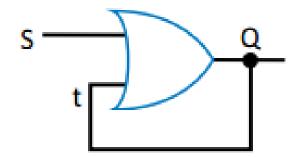
Latch & Flip flop

Sequential Circuit

- Output depends on both present inputs and history of inputs and outputs.
- It is combinational circuit with feedback



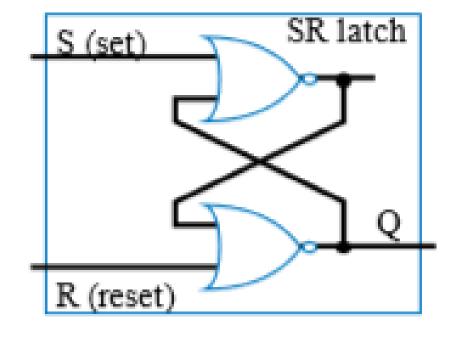
Latch & Flip-flop

- Building blocks that can store value of a bit
- Latch is level sensitive
- Filp-flop is edge sensitive

SR Latch

- S=1: set Q to 1
- R=1: reset Q to 0

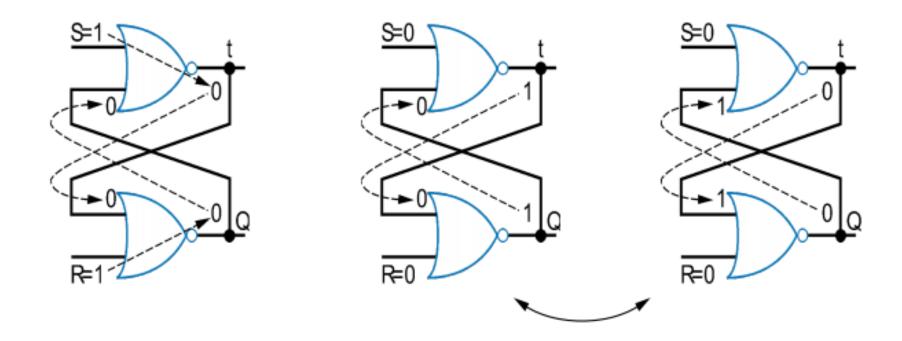
S(t)	R(t)	Q(t)	Q(t	+Δ) — Q+
0	0	0	0	hold
0	0	1	1	noid
0	1	0	0	reset
0	1	1	0	reset
1	0	0	1	set
1	0	1	1	301
1	1	0	Х	not allowed
1	1	1	X	



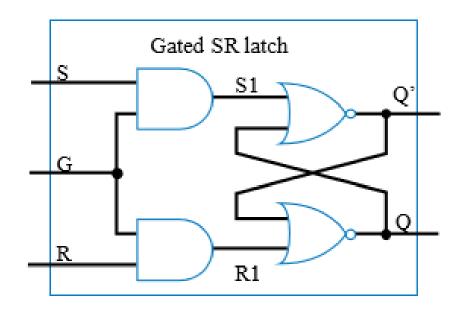
$$Q^+ = S + R'Q$$

SR Latch

• When S=1 & R=1, Q may oscillate when they both return to 0 simultaneously.



Gated SR Latch

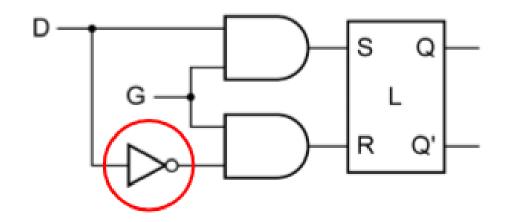


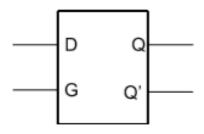
Characteristic Table

G	S	R	Q⁺
0	Х	Х	Q; Latch locked
1	0	0	Q; Hold state
1	0	1	0; Reset state
1	1	0	1; Set state
1	1	1	not allowed

Gated D Latch

- No unstable state as in SR latch
- Q = D when G = 1
- Q holds when G = 0





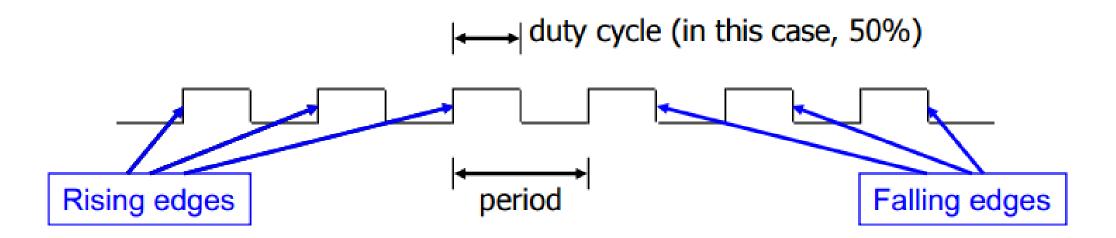
D latch symbol

Characteristic Table

G	D	Q+
1	0	0
1	1	1
0	Χ	Q

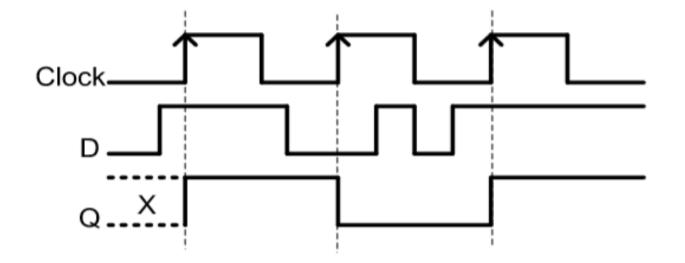
Clock Signal

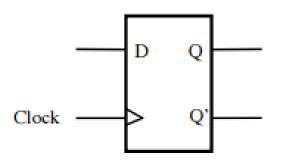
- Periodic pulse train
- Clock period: time interval between pulses



Rising-Edge Triggered D Filp Flop

Edge sensitive

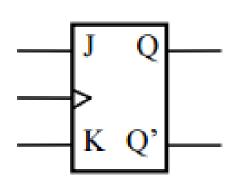




clock	D	Q ⁺
	0	0
	1	1
0	Χ	Q
1	Χ	Q

Characteristic equation: Q+ = D (at active clock edges)

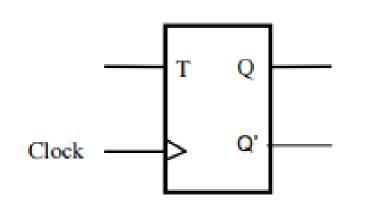
Rising-Edge Triggered J-K Filp Flop



J	K	Q+
0	0	Q
0	1	0
1	0	1
_1	1	Q'

Characteristic equation: Q+ = JQ' + K'Q

Rising-Edge Triggered T Filp Flop



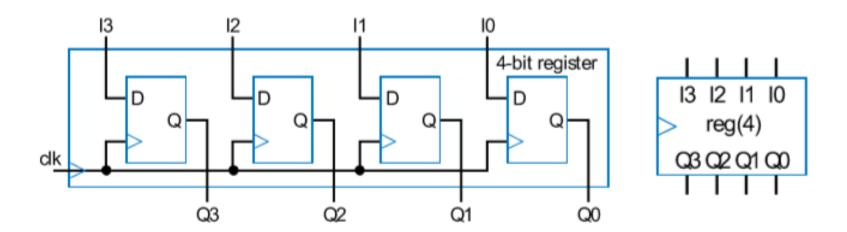
clock	Т	Q⁺	
	0	Q	
	1	Q'	

Characteristic equation:

$$Q^+ = T'Q + TQ' = T \oplus Q$$

Register

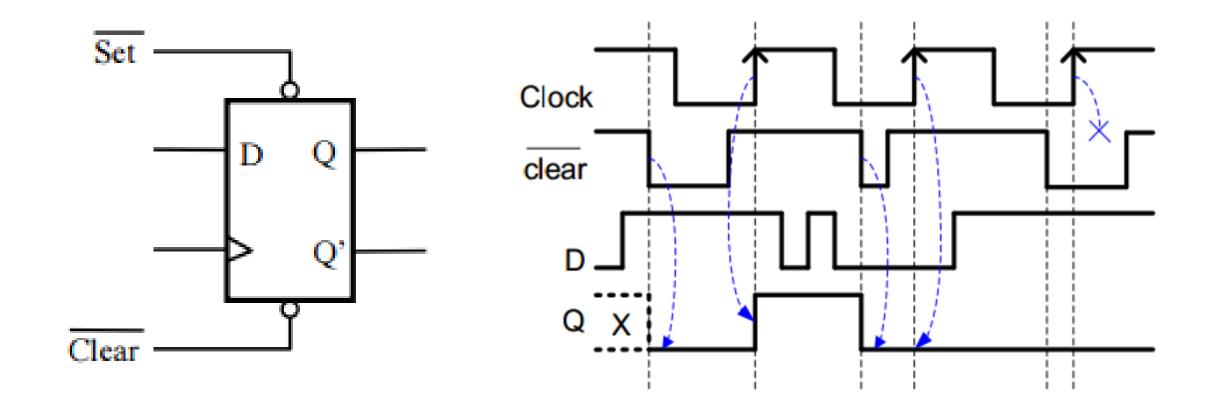
- Multiple filp-flops sharing clock signal
- Store muliple bits
- Load values simultaneously at rising edge



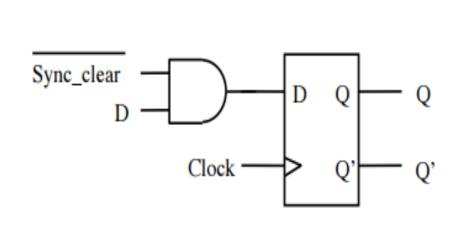
Control Inputs for Flip Flops

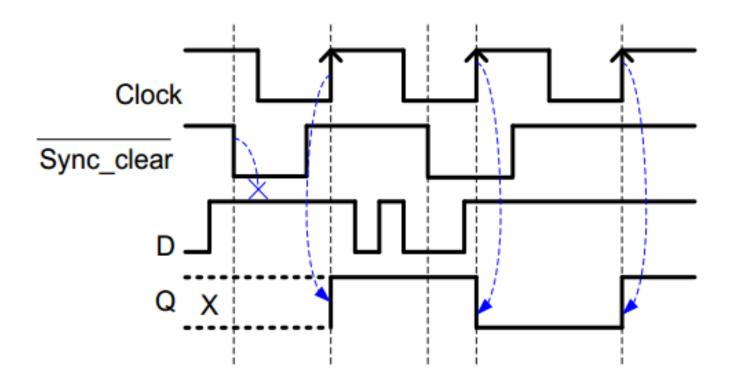
- Synchronous/asynchronous: depends on the clock signal or not
- Active low/high: controls when it's low/high

D flip flop with active low asynchronous Clear



D flip flop with active low synchronous Clear





Verilog

Verilog HDL

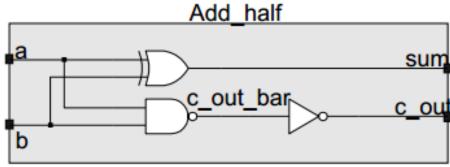
Syntax:

- module
- always
- if···else···
- switch
- assignment

Tips:

- [N-1:0] represents N bits
- Destination variables inside always must be reg type

Module



```
module ports
  module name
module Add half (sum, c out, a, b);
 input
            a, b;
                                   declaration of port modes
 output sum, c out;

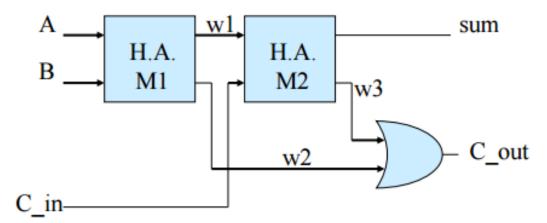
    declaration of internal signal

            c out bar; -
 wire
             (sum, a, b);
 xor
                                             instantiation of pre-defined
             (c out bar, a, b);
 nand
                                             primitive gates
             (c out, /c out bar);
 not
endmodule
                Same variable indicates connection
```

Module

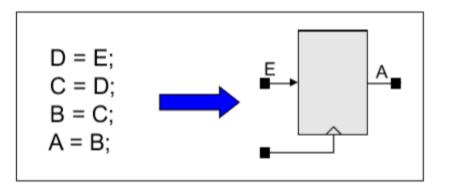
```
module Add_full (sum, c_out, a, b, c_in); // parent module
input a, b, c_in;
output c_out, sum;
wire w1, w2, w3;

Add_half M1 (w1, w2, a, b); // child module
Add_half M2 (sum, w3, w1, c_in); // child module
or (c_out, w2, w3); // primitive instantiation
endmodule
```

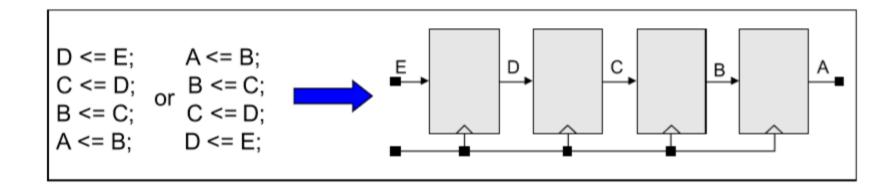


Assignment

Assign one statement by one statement



<= Assign simultaneously



```
Testbench
                           module Test Banch;
                             parameter half period = 50;
                             parameter counter size = 4;
Outputs of the module •
                             wire [counter size-1:0] Q;
                                  [counter size-1:0] Din;
                             req
Inputs of the module 

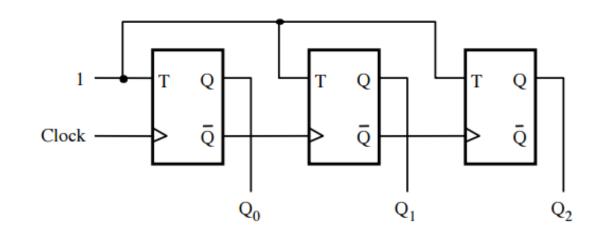
                                  clock, load, reset, CE;
                             rea
Module to be tested
                             counter N bit #(counter size) UUT (clock, reset, load, CE, Din, Q, CEO);
                             initial begin
                               #0
                                     clock = 0; Din = 0; load = 0; CE = 1; reset = 1;
                              #100 reset = 0;
                              #200 Din = 8; load = 1;
                              #100 load = 0;
Changes of inputs
                              #300 CE = 0;
                              #200 CE = 1;
                             end
                             always #half period clock = ~clock;
                             initial #2000 $stop;
```

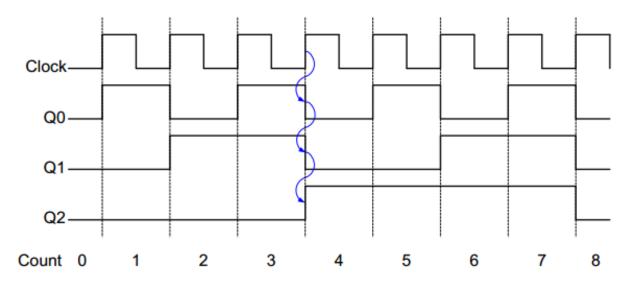
endmodule

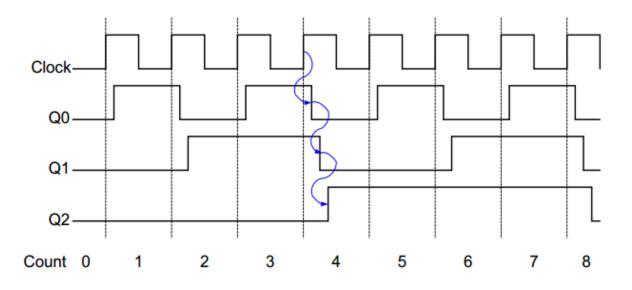
Counter

Asynchronous Counter

- Flip flops are cascaded together
- Delays → asynchronous output







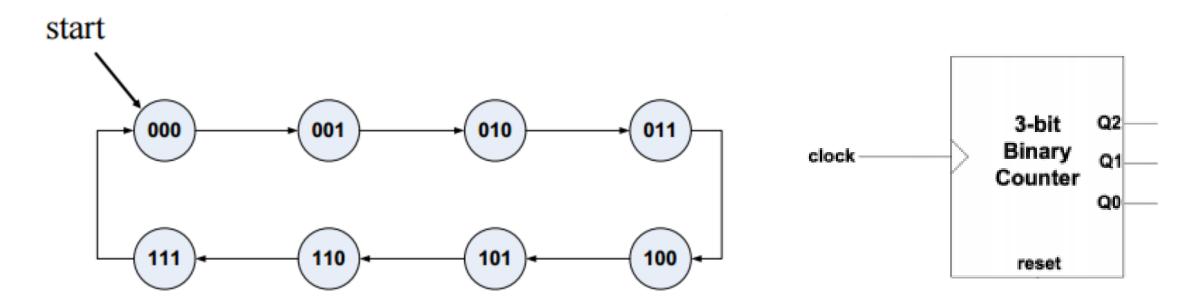
Ideal

Practical

Synchronous Counter

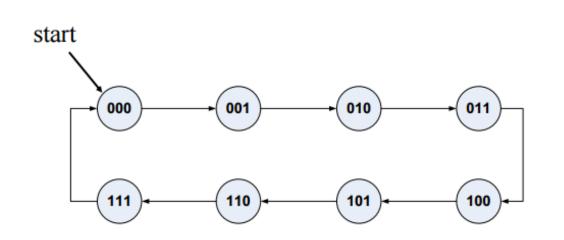
- All the filp-flops are triggered by the same clock
- May be implemented by different type of filp flops

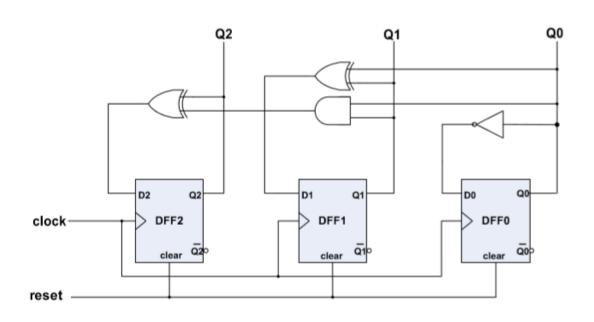
Synchronous 3-bit Binary Counter



How to design a synchronous counter given a sequence of numbers?

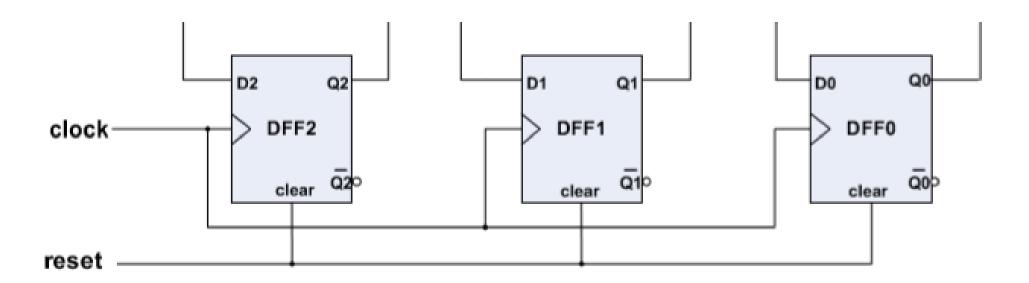
- Combine combinational circuit and sequential circuit
- Steps:
 - 1. Decide what kinds of flip flops to use and how many
 - 2. Draw truth table
 - 3. Derive logic equations
 - 4. Draw combinational circuit



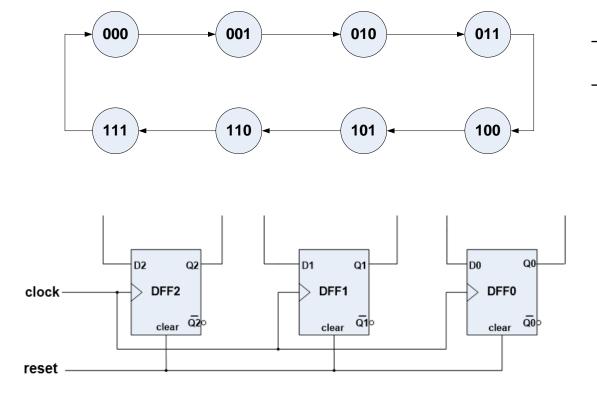


Step 1: FF

- 3-bit counter → 3 flip flops
- Here we choose D filp flops



Step 2: Draw Truth Table



Present State		Next State			D flip flop input			
Q2	Q1	Q0	Q2+	Q1+	Q0+	D2	D1	D0
0	0	0	0	0	1	0	0	1
0	0	1	0	1	0	0	1	0
0	1	0	0	1	1	0	1	1
0	1	1	1	0	0	1	0	0
1	0	0	1	0	1	1	0	1
1	0	1	1	1	0	1	1	0
1	1	0	1	1	1	1	1	1
1	1	1	0	0	0	0	0	0

 $Q^+ = D$ upon active edge

Step 3: Derive Logic Equations

Present State		Next State			D flip flop input			
Q2	Q1	Q0	Q2+	Q1+	Q0+	D2	D1	D0
0	0	0	0	0	1	0	0	1
0	0	1	0	1	0	0	1	0
0	1	0	0	1	1	0	1	1
0	1	1	1	0	0	1	0	0
1	0	0	1	0	1	1	0	1
1	0	1	1	1	0	1	1	0
1	1	0	1	1	1	1	1	1
1	1	1	0	0	0	0	0	0

$$D1 = Q1'Q0+Q1Q0'$$
$$= Q1 \oplus Q0$$

$$D0 = Q0$$

Step 3: Draw Combinational Circuit

$$D2 = Q2Q1'+Q2Q0'+Q2'Q1Q0$$

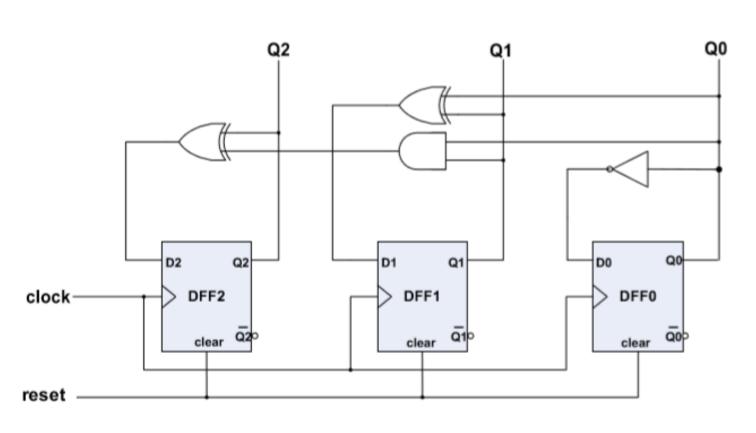
$$= Q2(Q1'+Q0')+Q2'Q1Q0$$

$$= Q2(Q1Q0)'+Q2'(Q1Q0)$$

$$= Q2 \oplus (Q1Q0)$$

$$D1 = Q1'Q0+Q1Q0'$$
$$= Q1 \oplus Q0$$

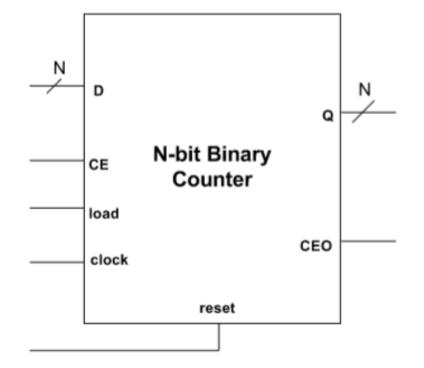
$$D0 = Q0$$



Control Signals

Priority of External Control Signals:

reset	load	CE	Action on the rising clock edge
1	X	×	Clear (Qn <= 0)
0	1	X	Load (Qn <= Dn)
0	0	1	Count
0	0	0	Hold (No Change)



Count Enable Output:

CEO = CE
$$\cdot Q_{N-1} \cdot Q_{N-2} \cdot ... \cdot Q_0$$

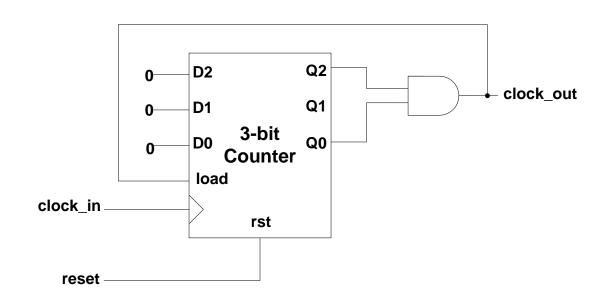
Clock Divider

Divide by N:

- $2^{n-1} < N < 2^n$, where n is the size of counter
- N-1 is the upper bound of the counter's sequence
- AND the bits that are 1 in the upper bound to the load

Divide by 6:

- $2^2 < 6 < 2^3 \rightarrow 3$ -bit counter
- Countering sequence: $000 \rightarrow 001 \rightarrow 010 \rightarrow 011 \rightarrow 100 \rightarrow 101 \rightarrow 000$
- Upper bound: 101 → load = Q0 AND A2



GOOD LUCK