

Hashing

The Idea

- The most basic but core idea:

lst = []

insert(12)

lst = [None, ..., None, 12]

index → 0, 1, 2, ..., 11, 12

insert(1)

lst = [None, 1, ..., 12]

index → 0, 1, ..., 12

Use Key to map item
so we can access the item
by the "index" $O(1)$

- Key, element = "my num", 5

[0, 9, 3, 420, 5, 40]



Hash Functions

```
def hashf(Key):
```

```
    hash_val = 0
```

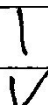
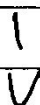
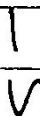
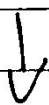
```
    for ix, c in enumerate(Key, start=2):
```

```
        hv += ix * ord(c)
```

```
    return hash_val
```

Notes:

- Maybe use a nice prime near a power of 2
- Maybe start at 0 since using the whole key gives more variation



• We want hash functions to evenly distribute the keys

Bins & Balls

- If rand toss rand ball equally likely to land in any bin

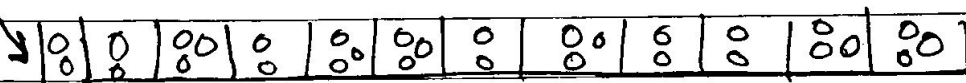
m bins n balls tossed, expect $\frac{n}{m}$ balls in each bin after awhile.

Expect two balls in the same bin after $\sim \sqrt{\pi \frac{n}{2}}$ tosses

Every bin has at least one ball after $\sim m \ln m$ tosses

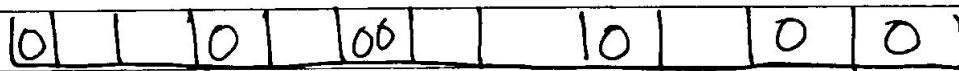
After m tosses the most loaded bin has $\Theta\left(\frac{\ln(m)}{\ln(\ln(m))}\right)$ balls.

After 30 tosses $\frac{n}{m}$ in each

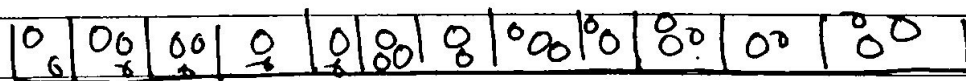


$12 = m$ bins, $n = 30$

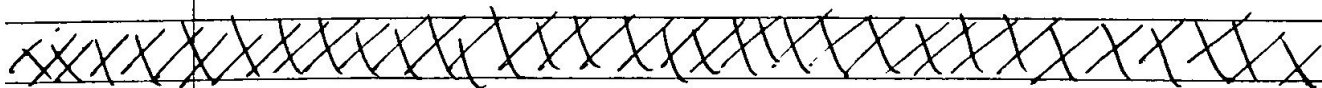
After 6.8 tosses $\sim \sqrt{\pi \frac{30}{2}}$



After 29.8 tosses $\sim 12 \ln 12$

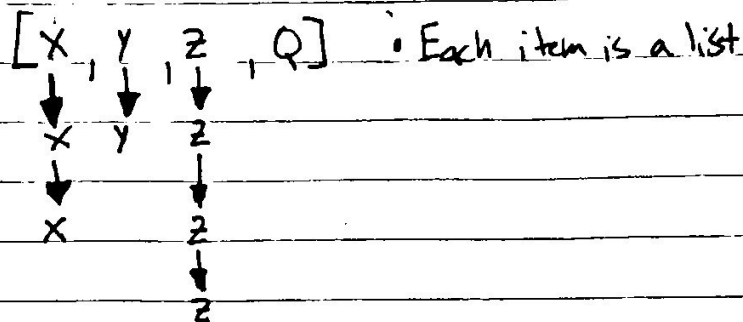


After 12 tosses $\sim \frac{\ln(12)}{\ln(\ln(12))} = 2.73$



Separate Chaining

$q, z, y, x = \text{any random element doesn't matter}$



* Doesn't have to use linked list.

* keep size of list to $n/5$ in order to say access is $O(1)$

Linear Probe

$[9, 3, 10, Nil, Nil, Nil, Nil]$

insert(420)

say the hash values here

$[9, 3, 10, 420, Nil, Nil, Nil]$

✓ ✓ ✓