PGMO Lecture: Vision, Learning and Optimization

2. Basic notion of convexity

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Overview

Convex functions

Legendre-Fenchel conjugate

Infimal convolution

Proximal map

Duality

Convexity

An extended real valued function $f: \mathcal{X} \to [-\infty, +\infty]$ is said to be *convex* if and only if its *epigraph*

$$\mathsf{epi}\, f := \{(x,\lambda) \in \mathcal{X} \times \mathbb{R} : \lambda \geq f(x)\}$$

is a convex set, that is, if when $\lambda \geq f(x)$, $\mu \geq f(y)$, and $t \in [0,1]$, we have $t\lambda + (1-t)\mu \geq f(tx+(1-t)y)$.

It is *proper* if it is not identically $+\infty$ and nowhere $-\infty$: in this case, it is convex if and only if, for all $x, y \in \mathcal{X}$ and $t \in [0, 1]$,

$$f(tx + (1-t)y) \le tf(x) + (1-t)f(y).$$

It is *strictly convex* if the above inequality is strict whenever $x \neq y$ and 0 < t < 1.

Subgradient

Given a convex, extended real valued, l.s.c. function $f: \mathcal{X} \to [-\infty, +\infty]$, we recall that its subgradient at a point x is defined as the set

$$\partial f(x) := \{ p \in \mathcal{X} : f(y) \ge f(x) + \langle p, y - x \rangle \ \forall y \in \mathcal{X} \}.$$

Fermat's stationary conditions for non-smooth convex functions f is hence generalized as:

 $x \in \mathcal{X}$ is a global minimizer of f if and only if $0 \in \partial f(x)$.

Smooth functions

Let us recall that a function f has a L-Lipschitz continuous gradient if there exists a constant L>0 such that for all $x,y\in\mathbb{E}$ one has

$$\left\|\nabla f(x) - \nabla f(y)\right\|_{*} \leq L \left\|x - y\right\|.$$

This implies that

$$\langle x - y, \nabla f(x) - \nabla f(y) \rangle \stackrel{C.S.}{\leq} \|x - y\|_2 \|\nabla f(x) - \nabla f(y)\|_2 \leq L \|x - y\|_2^2$$

Theorem

Let f be a continuously differentiable function over \mathbb{E} with L-Lipschitz continuous gradient. Then, for any $x,y\in\mathbb{E}$

$$f(y) \le f(x) + \langle \nabla f(x), y - x \rangle + \frac{L}{2} ||x - y||^2,$$

which is known under the name "descent lemma". Moreover, we also have the reverse inequality

$$f(y) \ge f(x) + \langle \nabla f(x), y - x \rangle + \frac{1}{2I} \|\nabla f(x) - \nabla f(y)\|^2$$

which is called the co-coercivity of the gradient.

Strong convexity

A function f is strongly convex or " μ -convex" if in addition, for $x, y \in \mathcal{X}$ and $p \in \partial f(x)$, we have

$$f(y) \ge f(x) + \langle p, y - x \rangle + \frac{\mu}{2} \|y - x\|^2$$

or, equivalently, if $x \mapsto f(x) - \frac{\mu}{2} \|x\|^2$ is convex.

Furthermore, it satisfies

$$f(tx + (1-t)y) \le tf(x) + (1-t)f(y) - \mu \frac{t(1-t)}{2} ||y-x||^2$$

for any x, y and any $t \in [0, 1]$.

Finally, if f is strongly convex and x^* is a minimizer of f(x), then we have

$$f(y) \ge f(x^*) + \frac{\mu}{2} \|y - x^*\|^2$$

for all $y \in \mathcal{X}$. For a *L*-smooth and μ strongly convex function f, the co-coercivity reads

$$\langle \nabla f(x) - \nabla f(y), x - y \rangle \ge \frac{L\mu}{L+\mu} \|x - y\|^2 + \frac{1}{L+\mu} \|\nabla f(x) - \nabla f(y)\|^2,$$

which is nothing else then the co-coercivity for the function $f(x) - \frac{\mu}{2} ||x||^2$, which is $L - \mu$ -smooth.

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Legendre-Fenchel conjugate

To any function $f: \mathcal{X} \to [-\infty, +\infty]$ one can associate the *Legendre–Fenchel conjugate* (or convex conjugate)

$$f^*(y) = \sup_{x \in \mathcal{X}} \langle y, x \rangle - f(x)$$

which, as a supremum of linear and continuous functions, is obviously convex and l.s.c.

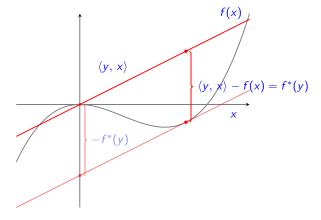


Figure: Illustration of the convex conjugate and its relation to the gradient of smooth functions.

Biconjugate and properties

The *biconjugate* f^{**} is the largest convex l.s.c. function below f. In general $f \ge f^{**}$ but if f is already convex and l.s.c. we have $f^{**} = f$.

Recall the celebrated Legendre-Fenchel identity:

$$y \in \partial f(x) \Leftrightarrow x \in \partial f^*(y) \Leftrightarrow f(x) + f^*(y) = \langle y, x \rangle$$
.

We have the following basic properties

- For f(x) = g(ax) with $a \neq 0$, the convex conjugate is given by $f^*(y) = g^*(y/a)$.
- ▶ For f(x) = g(x + b), the convex conjugate is given by $f^*(y) = g^*(y) \langle b, y \rangle$.
- for f(x) = ag(x) with a > 0, the convex conjugate is given by $f^*(y) = ag^*(y/a)$.
- ▶ Note that we also have the property:

$$f(x) \ge g(x) \Longleftrightarrow g^*(y) \ge f^*(y)$$

Monotone operator

The subgradient of a convex function is a monotone operator, that is it satisfies

$$\langle p-q, x-y \rangle \geq 0 \quad \forall (x,y) \in \mathcal{X}^2, \ p \in \partial f(x), \ q \in \partial f(y),$$

It is *strongly monotone* if *f* is strongly convex:

$$\langle p-q, x-y \rangle \ge \mu \|x-y\|^2 \quad \forall (x,y) \in \mathcal{X}^2, \ p \in \partial f(x), \ q \in \partial f(y).$$

An important remark is that f is μ -strongly convex if and only if its conjugate f^* is continuously differentiable with $1/\mu$ -Lipschitz gradient, i.e.

$$\|
abla f^*(p) -
abla f^*(q)\| \leq rac{1}{\mu} \|p - q\|$$
 .

Hence, duality allows to trade strong convexity with smoothness.

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The infimal convolution

Definition

Let $f, g : \mathcal{X} \to (-\infty, \infty]$ be two proper functions. The infimal convolution between f and g is defined by

$$(f\Box g)(x) = \min_{u} f(u) + g(x - u)$$

- ▶ Interpretation: Take the function f and convolve it with the minimum of the function g. The function $f \square g$ is given by the resulting envelope function.
- ▶ Remark: In case f and g are positive and one-homogeneous functions (e.g. norms or semi-norms) the infimal convolution is equivalent to the convex envelope (i.e. biconjugate) of the minimum of the two functions

$$(f\Box g) = (\min\{f(x), g(x)\})^{**}.$$

Graphical interpretation

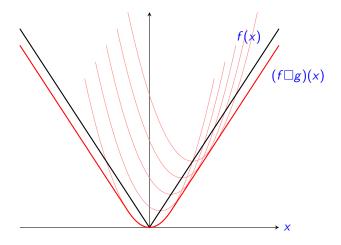


Figure: Illustration of the infimal convolution of the function f(x) = |x| and the quadratic function $g(x) = \frac{1}{2}x^2$.

Properties of the infimal convolution

► The infimal convolution of two convex functions enjoys the following two fundamental properties.

Proposition

Let f and g be two proper convex functions, then the following holds:

- $ightharpoonup f \Box g$ is a convex function.
- ▶ Proof: The convexity follows from the convexity of partial minimization. The symmetry follows from the fact that the infimal convolution can also be written as

$$(f \square g)(x) = \min_{u} f(u) + g(x - u) = \min_{x = u + v} f(u) + g(v).$$

Infimal convolution and convex conjugates

There are interesting connections between the convex conjugate and the infimal convolution.

Theorem

Let $f, g : \mathcal{X} \to (-\infty, \infty]$ be two proper functions and let f^*, g^* be their convex conjugates. It holds that

$$(f\Box g)^* = f^* + g^*.$$

The second direction requires additional assumptions such as convexity.

Theorem

Let $f: \mathcal{X} \to (-\infty, \infty]$ be a proper convex function and let $g: \mathcal{X} \to \mathbb{R}$ be a real-valued convex function. Then

$$(f+g)^*=f^*\Box g^*.$$

The Moreau envelope

- An interesting special case of the infimal convolution is the Moreau envelope.
- ▶ It is obtained by choosing one function in the infimal convolution as a quadratic function

$$x \mapsto f_{\lambda}(x) = (f \Box \frac{1}{2\lambda} \|\cdot\|_{2}^{2})(x) = \min_{u} f(u) + \frac{1}{2\lambda} \|x - u\|^{2},$$

where $\lambda > 0$ is a "smoothing" parameter.

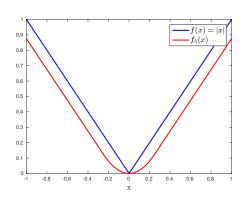
It can be used to "smooth" a non-smooth convex function without destroying its minimizer.

Example

▶ Let f(x) = |x|. The Moreau-envelope $f_{\lambda}(x)$ is given by

$$f_{\lambda}(x) = \begin{cases} rac{x^2}{2\lambda} & \text{if } |x| \leq \lambda \\ |x| - rac{\lambda}{2} & \text{else.} \end{cases}$$

▶ This is exactly the Huber function from robust statistics!



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The proximal map

Let us consider again the Moreau envelope of a function f:

$$\min_{x} f(x) + \frac{1}{2\lambda} \left\| x - y \right\|^2$$

▶ The point that attains the minimum in the Moreau envelope is called the proximal map.

Definition

Let $\lambda > 0$ be a parameter. The proximal map with parameter λ is defined as

$$\operatorname{prox}_{\lambda f}(y) = \arg \min_{x} f(x) + \frac{1}{2\lambda} \|x - y\|^{2}$$

▶ It can be seen as a generalization of the Euclidean projection of a point *y* on a convex set *C*.

Optimality condition of the proximal map

▶ The optimality condition of the proximal map is given by

$$\lambda \partial f(x) + x - y \ni 0 \iff y - x \in \lambda \partial f(x)$$

▶ In operator notation, the proximal map is written as

$$x - y + \lambda \partial f(x) \ni 0$$

$$\Leftrightarrow x + \lambda \partial f(x) \ni y$$

$$\Leftrightarrow (I + \lambda \partial f)x \ni y$$

$$\Leftrightarrow x = (I + \lambda \partial f)^{-1}(y)$$

$$\Leftrightarrow x = \operatorname{prox}_{\lambda f}(y)$$

Examples

▶ Let $f(x) = \frac{1}{2}x^TAx + b^Tx + c$. The proximal map is given by

$$\operatorname{prox}_{\lambda f}(y) = (I + \lambda A)^{-1}(y - \lambda b)$$

▶ Let $f(x) = ||x||_1$, the proximal map is given by

$$(\operatorname{prox}_{\lambda f}(y))_i = \max(0, |y_i| - \lambda)\operatorname{sgn}(y_i)$$

Let $f(x) = -\sum_{i=1}^{n} \log x_i$, the proximal map is given by

$$(\operatorname{prox}_{\lambda f}(y))_i = \frac{y_i + \sqrt{y_i^2 + 4\lambda}}{2}$$

Let $f(x) = \delta_{\|\cdot\|_{\infty} < 1}(x)$, the proximal map is given by

$$(\mathsf{prox}_f(y))_i = \frac{y_i}{\mathsf{max}(1,|y_i|)}$$

Proximal map calculus

Let $f(x,y) = f_1(x) + f_2(y)$, then the proximal map with respect to f is given by

$$\mathsf{prox}_f(u,v) = (\mathsf{prox}_{f_1}(u), \mathsf{prox}_{f_2}(v)).$$

▶ If $f(x) = \alpha g(x) + b$, with $\alpha > 0$, then

$$prox_f(u) = prox_{\alpha f}(u)$$

▶ If $f(x) = g(\alpha x + b)$, with $\alpha \neq 0$ then

$$\operatorname{prox}_f(u) = \frac{1}{\alpha} \left(\operatorname{prox}_{\alpha^2 g} (\alpha u + b) - b \right)$$

▶ If f(x) = g(Qx), with Q orthogonal (such that $QQ^T = Q^TQ = I$), then

$$\operatorname{prox}_f(u) = Q^T \operatorname{prox}_g(Qu)$$

Proximal map calculus

$$prox_f(u) = prox_g(u - a)$$

► If
$$f(x) = g(x) + \frac{\gamma}{2} ||x - a||^2$$
, then

$$\operatorname{prox}_{f}(u) = \operatorname{prox}_{\tilde{\gamma}g} \left(\tilde{\gamma}u + \tilde{\gamma}\gamma a \right),$$

with
$$\tilde{\gamma} = 1/(1+\gamma)$$
.

Moreau identity

▶ The Moreau identity [Moreau '65] connects the proximal map of a function with the proximal map of its convex conjugate.

Theorem

Let f be a convex function and let f^* be its convex conjugate. Then

$$x = \operatorname{prox}_{f}(x) + \operatorname{prox}_{f^{*}}(x)$$

One also has that

$$x = \operatorname{prox}_{\tau f}(x) + \tau \operatorname{prox}_{\frac{1}{\tau}f^*} \left(\frac{x}{\tau}\right)$$

- \triangleright From a practical point of view, it shows that the proximal map of a function f is as easy to compute as the proximal map of its convex conjugate.
- ► Example: f(x) = |x|. Its convex conjugate is $f^*(y) = \delta_{[-1,1]}(y)$ and $\text{prox}_{f^*}(y) = \max(-1, \min(1, x))$, hence

$$\operatorname{prox}_f(x) = x - \max(-1, \min(1, x)).$$

The gradient of the Moreau envelope

▶ It turns out that the gradient of the Moreau envelope of a convex function *f* has a particularly appealing structure.

Proposition

Let f be a convex function and let $\lambda > 0$. Then the Moreau envelope of f defined through

$$(f \Box \frac{1}{2\lambda} \|\cdot\|_2^2)(x) = \min_{y} f(y) + \frac{1}{2\lambda} \|x - y\|_2^2$$

is continuously differentiable and its gradient

$$\nabla\left(f\Box\frac{1}{2\lambda}\left\|\cdot
ight\|_{2}^{2}
ight)=rac{1}{\lambda}(I-\mathsf{prox}_{\lambda f})$$

is λ^{-1} Lipschitz continuous.

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Fenchel-Rockafellar Duality

Let us consider the following (primal) optimization problem

$$\min_{x} f(Kx) + g(x),$$

where f, g are closed convex functions and $K \in \mathbb{R}^{m \times n}$ is a matrix. Using the fact that $f = f^{**}$ we have

$$\min_{x} f(Kx) + g(x) = \min_{x} \max_{y} \langle y, Kx \rangle - f^{*}(y) + g(x).$$

Under very mild conditions, we can swap the min and max which yields

$$\max_{y} \min_{x} \langle y, Kx \rangle - f^{*}(y) + g(x) = \max_{y} - \max_{x} - \langle x, K^{*}y \rangle + f^{*}(y) - g(x).$$

ightharpoonup Using the definition of g^* we obtain the Rockafellar-Fenchel dual

$$\max_{y} -f^{*}(y) - g^{*}(-K^{*}y),$$

Strong duality

Let us denote by

$$\mathcal{P}(x) = f(Kx) + g(x), \quad \mathcal{D}(y) := -f^*(y) - g^*(-K^*y),$$

the primal and dual problems.

From the Fenchel-Rockafellar duality it follows that strong dualty holds, i.e.

$$\max_{y} \mathcal{D}(y) = \min_{x} \mathcal{P}(x)$$

Moreover.

$$\mathcal{D}(y) \leq \mathcal{P}(x), \quad \forall (x, y).$$

Hence, it is natural to define the primal-dual gap

$$\mathcal{G}(x,y) = \mathcal{P}(x) - \mathcal{D}(y) \ge 0,$$

which vanishes if and only if (x, y) is an optimal solution pair of the primal and dual problems.

Saddle point formulation

Let us define the Lagrangian function

$$\mathcal{L}(x,y) := \langle y, Kx \rangle - f^*(y) + g(x)$$

If x^* is a solution to the primal problem and y^* is the solution of the dual problem, then

$$\max_{y} \mathcal{L}(x^*, y) = \mathcal{P}(x^*) = \mathcal{L}(x^*, y^*) = \mathcal{D}(y^*) = \min_{x} \mathcal{L}(x, y^*)$$

Observing that

$$\mathcal{L}(x^*, y) \le \mathcal{P}(x^*), \quad \mathcal{D}(y^*) \le \mathcal{L}(x, y^*),$$

we obtain

$$\mathcal{L}(x^*, y) \le \mathcal{L}(x^*, y^*) \le \mathcal{L}(x, y^*),$$

hence the optimal solution pair (x^*, y^*) is a saddle point of the Lagrangian.

Minimax Theorem

The mixed primal-dual formulation that appeared in the derivation of the Fenchel-Rockafellar dual problem belongs to the class of convex-concave minimax problems.

Theorem

Let $X \subset \mathcal{X}$ and $Y \subset \mathcal{X}^*$ be convex and compact sets. Moreover, let $f(\cdot, y)$ be convex for all fixed y and $f(x, \cdot)$ concave for all fixed x. Then, we have that:

$$\min_{x \in X} \max_{y \in Y} f(x, y) = \max_{y \in Y} \min_{x \in X} f(x, y)$$

- ▶ Proved by John von Neumann in 1935: As far as I can see, there could be no theory of games . . . without that theorem . . . I thought there was nothing worth publishing until the Minimax Theorem was proved.
- Extended to quasi-convex/concave functions by Sion in 1958.

Non-convex minimax theorem

Finally, we present a result on general minimax problems:

Theorem

Let $f: X \times Y \mapsto \mathbb{R}$ be a saddle point problem, where X and Y are some subsets of \mathcal{X} and \mathcal{X}^* , then

$$\min_{x \in X} \max_{y \in Y} f(x, y) \ge \max_{y \in Y} \min_{x \in X} f(x, y).$$

- Interpretation: It matters who plays first in games.
- ▶ Proof: First note that we have

$$f(x,y) \ge \min_{\xi} f(\xi,y), \quad \forall (x,y) \in X \times Y$$

► Taking the maximum wrt *y* on both sides

$$\max_{y} f(x, y) \ge \max_{y} \min_{\xi} f(\xi, y), \quad \forall x \in X$$

Finally, taking the minimum wrt x on both sides yields the desired result.

Example: Dual formulation of the ROF model

Recall that the primal ROF model is given by

$$\min_{u} \| \mathrm{D} u \|_{p,1} + \frac{1}{2} \| u - d \|^{2}$$

The dual ROF model is given by

$$\max_{\mathbf{p}} -\delta_{\{\|\cdot\|_{q,\infty} \leq \lambda\}}(\mathbf{p}) - \frac{1}{2} \|\mathbf{D}^*\mathbf{p}\|^2 + \langle \mathbf{D}^*\mathbf{p}, f \rangle,$$

where $u = d - D^*p$. The primal-dual gap is given by

$$\mathcal{G}(u,\mathbf{p}) = \|\mathrm{D}u\|_{p,1} + \delta_{\{\|\cdot\|_{q,\infty} \le \lambda\}}(\mathbf{p}) - \langle \mathbf{p}, \, \mathrm{D}u \rangle + \frac{1}{2} \|d - \mathrm{D}^*\mathbf{p} - u\|^2 \ge 0$$

It turns out that for u^* being the minimizer of the primal ROF model,

$$\mathcal{G}(u,\mathbf{p}) \geq \frac{1}{2} \|u - u^*\|^2 + \frac{1}{2} \|d - D^*\mathbf{p} - u^*\|^2.$$