

1 Definitions

Let $m = \left\lceil \left(\frac{N}{93} \right)^{\frac{1}{d}} \right\rceil$, base of the counter

MSR = most significant digit region

\mathcal{C}_0 = starting value of counter

$d = \lceil \log_m(\mathcal{C}_0) \rceil$, number of digits per row

$\mathcal{DR} = \left\lceil \frac{d}{3} \right\rceil$, digit regions per row

$\mathcal{C}_f = m^d - 1$, final value of the counter

$\mathcal{C}_\Delta = \mathcal{C}_f - \mathcal{C}_0$, number of rows

$l = \lceil \log m \rceil + 2$, bits needed to represent each digit plus 2 for MSR and MSD

$\mathcal{DR}_{height} = 3 \cdot (l + 30)$, height of a row

$h = c_\Delta \cdot \mathcal{DR}_{height}$, height of constuction without roof tiles

In order to determine the number of digits per row, given some value k , let $R = k \bmod 6$. There are 3 distinct cases to consider.

- if $0 \leq R \leq 1$, then $\{ d \in \mathbb{N} \mid 3 \text{ divides } d \}$
MSR \leftarrow 3 digits
- if $2 \leq R \leq 3$, then $\{ d \in \mathbb{N} \mid 3 \text{ divides } d - 1 \}$
MSR \leftarrow 1 digit
- if $4 \leq R \leq 5$, then $\{ d \in \mathbb{N} \mid 3 \text{ divides } d - 2 \}$
MSR \leftarrow 2 digits

2 Determining the starting value \mathcal{C}_0

...therefore, let $d = \lfloor \frac{k}{2} \rfloor$, $m = \left\lceil \left(\frac{N}{93} \right)^{\frac{1}{d}} \right\rceil$, $l = \lceil \log m \rceil + 2$, $\mathcal{C}_0 = m^d - \left\lfloor \frac{N-3l-76}{3l+90} \right\rfloor$, where d is the number of digits per row of the counter, m is the base of the counter, l is the number of bits needed to encode each digit in binary plus 2 for indicating whether a digit is in the MSR and is the MSD in that region, and \mathcal{C}_0 is the start of the counter in decimal.

In general, the height of a digit region is $3l + 90$. There are two cases when this is false, in the first digit region the height is $3l + 91$, and the last digit region has a height of $3l + 75$. To ensure the rectangle ends with height N , let h be the height of the construction before any filler/roof tiles are added. If we define \mathcal{C}_Δ as the number of **C**ounter unit rows, then $h = n(3l + 90) + (3l + 91) + (3l + 75)$, which we will re-write as $n(3l + 90) + 3l + 76$. So then the maximum height of the counter is $m^d(3l + 90) + 3l + 76$. Since our goal is to end with with a rectangle of height N , we need to pick a base such that the counter can increment so many times that when it stops, it is at least N .

Lemma 1. $N \leq m^d(3l + 90) + 3l + 76$.

Proof.

$$\begin{aligned}
N &= 93 \left(\frac{N}{93} \right) = 93 \left(\left(\frac{N}{93} \right)^{\frac{1}{d}} \right)^d \leq 93 \left\lceil \left(\frac{N}{93} \right)^{\frac{1}{d}} \right\rceil^d \\
&= 93m^d \leq 3lm^d + 90m^d \leq 3lm^d + 90m^d + 3l + 76 \\
&= m^d(3l + 90) + 3l + 76
\end{aligned}$$

□

2.1 Filling in the gaps

This means that the number of **Counter** unit rows \mathcal{C}_Δ is $m^d - \mathcal{C}_0$, where we have defined \mathcal{C}_0 as the starting value of the counter. To choose the best starting value, we find the value for \mathcal{C}_Δ that gets h as close to N without exceeding N . It follows from the equation $h = \mathcal{C}_\Delta(3l + 90) + 3l + 76$, that $\mathcal{C}_\Delta = \left\lfloor \frac{N-3l-76}{3l+90} \right\rfloor$. Thus, $\mathcal{C}_0 = m^d - \left\lfloor \frac{N-3l-76}{3l+90} \right\rfloor$. As a result of each digit requiring a width of 2 tiles, if k is odd, one additional tile column must be added. The number of filler tiles needed for the width is $k \bmod 2$, and the number of filler tiles for the height is $N - h$, which is $N - 3l - 76 \bmod 3l + 90$.

3 General counter

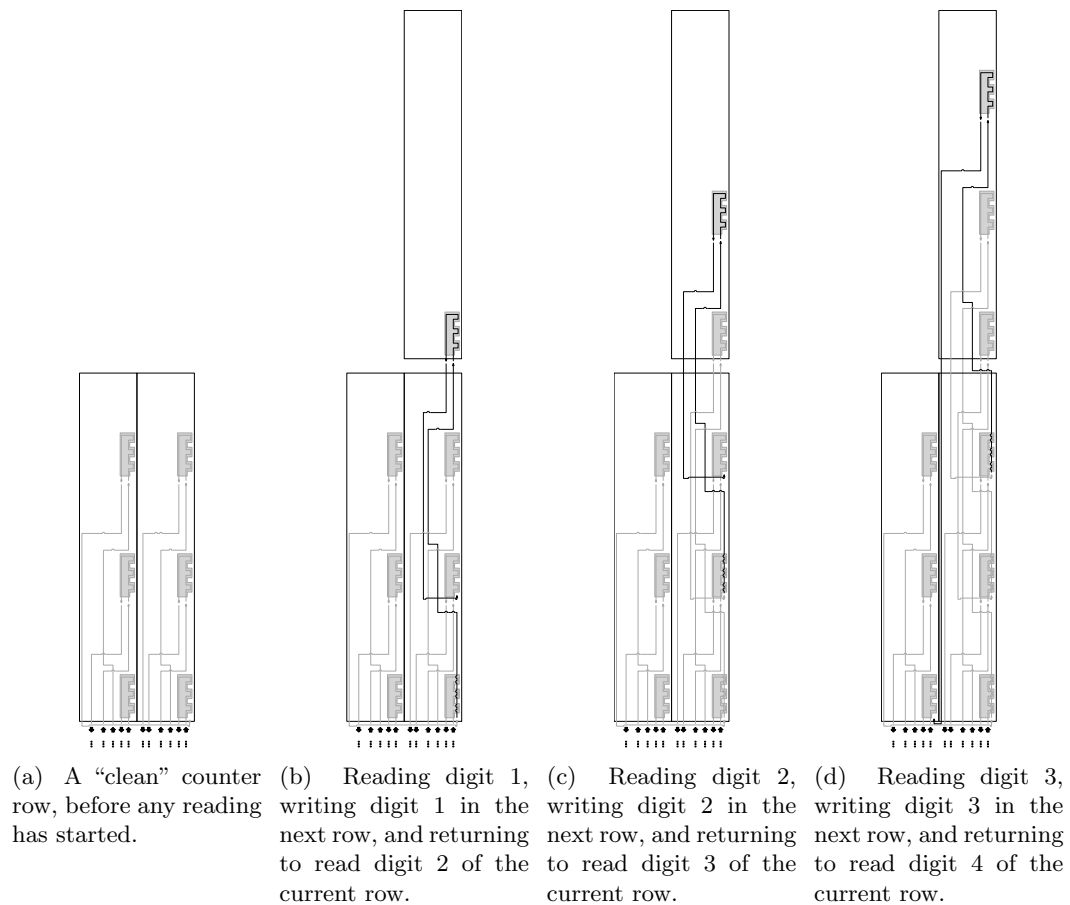


Figure 1: Progression of the counter as it reads the 3 least significant digits in one value, and writes the corresponding digits in the next row/value.