## **Identity Management**

### **Overview**

is used to retrieve the list of delegates in the current epoch. Asymmetric public keys of other gates are used for encryption of the delegate's address advertisement. The delegate enticates the message by signing it with the private bls key.

### **ass Diagram**

is the delegate's address advertisement message. The message has ip and port encrypted with ECIES public key of every delegate except for the sender and the hash of the encrypted data and other data is signed with the sender's private BLS key. AddressAdMessageTxAcceptor is the delegate's tx acceptor advertisement message. The hash of the message is signed with the sender's private key. P2pConsensusHeader precedes P2p consensus message. The header identifies receiver of the consensus message. - P2pHeader is the main p2p message header. If app\_type is Consensus then it is followed by P2pConsensusHeader, which in turn is followed by corresponding consensus message. If app\_type is AddressAd then it is followed by AddressAdMessage. If app\_type is AddressAdTxAcceptor then it is followed by AddressAddMessageTxAcceptor. - DelegateIdentityManager handles all identity manager functionality related to advertisement and persistence of the delegate and tx acceptor addresses.

The class diagram is shown on Figure 1.

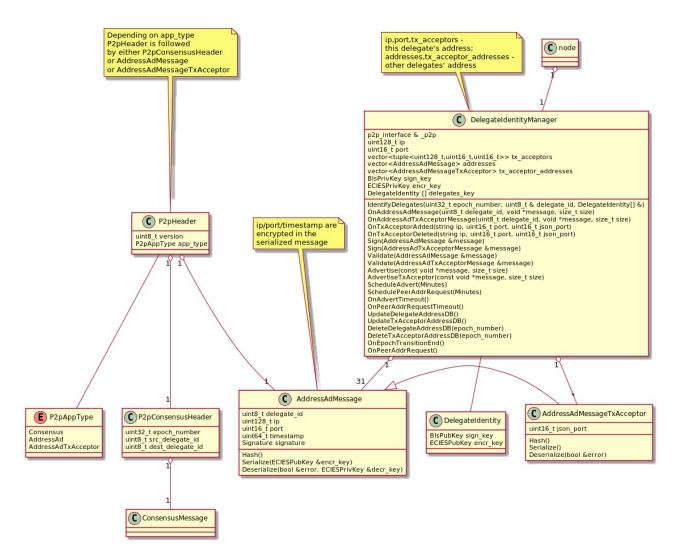


Figure 1. Class diagram.

## Sequence Diagram Delegate and Tx Acceptor Address Advertisement

On start up the node ascertains if it is the delegate in the current epoch (IdentifyDelegates()). If it is the delegate then it advertises via P2p substym its encrypted address and all of its tx acceptor addresses (Advertise(), AdvertiseTxAcceptor()). If the node is the delegate in the current time is greater than one hour before the epoch s<del>tart </del>then the node <mark>schedules a timer</mark> eduleAdvert) at (epoch start - now - 60min). On timeout 🚃 AdvertTimeout()), the node advertises via P2p substym its encrypted address and all of its tx acceptor addresses and schedules a timer (epoch start - now - 30min). On timeout, the node advertises via P2p subsystem its encrypted address and all of its tx acceptor addresses. The node also executes CheckDelegateAddress (Figure 3). If the node is the delegate in the next epoch and the current time is less than one hour before the epoch start and greater than 30min before epoch start then the node advertises via P2p subsystem its encrypted address and all of its tx acceptor addresses and schedules a timer at (epoch start - now - 30min). On timeout, the node advertises via P2p substym its encrypted address and all of its tx acceptor addresses. If the node is the delegate in the next epoch and the current time is less than 30min before the epoch start then the node advertises via P2p substym its encrypted address and all of its tx acceptor addresses. If the node is not the delegate in the current or next epoch then it doesn't execute and action and on

characteristic event it repeats the advertisement sequence (AdvertCheck). When a node ives tx acceptor address via rps call (OnTxAcceptorAdded), it saves the address to the database (UpdateTxAcceptorAddressDB()). If the node is the delegate in the current epoch then it expertises tx acceptor address. When a node receives tx acceptor address deletion rpc call then node advertises tx acceptor deletion via P2p subsystem. The node deletes tx acceptor from the database. Sequence diagram for the address advertisement is shown on Figure 2.

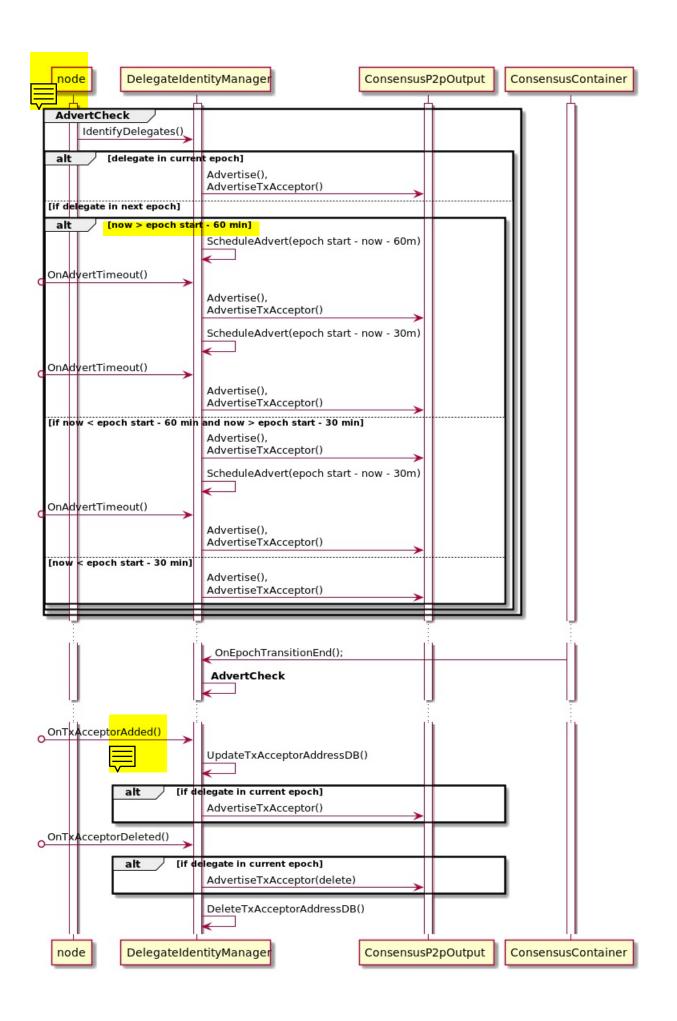


Figure 2. Sequence diagram, address advertisement

# Sequence Diagram Delegate Address and Tx Acceptor Address Receiving

On start up the node ascertains if it is the delegate in the current epoch or it is the delegate in the next epoch and current time is less than epoch start - 1 hour (IdentifyDelegates()). If it is the delegate and it dogsp\; have addresses of all other delegates then the delegate requests all missing addresses this peers via P2p subsystem (RequestDelegateAddresses()) and schedules 10 min timer (SchedulePeerAddrRequest()). When the delegate receives the response back from its peers (OnPeerAddrRequest()), it updates delegates address (UpdateDelegateAddressDB()) and connects to the received peers address (ConnectToPeer()) if it is the delegate in the current epoch. On timeout (OnPeerAddrRequestTimeout()) the delegate repeats the sequence (CheckDelegateAddress). When a node receives a delegate's or tx acceptor address advertisement (OnAddressAdMessage(), OnAddressAdTxAcceptorMessage()), and if the node is the delegate in current or next epoch then it updates the address database. If the node is the delegate in the current epoch then it connects to the received delegate address (ConnectToPeer()). On epoch transition end (OnEpochTransitionEnd()) the node telegate and tx acceptor advertisment messages that are older than the current epoch (DeleteDelegateAddressDB(), DeleteTxAcceptorAddressDB()). Sequence diagram for the address receiving is shown on Figure 3.

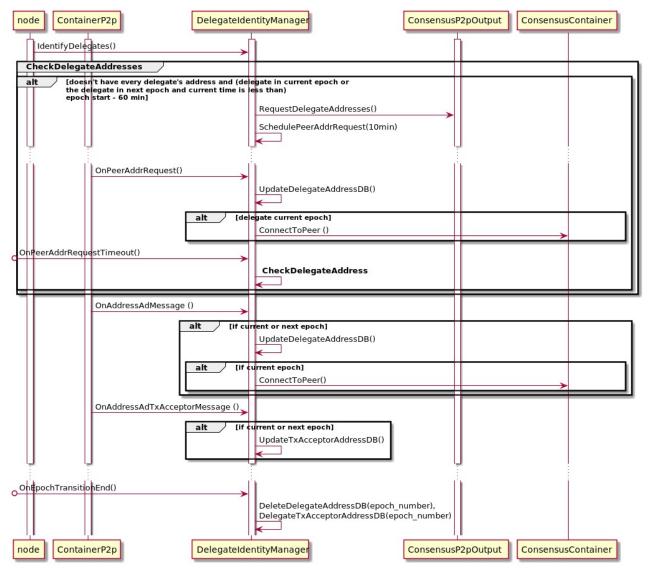


Figure 3.Sequence diagram, address receiving.

#### **Database**

database id to store the delegate's address is delegate\_address\_db and is keyed by the gate's ip, port, and epoch number. The value is blank. The database id to store the delegate's cceptor is txacceptor\_address\_db is keyed by the delegate's id, tx acceptor ip, json port, binary port, and epoch number. The value is blank. The keys serialization is described in the Table 1 and 2.

| Field Name |Size (Byte)| Description | | --- | -------| ------| | delegate id | 1 | Delegate id | IP | 39 | ip as the string | Port | 2 | Binary port | Epoch number | 4 | Binary epoch number |

Table 1. Delegate address key.

| Field Name |Size (Byte)| Description | | --- | -------| ------- | | delegate id | 1 | Delegate id | IP | 39 | ip as the string | Json port | 2 | Binary formatted port | Binary port | 2 | Binary formatted port | Epoch number | 4 | Binary formatted epoch number |

Table Acceptor address key.