

UTA I.D :- 1002229720

Namra Patel

## Hands On - 3

1. Find the runtime of algorithm mathematically

Ans → So here outer runs  $n$  time

$$T(n) = 1 + \sum_{i=1}^n + \sum_{i=1}^n \sum_{j=1}^{n+1} + \sum_{i=1}^n \sum_{j=1}^1 1$$

$$T(n) = 1 + (n+1) (n^2 + n) + n^2$$

$$T(n) = 2 + 3n + 3n^2$$

2. In github code.

3. Find polynomials that are upper and down bounds on your curve from #2

Ans →

Big-O → the upper bound on graph shown by the graph dashed line, which slightly exceeds the other curve.

Big-omega → So this represents the lower bound on time complexity. So it grows at least guarantichy repeat to  $n$  so  $\sim (n^2)$



Big theta: Since both the upper and closely follow the actual timing data runtime is  $\Theta(n^2)$ .

4. In github

$x = f(n)$

$x = 1;$

$y = 1;$

for  $j = 1 : n$

for  $j = 1 : n$

$x = x + 1;$

$y = 1 + y;$

4. will this increase now why it takes the algorithm to run.

Ans. Yes the modifier function will take slightly more time to run because of operation  $y = 1 + y$  within inner loop. However, this operation is constant ~~time~~ <sup>time</sup>  $\Theta(1)$ , and since it is executed same time as the other operation.

5. Will it effect your result from #1.

→ So it will not effect the results as the time ~~so~~ complexity will remain the same  $O(n^2)$   
So the array does not affect the complexity.