### Haskell

A functional paradigm exemplar and other aspects (cont'd)

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### **Typeclasses**

Types can be part of typeclasses (members) $^1$ :

```
Prelude> : t elem elem :: Eq a \Rightarrow a \rightarrow [a] \rightarrow Bool
```

- Eq: type supports == en /=
- Ord: type is ordered
- Show: type can be represented as a string
- Read: string can be represented as the type
- Enum: type is sequentially ordered
- Bounded: type has lower- and upper-bounds
- Num: type acts as a number
- Integral: type is an integer number (Int/Integer)
- Floating: type is a floating point number (Float/Double)



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<sup>1</sup>https://www.haskell.org/tutorial/classes.html

### The Eq typeclass

```
module MyEq where
class Eq a where
(==) :: a -> a -> Bool
(/=) :: a -> a -> Bool
```

```
*MyEq> [1,2] MyEq.== [3,4]

<interactive >:52:1: error:
    No instance for (MyEq.Eq [Integer])
    arising from a use of MyEq.==
    In the expression: [1, 2] MyEq.== [3, 4]
    In an equation for it : it = [1, 2] MyEq.== [3, 4]
```

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### Type system

- Type checker:
  - $\mathsf{static} o \mathsf{at} \ \mathsf{compile}\mathsf{-time} \quad \mathsf{strong} o \mathsf{No} \ \mathsf{unchecked} \ \mathsf{runtime} \ \mathsf{type} \ \mathsf{error}$
- Type inference: algorithmic reasoning about types of variables and parameters

```
Prelude> [1,'a']
<interactive>:62:2: error:

No instance for (Num Char) arising from the literal 1
In the expression: 1
In the expression: [1, 'a']
In an equation for it : it = [1, 'a']
```

• Type classes: introduce overloading as a principle  $\rightarrow$  ad-hoc polymorphism <sup>2</sup>

```
class (Eq a) => Ord a where
(<), (<=), (>=), (>) :: a -> a -> Bool
max, min :: a -> a -> a
```

multiple inheritance

```
class (Eq a, Show a) => ShowOrd a where ...
```

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<sup>2</sup>https://www.researchgate.net/publication/2710954\_How\_to\_Make\_Ad-Hoc\_Polymorphism\_Less\_Ad\_Hoc\_

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# Type Synonyms

We can define aliases using the type keyword:

```
Lecture.hs
```

```
type Feedback = (Int, Int)
```

# Newtypes

We can define new types using the **newtype** keyword  $^3$ . These will not have the same typeclasses inherited, but the same representation.

```
module MN where
newtype Natural = MakeNatural Integer
toNatural :: Integer -> Natural
to Natural x \mid x < 0 = error "Can't create negative naturals!"
           | otherwise = MakeNatural x
from Natural :: Natural -> Integer
from Natural (Make Natural i) = i
instance Num Natural where
  fromInteger = toNatural
 \times + v
                 = toNatural (fromNatural x Prelude.+ fromNatural v)
 x - y
                     = let r = fromNatural x Prelude. - fromNatural y in
                           if r < 0 then error "Unnatural subtraction"
                                    else toNatural r
                     = toNatural (fromNatural x Prelude.* fromNatural y)
  x * y
  ahs x
                     = x
                     = toNatural (Prelude.signum $ fromNatural x)
  signum x
```

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### Custom datatypes

We can define our own types using the **data** keyword:

```
Lecture.hs

data Color = Red | Yellow | Blue | Green | Orange | Purple
deriving (Eq, Show, Bounded, Enum)
```

```
Prelude> [minBound..maxBound] :: [Color]
[Red, Yellow, Blue, Green, Orange, Purple]
Prelude> let foo = Red
Prelude> foo
Red
Prelude> :t foo
foo :: Color
```

# Playing with Types<sup>4</sup>

"casting"

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# Playing with Types

#### "inferring"

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## Playing with Types

"inferring" cont'd

# Playing with Types

"inferring" cont'd

```
|Prelude> elem'a = foldr((||). (==a)) False
```

### Type and data constructors

Tree is a type constructor, while Leaf and Node are data constructors:

```
Lecture.hs
data Tree a = Leaf | Node (Tree a) (Tree a)
```

## Using Type constructors and functors

#### You make the Tree datatype a functor by:

```
instance Functor Tree where fmap f (Leaf x) = Leaf (f x) fmap f (Node t1 t2) = Node (fmap f t1) (fmap f t2)
```

Functors are also known as <\$>

### Monads

For a gentle introduction to Functors, Applicatives and Monads go to http://adit.io/posts/2013-04-17-functors,\_applicatives,\_and\_monads\_in\_pictures.html

```
class Monad m where

(>>=) :: m a -> ( a -> m b) -> m b -- the bind operator

(>>) :: m a -> m b -- sequence, but discard intermediate

return :: a -> m a -- wrap it back for purity

fail :: String -> m a
```

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## The Maybe Monad

```
Maybe a = Nothing | Just a
instance Monad Maybe where
Nothing >>= _ = Nothing
(Just x) >>= f = f x
```

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#### Let's revisit Natural

```
module MN where ... instance Show Natural where show x = show (fromNatural x)
```