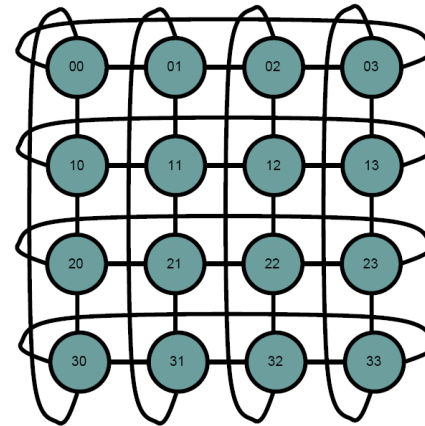


Direct and Indirect Networks

■ Direct Network

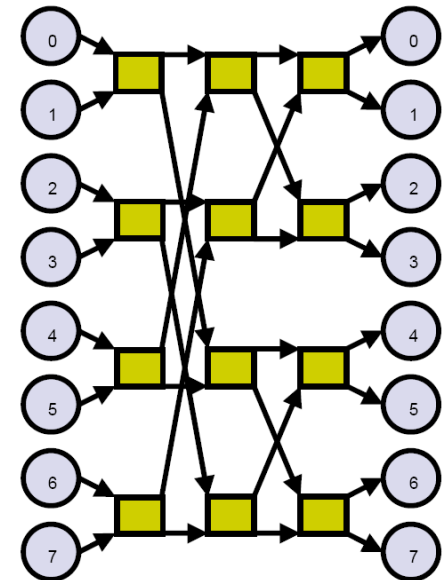
→ Every Node in the network is both a terminal and a switch



Direct Network

■ Indirect Network

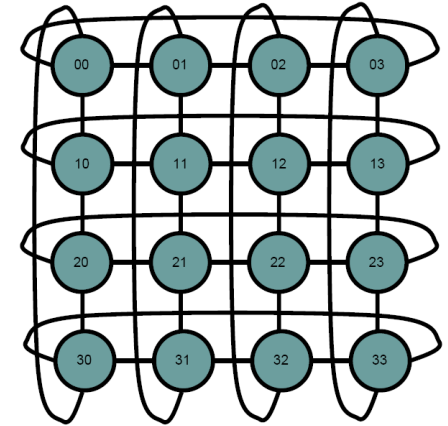
→ Nodes are either switches or terminal



Indirect Network

Direct Networks

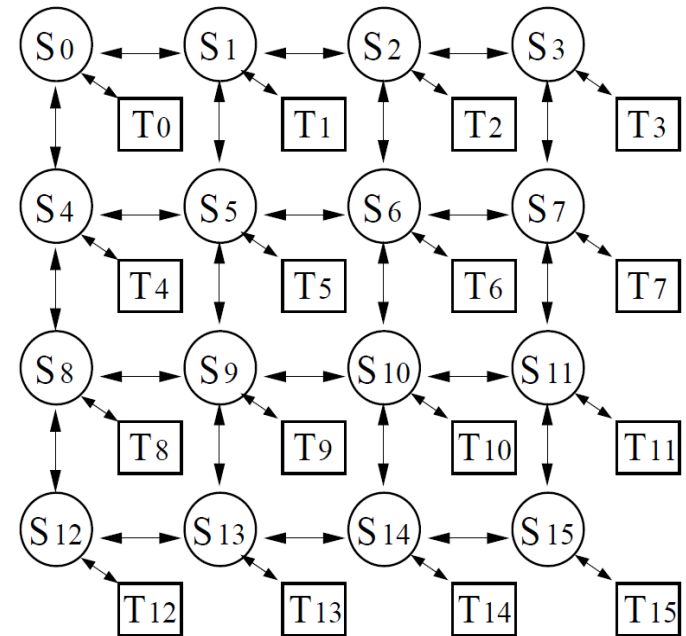
- aka point-to-point network
- Consists of a set of nodes, each one being directly connected to a (usually small) subset of other nodes in the network
 - These nodes may have different functional capabilities
 - ✓ *E.g., vector processors, graphics processors, I/O processors, etc.*



Direct Network

Direct Networks - Router

- A common component of the node is *the router*
 - It handles message communication among nodes
 - ✓ For this reason, direct networks are also known as router-based networks
 - Each router has direct connections to the router of its neighbors



Direct Networks - Links

- Two neighboring nodes are connected by a pair of unidirectional channels in opposite directions
- A bidirectional channel may also be used to connect two neighboring nodes

Direct Networks - Scalability

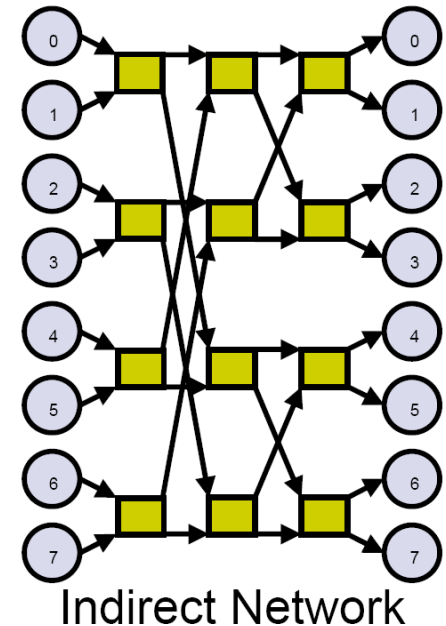
- As the number of nodes in the system increases, the total communication bandwidth also increase
 - ➔ Thus, direct networks have been a popular interconnection architecture for constructing large-scale parallel computers

Direct Networks - Topologies

- Many network topologies have been proposed in terms of their **graph-theoretical properties**
 - Very few of them have ever been implemented
 - Most of the implemented networks have an *orthogonal topology*

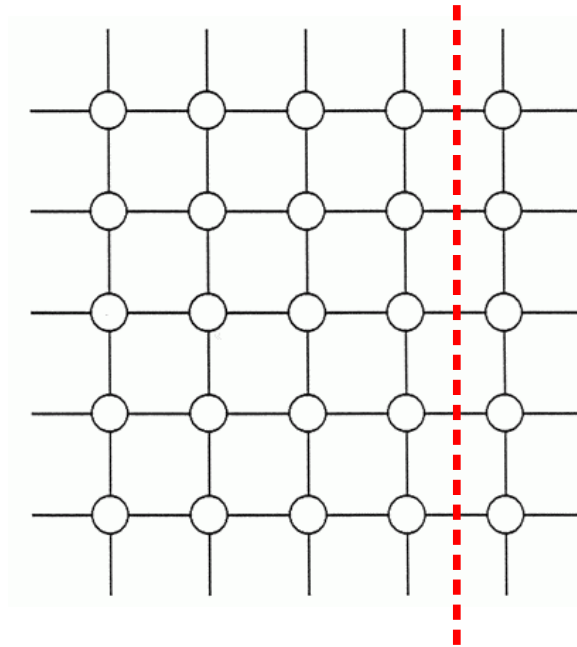
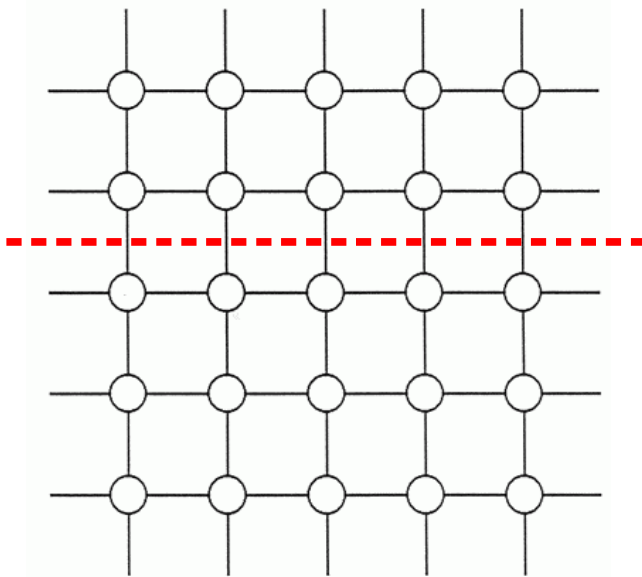
Indirect Networks

- The communication between any two nodes is carried through some switches
- Each node has a network adapter that connects to a network switch
- The interconnection of those switches defines various network topologies



Cuts

- A **cut** of a network, $C(N_1, N_2)$, is a set of channels that partitions the set of all nodes into two disjoint sets, N_1 and N_2
 - Each element in $C(N_1, N_2)$ is a channel with a source in N_1 and destination in N_2 or vice versa



Bandwidth of the Cut

- Total *bandwidth of the cut* $C(N_1, N_2)$

$$B(N_1, N_2) = \sum_{c \in C(N_1, N_2)} b_c$$

Bisection

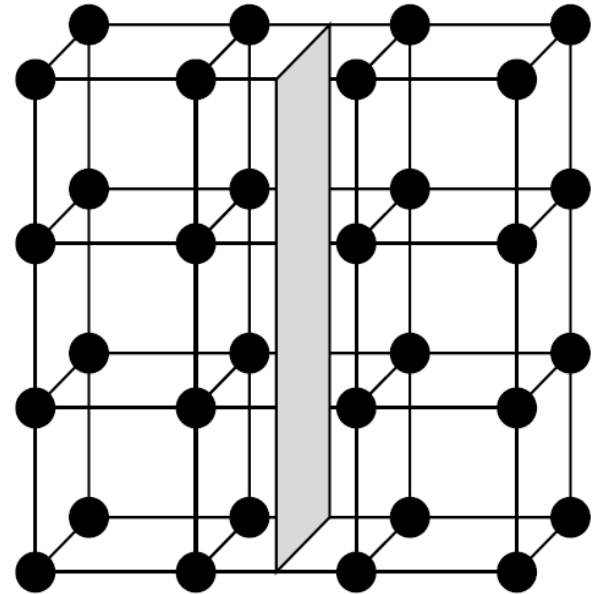
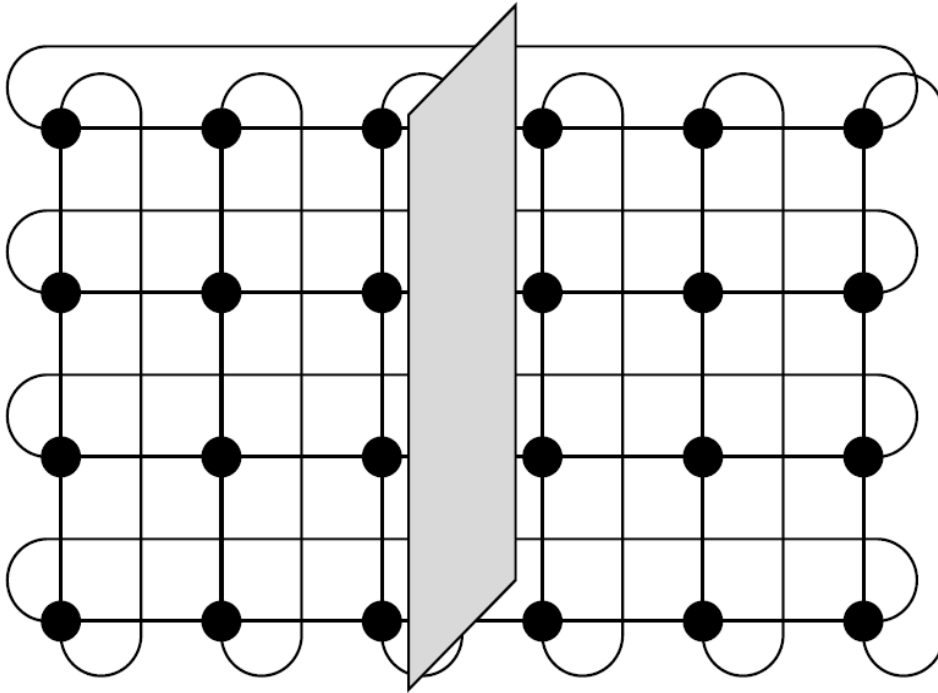
- The **bisection** is a cut that partitions the entire network nearly in half
- The **channel bisection** of a network, B_C , is the minimum channel count over all bisections

$$B_C = \min_{\text{bisections}} |C(N_1, N_2)|$$

- The **bisection bandwidth** of a network, B_B , is the minimum bandwidth over all bisections

$$B_B = \min_{\text{bisections}} |B(N_1, N_2)|$$

Bisection Examples



Diameter

- The **diameter** of a network, H_{\max} , is the largest, minimal hop count over all pairs of terminal nodes

$$H_{\max} = \max_{x,y \in N} |H(x,y)|$$

For a fully connected network with N terminals built from switches with out degree δ_o , H_{\max} is bounded by

$$H_{\max} \geq \log_{\delta_o} N \quad (1)$$

Each terminal can reach at most δ_o other terminals after one hop

At most δ_o^2 after two hops, and at most δ_o^H after H hops

If we set $\delta_o^H = N$ and solve for H , we get (1)

Average Minimum Hop count

- The ***average minimum hop count*** of a network, H_{\min} , is defined as the average hop count over all sources and destinations

$$H_{\min} = \frac{1}{N^2} \sum_{x, y \in N} H(x, y)$$