

Namra Patel

1002229720

DAA

H.W-17

① Inserting n elements using

a) Aggregate Method.

→ The table doubles in size when it runs out of space.

→ So, if the original size is s , after insertion it doubles to $2s$ etc.

→ In general, after k doublings the size is 2^k .

Pseudo code

→ initialize table with capacity = 1
for i 1 to n :

if table is full:

new table = create new table
with size 2 to current size of element
from old to new table.

table = new table

'insert element i ' into table let $k = \log(n+1)$
- 1

$$\begin{aligned}\text{Total cost} &= O(n) * k \\ &= O(n \log n)\end{aligned}$$

A worst case cost per insertion = $O(\log n)$

Runtime per insertion is $O(\log n)$

Total time is $O(n)$ to $\log(n+1)$

b) Accounting method.

* charge 2 units per each insertion

→ when the table double in size from m to $2m$ credit m units

→ The credit executes pay for copy cost at $O(n)$.

→ Total credit is $m + 2m + nm + \dots \cdot \frac{n}{2}m = O(n)$

Pseudo code:

→ initialise table with capacity = 1

for $i = 0$ to n :

if table is full:

newtable = create with size $2 \times$

current size

copy element from old table to new table.

table = newtable.

insert element i into table.

initialise charges = 0

initialize credits = 0

for $i = 1$ to n ;

charges $+= 2$

if table doubled in size
from m to $< 2m$:

credits $+= m$

total charges = $2n = O(n)$

total credits = $m + 2m + \dots + n/2 \cdot m = O(n)$

Amortized cost per insertion -

total $(n, O(n)) / n$

= 1

Runtime per insertion = $O(1)$

total time = $O(n)$