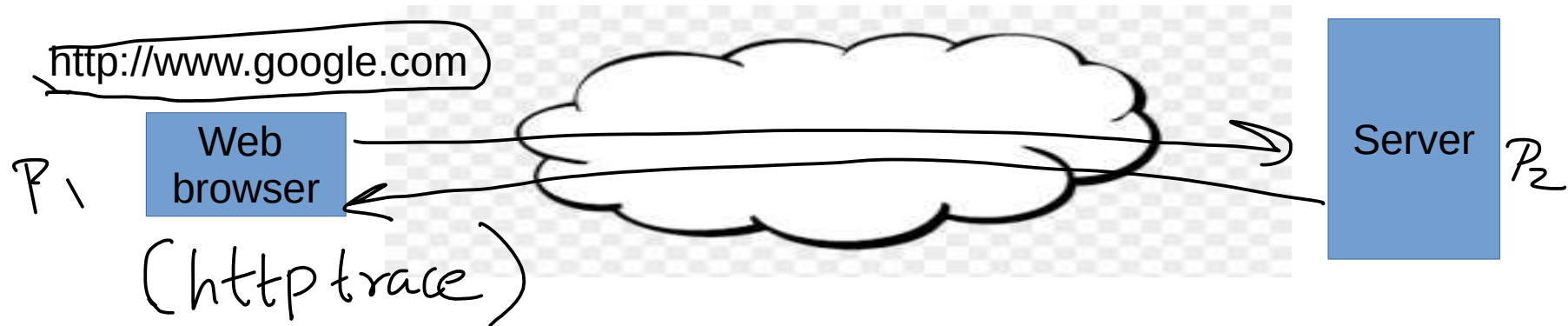

Introduction to Communicating Distributed Processes

Lecture 1

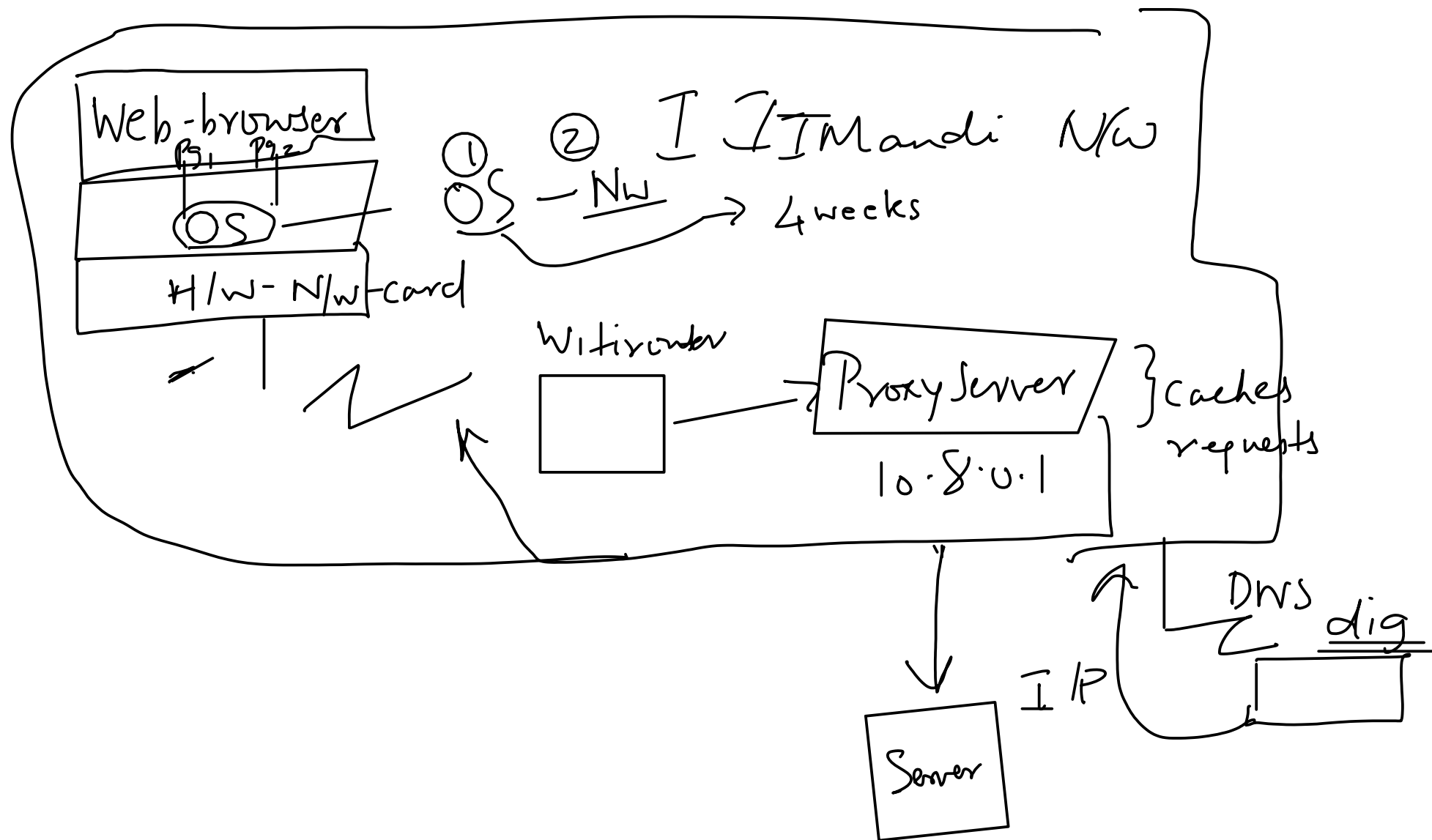
Motivating Example

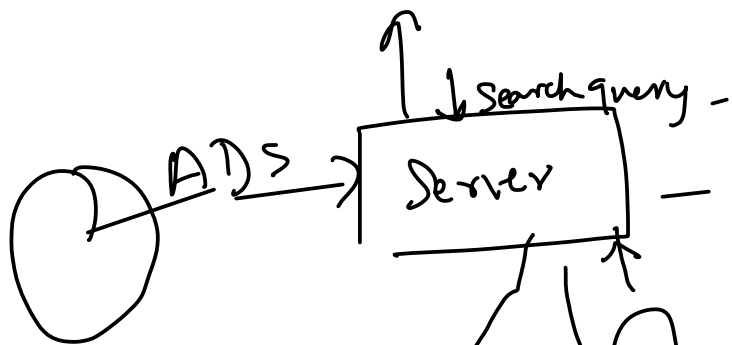


Get - Appⁿ layer protocol

- Flood the msg to all n/c
(Broadcast) - infeasible in
Internet-scale

Content-Delivery Network - latencies ↓





1 rep — $\frac{10005 \text{ rep/s}}{\text{Sequential}}$

Google data
100005 m/c

↓ multithreading

Concurrent

ACID

Threads

web
crawling

map $\begin{matrix} w_1, w_2, \dots \\ P_1 \end{matrix}$

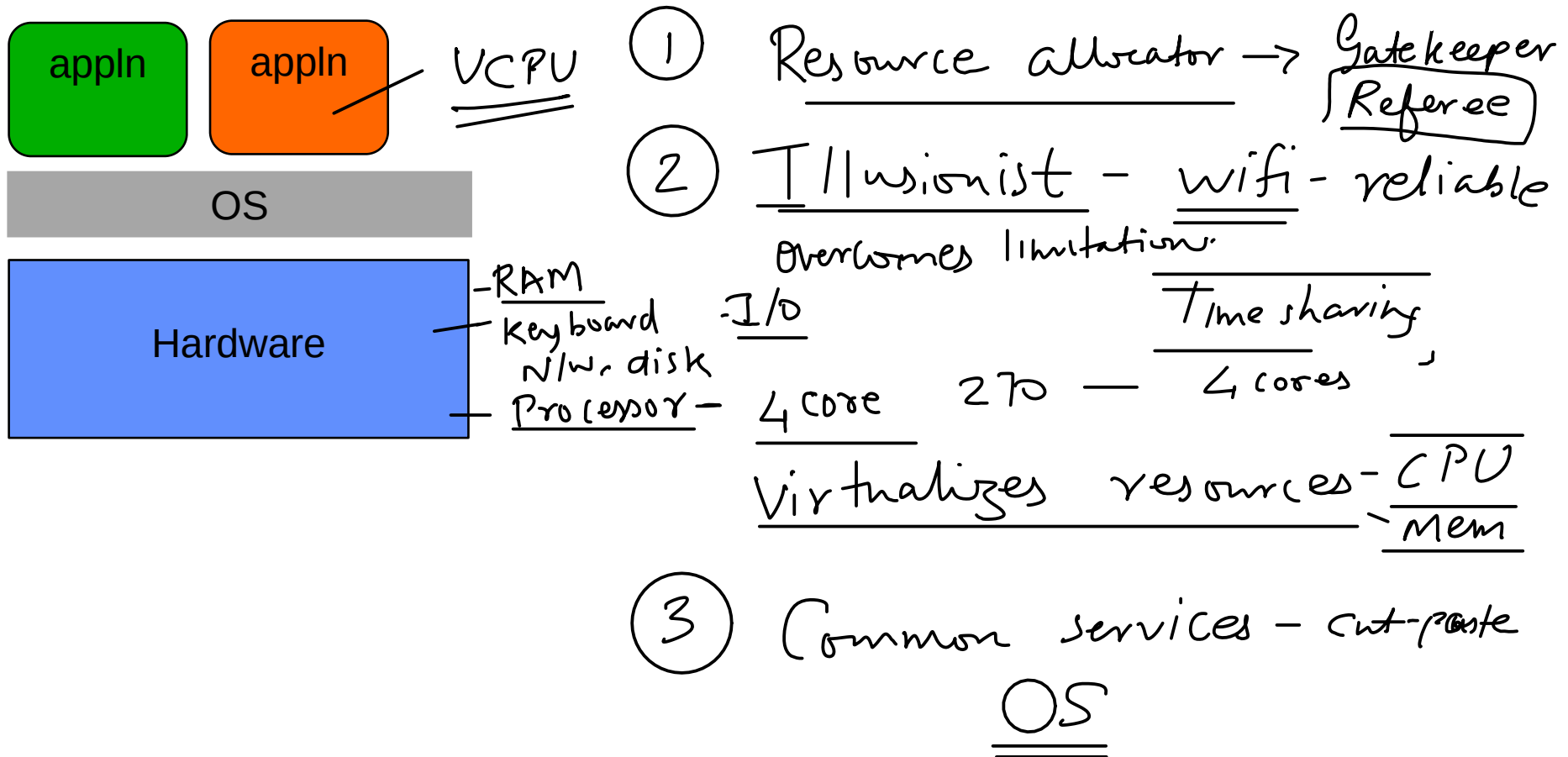
↓ (w_1, P_1)
 (w_2, P_1)
⋮

Index: word → P_1, P_2, \dots, P_n

$[w_1 \rightarrow P_1, P_2, \dots, P_n] \cap w_3$

What is an operating system?

- Special layer of software that manages a computer's resources for its users and applications



Roles played by OS

- Referee

- Resource allocation
- Isolation

```
def test():  
    while True:  
        print 'x=1'
```

- Communication -

- Illusionist

- Virtualize resources - illusion of reliable service using Wifi, infinite memory, ability to deal with evolving hardware

- Common services

- Cut, copy and paste across different applications

Handwritten notes:

- Timer Interrupt
- OS - Ctrl-C
- segmentation fault
- $2^{32} - 4\text{GB} \times 4\text{GB}$
- 64-bit processor: 2^{64}

Evolution of OS: a brief history

- 1945: ENIAC- loading programs / data was difficult and time-consuming
- 1949: EDVAC - Computers got memory
- 1949: BINAC - Programming language
- 1951: UNIVAC - Reusable code
- 1952: Shared IO routines and True assembler
- 1956: Interrupts
- 1954 - 1957: Batch processing, Fortran [architecture independent, high level language]
- 1960s - Multi-programming, multi-processing, virtual memory to make OS portable
- 1960s - disks became mainstream
- 1966 - mini-computers became cheaper [time-sharing system] - Interactive use
- 1969 - Unix Operating System - Assembly lang
- 1972 - Virtual machine operating system
- 1973 - Unix written in C [portable]
- Graphical user Interface and then ubiquitous devices

ml bit-bit

CPU - idle

huge - expensive
m/c

humans
cheap

fn.

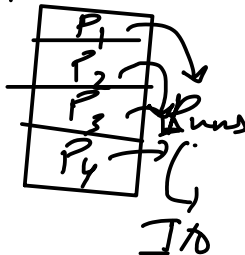
error-handling

CPU

Mem

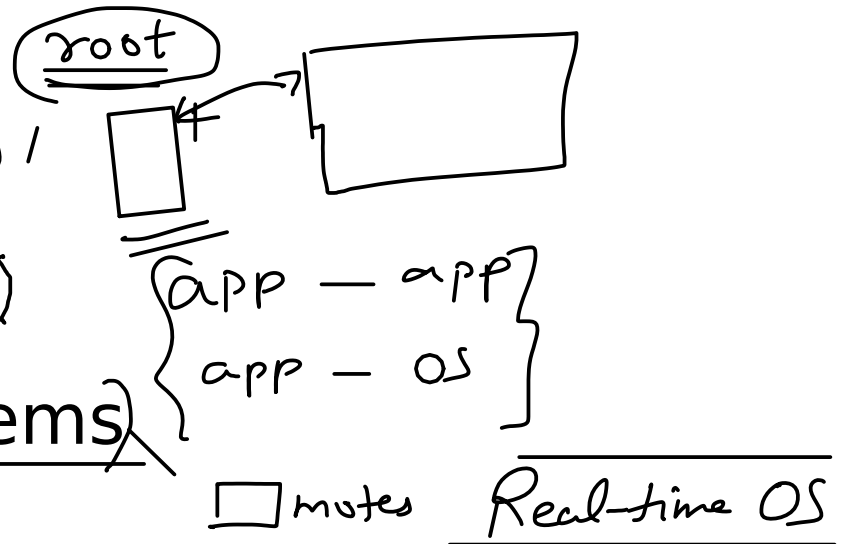
I/O

DMA transfer



improve utilization

Modern Operating Systems

- Desktop / Sudo Su - (root)
 - Smartphone / portability, samsung
 - Cloud Operating Systems
 - Embedded operating systems
- 
- notes Real-time OS

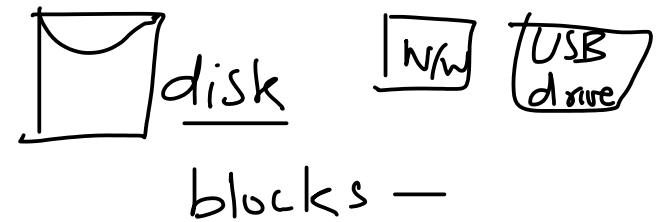
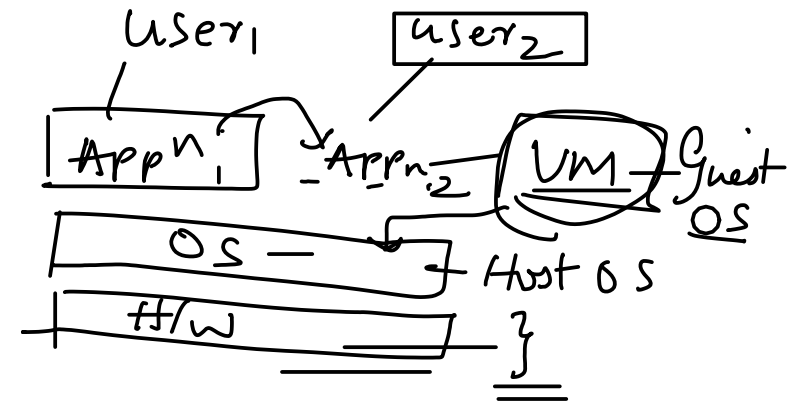
Where is it headed?

- Very large-scale data-centers, very heterogeneous hardware, multi-core machines, large storage

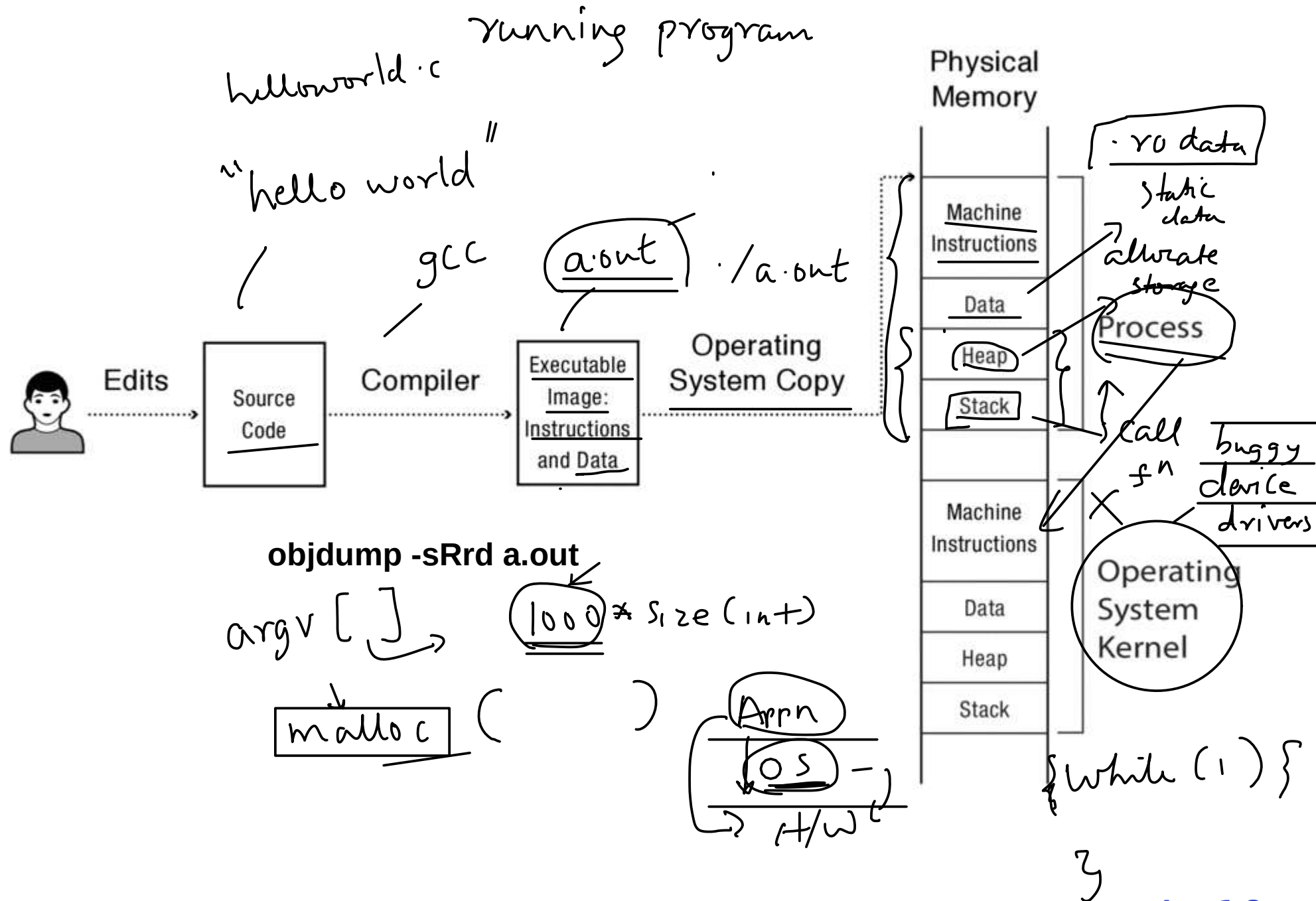
Recap: Sept. 17, 2020

OS plays 3 roles

- Referee - protection
- Illusionist - virtualizes res.
convenient abstr.
 - infinite mem.
- Common service
 - VMs
 - ↳ Windows OS on VM
 - Linux OS



Program to Process



C program

```
#include <stdio.h>
```

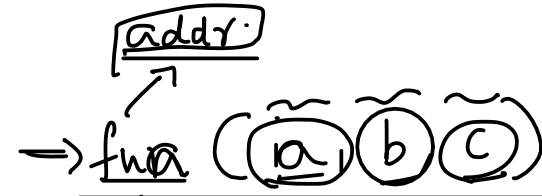
```
int main(int argc, char const *argv[])
```

```
{
```

```
    printf("Hello world\n");
```

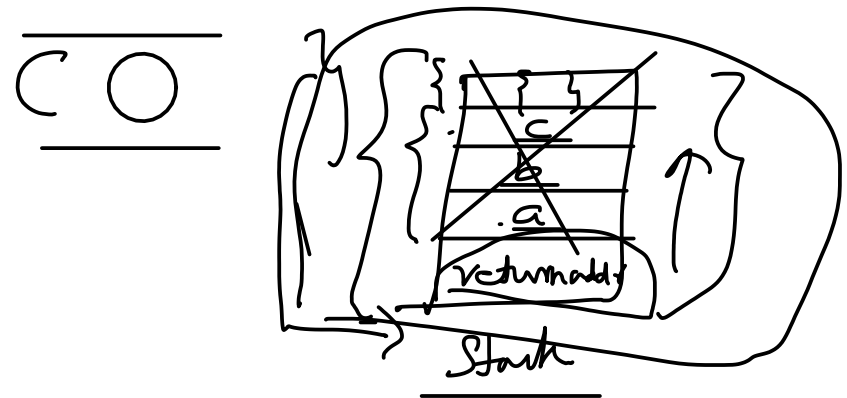
```
    return 0;
```

```
}
```

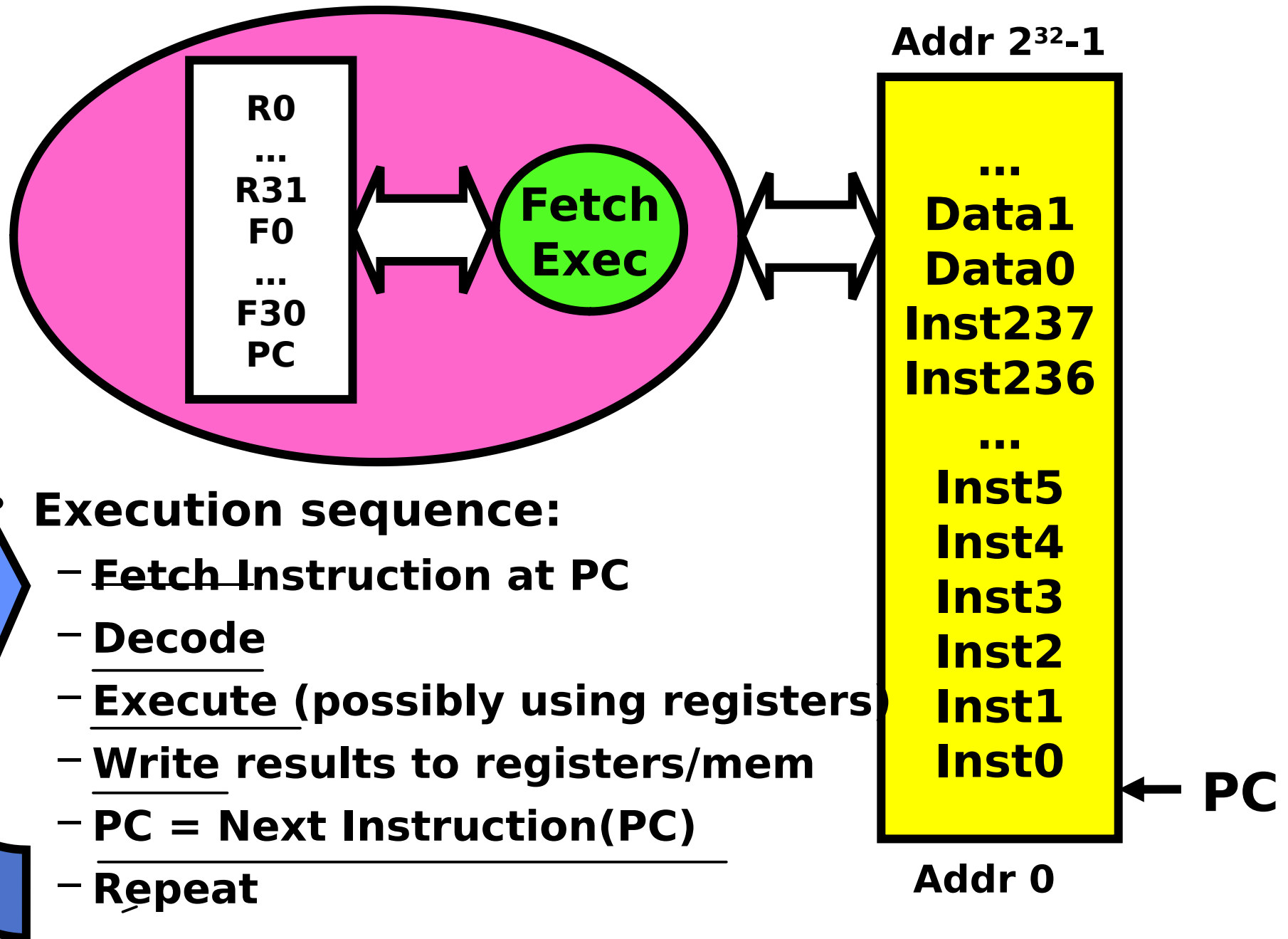


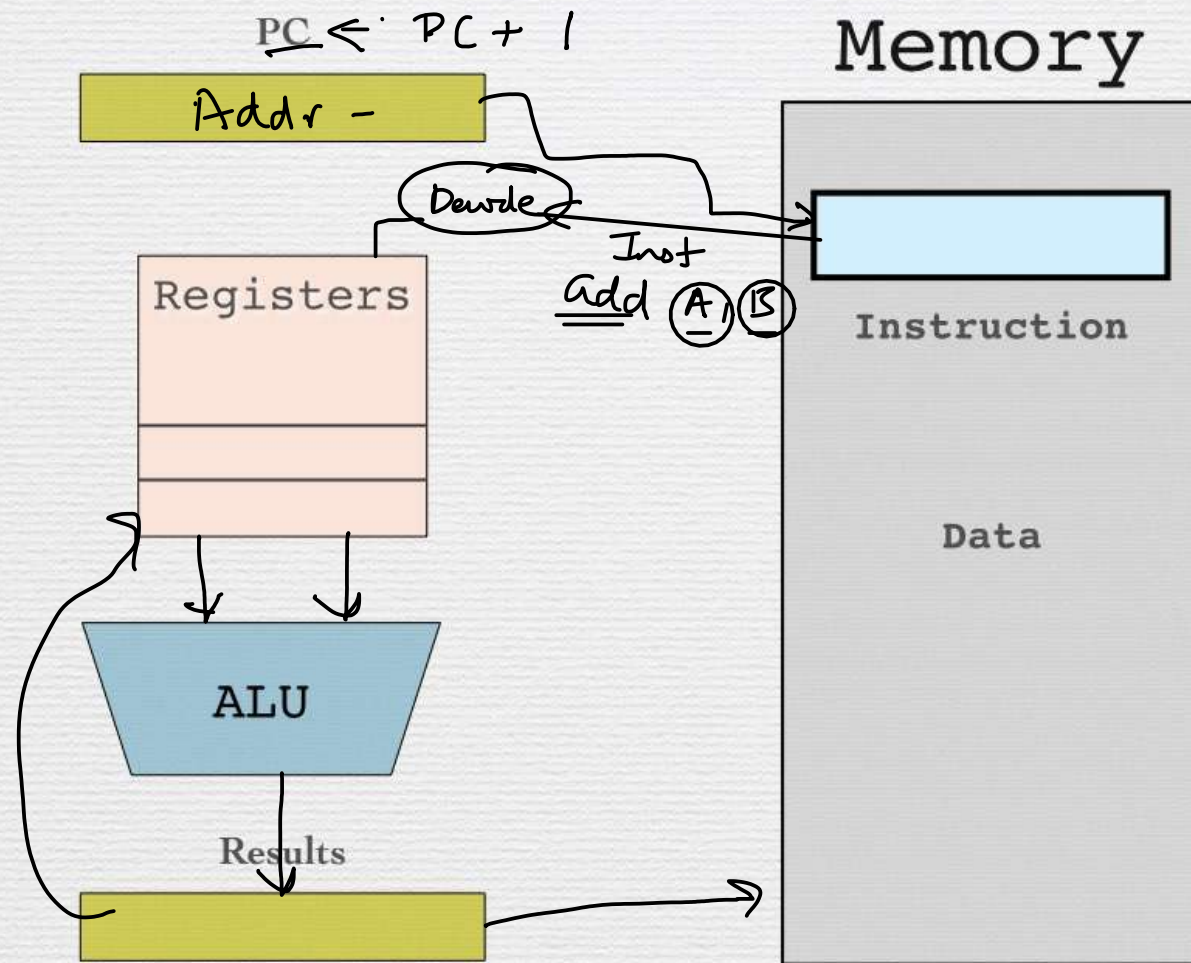
- return addr { int d; local var.

```
objdump -sRrd a.out
```



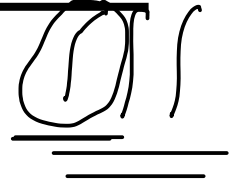
What happens during program execution?





First OS Concept: Thread of Control Process

- **Thread: Single unique execution context**



- Program Counter, Registers, Execution Flags, Stack, State in memory for that thread

- **PC: holds the address of executing instruction in the thread.**

- **Certain registers hold the *context* of thread**

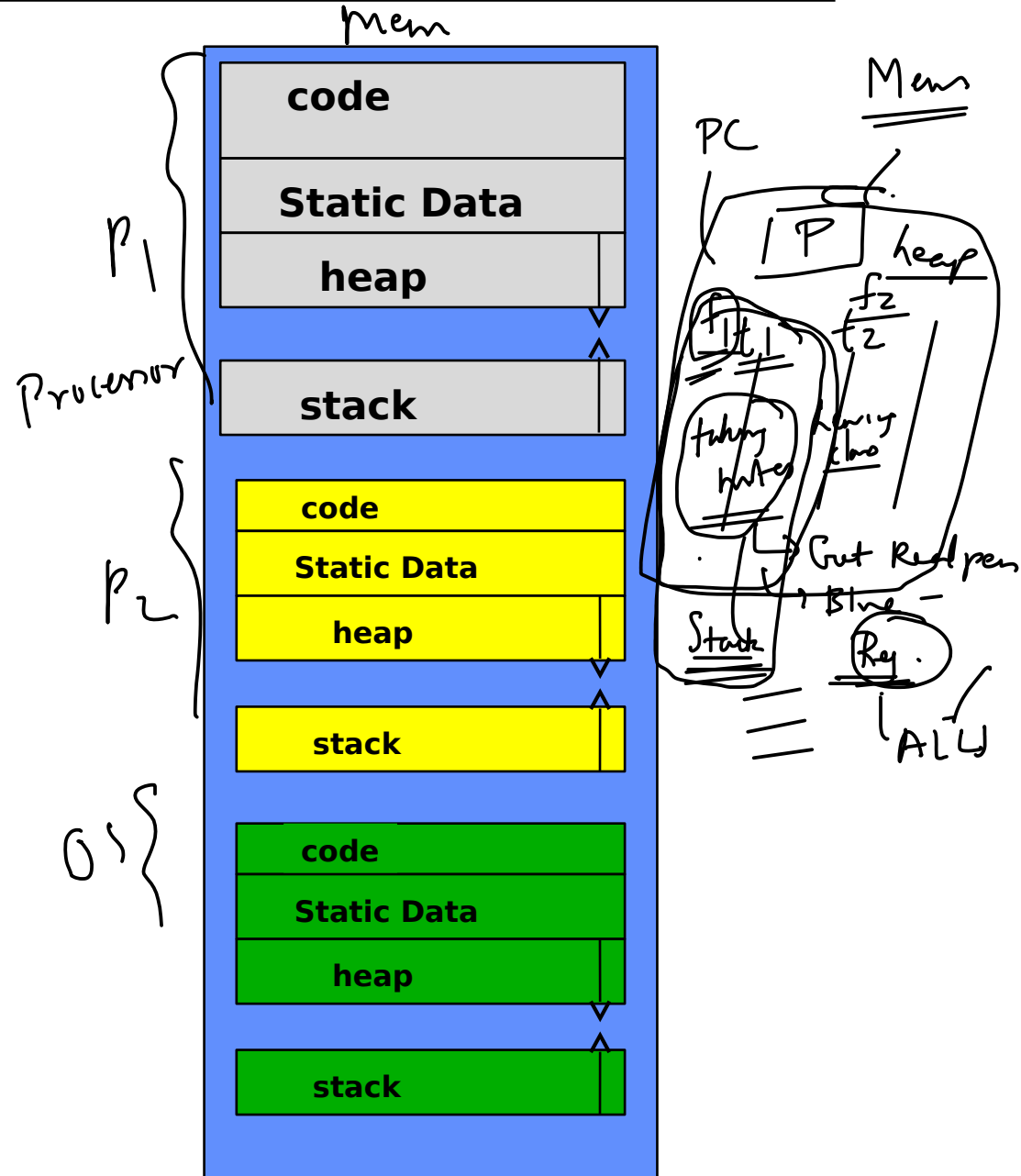
- Stack pointer, Heap Pointer, Data

- **Registers hold the root state of the thread.**

- The rest is “in memory”

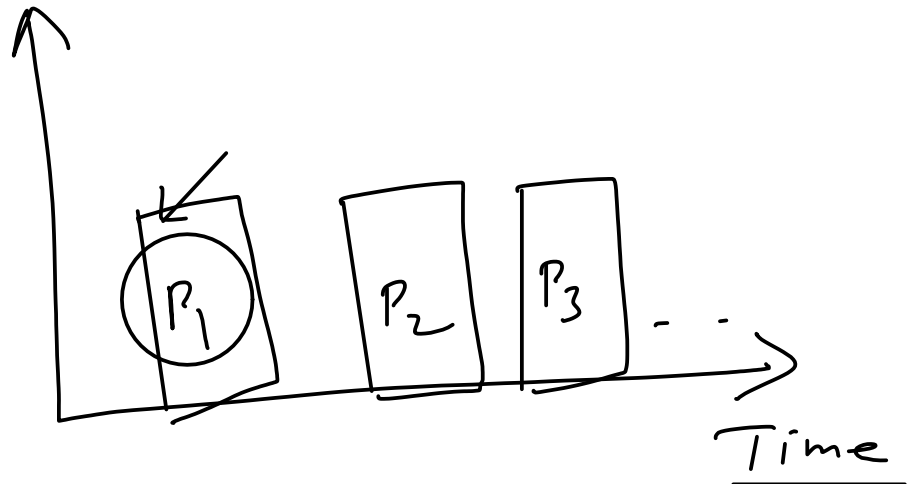
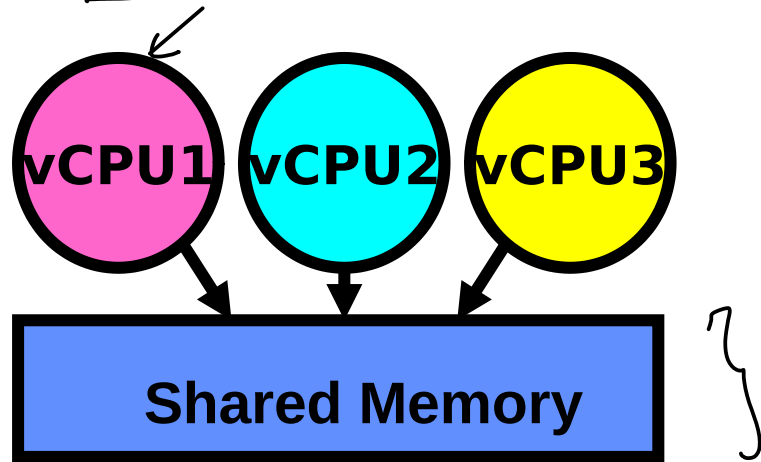


Multiprogramming



Virtualization of Resources

- Virtual CPU**



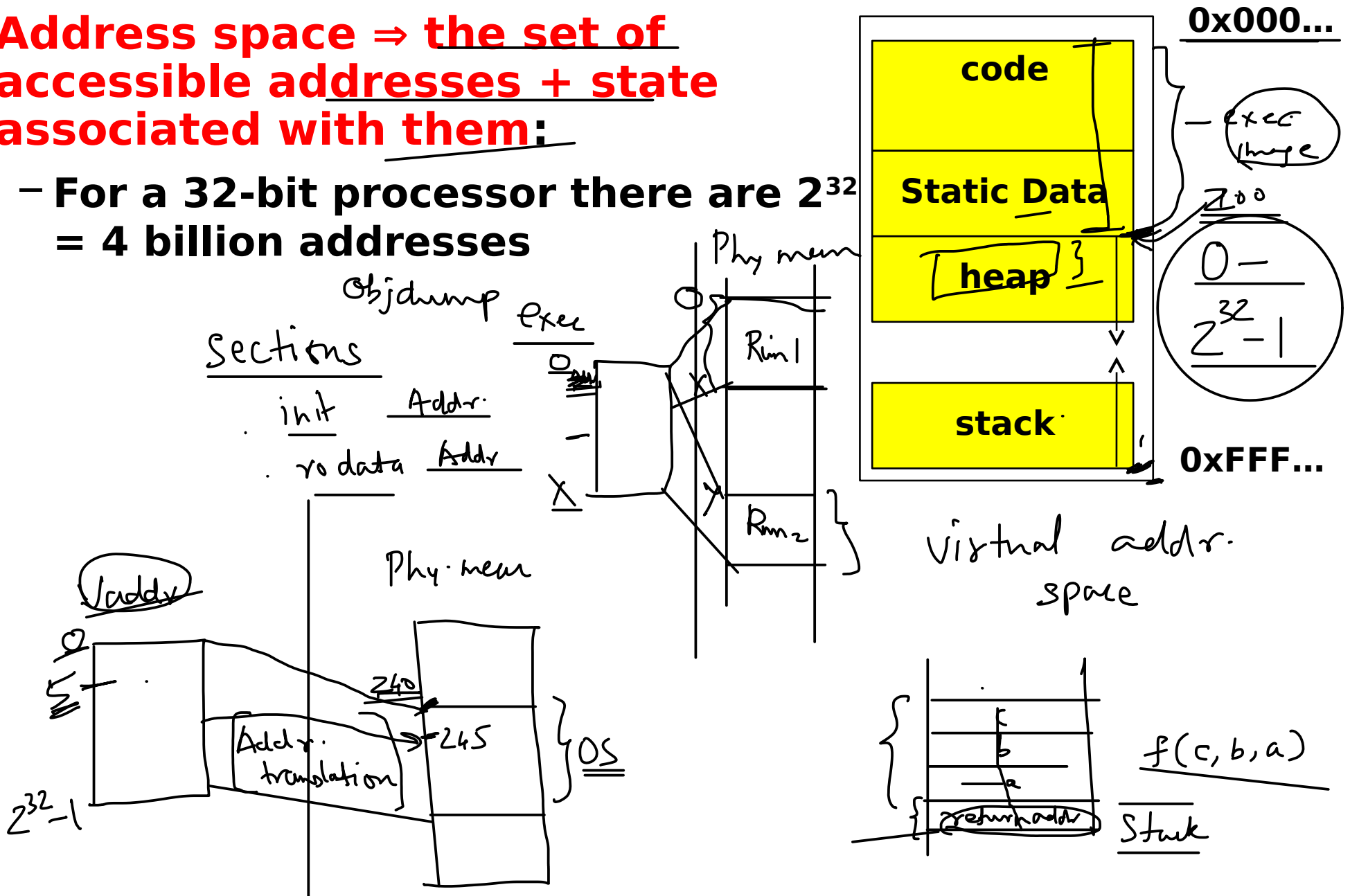
Recap Sept 18

- program → process
- Thread → unique ^{single} execution context
PC, Registers, Exec. flags.

Second OS Concept: Program's Address Space

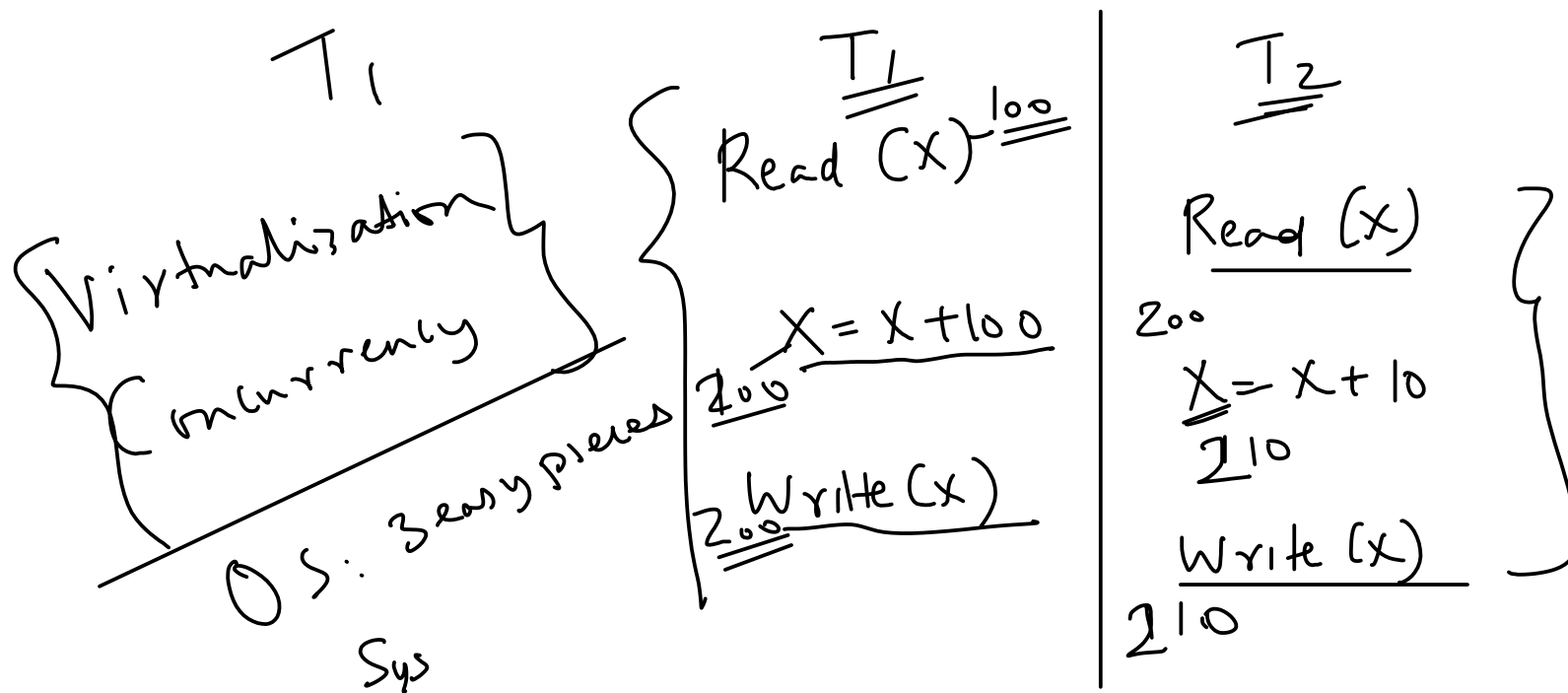
- **Address space \Rightarrow the set of accessible addresses + state associated with them:**

- For a 32-bit processor there are 2^{32} = 4 billion addresses

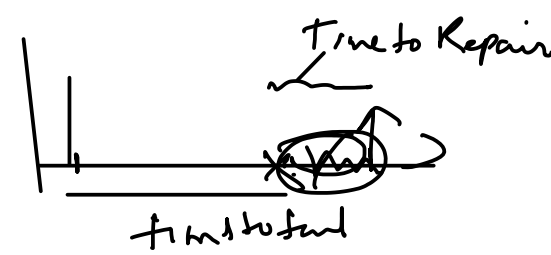


Demos

- **Virtualization of CPU and memory**
- **Concurrency issues lead to non-reproducible and non-deterministic output**



How to evaluate an Operating System?

- **Reliability** — correct — write 3 — 2 to 50 million lines —
 - **Availability** — crashing few sec. — $MTTR \downarrow$ $MTTF \uparrow$
 - **Adoption** —
 - **Security and Privacy** — unauthorized
 - **Portability** — $\left. \begin{array}{l} \nearrow \text{Newer Apps} \\ \searrow \text{H/w Abstraction layer} \end{array} \right\} \text{— } \underline{\underline{API}}$
 - **Performance** — \downarrow $\left\{ \begin{array}{l} \text{fast} \\ \text{efficiency} \end{array} \right.$
- 

Summary

- This course covers concepts from OS and Networks.
- OS is a layer of software that manages computer resources for its users and applications.
- Evolution of OS: IO routines, Batch processing, Multi-programming, Interactive processing ...
- Roles played by OS: referee, illusionist and common services
- Two concepts: threads and address space
- Metrics for evaluating the OS