

# Lab05-DynamicProgramming

CS214-Algorithm and Complexity, Xiaofeng Gao, Spring 2021.

\* If there is any problem, please contact TA Haolin Zhou.

\* Name: Zilong Li Student ID: 518070910095 Email: logcreative-lzl@sjtu.edu.cn

1. *Optimal Binary Search Tree.* Given a sorted sequence  $K = \langle k_1, k_2, \dots, k_n \rangle$  of  $n$  distinct keys, and we wish to build a binary search tree from these keys. For each key  $k_i$ , we have a probability  $p_i$  that a search will be for  $k_i$ . Some searches may be for values not in  $K$ , and so we also have  $n + 1$  *dummy keys*  $d_0, d_1, d_2, \dots, d_n$  representing values not in  $K$ . In particular,  $d_0$  represents all values less than  $k_1$ , and  $d_n$  represents all values greater than  $k_n$ . For  $i = 1, 2, \dots, n - 1$ , the dummy key  $d_i$  represents all values between  $k_i$  and  $k_{i+1}$ . For each dummy key  $d_i$ , we have a probability  $q_i$  that a search will correspond to  $d_i$ . Each key  $k_i$  is an internal node, and each dummy key  $d_i$  is a leaf. Every search is either successful (finding some key  $k_i$ ) or unsuccessful (finding some dummy key  $d_i$ ), and so we have  $\sum_{i=1}^n p_i + \sum_{i=0}^n q_i = 1$ .
  - (a) Prove that if an optimal binary search tree  $T$  ( $T$  has the smallest expected search cost) has a subtree  $T'$  containing keys  $k_i, \dots, k_j$ , then this subtree  $T'$  must be optimal as well for the subproblem with keys  $k_i, \dots, k_j$  and dummy keys  $d_{i-1}, \dots, d_j$ .
  - (b) We define  $e[i, j]$  as the expected cost of searching an optimal binary search tree containing the keys  $k_i, \dots, k_j$ . Our goal is to compute  $e[1, n]$ . Write the state transition equation and pseudocode using **dynamic programming** to find the minimum expected cost of a search in a given binary tree. (**Remark:** You may use  $w(i, j) = \sum_{l=i}^j p_l + \sum_{l=i-1}^j q_l$ ).
  - (c) Implement your proposed algorithm in C/C++ and analyze the time complexity. ([The framework Code-OBST.cpp is attached on the course webpage](#)). Give the minimum search cost calculated by your algorithm. The test case is given as following:

$i$	0	1	2	3	4	5	6	7
$p_i$		0.04	0.06	0.08	0.02	0.10	0.12	0.14
$q_i$	0.06	0.06	0.06	0.06	0.05	0.05	0.05	0.05

- (d) Please draw the structure of the optimal binary search tree in the test case, and explain the drawing process.
2. *Dynamic Time Warping Distance.* **DTW** stretches the series along the time axis in a dynamic way over different portions to enable more effective matching. Let  $DTW(i, j)$  be the optimal distance between the first  $i$  and first  $j$  elements of two time series  $\bar{X} = (x_1 \dots x_n)$  and  $\bar{Y} = (y_1 \dots y_m)$ , respectively. Note that the two time series are of lengths  $n$  and  $m$ , which may not be the same. Then, the value of  $DTW(i, j)$  is defined recursively as follows:

$$DTW(i, j) = |x_i - y_j| + \min(DTW(i, j - 1), DTW(i - 1, j), DTW(i - 1, j - 1))$$

- (a) Implement the proposed DTW algorithm in C/C++ and analyze the time complexity of your implementation. ([The framework Code-DTW.cpp is attached on the course webpage](#)). Two test cases have been given in the source code.
- (b) The window constraint imposes a minimum level  $w$  of positional alignment between matched elements. The window constraint requires that  $DTW(i, j)$  be computed only when  $|i - j| \leq w$ . Modify your code to add a window constraint and give the results of  $w = 0$  and  $w = 1$  on the two test cases.

**Remark:** You need to include your .pdf and .tex and 2 source code files in your uploaded .rar or .zip file. Screenshots of test case results are acceptable.