

## 01-D-2 级数

#数据结构邓神

### 算数级数

$$T(n) = 1 + 2 + \dots + n = n(n+1)/2 = O(n^2) \quad // \text{ BigO}$$

最后的复杂度等于末项的平方

### 幂方级数:比幂次高一阶

$$T(n) = 1^2 + 2^2 + \dots + n^2 = n(n+1)(2n+1)/6 = O(n^3)$$

$$T(n) = 1^3 + 2^3 + \dots + n^3 = n^2(n+1)^2/4 = O(n^4)$$

...

### 严格证明

$$\sum_{k=0}^n k^d \approx \int_0^n x^{d+1} dx = \frac{1}{d+1} x^{d+1} \Big|_0^n = \frac{1}{d+1} n^{d+1} = O(n^{d+1})$$

### 几何级数 ( $a > 1$ ): 与末项同阶

$$T_a(n) = a^0 + a^1 + \dots + a^n = (a^{n+1} - 1)/(a - 1) = O(a^n)$$

$$1 + 2 + 4 + \dots + 2^n = 2^{n+1} - 1 = O(2^n)$$

### 收敛级数

$$1/1/2 + 1/2/3 + 1/3/4 + \dots + 1/(n-1)/n = 1 - 1/n = O(1)$$

$$1 + 1/n^2 + \dots + 1/n^2 < 1 + 1/2^2 + \dots = \pi^2/6 = O(1)$$

$$1/3 + 1/7 + 1/8 + \dots = 1 = O(1)$$

有必要吗, 每一项都是分数?

难道基本操作次数和存储单元数可能是分数吗？

某种意义上是！

$$(1-\lambda) * [1 + 2\lambda + 3\lambda^2 + \dots] = 1/(1 - \lambda) = O(1) \quad 0 < \lambda < 1 \quad // \text{ 几何分布}$$

未必收敛，但是长度有限

$$h(n) = 1 + 1/2 + 1/3 + \dots + 1/n = O(\log n) \quad // \text{ 调和级数}$$

$$\log 1 + \log 2 + \log 3 + \dots + \log n = \log(n!) = O(n \log n) \quad // \text{ 对数级数}$$