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# **Preface**

# I. Project Description

This project builds an Automatic Car System using FreeRTOS. This automatic car is capable of multitasking such as detecting obstacles, tracking lines and recieving WiFi and Bluetooth command concurrently. This project also invloves building a control panel for the main controller to override the functions with command send through WiFi.

# II. Project Goals

- 1. Build an automatic car which is capable of multitasking such as detecting obstacles, tracking lines and recieving WiFi and Bluetooth command concurrently.
- 2. Build a control panel for the main controller to override the functions with command send through WiFi.

# **Background Knowledge**

## I. Introduction to RTOS

#### **Overview of RTOS**

• Real-Time Operating System (RTOS): A specialized operating system designed for embedded systems requiring deterministic, timely responses to events. It's often lightweight, making it suitable for devices with constraints on memory, compute, and power.

#### **Multitasking and Concurrency**

- **Multitasking**: RTOSes, like FreeRTOS, can execute multiple tasks concurrently by rapidly switching between them, although a single-core processor can only run one task at a time.
- Tasks vs. Threads: RTOSes often refer to tasks instead of threads, as they usually don't have virtual memory.
- **Benefits of Multitasking**: Partitioning applications into manageable tasks allows for easier testing, team collaboration, and code reuse. The RTOS kernel manages timing and sequencing.

#### Scheduling

- **Scheduler**: Part of the RTOS kernel that decides which task to execute. Tasks can be paused and resumed, allowing for prioritization based on real-time requirements.
- Task States: Tasks can yield, sleep, or block, allowing the kernel to manage and select the appropriate task to run according to the scheduling algorithm.

## **Real-Time Scheduling and Priority**

- **Scheduling Policy**: Real-time systems use priority-based scheduling, ensuring the highest-priority task gets processing time to meet deadlines.
- Task Prioritization: Tasks are assigned priorities based on deadlines and the consequences of missing them. For example, a control task with strict timing is prioritized over a key handler task with a more lenient timing requirement.

#### **Example Scenario**

- In a system with a **keypad and control function**, a key handler task gives feedback on user input, while a control task performs time-sensitive operations.
- The control task, needing timely execution, is assigned higher priority. The RTOS switches between tasks
  to ensure deadlines are met, pausing lower-priority tasks like the idle task or the key handler if the control
  task is ready.

This setup ensures that the system can respond to real-world events within specified time constraints, fulfilling the requirements of a real-time embedded system.

#### II. Introduction to FreeRTOS

FreeRTOS is a lightweight Real-Time Operating System (RTOS) designed to run on microcontrollers, which are compact and resource-limited processors typically used in embedded applications. These applications often involve specific, dedicated tasks, making a full-featured RTOS unnecessary or impractical. FreeRTOS provides essential real-time scheduling, inter-task communication, timing, and synchronization primitives, making it a real-time kernel rather than a full RTOS. Additional features, like networking or a command console interface, can be added as needed.

# **III. FreeRTOS Coding Tips**

## xTaskCreate() & vTaskStartScheduler() 建立任務

- 1. Each task has:
  - a. Register set
  - b. Stack area
  - c. Priority
- 2. 建立任務,沒建立task會執行idle task
  - a. 把AppTaskCreate\_Handle當作該Task控制點

3. 建立子任務

```
30  //AppTaskCreate_Handle
31  static TaskHandle t AppTaskCreate_Handle = NULL;
32  static TaskHandle_t LED_Task_Handle = NULL;
33  static TaskHandle_t LED1_Task_Handle = NULL;
34  static TaskHandle_t LED2_Task_Handle = NULL;
```

也可用 xTaskHandle , xTaskHandle 跟 TaskHandle\_t 一樣

```
106 static void AppTaskCreate(void)
107 □ {
     BaseType_t xReturn = pdPASS;/* 定義一個建立資訊返回值,預設為pdPASS */
109
     taskENTER CRITICAL();
                           //進入臨界區
110
     /* 建立LED Task任務 */
112
     113
114
115
116
117
                     (TaskHandle_t* ) & LED_Task_Handle);/* 任務控制塊指標 */
118
119
     if(pdPASS == xReturn)
     printf("LED_Task!\r\n");
/* 建立LED1_Task_Task任務 */
120
121
     122
123
124
125
126
127
128
     if(pdPASS == xReturn)
129
     printf("LED1_Task!\r\n");
/* 建立LED2_Task_Task任務 */
130
     131
132
133
134
135
136
137
     if(pdPASS == xReturn)
138
      printf("LED2_Task!\r\n");
139
     //AppTaskCreate_Handle
140
     wTaskDelete(AppTaskCreate_Handle); //刪除AppTaskCreate任務
141
142
     taskEXIT_CRITICAL();
                           //退出臨界區
144 }
```

#### 4. 撰寫各任務內容

a. sample code

```
153 static void LED_Task(void* parameter)
154 🚍 (
        while (1)
155
           LEDO( ON );
//printf("LED_Task Running, LED1_ON\r\n");
158
          vTaskDelay(5000); /* 延時500個tick */
159
160
161
         LEDO( OFF );
//printf("LED_Task Running, LED1 OFF\r\n");
162
163
164
165
    [ }
           vTaskDelay(5000); /* 延時500個tick */
166
      static void LED1_Task(void* parameter)
        while (1)
169
LED1( ON );
//printf("LED_Task Running,LED1 ON\r\n");
vTaskDelay(1000); /* 延時500個tick */
174
175
176
          LED1( OFF );
//printf("LED_Task Running, LED1_OFF\r\n");
           vTaskDelay(1000); /* 延時500個tick */
```

- b. 任務中通常有:
  - i. 一個無窮迴圈
  - ii. 呼叫OS Service

# taskENTER\_CRITICAL() & taskEXIT\_CRITICAL()

- 1. 在臨界區中執行的程式不會被打擾
- 2.  $taskenter_critical() \Rightarrow portenter_critical() \Rightarrow vPortenterCritical() \Rightarrow portDISABLe_INTERRUPTS()$

### vTaskDelete()

- vTaskDelete(AppTaskCreate\_Handle); //Delete AppTaskCreate
- 2. 不會馬上刪除,需回收時間 (遇return)

### vTaskPrioritySet() & uxTaskPriorityGet()

- 1. vTaskPrioritySet(LED\_Task\_Handle,11)
- 2. Task\_pri=uxTaskPriorityGet(LED\_Task\_Handle); //get priority
- 3. printf("New priority is %ld", Task\_pri);
- 4. 數字越大,優先權越高

#### vTaskSuspend() & vTaskResume()

- 1. suspend task and resume task
- vTaskSuspend(LED\_Task\_Handle);
- 3. vTaskResume(LED\_Task\_Handle);

#### vTaskSuspendAll() & vTaskResumeAll()

1. suspend所有除了自己的Task, OS通常也直接不作用(eg. delay won't work)

## xTaskGetTickCount () & vTaskDelay() & vTaskDelayUntil() & vTaskSuspend( NULL )

1. sample code 1: xTaskGetTickCount () & vTaskDelay()

```
96 void vTaskl( void *pvParameters )
97 🖯 (
          portTickType xDelay, xNextTime;
99
          /st As per most tasks, this task is implemented in an infinite loop. st/
100
          while(1)
L01 🛱
            LEDO_TOGGLE;
L03
            // printf("Taskl is running \r\n");
xDelay = xNextTime - xTaskGetTickCount ();
L05
            xNextTime += ( portTickType ) 1000;
            if( xDelay <= ( portTickType ) 1000 )</pre>
108
              vTaskDelay( xDelay );
110
       }
```

- i. portTickType 是 xTaskGetTickCount () 回傳值的data type
- 2. sample code 2: xTaskGetTickCount () & vTaskDelayUntil()

```
114 void vTask2 ( void *pvParameters )
116
         portTickType xLastWakeTime;
117
118
         const portTickType xFrequency = 100;
          // Initialise the xLastWakeTime variable with the current time.
120
121
122
          xLastWakeTime = xTaskGetTickCount();
123
         /* As per most tasks, this task is implemented in an infinite loop. */
124
         while(1)
125
             //printf("Task2 is running\r\n");
126
             vTaskDelayUntil( &xLastWakeTime, xFrequency );
128
130
131 }
```

- vTaskDelayUntil( &xLastWakeTime, xFrequency );
- ii. 從 xLastWakeTime 開始向後數 xFrequency
- 3. sample code 3: vTaskSuspend( NULL )
  - i. vTaskSuspend( NULL ) 把自己暫停掉

#### uxTaskGetNumberOfTasks()

1. 取得目前存在多少Task

## vTaskList() 取得CPU Usage

1. Table produced

Name *********	State	Priority	Stack	Num
Print	R	4	331	29
Math7	R	0	417	7
Math8	R	0	407	8
QConsB2	R	0	53	14
QProdB5	R	0	52	17
QConsB4	R	0	53	16
SEM1	R	0	50	27
SEM1	R	0	50	28
IDLE	R	0	64	0
Math1	R	0	436	1
Math2	R	0	436	2

- 2. State Types:
  - i. B: Block (delay會進入block)
  - ii. R: Ready
  - iii. S: Suspend
  - iv. D: has been deleted, but memory has not yet been freed
- 3. Stack: 剩下的Stack容量
- 4. Num: 隨機編號
- 5. sample code

#### vTaskGetRunTimeStats()

- vTaskGetRunTimeStats((char \*)&CPU\_RunInfo);
- 2. CPU排程執行此Task的次數
- 3. produced content:

Task	Abs Time	% Time
*******	*****	*****
uIP	12050	<1%
IDLE	587724	24%
QProdB2	2172	<1%
QProdB3	10002	<1%
QProdB5	11504	<1%
QConsB6	11671	<1%
PolSEM1	60033	2%
PolSEM2	59957	2%
IntMath	349246	14%
MuLow	36619	1%
GenO	579715	2.4%

## xQueueHandle & xQueueCreate() & xQueueReceive() & xQueueSend()

- xQueueHandle xQueue;
- 2. xQueue = xQueueCreate(5, sizeof(uint8\_t));
- 3. 這個Queue建立在同步化的機制
- 4. sample code
  - a. 收資料
    - i. 也可用BaseType\_t來收回傳值,0代表收不到資料

```
87 void vTask3( void *pvParameters)
                        uint8_t Count[5];
                         while(1)
  91
  92
  93
94 =
                         if( xQueue != 0 )
                             //put xQueue's data to &Count[0], wait at most portMAX_DELAY
xQueueReceive( xQueue, &Count[0], portMAX_DELAY);
xQueueReceive( xQueue, &Count[1], portMAX_DELAY);
xQueueReceive( xQueue, &Count[2], portMAX_DELAY);
xQueueReceive( xQueue, &Count[3], portMAX_DELAY);
xQueueReceive( xQueue, &Count[4], portMAX_DELAY);
  95
  96
97
98
99
100
 101
 102
                                  printf("Count0 = %d\r\n", Count[0]);
printf("Count1 = %d\r\n", Count[1]);
printf("Count2 = %d\r\n", Count[2]);
printf("Count3 = %d\r\n", Count[3]);
printf("Count4 = %d\r\n", Count[4]);
103
 104
 105
 106
 107
109 }
```

#### b. 塞資料

```
111 void vTask4 ( void *pvParameters )
                 portTickType xLastWakeTime;
                 const portTickType xFrequency = 1000;
uint8_t ucVar[5] = {0};
 115
116
117
118
119
                   // Initialise the xLastWakeTime variable with the current time. xLastWakeTime = xTaskGetTickCount();
                      while(1)
 120
121
122
123
124
125
126
127
128
                      ucVar[1]++;
ucVar[2]++;
                      ucVar[3]++;
ucVar[4]++;
 129
130
131
132
                      if( xQueue != 0 )
                       {
//put &ucVar[0]'s data to xQueue, wait at most 10
xQueueSend( xQueue, ( void * ) &ucVar[0], ( portTickType ) 10 );
xQueueSend( xQueue, ( void * ) &ucVar[1], ( portTickType ) 10 );
xQueueSend( xQueue, ( void * ) &ucVar[2], ( portTickType ) 10 );
xQueueSend( xQueue, ( void * ) &ucVar[3], ( portTickType ) 10 );
xQueueSend( xQueue, ( void * ) &ucVar[4], ( portTickType ) 10 );
 133
 134
135
136
137
                                vTaskDelayUntil( &xLastWakeTime, xFrequency );
 138
139
140
```

# Mutex: xSemaphoreHandle & xSemaphoreCreateMutex() & xSemaphoreTake() & xSemaphoreGive()

- 1. xSemaphoreHandle xSemaphore;
- 2. xSemaphore = xSemaphoreCreateMutex();
- 3. sample code:
  - a. printf is mutual exclusive
  - b. vTask3 跟 vTask4 會輪流取得資源
  - C. 可取得 xSemaphoreTake() 回傳值,來確認能不能取得資源

```
87 void vTask3( void *pvParameters )
88 🖵 {
         portTickType xLastWakeTime;
89
90
         const portTickType xFrequency = 1000;
       // Initialise the xLastWakeTime variable with the current time.
92
93
         xLastWakeTime = xTaskGetTickCount();
94
95
          while(1)
96
           // from xSemaphore take data for at most portMAX DELAY
97
98
           xSemaphoreTake( xSemaphore, portMAX DELAY);
                   == Critical Section Start ==
99
100
           //this printf can be seen as a shared resource
          printf("Task3 is running\r\n");
//===== Critical Section End ==
101
102
103
           // return resource
104
           xSemaphoreGive( xSemaphore );
105
           vTaskDelayUntil( &xLastWakeTime, xFrequency );
106
106 |-
109 void vTask4( void *pvParameters )
 110 □ {
 111
          portTickType xLastWakeTime;
 112
          const portTickType xFrequency = 1000;
 113
 114
         // Initialise the x {\tt LastWakeTime} variable with the current time.
 115
           xLastWakeTime = xTaskGetTickCount();
 116
 117
118 =
          while(1)
 119
            // from xSemaphore take data for at most portMAX_DELAY
 120
121
            // finite waiting cus the tasks are arranged in waiting Queue
            xSemaphoreTake( xSemaphore, portMAX DELAY);
           //===== Critical Section Start =====
//this printf can be seen as a shared resource
 122
 123
           printf("Task4 is running\r\n");
 124
 125
126
           //===== Critical :
// return resource
                    == Critical Section End =
 127
            xSemaphoreGive( xSemaphore );
 128
             vTaskDelayUntil( &xLastWakeTime, xFrequency );
 129
```

# Binary Semaphore: xSemaphoreHandle & vSemaphoreCreateBinary() & xSemaphoreTake() & xSemaphoreGive()

- 1. xSemaphoreHandle xSemaphore;
- 2. vSemaphoreCreateBinary( xSemaphore );
- 3. sample code:
  - a. vTask3 要等 vTask4 release xsemaphore 才可再存取 ⇒ vTask3 要等 vTask4 執行完才可再執行

```
86 void vTask3( void *pvParameters )
87 🗏 {
88
        uint8 t Count = 100;
89
90
91
92
            xSemaphoreTake( xSemaphore, portMAX_DELAY); // Take resource
93
           printf("Count = %d\r\n", Count);
95
96 }
98 void vTask4 ( void *pvParameters )
         portTickType xLastWakeTime;
101
         const portTickType xFrequency = 1000;
102
103
        // Initialise the xLastWakeTime variable with the current time.
104
          xLastWakeTime = xTaskGetTickCount();
105
106
         while(1)
107
           printf("Task4 is running\r\n");
108
     xSemaphoreGive( xSemaphore ); // Release resource
109
110
           vTaskDelayUntil( &xLastWakeTime, xFrequency );
111
112 }
```

# Counting Semaphore: xSemaphoreHandle & xSemaphoreCreateCounting() & xSemaphoreTake() & xSemaphoreGive()

- 1. xSemaphoreHandle xSemaphore;
- 2. xSemaphore = xSemaphoreCreateCounting(5,2);
- 3. first num: maximum, second num: initial
- 4. delay少的Task可能可以一次搶得多個資源

# **Development Tools**

# I. Development Platform

- 1. Arduino IDE
- 2. Qt Creator
- 3. Visual Studio Code

# II. Project Package

- 1. ESP32 Wrover Module (with WiFi and Bluetooth functions)
- 2. Line Tracking Module (for line tracking)
- 3. HC-SR04 Ultrasonic Module (for obstacle avoidance)
- 4. Dot Matrix Module
- 5. RGB LED (on board module)
- 6. Buzzer (on board module)

# III. Technologies Used

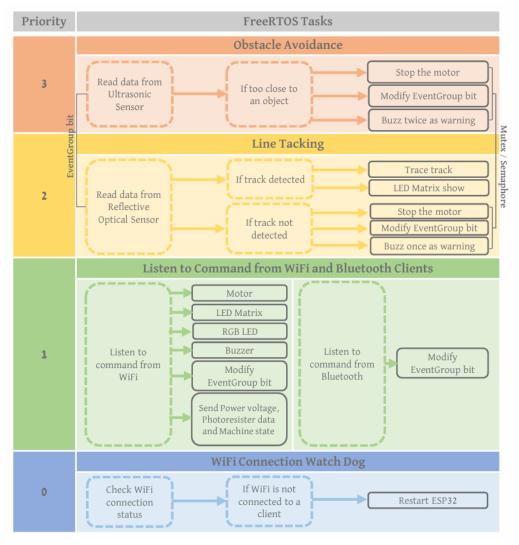
- 1. MCU (Microcontroller Unit)
- 2. WiFi
- 3. Bluetooth
- 4. FreeRTOS

# IV. Programming Language Used

- 1. C Programming Language
- 2. Python Programming Language

# **Project Scenario**

The automatic car has three STATE: *IN MOTION : AUTO, IN MOTION : TRACK, MOTIONLESS, OVERRIDE 1* and *OVERRIDE 2*, which are mainly controlled by Event Group. In these STATEs, the automatic car will perform multiple tasks that have different priorities and are sychronized through Mutex and Event Group.



**Priority Chart** 

### I. IN MOTION: AUTO

In this state, Task 1-4 are all working simultaneously. Obstacle Avoidance and Line Tracking tasks are triggered by setting the first EventGroup bit to 1 after recieving command from Control Panel. Line Tracking task will start after Obstacle Avoidance has been executed one time to avoid the automatic car bumping into an obstacle in sight due to the later execution of Obstacle Avoidance task. The late execution of Line tracking task is controlled by the seventh bit of EventGroup value (Track Begin, TB). When the Obstacle Avoidance task has been executed once, the TB bit will be set to 1. In other word, the Line Tracking task will not be executed until the Obstacle Avoidance task has been executed once.

### II. IN MOTION: TRACK

In this state, only Task 2-4 will be working. Obstacle Avoidance will be blocked right after entering the while loop as the first EventGroup bit has been cleared. This state is triggered by setting both the second and the TB bit of EventGroup to 1.

## **III. MOTIONLESS**

The cause of Motionless state can be due to several reasons: 1) Obstacle detected, 2) No track found 3) Bluetooth stop instruction, and 4) WiFi override instruction. The cause of Motionless state will be record with the third to sixth bit of the EventGroup.

#### IV. OVERRIDE 1

In this state, only RGB LED and LED matrix functionalities will be override, which means only these functionalities can be controlled through control panel. The rest of the functionalities will be contolled through its original state.

#### V. OVERRIDE 2

In this state, all the functionalities will be override by the control panel. Before entering this state, the automatic car will be forced to a stop, which correspond to the fourth reason of Motionless state: WiFi override instruction.

# **Main Progam**

# I. ESP32 Program

#### **Module Setup**

#### **FreeRTOS Tasks Creation**

#### Main Loop: Listen to WiFi command

```
if (CmdArray[0] == CMD_OVERRIDE) { //Override Mode

xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit

xEventGroupSetBits(xEventGroup, E6_BIT); // Set E6 bit

227

} //Clears the command array and parameter array

memset(CmdArray, 0, sizeof(CmdArray));

memset(paramters, 0, sizeof(paramters));

}

Emotion_Show(emotion_task_mode);//Led matrix display function

WS2812_Show(ws2812_task_mode);//Car color lights display function

Car_Select(carFlag);//ESP32_Car_mode_selection function
```

#### Main Loop: Listen to Bluetooth command

```
// Bluetooth Cmd
if (SerialBT.available()){
    message=SerialBT.read();
    Serial.write(message);

if(message=turnON){
    W52812_Show(4); //Car color lights display function
    Serial.println(F(":LED ON"));
    SerialBT.println(F("LED ON"));
}
else if(message=turnOFF){
    W52812_Show(0); //Turn LED Off
    Serial.println(F(":LED OFF"));
    SerialBT.println(F("LED OFF"));
    SerialBT.println(F("LED OFF"));
}
else if(message=restart){
    xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
    xEventGroupSetBits(xEventGroup, E1_BIT); // Set E1 bit
    Serial.println(F(":RESTART"));
    SerialBT.println(F("RESTART"));
}
else if(message=stop){
    xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
    xEventGroupSetBits(xEventGroup, ALL_BIT); // Clear all bit
    xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
    xEventGroupClearBits(xEventGroup, E5_BIT); // Set E5 bit
    Serial.println(F(":STOP"));
    SerialBT.println(F(":STOP"));
    SerialBT.println(F("STOP"));
}
```

#### Main Loop: Restart EPS32 if WiFi Client Not Connected

#### **Obstacle Avoidance Task**

```
case 7: //111
Emotion_SetWode(@);
vTaskbelay(100);

Motor_Move(@, 0, 0, 0, 0);
xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
xEventGroupSetBits(xEventGroup, E4_BIT); // Set E4 bit
Buzzer_Alert(1, 1);
Emotion_SetWode(0);
Emotion_SetWode(0);
Emotion_SetWode(0);
Emotion_SetWode(0);

Emotion_SetWode(0);
Emotion_SetWode(0);
Emotion_SetWode(0);
Emotion_SetWode(0);
Emotion_SetWode(0);
Emotion_SetWode(0);
Emotion_SetWode(0);

Emotion_SetWode(0);
Emotion_SetWode(0);

Motor_Move(-SPEED_LV3, -SPEED_LV4, SPEED_LV4); //Turn Left

break;

case 4: //100
case 6: //110
Emotion_SetWode(5);
Motor_Move(-SPEED_LV4, - SPEED_LV3, -SPEED_LV3); //Turn Right
break;

default:
break;

}

Emotion_Show(emotion_task_mode); //Led matrix display function
}

xSemaphoreGive(mutex);
}else{
continue;
}

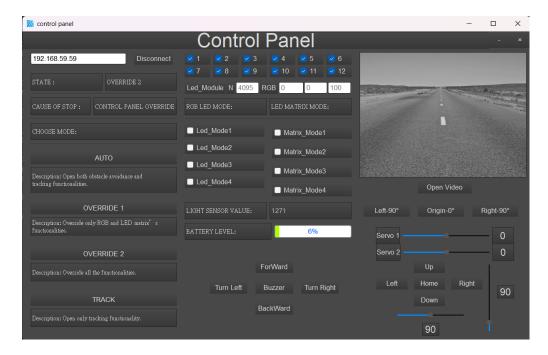
330
}

}
```

# WiFi Connection Watch Dog

# II. Contol Panel (run on PC)

## **Control panel in OVERRIDE 2 STATE**

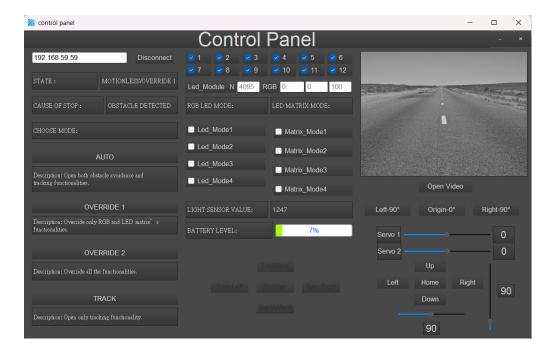


## **Control panel in OVERRIDE 1 STATE**

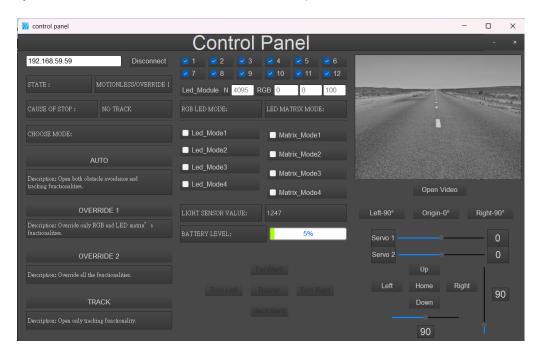


Control panel in MOTIONLESS/OVERRIDE 1 STATE, Obstacle Detected

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#### Control panel in MOTIONLESS/OVERRIDE 1 STATE, No Track Detected



# **Problem Discussion**

#### I. Automatic Car Motor Control Race Condition

If we did not protect the automatic car motor with Mutex or Semaphore, race condition can occur. A race condition is a problematic situation that arises when a device or system tries to perform multiple operations simultaneously, but the correct outcome depends on a specific order of execution. In programming, race conditions occur when two processes or threads try to access the same resource at the same time, potentially causing system errors or unexpected behavior. In our case, race condition occurs when both Line Tracking task and Obstacle Avoidance task try to control the wheel motor.

#### Without Mutex or Semaphore

In the following code, because the marked code block is not protected with Mutex or Semaphore, the automatic car might not be able to stop in time when detecting obstacles.

#### **Obstacle Avoidance Task**

```
void start_Line_Track(void *pvParameters){

Serial.println("start_Line_Track Created");

// define a variable which holds the state of events

const EventBits_t xBitsToWaitFor = (E1_BIT|E2_BIT);

EventBits_t xEventGroupValue;

while(1){
```

From the following output, we can discover that when the automatic car detect obstacles, it does not stop directly due to two reasons:

- 1. The Line Tracking task can still be running concurrently on another core (ESP32 has two cores)
- 2. Context switch can occur after obstacle is detected, or before clearing up the EventGroup bits.

In other words, by stopping the car and clearing up the EventGroup bits sequentially is not enough to stop the car upon detecting the obstacles.

```
17:42:49.912 -> Enter start Line Track

17:42:49.956 -> Enter start_Obst_Avoid

17:42:50.034 -> Enter start_Line_Track

17:42:50.122 -> Enter start_Line_Track

17:42:50.228 -> Enter start_Line_Track

17:42:50.362 -> Enter start_Line_Track

17:42:50.472 -> Enter start_Line_Track

17:42:50.583 -> Enter start_Line_Track

17:42:50.676 -> Enter start_Line_Track

17:42:50.676 -> Enter start_Line_Track

17:42:50.768 -> Enter start_Line_Track

17:42:50.768 -> Enter start_Line_Track

17:42:51.712 -> CMD_POWER#

17:42:51.172 -> CMD_POWER#
```

# Use vTaskSuspend() and vTaskResume() to Stop Line Tracking Task Temporarily

In the following code, using vTaskSuspend() and vTaskResume() to suspend the line tracking task until the EventGroup bits are cleared can potentially stop the automatic car in time. Yet, problem can still occure when the Line Tracking task is resumed.

```
void(void *pvParameters){
Serial.println("start_Obst_Avoid Created");
// define a variable which holds the state of
const EventBits_t xBitsToWaitFor = (E1_BIT);
EventBits t xEventGroupValue;
   xEventGroupValue = xEventGroupWaitBits(xEventGroup,
                                                              xBitsToWaitFor,
                                                              pdFALSE.
                                                             pdTRUE,
  if((xEventGroupValue & E1 BIT) != 0){
    Serial.println("Enter start_Obst_Avoid");
xEventGroupSetBits(xEventGroup, TB_BIT); // Set TB bit
    float distance = get_distance();
if (distance <= OBSTACLE DISTANCE)</pre>
       if (line_Track_Handle != NULL){
          vTaskSuspend(line_Track_Handle);
       Motor_Move(0, 0, 0, 0); //Stop the car to judge the situation
Buzzer_Alert(2, 1);
Emotion_SetMode(0);
        Emotion_Seconde(0);
Emotion_Show(emotion_task_mode);
xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
xEventGroupSetBits(xEventGroup, E3_BIT); // Set E3 bit
Serial.println("STOP");
    vTaskResume(line_Track_Handle);
```

From the following output, we can discover that the automatic car stops upon detecting an obstacle and the EventGroup bits are successfully cleared. However, when the Line Tracking task is resumed, The Line Tracking task will not be executed from the beginning of the task where the EventGroup bits will be checked. Instead, the Line Tracking task will start from where it had left before being suspended. Accordingly, the automatic car might move unexpectedly after the Line Tracking task has resumed.

```
17:59:01.037 -> Enter start_Obst_Avoid
17:59:01.068 -> Enter start_Line_Track
17:59:01.206 -> Enter start_Line_Track
17:59:01.299 -> Enter start_Line_Track
17:59:01.378 -> Enter start_Line_Track
17:59:01.501 -> Enter start_Dost_Avoid
17:59:02.390 -> STOP
17:59:04.020 -> CMD_POWER#
17:59:07.013 -> CMD_POWER#
```

# With Mutex or Semaphore

In the following code, mutex is used to protect the code block in which wheel motors have been controlled in both the Line Tracking and Obstacle Avoidance tasks.

**Obstacle Avoidance Task** 

```
t Avoid(void *pvParameters){
         Serial.println("start_Obst_Avoid Created");
         // define a variable which holds the state of events
         const EventBits t xBitsToWaitFor = (E1 BIT):
         EventBits t xEventGroupValue;
           xEventGroupValue = xEventGroupWaitBits(xEventGroup,
                                                        xBitsToWaitFor,
                                                        pdTRUE,
                                                        portMAX_DELAY
           if((xEventGroupValue & E1_BIT) != 0){
                      coupSetBits(xEventGroup, TB_BIT); // Set TB bit
             float distance = get_distance();
             if (distance <= OBSTACLE DISTANCE)
            if (distance <= OBSTACLE DISTANCE);

xSemaphoreTake(mutex, portMAX_DELAY);

(0 0 0 0 0). //Stop the car to judge
               Motor_Move(0, 0, 0, 0);
Buzzer_Alert(2, 1);
               Emotion_SetMode(0);
               Emotion Show(emotion_task_mode);
               xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
               xEventGroupSetBits(xEventGroup, E3_BIT); // Set E3 bit
Serial.println("STOP"):
            xSemaphoreGive(mutex);
456
             }
              vTaskDelay(450);
```

From the following output, we can discover that before stopping, the Line Tracking task is still running. Yet, when it tries to gain access to the wheel motor, it is blocked outside of the code block as the key has already been taken by the Obstacle Avoidance task. Afterward, when the key is retruned, the Line Tracking Task is allowed to execute the protected code block. However, the EventGroup bits have already been cleared. Thus, the Line Tracking task is still unable to control the wheel motor. As a result, the automatic car can be stopped in time when an obstacle is detected and avoid unexpected movement caused by the Line Tracking task.

```
18:34:25.734 -> Enter start_Obst_Avoid

18:34:25.734 -> Enter start_Line_Track

18:34:25.870 -> Enter start_Line_Track

18:34:25.964 -> Enter start_Line_Track

18:34:26.051 -> Enter start_Line_Track

18:34:26.181 -> Enter start_Line_Track

18:34:26.181 -> Enter start_Line_Track

18:34:26.261 -> Enter start_Line_Track

18:34:27.201 -> STOP
```

# II. Late Start of Line Tracking

Another problem we face in this project is the late start problem. When the automatic car enter AUTO STATE, we want to garantee that the Obstacle Avoidance task starts first and the Line Tracking task starts after the Obstacle Avoidance task has finished its first round. By doing so, we can prevent the automatic car from bumping into an obstacle in sight because the Line Tracking task has already been working while the Obstacle Avoidance task has not. The following approaches are some potential solutions.

### **Using Binary Semaphore**

In FreeRTOS, a binary semaphore is a synchronization mechanism used to manage resource access between tasks or signal task completion. It operates with two states: "taken" (0) or "available" (1). A task can "take" the semaphore if it's available, indicating exclusive access to a shared resource, and other tasks must wait until it's released. Binary semaphores are often used for simple signaling, such as notifying a task that an event has occurred.

By letting the Obstacle Avoidance task to release the resource after finish executing the first round, and letting the Line Tracking task to take the resource in the beginning of the task, the Line Tracking task is forced to wait until the Obstacle Avoidance task releases the resource. The late start of Line Tracking task is then fulfilled.

The following is the code.

#### **Obstacle Avoidance Task**

```
| start_Line_Track(void *pvParameters){
rial.println("start Line Track Created
xSemaphoreTake(binary_sem, portMAX_DELAY);
const EventBits_t xBitsToWaitFor = (E1_BIT|E2_BIT);
EventBits_t xEventGroupValue;
  if (xSemaphoreTake(mutex, 1/portTICK_PERIOD_MS) == pdTRUE){
     xEventGroupValue = xEventGroupWaitBits(xEventGroup,
                                                xBitsToWaitFor.
                                                pdFALSE.
                                                pdFALSE,
                                                portMAX DELAY
     Serial.println("Enter start_Line_Track");
Track_Read();
     switch (sensorValue[3])
       Emotion_SetMode(3);
Motor_Move(SPEED_LV1, SPEED_LV1, SPEED_LV1, SPEED_LV1); //Move Forward
      case 0: //000
case 7: //111
Emotion_SetMode(6);
vTaskDelay(100);
        break;
case 1: //001
case 3: //011
         Emotion_SetMode(4);
         Motor_Move(-SPEED_LV3, -SPEED_LV3, SPEED_LV4); //Turn Left
```

```
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Emotion_Show(emotion_task_mode);//Led matrix display function
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xSemaphoreGive(mutex);
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}else{
continue;
}
388
}
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}
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case 4: //100
case 6: //110
Emotion_SetMode(5);
Motor_Move(SPEED_LV4, - SPEED_LV3, -SPEED_LV3);//Turn Right
break;

default:
break;
}

Emotion_Show(emotion_task_mode);//Led matrix display function

xSemaphoreGive(mutex);
}else{
continue;
}

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}

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}

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}
```

However, when using binary semaphore, there are some undesired effects for our scenario. After the first round of execution, the Line Tracking task is still affected by this binary semaphore, which causes the Line Tracking task to wait for the Obstacle Avoidance task releasing its resource from time to time. Yet, after the first round of execution, we want the two tasks to run non-blockingly. Thus, we change the code to the following.

### **Using Two EventGroup Bits**

Override Stop	Bluetooth Stop	No Track Stop	Obstacle Stop	Line Tracking (T)	Obstacle
(OS)	(BS)	(TS)	(OS)		Aoidance (O)
1	1	1	1	1	1

Here, we use the first two bits of the EventGoup to trigger the execution of the Obstacle avoidance task and the Line Tracking task. When the automatic car enter AUTO STATE, (O) will first be set to 1 while (T) will be set to 1 after the Obstacle Avoidance task had finished the first round of execution. By doing so, the Line Tracking task will be triggered later than the Obstacle Avoidance task.

The following is the code.

#### **Obstacle Avoidance Task**

```
ost_Avoid(void *pvParameters){
           Serial.println("start_Obst_Avoid Created");
           const EventBits_t xBitsToWaitFor = (E1_BIT);
           EventBits_t xEventGroupValue;
             xEventGroupValue = xEventGroupWaitBits(xEventGroup,
                                                                     xBitsToWaitFor,
                                                                     pdFALSE,
                                                                     pdTRUE,
                                                                     portMAX DELAY
             if((xEventGroupValue & E1_BIT) != 0){
            Serial.println("Enter start Obst Avoid");
xEventGroupSetBits(xEventGroup, E2_BIT); // Set E2 bit
441
                 float distance = get_distance();
                if (distance <= OBSTACLE_DISTANCE){</pre>
                  xSemaphoreTake(mutex, portMAX_DELAY);
Motor_Move(0, 0, 0, 0); //Stop the
Buzzer_Alert(2, 1);
Emotion_SetMode(0);
Emotion_Show(emotion_task_mode);
                                                      //Stop the car to judge the situation
                   Serial.println("STOP");
                   xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
xEventGroupSetBits(xEventGroup, E3_BIT); // Set E3 bit
xSemaphoreGive(mutex);
                 xSemaphoreGive(binary_sem, portMAX_DELAY);
vTaskDelay(450);
```

```
Serial.println("start_Line_Track Created");
xSemaphoreTake(binary_sem, portMAX_DELAY);
const EventBits_t xBitsToWaitFor = (E2_BIT);
   entBits_t xEventGroupValue;
   if (xSemaphoreTake(mutex, 1/portTICK_PERIOD_MS) == pdTRUE){
     xEventGroupValue = xEventGroupWaitBits(xEventGroup)
                                                   xBitsToWaitFor.
                                                   pdFALSE,
                                                   pdFALSE.
                                                    portMAX DELAY
    if((xEventGroupValue & (E2_BIT)) != 0){
       Serial.println("Enter start_Line_Track");
Track_Read();
        switch (sensorValue[3])
           ase 5: //10.1

Emotion_SetMode(3);

Motor_Move(SPEED_LV1, SPEED_LV1, SPEED_LV1); //Move Forward
          case 0:
            Emotion_SetMode(6);
vTaskDelay(100);
            Motor_Move(0, 0, 0, 0);
            xEventGroupClearBits(xEventGroup, ALL_BIT); // Clear all bit
xEventGroupSetBits(xEventGroup, E4_BIT); // Set E4 bit
            Buzzer_Alert(1, 1);
                tion_Show(emotion_task_mode);
```

However, when using only two bits of the EventGoup to control the sequence of execution, some undesired side effects will occur. For instance, when the automatic car stops due to no track detected in the Line Tracking task (marked with blue in the above code), both (O) and (T) will be cleared to stop the two tasks from entering the next round. Nonetheless, the Obstacle Avoidance task can still be executing its current round, which include setting (T) to 1. This will cause the Line Tracking task starts working again, which isn't what we want. Therefore, we change the code to the following.

# **Using Three EventGroup Bits**

Track Begin (TB)	Override Stop (OS)	Bluetooth Stop (BS)	No Track Stop (TS)	Obstacle Stop (OS)	Line Tracking (T)	Obstacle Aoidance (O)
1	1	1	1	1	1	1

By adding another Track Begin bit (TB), we can avoid the aforementioned problems to occur. The Line Tracking task will be triggered only when either (T) and (TB) or (O) and (TB) are set to 1. When the automatic car enter

AUTO STATE, (O) will first be set to 1 while (TB) will be set to 1 after the Obstacle Avoidance task had finished the first round of execution.

The following is the code.

#### **Obstacle Avoidance Task**

```
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```

When the automatic car stops due to no track detected in the Line Tracking task, both (O) and (T) will be cleared to stop the two tasks from entering the next round. Although the Obstacle Avoidance task can still be executing its current round, which include setting (TB) to 1, it is not enough to cause the Line Tracking task to enter the next round as both (O) and (TB) need to be set. As for the TRACK STATE in which only Line Tracking task is working, (T) and (TB) should be set together at once so that the Line Tracking task can be triggered directly. At this point, all the requirements are fulfilled.