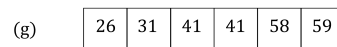
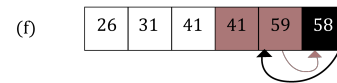
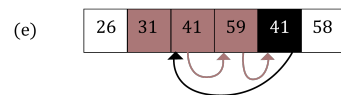
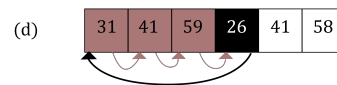
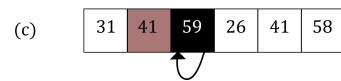
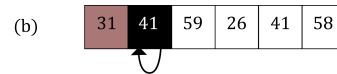
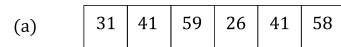


Chapter 2 Part 1 Exercises

2.1-1

The steps are the following:



2.1-2

The new instertion-sort will be the following:

```

REVERSE-INSERTION-SORT( $A$ )
1  for  $j = 2$  to  $A.length$ 
2      do
            $key = A[j]$ 
3       $i = j - 1$ 
4      while  $i > 0$  and  $A[i] < key$ 
5          do
                $A[i + 1] = A[i]$ 
6               $i = i - 1$ 
7       $A[i + 1] = key$ 

```

2.1-3

The linear search algorithm will be:

```

LINEAR-SEARCH( $A, v$ )
1  while  $i \leq A.length$ 
2      do
           if  $A[i] = v$ 
3           then
               return  $i$ 
4       $i = i + 1$ 
5  return NIL

```

If the linear search finishes, then v has not been found in the array, so we return NIL. Otherwise, the loop stops when the first occurrence of v is found.

The **loop invariant** of the algorithm is:

At the start of each iteration of the loop, the subarray $A[1..i - 1]$ does not hold the value v .

Let us see now how the loop invariant properties hold now.

Initialization: We start by showing that the loop invariant holds before the first loop operation, when $i = 1$. The subarray of A is empty, so by definition it does not contain v .

Maintenance: Informally, the body of the while loop compares v with $A[i]$ and updates exits the loop if they are equal. The subarray $A[1..i]$ consists of elements that are not equal to v , as otherwise the loop would have ended.

Termination: We examine what happens when the loop terminates. When the loop terminates the value of $i = A.length + 1 = n + 1$. Then, the whole array A is the left subarray $A[1..i]$, so we have gone through the whole array and not found the value v , so we return NIL.

2.1-4

Stating the problem formally:

Input: Two arrays A, B of size n containing binary digits.

Output: An array C, which is of size $n + 1$ and contains the binary sum of A and B.

Code:

BINARY-ADDITION(A, B, C)

```
1  carry = 0
2  for  $i = A.length$  downto 1
3      do
4           $C[i + 1] = A[i] + B[i] + carry$ 
5          if  $C[i + 1] = 2$ 
6              then
7                   $C[i + 1] = 0$ 
8                   $carry = 1$ 
9          elseif  $C[i + 1] = 3$ 
10             then
11                  $C[i + 1] = 1$ 
12                  $carry = 1$ 
13         else
14              $carry = 0$ 
15      $C[1] = carry$ 
```