Algorithms: FIRST and FOLLOW Sets

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Terminology

Before discussing the algorithms, let's review terminology. The set T defines all terminals of the grammar. The nonterminals are defined either by the set N or by looking at the vocabulary (V) and removing all terminals. Everything in the vocabulary is called a symbol. Thus, all nonterminals and terminals are symbols. The set P defines all productions. A production consists of the left-hand side and right-hand side of the production arrow (\rightarrow) . In the grammars we care about—context-free grammars—the left-hand side will only contain a single, nonterminal. Here is an example production: $A \rightarrow aB$

A is a nonterminal because it is on the left-hand side of the production. aB is on the right-hand side of the production. For this document, we will use the convention that upper-case letters are nonterminals and lower-case letters are terminals. When in doubt, use T, V or N, and the left-hand side of productions to determine which symbols are terminals and which are nonterminals.

Algorithm 1: Computing FIRST sets

We compute FIRST sets on productions. When computing the FIRST set of a the production use these steps (let's say A is the nonterminal on the right-hand side of the production):

- 1) Look at the leftmost symbol on the right-hand side of the production
- 2.1) If the leftmost symbol is a terminal then that terminal is the only symbol in the FIRST set of the production $A \to \dots$
- 2.2) Else, if the leftmost symbol is a nonterminal (let's call it B) then add all terminals in the FIRST sets of all productions with B on the *left-hand* side into the FIRST set of the original production $A \to B \dots$

Algorithm 2: Computing FOLLOW sets

When computing the FOLLOW set of a nonterminal A use these steps:

- 1) Find all the productions that *produce* the nonterminal A
- 2) For each production in step 1, find all places where the nonterminal A is produced and what single symbol is produced immediately afterward
- 3.1) For each symbol found in step 2, If the symbol is a terminal, then add that terminal into the FOLLOW set of A
- 3.2) Else, if the symbol is a nonterminal (let's call it B), then add each element in the FIRST set of B into the FOLLOW set of A
- 3.3) Else, if there is no symbol, go up another level, by repeating steps 1-3 for B (if we go up to the starting nonterminal then add the end of input symbol, #, to the FOLLOW set of A)