} \\
 &= \sum_x P(X_{n-1} = x | X_n, Y_1, Y_2, \dots, Y_{n-1}) P(Y_n | X_{n-1} = x, X_n, Y_1, Y_2, \dots, Y_{n-1}) \quad \text{(product rule)} \\
 &= \sum_x P(X_{n-1} = x | Y_1, Y_2, \dots, Y_{n-1}) P(Y_n | X_{n-1} = x, X_n) \quad \text{(conditional independence)}
 \end{aligned}$$

(e) **Denominator (2 pts)**

$$\begin{aligned}
 P(Y_n | Y_1, Y_2, \dots, Y_{n-1}) &= \sum_{x'} P(Y_n, X_n = x' | Y_1, Y_2, \dots, Y_{n-1}) \quad \text{(marginalization)} \\
 &= \sum_{x'} P(Y_n | X_n = x', Y_1, Y_2, \dots, Y_{n-1}) P(X_n = x' | Y_1, Y_2, \dots, Y_{n-1}) \quad \text{(product rule)} \\
 &= \sum_{x, x'} P(X_{n-1} = x | Y_1, Y_2, \dots, Y_{n-1}) P(Y_n | X_{n-1} = x, X_n = x') P(X_n = x') \quad \text{(substitution from c&d)}
 \end{aligned}$$

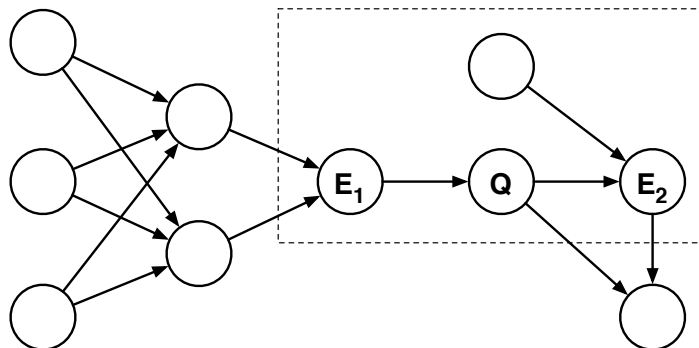
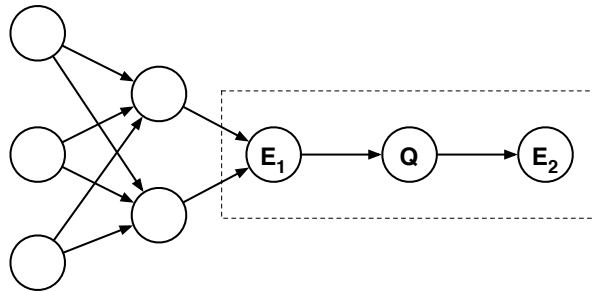
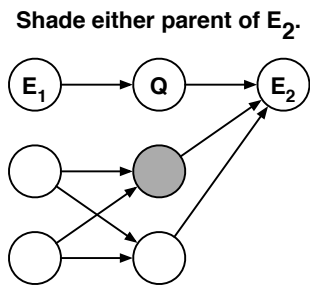
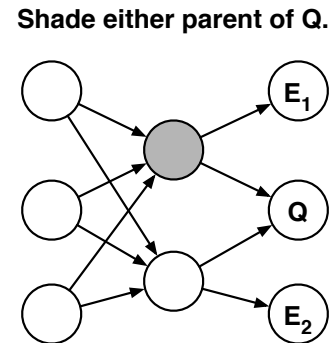
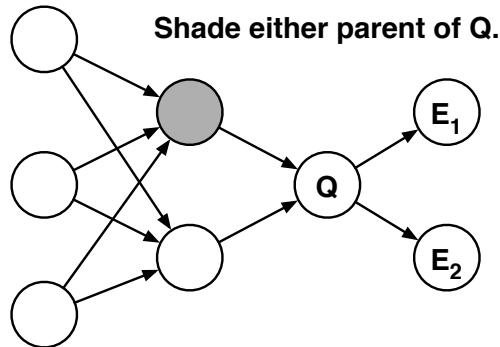
3.4 Node clustering and polytrees (5 pts)

In the figure below, *circle* the DAGs that are polytrees. In the other DAGs, shade **two** nodes that could be *clustered* so that the resulting DAG is a polytree.



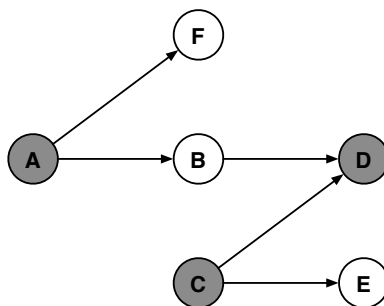
3.5 Cutsets and polytrees (5 pts)

For each of the five loopy belief networks shown below, consider how to compute the posterior probability $P(Q|E_1, E_2)$. If the inference can be performed by running the polytree algorithm on a subgraph, enclose this subgraph by a dotted line as shown on the previous page. On the other hand, if the inference cannot be performed in this way, shade **one** node in the belief network that can be instantiated to induce a polytree by the method of cutset conditioning.



For the last belief network, note that $P(Q|E_1, E_2)$ can be computed in terms of the CPTs for $P(Q|E_1)$, $P(\text{pa}(E_2))$, and $P(E_2|Q, \text{pa}(E_2))$.

3.6 Even more inference (5 pts)



(a) **Markov blanket** (3 pts)

$$\begin{aligned}
 P(B|A, C, D) &= \frac{P(D|B, A, C) P(B|A, C)}{P(D|A, C)} \quad (\text{conditionalized Bayes rule}) \\
 &= \frac{P(D|B, C) P(B|A)}{P(D|A, C)} \quad (\text{conditional independence}) \\
 &= \frac{P(D|B, C) P(B|A)}{\sum_b P(B=b, D|A, C)} \quad (\text{marginalization}) \\
 &= \frac{P(D|B, C) P(B|A)}{\sum_b P(B=b|A, C) P(D|A, B=b, C)} \quad (\text{product rule}) \\
 &= \frac{P(B|A) P(D|B, C)}{\sum_b P(B=b|A) P(D|B=b, C)} \quad (\text{conditional independence})
 \end{aligned}$$

(b) **Conditional independence** (1 pt)

The nodes A, C, and D form the Markov blanket of B.
Hence $P(B|A, C, D, E, F) = P(B|A, C, D)$.
The answer is the same as part (a).

(c) **More conditional independence** (1 pts)

$$\begin{aligned}
 P(B, E, F|A, C, D) &= P(B|A, C, D) P(E|A, B, C, D) P(F|A, B, C, D, E) \quad (\text{product rule}) \\
 &= P(B|A, C, D) P(E|C) P(F|A) \quad (\text{conditional independence})
 \end{aligned}$$

The first term on the right side is given by the answer in part (a).

3.7 Inference in a chain (11 pts) (250A only)

(a) **Powers of transition matrix** (3 pts)

We will prove the claim by induction. First we note that the claim is trivial at node $t = 1$. Next, assume that the claim is true at node $t - 1$. Then at node t :

$$\begin{aligned}
 P(X_{t+1} = j | X_1 = i) &= \sum_k P(X_{t+1} = j, X_t = k | X_1 = i) && \boxed{\text{marginalization}} \\
 &= \sum_k P(X_{t+1} = j | X_t = k, X_1 = i) P(X_t = k | X_1 = i) && \boxed{\text{product rule}} \\
 &= \sum_k P(X_{t+1} = j | X_t = k) P(X_t = k | X_1 = i) && \boxed{\text{conditional independence}} \\
 &= \sum_k P(X_t = k | X_1 = i) P(X_{t+1} = j | X_t = k) \\
 &= \sum_k [A^{t-1}]_{ik} A_{kj} && \boxed{\text{induction}} \\
 &= [A^t]_{ij}.
 \end{aligned}$$

(b) **Linear-time complexity** (2 pts)

Note that computing any entry in the j th column of A^t requires only the j th column of A^{t-1} :

$$[A^t]_{ij} = \sum_k A_{ik} [A^{t-1}]_{kj}.$$

Inductively, we need only compute the j th column of each power of A . Each column computation takes $O(m^2)$ time; thus the total time to compute $[A^t]_{ij}$ is $O(m^2 t)$.

(c) **Log-time complexity** (2 pts)

Alternatively, we can compute A, A^2, A^4, \dots, A^k , where $k = 2^{\lfloor \log_2 t \rfloor}$ is the largest power of 2 less than t , by repeated squaring in time $O(m^3 \log_2 t)$. A^t can then be computed by multiplying the matrices corresponding to ones in the binary expansion of t . (For example, $A^{11} = A^8 A^2 A$.) This process takes time $O(m^3 \log_2 t)$, since matrix multiplication involves $O(m^3)$ computations. Hence the total computation is $O(m^3 \log_2 t)$.

(d) **Sparse transition matrices** (1 pt)

If the transition matrix A_{ij} is sparse, then the sum in (b) can be restricted to the $s \ll m$ non-zero elements of each row. This reduces the computation to $O(smt)$.

(e) **Posterior probability** (3 pts)

$$\begin{aligned}P(X_1 = i | X_{T+1} = j) &= \frac{P(X_{T+1} = j | X_1 = i) P(X_1 = i)}{P(X_{T+1} = j)} && \boxed{\text{Bayes rule}} \\&= \frac{P(X_{T+1} = j | X_1 = i) P(X_1 = i)}{\sum_k P(X_1 = k, X_{T+1} = j)} && \boxed{\text{marginalization}} \\&= \frac{P(X_{T+1} = j | X_1 = i) P(X_1 = i)}{\sum_k P(X_1 = k) P(X_{T+1} = j | X_1 = k)} && \boxed{\text{product rule}} \\&= \frac{[A^T]_{ij} P(X_1 = i)}{\sum_k [A^T]_{kj} P(X_1 = k)} && \boxed{\text{substitution}}\end{aligned}$$
