# **OS Lab6**

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# 实验步骤

## 环境搭建



### 管理空闲内存空间

#### 环境变量

```
#define P_PAGE 0x1000
#define parent(x) (((x + 1) >> 1) - 1)
#define lson(x) (((x + 1) << 1) - 1)
#define rson(x) ((x + 1) << 1)

struct buddy buddy;
unsigned ins_bitmap[2 * P_PAGE - 1];
unsigned long page_offset = 0x80000000ul;

extern unsigned long *_end;</pre>
```

P\_PAGE为总页数, 16M/4K = 0x1000

parent、Ison、rson分别为计算父亲、左儿子和右儿子的标号

ins\_bitmap是定义在程序字段的bitmap, buddy.bitmap指向这个数组的开头

page\_offset是决定alloc\_pages返回的地址的偏移量,起始置为0x80000000,alloc\_pages返回物理地址,当paging\_init结束后置page\_offset为0xffffffe000000000,alloc\_pages返回虚拟地址。

\_end为从lds中引入的标签

### init\_buddy\_system

```
void init_buddy_system(void) {
   buddy.size = P_PAGE;
   buddy.bitmap = ins_bitmap;
   unsigned long x, y, i;
   for (x = buddy.size, y = 1, i = 0; x > 0; x >>= 1, i += y, y <<= 1) {
      int j;
      for (j = 0; j < y; ++j) {
         ins_bitmap[i + j] = x;
      }
   }
   alloc_pages((((unsigned long)&_end - 0x80000000ul) >> 12));
}
```

### alloc\_pages

```
void *alloc_pages(int n) {
   int i;
   for (i = 0; ; ++i) {
      if ((1 << i) >= n) {
            n = (1 << i);
            break;
      }
}</pre>
```

```
if (ins_bitmap[0] < n) {</pre>
       return 0;
    }
    i = 0;
    int full_size = buddy.size;
    unsigned t bitmap;
    while (i < 2 * P_PAGE - 1) {
        if (ins_bitmap[i] == n && ins_bitmap[i] == full_size) {
            t_bitmap = ins_bitmap[i];
            ins_bitmap[i] = 0;
            int ii = i;
            while (ii > 0) {
                ii = parent(ii);
                ins_bitmap[ii] = ins_bitmap[rson(ii)] > ins_bitmap[lson(ii)] ?
                         ins_bitmap[rson(ii)] : ins_bitmap[lson(ii)];
            }
            break;
        } else {
            if (ins_bitmap[lson(i)] >= n) {
                full_size >>= 1;
                i = lson(i);
            } else if (ins_bitmap[rson(i)] >= n) {
                full size >>= 1;
                i = rson(i);
            }
        }
    }
    unsigned long addr = (((i + 1) * t_bitmap - buddy.size) << 12) +</pre>
page_offset;
    if (addr >= 0xffffffe000000000) {
        puts("[S] Buddy allocate addr: ");
        puti64(addr);
        puts("\n");
    }
    return addr;
}
```

注意使用page\_offset来控制返回的地址是物理地址还是虚拟地址。

### free\_pages

```
void free_pages(void* x) {
  unsigned long index = x;
  index -= page_offset;
  index >>= 12;
  index += buddy.size;
```

```
int i = 0;
    int full_size = buddy.size;
    while (i < 2 * P PAGE - 1) {
        if ((i + 1) * ins bitmap[i] == index) {
            if (ins_bitmap[i] == 0) {
                ins_bitmap[i] = full_size;
                int ii = i;
                int t full size = full size;
                while (ii > 0) {
                    ii = parent(ii);
                    t full size <<= 1;
                    if (ins_bitmap[lson(ii)] == (t_full_size >> 1)
                        && ins_bitmap[rson(ii)] == (t_full_size >> 1)) {
                        ins_bitmap[ii] = t_full_size;
                    } else {
                        ins_bitmap[ii] = ins_bitmap[rson(ii)] >
ins_bitmap[lson(ii)] ?
                                ins_bitmap[rson(ii)] : ins_bitmap[lson(ii)];
                    }
                }
            } else {
                full size >>= 1;
                i = lson(i);
            }
        } else {
            full_size >>= 1;
            i = rson(i);
        }
    }
}
```

## 实现统一内存分配接口

```
void *kmalloc(size_t size)
{
  int objindex;
  void *p;

  if(size == 0)
    return NULL;

// size 若在 kmem_cache_objsize 所提供的范围之内,则使用 slub allocator 来分配内存
for(objindex = 0; objindex < NR_PARTIAL; objindex ++) {
  if (size <= kmem_cache_objsize[objindex]) {
    p = kmem_cache_alloc(slub_allocator[objindex]);
    break;
  }
  // YOUR CODE HERE</pre>
```

```
// size 若不在 kmem_cache_objsize 范围之内,则使用 buddy system 来分配内存
if(objindex >= NR_PARTIAL){
    // YOUR CODE HERE
    p = alloc_pages(size >> 12);
    set_page_attr(p, (size-1) / PAGE_SIZE, PAGE_BUDDY);
}
return p;
}
```

```
void kfree(const void *addr)
{
    struct page *page;
   if(addr == NULL)
       return;
    // 获得地址所在页的属性
   // YOUR CODE HERE
   page = ADDR_TO_PAGE(addr);
    // 判断当前页面属性
   if(page->flags == PAGE_BUDDY){
       // YOUR CODE HERE
       free_pages(addr);
       clear_page_attr(ADDR_TO_PAGE(addr)->header);
    } else if(page->flags == PAGE_SLUB){
       // YOUR CODE HERE
       kmem cache free(addr);
    return;
}
```

## 为 mm\_struct 添加 vm\_area\_struct 数据结构

```
typedef struct { unsigned long pgprot; } pgprot_t;
struct vm_area_struct {
    /* Our start address within vm_area. */
    unsigned long vm_start;
    /* The first byte after our end address within vm_area. */
    unsigned long vm_end;
    /* linked list of VM areas per task, sorted by address. */
    struct vm_area_struct *vm_next, *vm_prev;
    /* The address space we belong to. */
    struct mm_struct *vm_mm;
    /* Access permissions of this VMA. */
    pgprot_t vm_page_prot;
```

```
/* Flags*/
unsigned long vm_flags;
};

struct mm_struct {
   unsigned long *pgtbl;
   unsigned long text_start;
   unsigned long text_end;
   unsigned long stack_top;
   unsigned long pa_for_stack;
   struct vm_area_struct *vm_area_head;
};
```

## 实现mmap/munmap/mprotect的SysCall

#### 实现mmap函数

```
#define off t unsigned long
unsigned long get_unmapped_area(size_t length) {
    unsigned long i = 0;
    for (i = 0; ; ++i) {
        struct vm_area_struct *ptr;
        int flag = 0;
        for (ptr = current->thread.mm->vm area head; ptr != NULL; ptr = ptr-
>vm_next) {
            if (ptr->vm_start < i + length && ptr->vm_end > i) {
                flag = 1;
                break;
            }
        if (flag == 0) {
           return i;
        }
    }
}
void *do_mmap(struct mm_struct *mm, void *start, size_t length, int prot) {
    struct vm_area_struct *ptr;
    for (ptr = mm->vm area head; ptr != NULL; ptr = ptr->vm next) {
        if (ptr->vm_start < (unsigned long)start + length && ptr->vm_end >
(unsigned long)start) {
            start = get_unmapped_area(length);
        }
    }
```

```
struct vm_area_struct *new_vm_area = kmalloc(sizeof(struct
vm area struct));
   new_vm_area->vm_start = (unsigned long)start;
   new vm area->vm end = start + length;
   new_vm_area->vm_flags = prot;
   new vm area->vm mm = mm;
   new_vm_area->vm_page_prot.pgprot = ((prot & PROT_READ ? 1 : 0) | (prot &
PROT_WRITE ? 2 : 0) | (prot & PROT_EXEC ? 4 : 0)) << 1;
    if (mm->vm_area_head == NULL) {
       mm->vm_area_head = new_vm_area;
    } else {
       for (ptr = mm->vm_area_head; ptr->vm_next != NULL; ptr = ptr->vm_next)
;
       ptr->vm_next = new_vm_area;
       new_vm_area->vm_prev = ptr;
    }
   return start;
}
void *mmap (void *__addr, size_t __len, int __prot,
            int flags, int fd, off t offset) {
    return do_mmap(current->thread.mm, __addr, __len, __prot);
}
```

#### 修改task init函数代码,更改为需求分页机制

```
sscratch top = kmalloc(PAGE SIZE) + PAGE SIZE;
task[1]->state = TASK_RUNNING;
#ifdef SJF
task[1]->counter = rand();
#endif
#ifdef PRIORITY
task[1]->counter = 8 - 1;
#endif
task[1]->priority = 5;
task[1]->blocked = 0;
task[1]->pid = 1;
task[1]->thread.sp = USER_END;
task[1]->thread.sepc = 0;
task[1]->thread.ra = (unsigned long)task_first_ret;
task[1]->thread.sscratch = sscratch_top;
task[1]->thread.mm = kmalloc(sizeof(struct mm struct));
task[1]->thread.mm->pgtbl = (unsigned long)kmalloc(PAGE_SIZE) - page_offset +
0x80000000ul;
task[1]->thread.mm->vm area head = NULL;
task[1]->thread.mm->pa for stack = 0;
unsigned long pgtbl_va = (unsigned long)task[1]->thread.mm->pgtbl;
```

```
initial_pgtbl(pgtbl_va);

do_mmap(task[1]->thread.mm, 0, PAGE_SIZE, PROT_READ | PROT_WRITE | PROT_EXEC);
do_mmap(task[1]->thread.mm, USER_END - PAGE_SIZE, PAGE_SIZE, PROT_READ |
PROT_WRITE);
```

#### 实现munmap函数

```
int munmap(void *start, size t length) {
    struct vm_area_struct *found = NULL;
    if (current->thread.mm->vm_area_head == NULL) {
        return -1;
    } else {
        struct vm_area_struct *ptr;
        for (ptr = current->thread.mm->vm area head; ptr != NULL; ptr = ptr-
>vm_next) {
            if (ptr->vm_start == (unsigned long)start && ptr->vm_end ==
(unsigned long)start + length) {
                found = ptr;
                break;
            }
        }
    }
    if (found == NULL) {
        return -1;
    free page tables(current->thread.mm->pgtbl, start, length, 1);
    if (found->vm_prev == NULL && found->vm_next == NULL) {
        current->thread.mm->vm area head = NULL;
    } else if (found->vm prev == NULL) {
        found->vm_next->vm_prev = NULL;
        current->thread.mm->vm area head = found->vm next;
    } else if (found->vm_next == NULL) {
        found->vm_prev->vm_next = NULL;
    } else {
        found->vm_prev->vm_next = found->vm_next;
        found->vm next->vm prev = found->vm prev;
   kfree(found);
   return 0;
}
```

#### 实现free\_page\_tables函数

```
unsigned long free_table_dir(unsigned long *tblptr, unsigned long va, unsigned
long right) {
  tblptr = (unsigned long)tblptr - 0x80000000ul + page_offset;
```

```
unsigned long tbl index = (va >> (unsigned long)right) & (unsigned
long) 0x1FF;
    if ((tblptr[tbl index]) & 1) {
        if (((tblptr[tbl_index] >> 1) & 0x7) == 0) {
            return free_table_dir(((((tblptr[tbl_index]) >> 10) & (unsigned)
long)0xFFFFFFFFFFF) << 12), va, right - 9);</pre>
        } else {
            tblptr[tbl_index] ^= 1;
            return (tblptr[tbl_index] >> 10) << 12;</pre>
        }
    } else {
        return NULL;
    }
}
void free page tables (unsigned long pgtbl, unsigned long va, unsigned long sz,
int free frame) {
    unsigned long i = 0;
    unsigned long va_aligned = va;// >> 12) << 12;</pre>
    for (i = 0; i < sz; i += 0x1000) {
        unsigned long pa = free_table_dir(pgtbl, va_aligned + i, 30);
        if (pa != NULL && free_frame) {
            kfree(pa + kmem_struct.virtual_offset);
        }
    }
}
```

### 实现mprotect函数

```
int mprotect (void *__addr, size_t __len, int __prot) {
    struct vm area struct *found = NULL;
    if (current->thread.mm->vm_area_head == NULL) {
        return -1;
    } else {
        struct vm_area_struct *ptr;
        for (ptr = current->thread.mm->vm area head; ptr != NULL; ptr = ptr-
>vm_next) {
            if (ptr->vm_start == (unsigned long)__addr && ptr->vm_end ==
(unsigned long)__addr + __len) {
                found = ptr;
                break;
            }
        }
    }
    if (found == NULL) {
       return -1;
    }
```

```
protect_page_tables(current->thread.mm->pgtbl, __addr, __len, __prot);
found->vm_flags = __prot;
found->vm_page_prot.pgprot = __prot << 1;
return 0;
}</pre>
```

## 实现 Page Fault 的检测与处理

```
void do_page_fault() {
    unsigned long scause;
    __asm__ _volatile__ ("csrr %0, scause\n\t" : "=r" (scause));
   unsigned long sepc;
    __asm__ _volatile__ ("csrr %0, sepc\n\t" : "=r" (sepc));
   unsigned long stval;
    asm volatile ("csrr %0, stval\n\t" : "=r" (stval));
   puts("[S] PAGE_FAULT: PID = ");
    puti(current->pid);
    puts("\n[S] PAGE_FAULT: scause: ");
   puti(scause);
   puts(", sepc: 0x");
   puti64(sepc);
   puts(", badaddr: 0x");
   puti64(stval);
   puts("\n");
   struct vm_area_struct *found = NULL;
    if (current->thread.mm->vm_area_head == NULL) {
        puts("Invalid vm area in page fault");
        return;
    } else {
        struct vm_area_struct *ptr;
        for (ptr = current->thread.mm->vm_area_head; ptr != NULL; ptr = ptr-
>vm_next) {
            if (ptr->vm_start <= stval && stval < ptr->vm_end) {
                found = ptr;
                break;
            }
        }
    if (found == NULL) {
        puts("Invalid vm area in page fault\n");
        return;
    if (scause == 12 && (found->vm_flags & PROT_READ) == 0 && (found->vm_flags
& PROT_EXEC) == 0 && (found->vm_flags & PROT_WRITE) == 0) {
```

```
puts("Invalid vm area in page fault\n");
        return;
    } else if (scause == 13 && (found->vm flags & (PROT READ)) == 0) {
        puts("Invalid vm area in page fault\n");
        return;
    } else if (scause == 15 && (found->vm flags & PROT READ) == 0 && (found-
>vm_flags & PROT_WRITE) == 0) {
        puts("Invalid vm area in page fault\n");
        return;
    }
    unsigned pa;
    unsigned long sz = (found->vm_end - found->vm_start + PAGE_SIZE - 1) /
PAGE_SIZE * PAGE_SIZE;
    if (current->thread.mm->pa_for_stack != 0 && found->vm_end == USER END) {
        pa = current->thread.mm->pa for stack;
        create mapping(current->thread.mm->pgtbl, found->vm start, current-
>thread.mm->pa_for_stack, sz, found->vm_page_prot.pgprot | (1 << 4));</pre>
    } else if (found->vm start == 0) {
        pa = 0x84000000ul;
        create_mapping(current->thread.mm->pgtbl, found->vm_start,
0x84000000ul, sz, found->vm page prot.pgprot | (1 << 4));
    } else {
        pa = kmalloc(sz) - page_offset + 0x80000000ul;
        create_mapping(current->thread.mm->pgtbl, found->vm_start / PAGE_SIZE *
PAGE SIZE, pa, sz, found->vm page prot.pgprot | (1 << 4));
    }
   puts("[S] mapped PA: ");
   puti64(pa);
    puts(" to VA: ");
   puti64(found->vm_start);
   puts(" with size: ");
   puti64(sz);
   puts("\n");
}
```

注意需要特判fork子进程的栈页和进程的代码段,这时不需要用kmalloc来获取物理页,前者通过映射到记录的已复制的子进程物理页,后者映射到0x84000000

### 实现 fork 系统调用

本次fork是实验的一大难点,增加了很多实验手册中所没有说明的代码。

fork

```
int fork() {
    ++task_num_top;
    task[task_num_top] = kmalloc(sizeof(struct task_struct));
    task[task_num_top]->state = TASK_RUNNING;
```

```
unsigned long sscratch top = kmalloc(PAGE SIZE) + PAGE SIZE;
    task[task_num_top]->state = TASK_RUNNING;
    #ifdef SJF
    task[task num top]->counter = rand();
    #endif
    #ifdef PRIORITY
    task[task_num_top]->counter = 8 - task_num_top;
    #endif
    task[task_num_top]->priority = 5;
    task[task_num_top]->blocked = 0;
    task[task num top]->pid = task num top;
    task[task_num_top]->thread.sp = USER_END;
   puts("[PID = ");
    puti(task[task_num_top]->pid);
    puts("] Process fork from [PID = ");
   puti(current->pid);
    puts("] Successfully! counter = ");
    puti(task[task_num_top]->counter);
    puts("\n");
    // initial_kernel_stack(sscratch_top, sscratch_top - 280);
    task[task_num_top]->thread.sepc = current->thread.sepc;
    task[task_num_top]->thread.ra = (unsigned long)forkret;
    task[task_num_top]->thread.sscratch = sscratch_top;
    task[task_num_top]->thread.sp = sscratch_top - 280;
    task[task_num_top]->thread.mm = kmalloc(sizeof(struct mm_struct));
    task[task_num_top]->thread.mm->pgtbl = (unsigned long)kmalloc(PAGE_SIZE) -
page offset + 0x80000000ul;
    task[task_num_top]->thread.mm->vm_area_head = NULL;
    task[task_num_top]->thread.mm->pa_for_stack = 0;
    unsigned long pgtbl va = (unsigned long)task[task num top]->thread.mm-
>pgtbl;
    initial_pgtbl(pgtbl_va);
    if (current->thread.mm->vm_area_head != NULL) {
        struct vm_area_struct *tail = NULL;
        for (struct vm_area_struct *ptr = current->thread.mm->vm_area_head; ptr
!= NULL; ptr = ptr->vm_next) {
            struct vm_area_struct *t = kmalloc(sizeof(struct vm_area_struct));
            t->vm_start = ptr->vm_start;
            t->vm end = ptr->vm end;
            t->vm_flags = ptr->vm_flags;
            t->vm_page_prot = ptr->vm_page_prot;
            t->vm_mm = task[task_num_top]->thread.mm;
            t->vm_prev = tail;
            t->vm_next = NULL;
            if (tail == NULL) {
                task[task num top]->thread.mm->vm area head = t;
```

```
} else {
                tail->vm next = t;
           tail = t;
       }
    }
    task[task num top]->thread.stack = task[task num top]->thread.sp;
    for (int i = 0; i < 280 / 8; ++i) {
        task[task num top]->thread.stack[i] = current->thread.stack[i];
    task[task_num_top]->thread.stack[9] = 0;
    char *new_stack = kmalloc(PAGE_SIZE);
    task[task_num_top]->thread.mm->pa_for_stack = (unsigned long)new_stack -
page offset + 0x80000000ul;
    for (int i = 0; i < PAGE_SIZE; ++i) {</pre>
        new_stack[i] = ((char*)(USER_END - PAGE_SIZE))[i];
    }
   return task num top;
}
```

这里的一大难点是,要在这里记录复制好的子进程的栈页表的物理地址。用于后续page fault的时候直接对复制好的物理地址做映射。

```
void forkret() {
   asm volatile ("csrrw x0, sscratch, %0" : : "r"(current->thread.sscratch));
   ret_from_fork(current->thread.stack);
}
```

子进程退出handler\_s时候的返回函数,主要为调用ret\_from\_fork

```
ret_from_fork:
    ld t0, 248(a0)
    csrw sepc, t0
    ld t0, 256(a0)
    csrw sstatus, t0
    ld t0, 264(a0)
    csrw stval, t0
    ld t0, 272(a0)
    csrw scause, t0
    ld x31, 240(a0)
    ld x30, 232(a0)
    ld x29, 224(a0)
    ld x28, 216(a0)
    ld x27, 208(a0)
    ld x26, 200(a0)
```

```
ld x25, 192(a0)
ld x24, 184(a0)
ld x23, 176(a0)
ld x22, 168(a0)
ld x21, 160(a0)
ld x20, 152(a0)
ld x19, 144(a0)
ld x18, 136(a0)
ld x17, 128(a0)
ld x16, 120(a0)
ld x15, 112(a0)
ld x14, 104(a0)
ld x13, 96(a0)
ld x12, 88(a0)
ld x11, 80(a0)
ld x9, 64(a0)
ld x8, 56(a0)
1d \times 7, 48(a0)
ld x6, 40(a0)
ld x5, 32(a0)
ld x4, 24(a0)
ld x3, 16(a0)
ld x2, 8(a0)
ld x1, 0(a0)
ld x10, 72(a0)
sret
```

这里从传入的stack参数中恢复寄存器。

## 编译及测试

## 遇到的问题

● create\_mapping的传入地址是虚拟地址还是物理地址

因为页表项的内容是物理地址,而我们kmalloc返回的地址均设为虚拟地址。这本来是很棘手的问题,但因为我们kmalloc返回的虚拟地址是内核地址,所以可以做简单的减法变换到物理变量。我通过在buddy.c中设置page\_offset,可以将alloc\_pages的返回值这个page\_offset相减并加上0x80000000来获得物理地址。

● fork的时候会发生嵌套异常,这需要特别处理sscratch来判断是否发生嵌套异常。(该思路来自于 linux的内核中关于嵌套处理的代码,由陈昱文同学提供思路)

## 运行结果

## 实验心得

本次实验,调的时候真的超级超级超级自闭,create\_mapping的bug卡了好久。最后也没找出是啥原因,重构了一遍代码后才解决bug。

好难好难好难,可能是因为工作量太大,我又拖到考试事周前才开始做,比较着急。

但是做下来真的学到了很多,感谢周亚金老师,感谢助教。

# 对实验指导的建议

fork提供的提示有些少,好多要自己摸索。可以提一下嵌套映射的问题,以及画一张图说明一下fork父 进程和子进程调用时的栈信息。