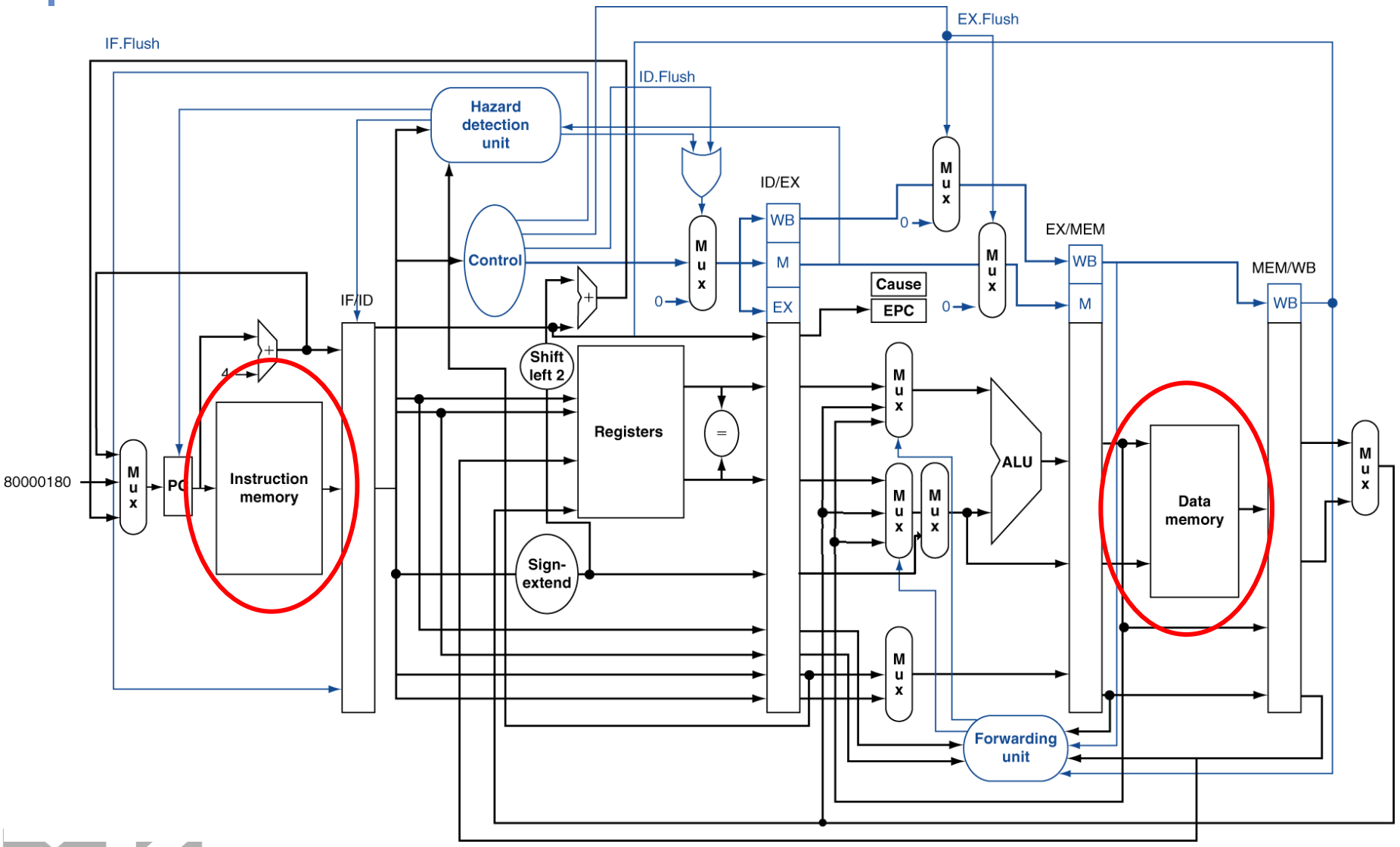
Topic 11

Memory Hierarchy

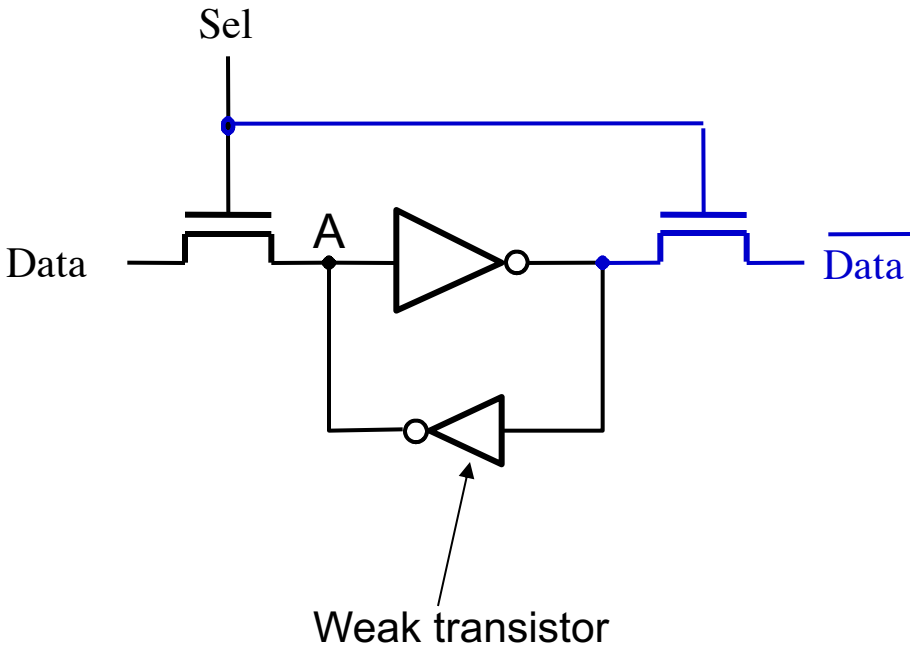
- Cache (1)

MIPS Pipeline Architecture



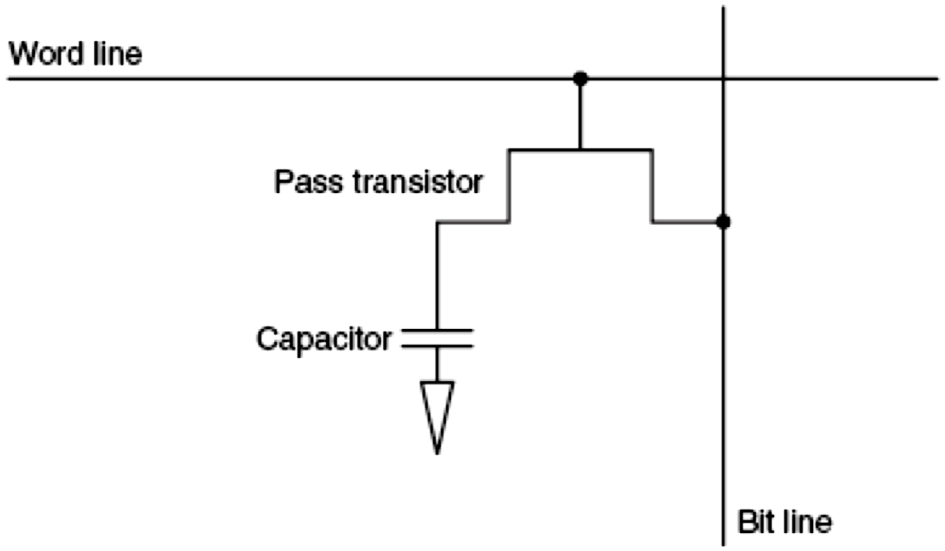
Static RAM (SRAM)

- When Sel = 1, Data is stored and retained in the SRAM cell by the feedback loop
- When Sel = 1, Data can be read out on the same port
- Point A is driven by both the Data transistor and the smaller inverter, but the Data transistor wins because the inverter is implemented using a weak transistor



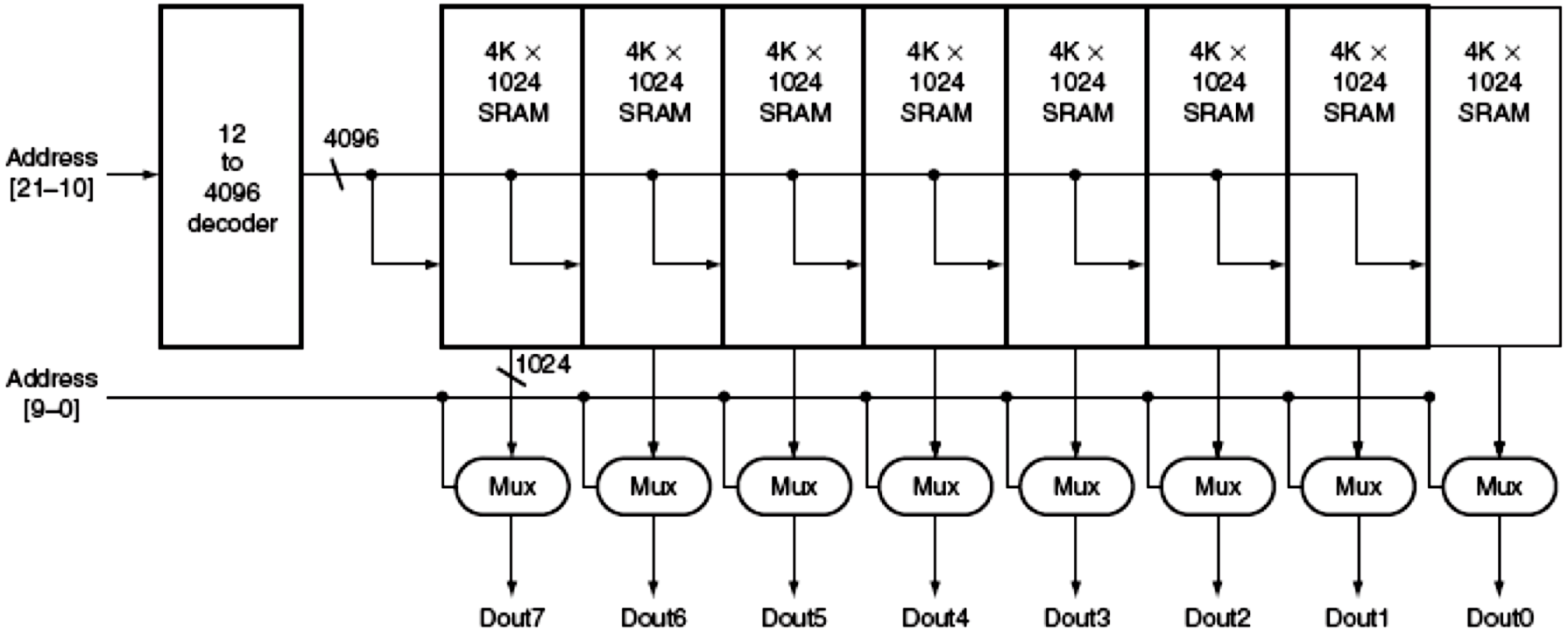
Dynamic RAM (DRAM)

- Write: turn on word line, charge capacitor through pass transistor by bit line
- Read: charge bit line halfway between high and low, turn on word line, then sense the voltage change on bit line
  - 1 if voltage increases
  - 0 if voltage decreases



Memory

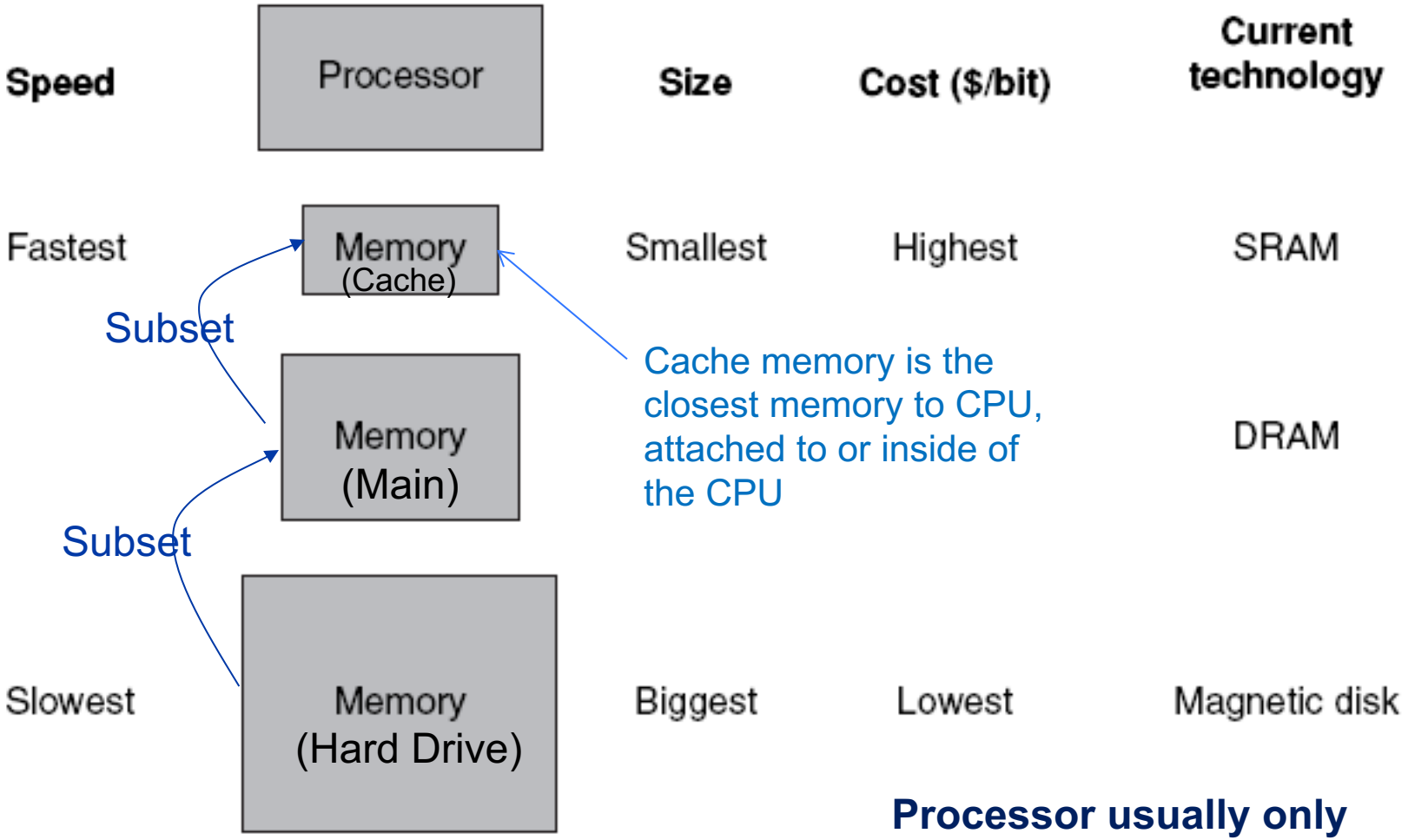
- Typical memory organization



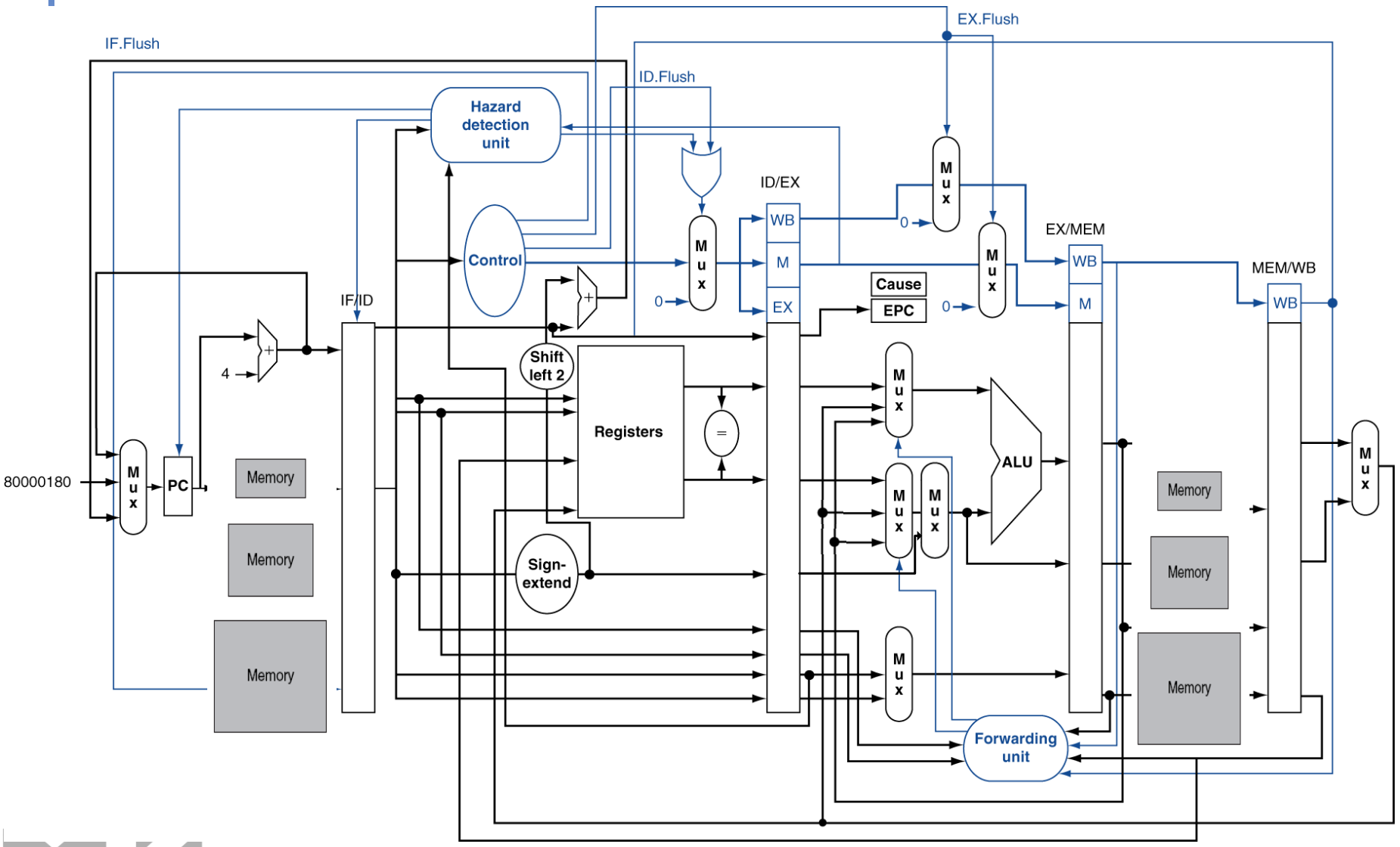
Memory Technology

- Static RAM (SRAM)
  - 0.5ns – 2.5ns, $2000 – $5000 per GB
- Dynamic RAM (DRAM)
  - 50ns – 70ns, $20 – $75 per GB
- Magnetic disk
  - 5ms – 20ms, $0.20 – $2 per GB
- Ideal memory
  - Access time of SRAM
  - Capacity and cost of disk

Memory Hierarchy



MIPS Pipeline Architecture



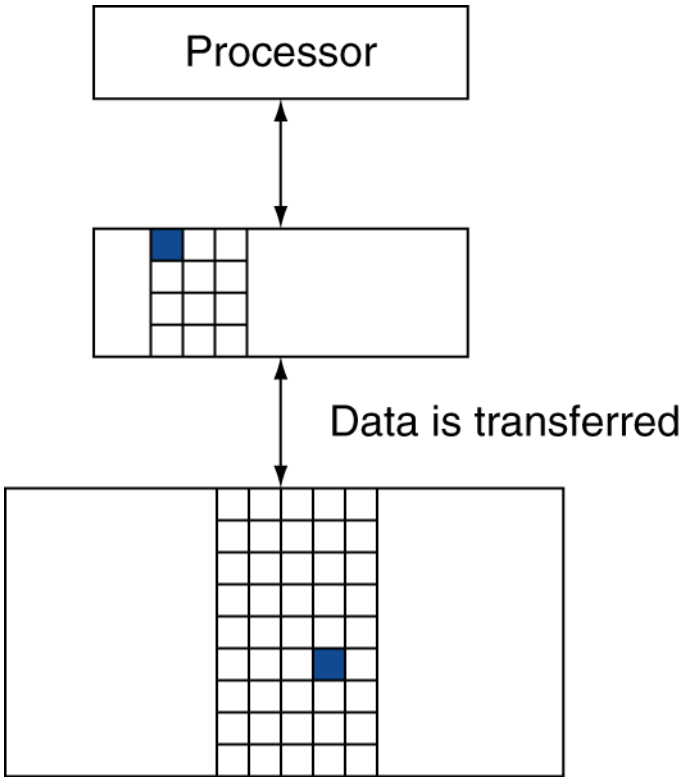
Principle of Locality

- Programs access a small proportion of their address space at any time
- Temporal locality
  - Items that are accessed recently are likely to be accessed again soon
  - e.g., instructions in a loop
- Spatial locality
  - Items near those that are accessed recently are likely to be accessed soon
  - E.g., sequential instruction access, array data

Principle of Memory Access

- Taking Advantage of Locality
- If a word of data/instruction is referenced
  - Copy this recently accessed word (temporal locality) and nearby items (spatial locality) together as a block from lower level memory to higher level memory (Hard Drive to Main memory or Main memory to Cache)

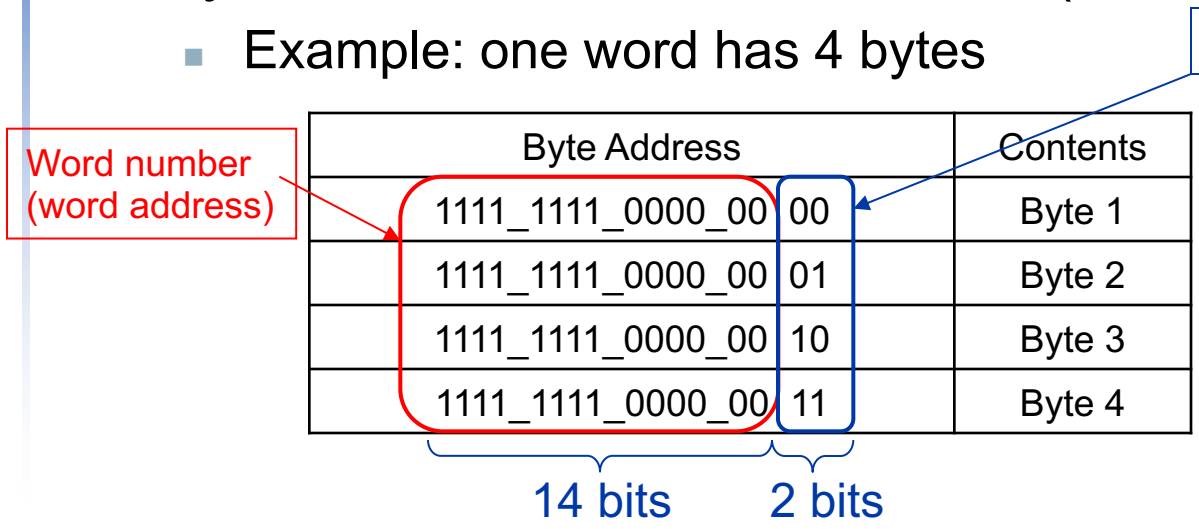
Access the Memory Hierarchy



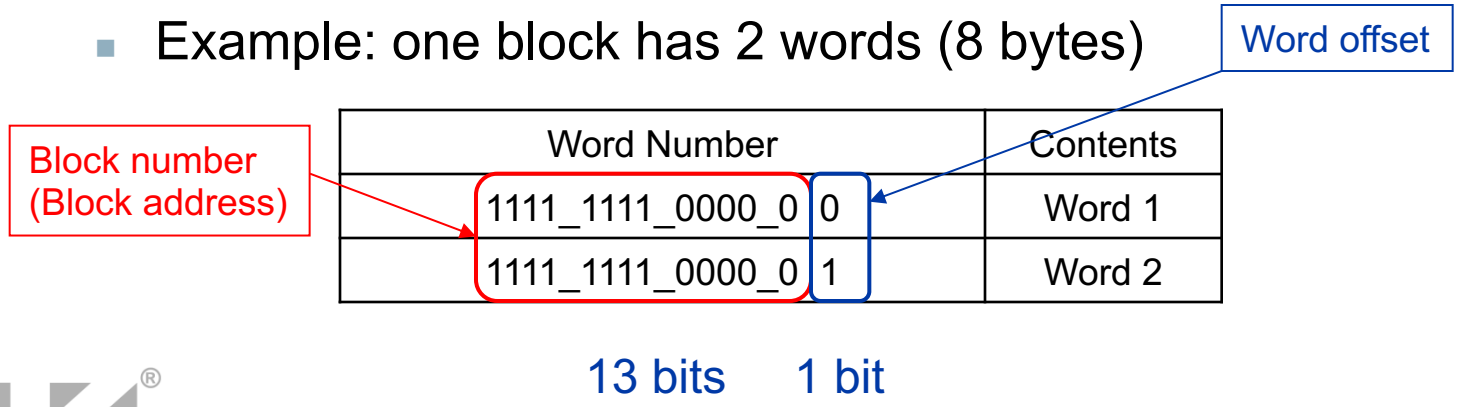
- Concepts:
  - Block (aka line)
    - Unit of data referencing to take advantage of temporal and spatial locality
    - May be one or multiple words
    - Block also has an address

Block, Word, Byte Addresses

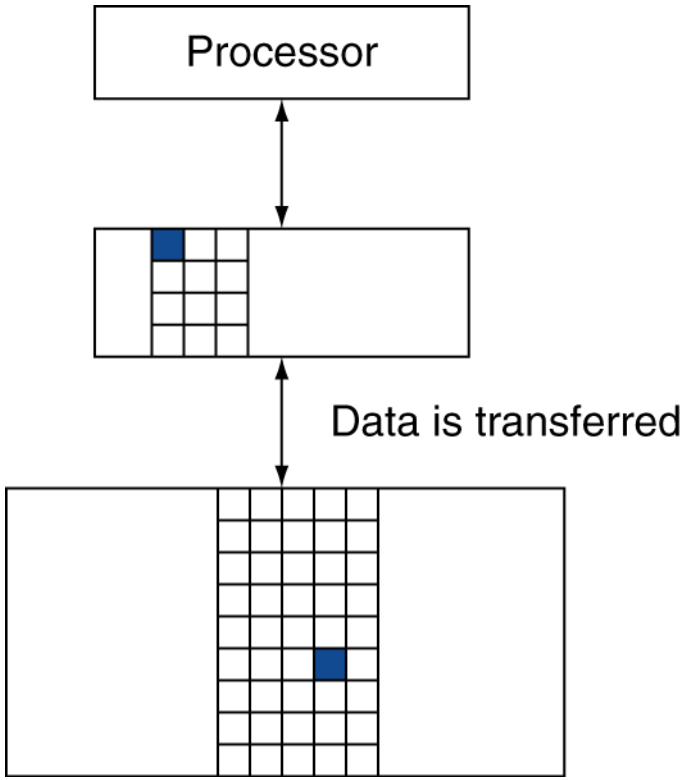
- Byte address vs. word address (assume 16-bit address)
  - Example: one word has 4 bytes



- Word address vs. block address (block number)
  - Example: one block has 2 words (8 bytes)

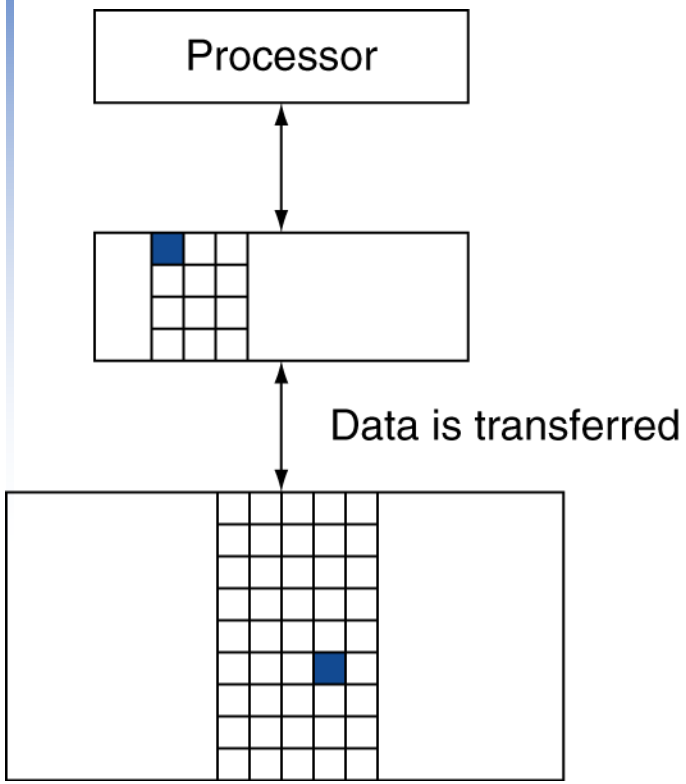


Access the Memory Hierarchy



- Concepts:
  - Hit
    - If accessed data is present in upper level, access satisfied by upper level
    - Hit rate: hits/accesses
  - Hit time: time to access a memory including
    - Time to determine whether a hit or miss
    - Time to pass block to requestor

Access the Memory Hierarchy



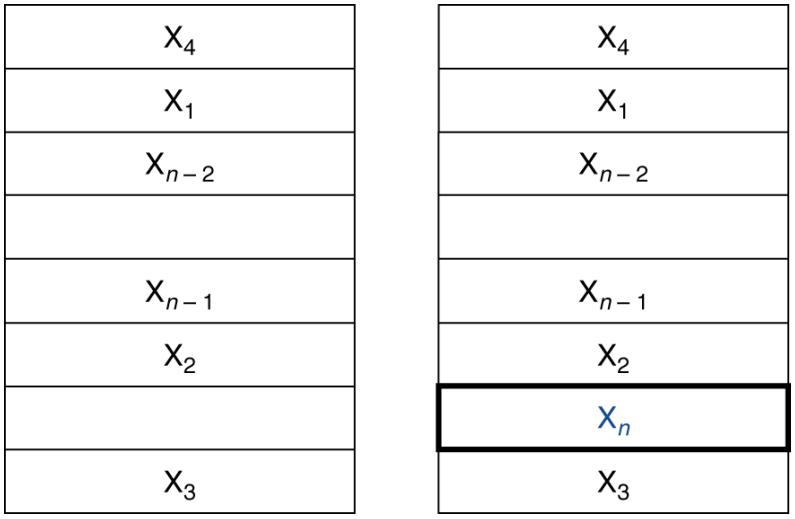
- Concepts:
  - Miss
    - If accessed data is absent
    - Block copied from lower to higher level
    - Then accessed data supplied from upper level
    - Time taken: miss penalty
    - Miss rate: misses/accesses = 1 – hit rate

Miss Penalty

- Miss (time) penalty
  - Time to fetch a block from lower level upon a miss including
    - Time to access the block
    - Time to transfer it between levels
    - Time to overwrite the higher level block
    - Time to pass block to requestor

Cache Miss

- Assume a cache with 1-word blocks X_1, …, X_{n-1}. Now CPU requests X_n



a. Before the reference to X_n b. After the reference to X_n

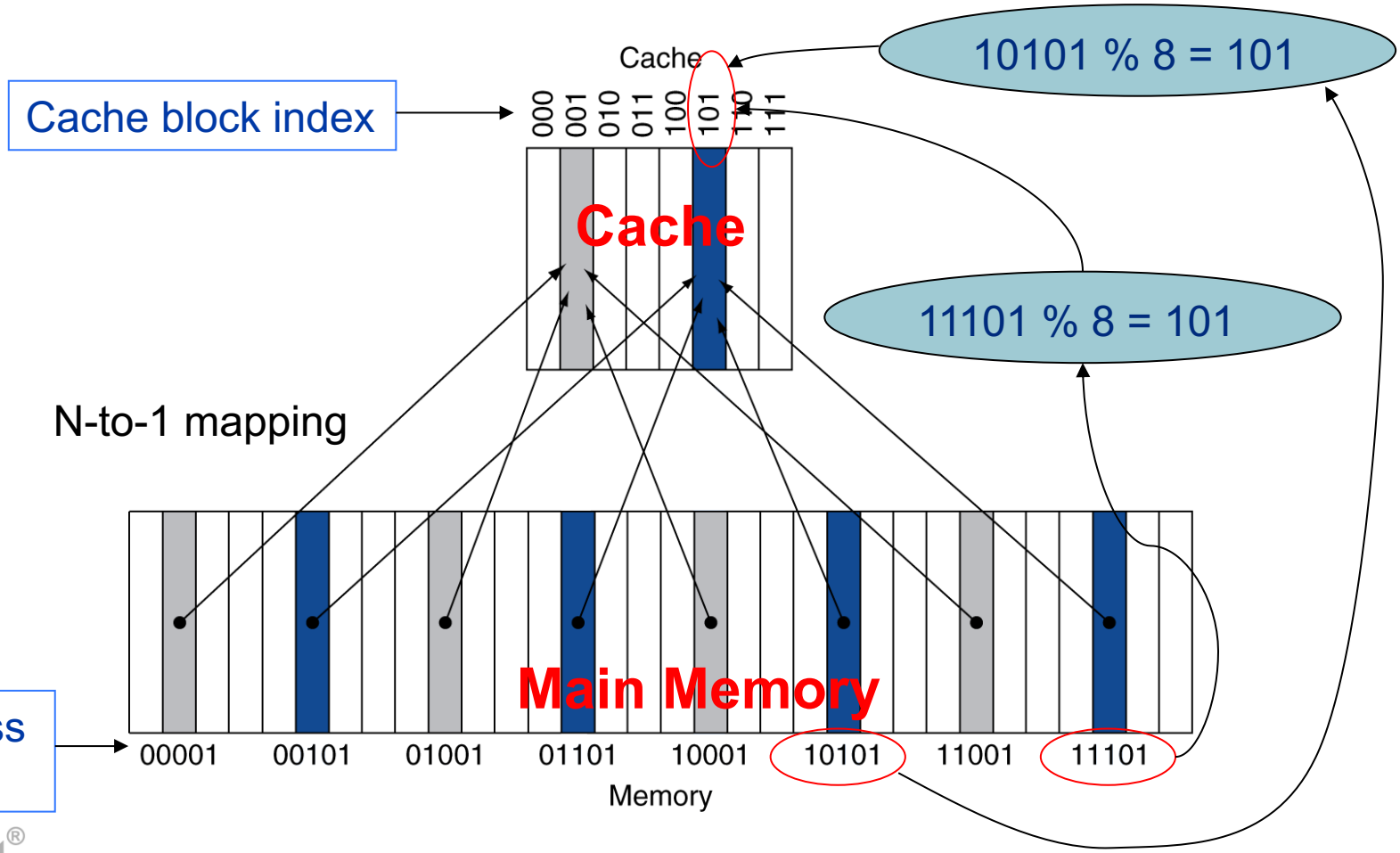
- Miss. Then X_n is brought in from lower level
- But
  - How do we know if there is a hit or miss?
  - Where do we look?

Direct Mapped Cache

- Direct mapped cache: each memory location corresponds to **one** choice in cache
- How to map?
- Location of a block in cache is determined by address of the requested word
  - Given word address, we can get block address (block number)
  - *Cache location (or cache block index) =*
    (Block address in memory) % (Number of blocks in cache) =
    lower $\log_2$(Number of blocks in cache) bits of block address
- **Multiple** memory locations correspond to **one** block in cache – n-to-1 mapping

Direct Mapped Cache

- Number of blocks in a cache is a power of 2
- Low-order bits of block address in memory = cache index

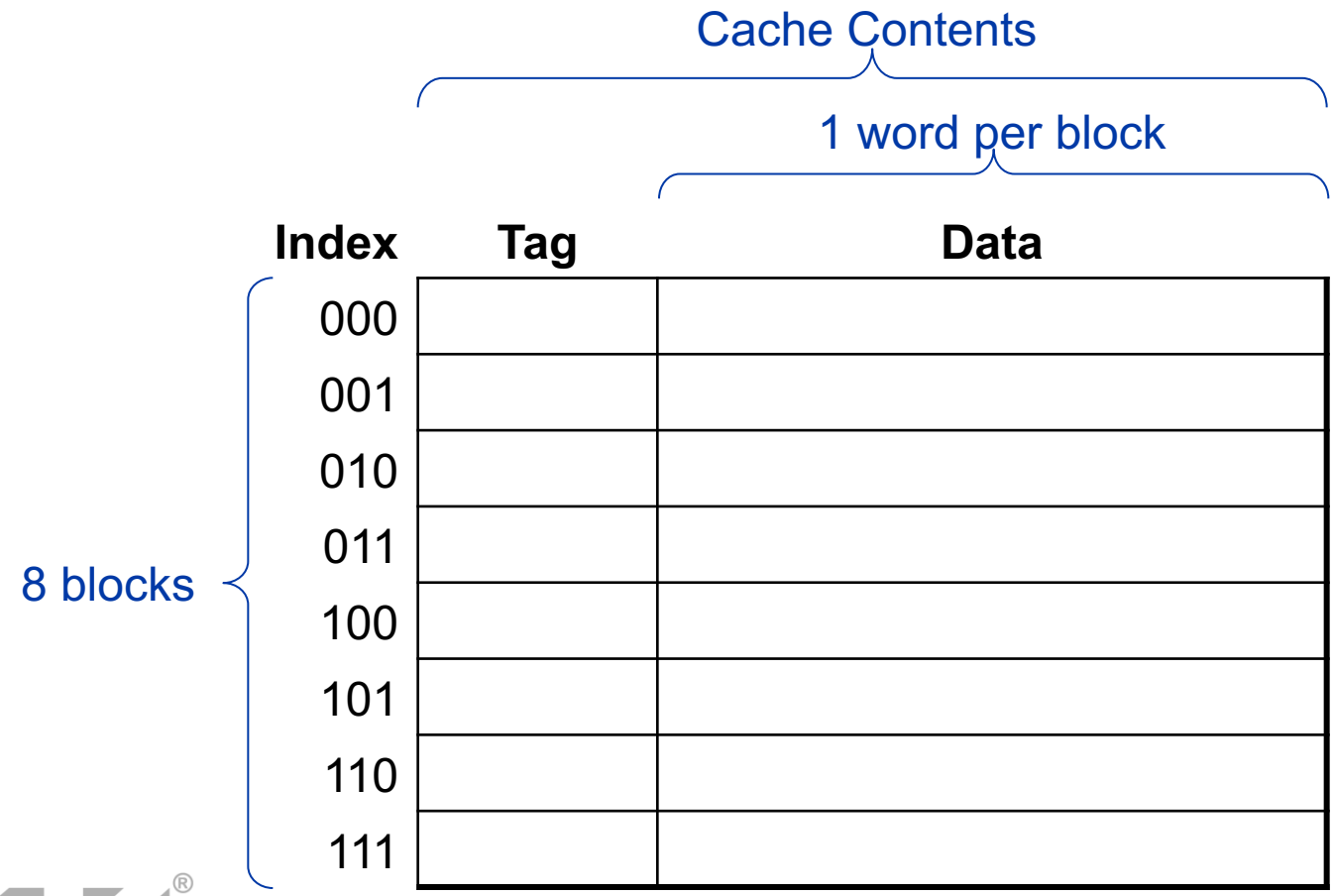


Tags in Cache

- How do we know which block of data is present in cache since it's n-to-1 mapping?
  - Store part of block address together with the data when the data is copied to higher level
    - Only need to store the high-order bits of memory block address, excluding the *lower log_2(Number of blocks in cache) bits of block address*
    - Called ***tag***

Direct Mapped Cache Example

- 8-block cache
- 1 word per block (word address = block address)

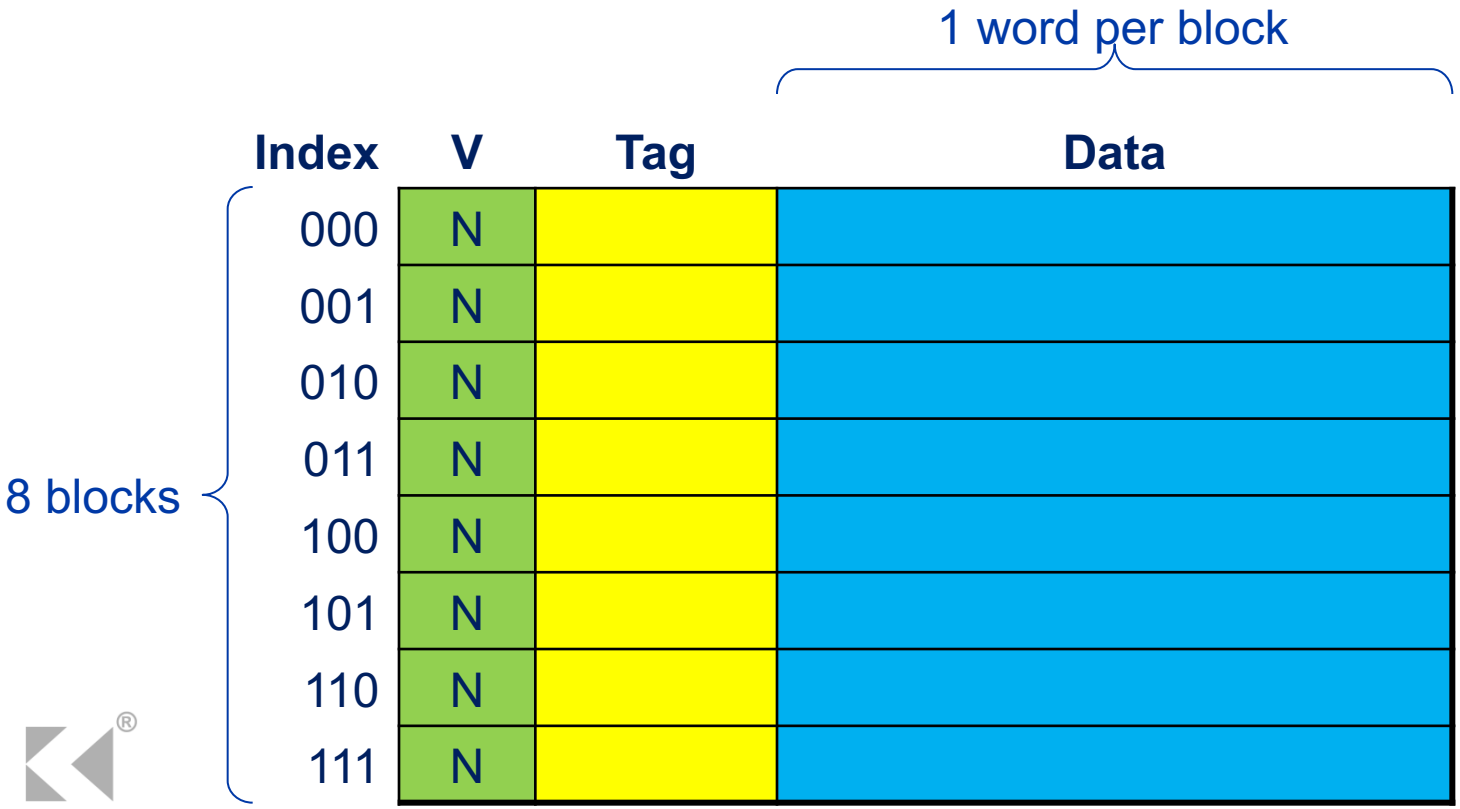


Valid Bits in Cache

- Waste searching time if there is no valid data in a location
  - Valid bit: 1 = present, 0 = not present
  - Initially 0 (N), to improve search speed

Cache Example 1

- 8-block cache
- 1 word per block (word address = block address)
- direct mapped
- Assuming 7-bit byte addresses



Cache Example 1 (cont.)

Word address without byte offset 00

lw R1 ← mem[22]			
Requested mem addr	Block (word) Addr	Hit/miss	Cache block
10110 00	10 110	Miss	110

Word addr (22)

Byte offset

Index	V	Tag	Data
000	N		
001	N		
010	N		
011	N		
100	N		
101	N		
110	**N**		
111	N		

Miss

Cache Miss

- On cache hit, CPU proceeds normally
- On cache miss
  - Stall the CPU pipeline
    - Using processor control unit and a cache controller
    - Freezes registers and waits for memory
  - Fetch block from next level of hierarchy

Cache Example 1 (cont.)

lw R1 ← mem[22]			
Requested mem addr	Block (word) Addr	Hit/miss	Cache block
10110 00	10 110	Miss	110

byte offset 00 are omitted

Index	V	Tag	Data
000	N		
001	N		
010	N		
011	N		
100	N		
101	N		
110	**Y**	**10**	**0x05ACF011**
111	N		

Cache memory

Word Addr	Data
00000(0)	0x81230431
00001(1)	0xABCD3305
...	...
10101(21)	0x12345678
10110(22)	**0x05ACF011**
...	...
11110(30)	0x00000000
11111(31)	0x00000000

Main memory

Cache Example 1 (cont.)

lw R2 ← mem[26]			
Requested mem addr	Block (word) Addr	Hit/miss	Cache block
11010 00	11 010	Miss	010

Index	V	Tag	Data
000	N		
001	N		
010	**Y**	**11**	**0x5DC60007**
011	N		
100	N		
101	N		
110	Y	10	0x05ACF011
111	N		

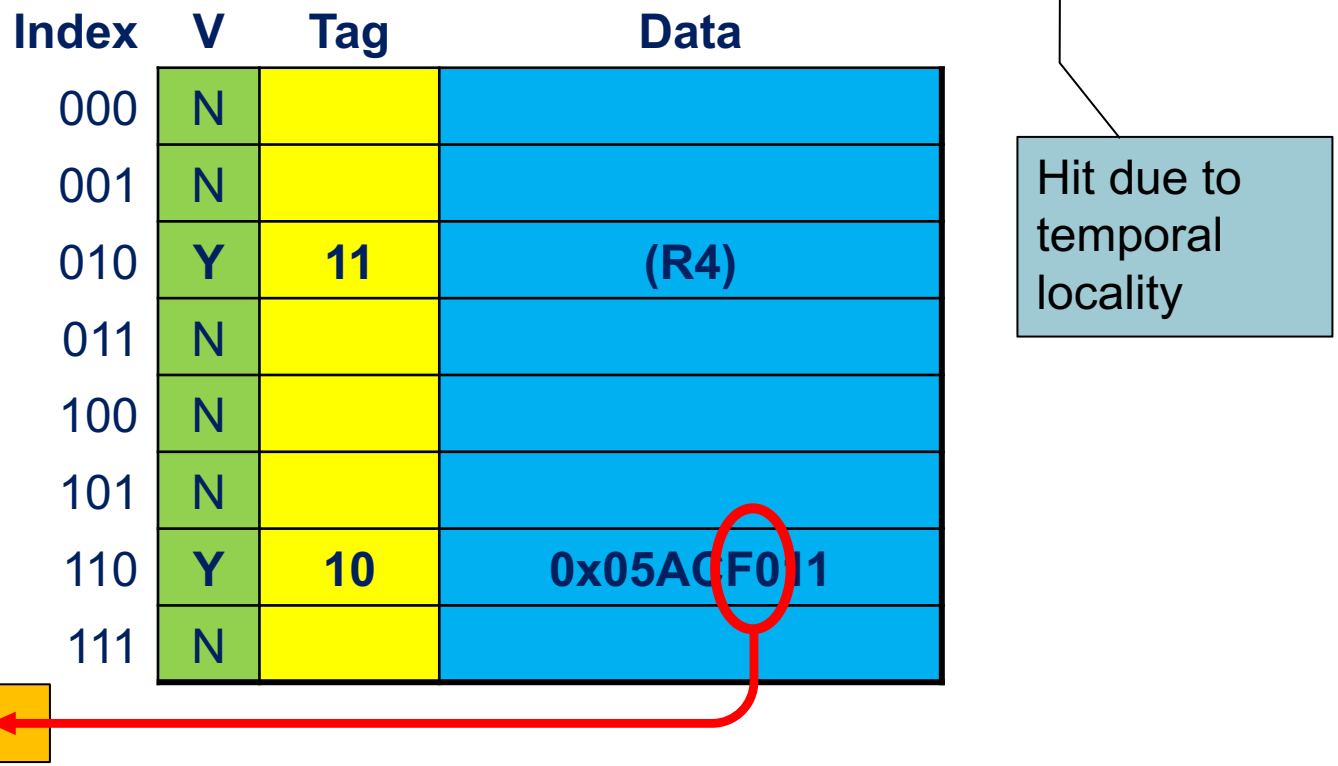
Cache memory

Word Addr	Data
00000 (0)	0x81230431
00001 (1)	0xABCD3305
…	…
10101 (21)	0x12345678
10110 (22)	0x05ACF011
…	…
11010 (26)	**0x5DC60007**
…	…
11110 (30)	0x00000000
11111 (31)	0x00000000

Main memory

Cache Example 1 (cont.)

	Requested mem addr	Block (word) Addr	Hit/miss	Cache block
`lb R3 ← mem[22]byte1`	10110 01	10 110	Hit	110
`sw R4 → mem[26]`	11010 00	11 010	Hit	010



Cache Example 1 (cont.)

Requested mem Addr	Block (word) Addr	Hit/miss	Cache block
1000000(lw)	10 000	Miss	000
0001100(sw R5)	00 011	Miss	011
1000000(lw)	10 000	Hit	000

Index	V	Tag	Data
000	**Y**	**10**	**0x8765ABCD**
001	N		
010	Y	11	(R4)
011	**Y**	**00**	**0xFFFF0002→(R5)**
100	N		
101	N		
110	Y	10	0x05ACF011
111	N		

Word Addr	Data
00000(0)	0x81230431
00001(1)	0xABCD3305
00010(2)	0xFFFF0001
00011(3)	**0xFFFF0002**
…	…
10000(16)	**0x8765ABCD**
…	…
10101(21)	0x12345678
10110(22)	0x05ACF011
…	…
11010(26)	0x5DC60007
…	…
11110(30)	0x00000000
11111(31)	0x00000000

Cache Example 1 (cont.)

Requested mem Addr	Block (word) Addr	Hit/miss	Cache block
1001000(lw)	10 010	Miss	010

Currently occupied by Mem[26]

Index	V	Tag	Data
000	Y	10	0x8765ABCD
001	N		
010	**Y**	**10**	**(R4)→0x0000FF00**
011	Y	00	0xFFFF0002→(R5)
100	N		
101	N		
110	Y	10	0x05ACF011
111	N		

Word Addr	Data
00000(0)	0x81230431
00001(1)	0xABCD3305
00010(2)	0xFFFF0001
00011(3)	0xFFFF0002
…	…
10000(16)	0x8765ABCD
10001(17)	0x12345678
10010(18)	**0x0000FF00**
…	…
11010(26)	0x5DC60007
…	…
11110(30)	0x00000000
11111(31)	0x00000000