

組別: _____ 簽名: _____

[group1]

1. Please order the sequence of the following steps for MIPS to invoke the exception handler.
 - a. Save PC in the exception program counter (EPC) register
 - b. Use mfc0 instruction to retrieve/copy EPC value to a general-purpose register
 - c. Jump to the predefined handler address
 - d. Return after the corrective action using jr instruction

a → c → b → d.

[group2]

2. Let two decimal numbers $A = 57$, $B = -84$, please answer the following questions.

1. Compute $A + B$ in 8-bit 2's complement
2. Compute $A - B$ in 8-bit 2's complement

Does overflow occur in the answers of question (1) and (2)? Why or why not? Please answer and explain them individually.

$$\begin{array}{r} 2 \overline{) 57} \\ \underline{2 \overline{) 28}} \\ 2 \overline{) 14} \\ \underline{2 \overline{) 7}} \\ 2 \overline{) 3} \\ \underline{2 \overline{) 2}} \\ 1 \end{array}$$

$$A = 57 = 00111001$$

$$B = -84 = 10101100$$

$$(1) A+B = 57-84 < -2^7$$

$$\begin{array}{r} 00111001 \\ + 10101100 \\ \hline 11100101 \end{array} \quad \text{not overflow}$$

$$\begin{array}{r} 2 \overline{) 84} \\ \underline{2 \overline{) 42}} \\ 2 \overline{) 21} \\ \underline{2 \overline{) 10}} \\ 2 \overline{) 5} \\ \underline{2 \overline{) 2}} \\ 1 \end{array}$$

$$+84 = 01010100$$

$$\Rightarrow 01010011$$

$$-84 = 10101100$$

$$(2) A-B = 57+84 = 141 > +2^7-1$$

$$\begin{array}{r} 00111001 \rightarrow 57 \\ + 01010100 \rightarrow 84 \\ \hline 10001101 \end{array} \quad \text{overflow}$$

[group5]

3. Which of the following statements are true?

- d.e.
- ☒ a. When we want to subtract b from a, we can add a to b', where b' is the action that 0 to 1, 1 to 0 in every bit. $\bar{b} + 1$
 - ☒ b. If we want to execute nor instruction, Ainvert, Bnegate, operation will be set to 1, 1, 01 respectively. 1100
 - ☒ c. When detecting overflow, we can use one exclusive ~~nor~~ gate to complete the task.
 - ☒ d. It is expensive to build a fully carry look-ahead adder.
 - ☒ e. It is impossible to occur overflow when add a to b, where a is positive and b is negative.

[group12]

4. Please select the correct options *a.c.d.*

- ☒ (a) When we implement 'set on less than' in ALU, the set value in ALU0 (LSB) is according to the result of adder in ALU31 (MSB).
- ☒ (b) We use 'XOR gate' to detect the carry into/out of most significant bit because the overflow occurs when the values in two bits are the same. $MSB \rightarrow sign$
- ☒ (c) 'Saturation' means that when a calculation overflows, the result is set to the largest positive number or the most negative number, rather than a modulo calculation as in 2's complement arithmetic. $R43$
- ☒ (d) When we implement 'subtraction' in ALU, we add 'one' because of 2's complement. C_n

[group7]

5. Choose the correct options *b.e.*

- ☒ a. all the languages require raising an exception of overflow *C not need*
- ☒ b. some languages will save PC in exception program counter register when overflow *ex. fortran*
- ☒ c. when adding operands with different signs, overflow ~~may~~ occur
- ☒ d. when carry into MSB ~~equal~~ to carry out of MSB means overflow
- ☒ e. add, addi will invoke exception handler during overflow

[group8]

a, c, d.

6. Carry look-ahead adder can diminish the carry delay which dominates the delay of ripple carry adder. Generate (g_i) and propagate (p_i) functions are two main operations of carry look-ahead adder. Assume a and b are two operands and c_{i+1} is the carry out of level i and carry in of level $i+1$, which is (are) correct?

✓ a. $g_i = a_i \cdot b_i$

✗ b. $p_i = (a_i + b_i) \cdot c_i$ $a_i \text{ XOR } b_i$

$c_{i+1} = g_i + p_i \cdot c_i$

✓ c. If g_i equals to 1, we can say the carry out of level i is 1.

✓ d. Carry look-ahead adder can be extended to multi-level style. The first group generate of a 3-bit group can then be defined as $G_0 = g^2 + (p^2 \cdot g^1) + (p^2 \cdot p^1 \cdot g^0)$

hint : 2 bit group $G_0 = g^1 + p^1 \cdot g^0$

[group14]

7. Suppose all numbers are 4-bits signed numbers, will the following formulas overflow?

- (a) $0111 + 0100 =$ **yes**
- (b) $0101 + 1011 =$ **No**
- (c) $1000 + 1111 =$ **yes**
- (d) $0010 + 0101 =$ **No**
- (e) $1101 + 1100 =$ **No**

以2補數計算

$$\begin{array}{r} 0111 \\ 0100 \\ \hline 1011 \end{array}$$

$$\begin{array}{r} 0101 \\ 1011 \\ \hline 10000 \end{array}$$

$$\begin{array}{r} 1000 \\ 1111 \\ \hline 10111 \end{array}$$

$$\begin{array}{r} 0010 \\ 0101 \\ \hline 0111 \end{array}$$

$$\begin{array}{r} 1101 \\ 1100 \\ \hline 11001 \end{array}$$

- (a) $0111 + 0100 = 1011 \rightarrow \text{overflow}$
- (b) $0101 + 1011 = 0110 \rightarrow \text{no overflow}$
- (c) $1000 + 1111 = 1111 \rightarrow \text{no overflow}$
- (d) $0010 + 0101 = 0111 \rightarrow \text{no overflow}$
- (e) $1101 + 1100 = 1001 \rightarrow \text{overflow}$

sign number $\rightarrow 7 \sim -7$

以sign number計算