# Implementación de los Tableaux Semánticos de la Lógica Proposicional

Semana 10

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Departmento de Matemáticas Aplicadas y Ciencias de la Computación



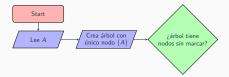
## Presentación

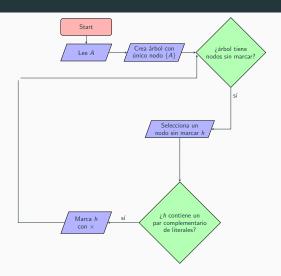
En esta sesión estudiaremos:

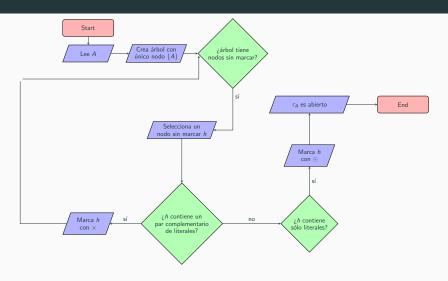
- 1 Elementos de la implementación
- 2 Algoritmo primero en anchura
- 3 Algoritmo primero en profundidad
- 4 Algoritmo backtracking

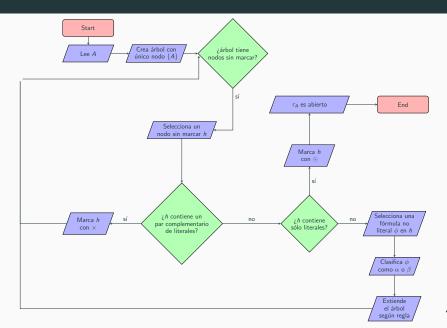
## Presentación

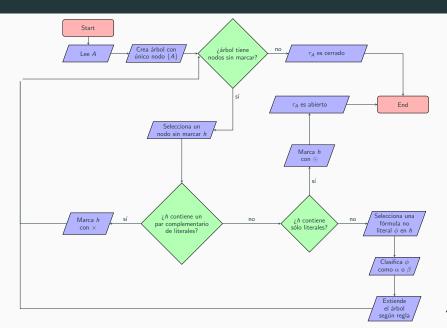
- 1 Elementos de la implementación
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- 4 Algoritmo backtracking











#### nodo\_tableaux — atributos

#### **Alfas**

self.alfas es una lista con las fórmulas tipo  $\alpha$ , su representación como cadena, el número de regla y la cadena 'alfa'.

Ejemplo:

$$n = \mathsf{nodo\_tableaux}(['-(p > q)', '(pOq)', '-p'])$$

$$n.\mathsf{alfas} = [(\mathit{Formula},' - (p > q)', 4,' \mathit{alfa'})]$$

5

#### nodo\_tableaux — atributos

#### **Betas**

self.betas es una lista con las fórmulas tipo  $\beta$ , su representación como cadena, el número de regla y la cadena 'beta'.

$$n = \mathsf{nodo\_tableaux}(['-(p > q)', '(pOq)', '-p'])$$

$$n.betas = [(Formula,'(pOq)', 2,'beta')]$$

#### nodo\_tableaux — atributos

#### Literales

self.literales es una lista con los literales, su representación como cadena, None y la cadena 'literal'.

$$n = \mathsf{nodo\_tableaux}(['-(p > q)', '(pOq)', '-p'])$$

$$n.$$
literales =  $[(Formula, '-p', None, 'literal')]$ 

## tiene\_lit\_comp

self.tiene\_lit\_comp() retorna True sii self.literales tiene un par complementario de literales.

$$n = \mathsf{nodo\_tableaux}(['-(p > q)', '(pOq)', '-p'])$$

$$n.tiene\_lit\_comp() = False$$

#### es\_hoja

self.es\_hoja() retorna:

Cerrada si tiene literales complementarios,

Abierta si no tiene literales complementarios y no tiene reglas ni alfa ni beta,

None en otro caso.

$$n = \mathsf{nodo\_tableaux}(['-(p>q)','(pOq)','-p'])$$
 
$$n.\mathsf{es\_hoja}() = \mathtt{None}$$

## interp

self.interp() retorna un diccionario que hace verdaderos a sus literales.

$$n = \mathsf{nodo\_tableaux}(['-(p > q)', '(pOq)', '-p'])$$

$$n.interp() = \{'p' : False\}$$

#### expandir

self.expandir() retorna un nodo\_tableaux que es el resultado de aplicar la primera regla  $\alpha$ . Si no hay reglas  $\alpha$ , retorna los dos nodo\_tableaux que son el resultado de aplicar la primera regla  $\beta$ . Si no hay reglas  $\beta$ , retorna None, None.

$$n = \mathsf{nodo\_tableaux}(['-(p > q)', '(pOq)', '-p'])$$

$$n.expandir() = nodo_tableaux(['(pOq)','p','-q','-p'])$$

First In First Out

 $\mathsf{FIFO} =$ 

First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{smallmatrix} \mathsf{A} \\ & \uparrow \\ 1 \end{smallmatrix} }$$

ADD(FIFO, A)

First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{A} & \mathsf{B} \\ \hline \uparrow & \uparrow \\ 1 & 2 \end{array} }$$

ADD(FIFO, B)

First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{tabular}{c|c} A & B & \cdots & Y \\ \hline $\uparrow$ & $\uparrow$ & $\uparrow$ \\ \hline $1$ & $2$ & $^{n-1}$ \\ \hline \end{tabular}}$$

#### First In First Out

$$\mathsf{FIFO} = \underbrace{ \begin{bmatrix} \mathsf{A} & \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow & \uparrow & \uparrow \\ 1 & 2 & & ^{n-1} & n \end{bmatrix} }_{\mathsf{h} = \mathsf{I}}$$

 $\mathsf{ADD}(\mathsf{FIFO},\,\mathsf{Z})$ 

#### First In First Out

$$\mathsf{FIFO} = \underbrace{ \begin{bmatrix} \mathsf{A} & \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow & \uparrow & \uparrow \\ 1 & 2 & {}^{n-1} & n \end{bmatrix}}_{n-1}$$

$$s \leftarrow \mathsf{POP}(\mathsf{FIFO})$$

#### First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & & ^{n-2} & ^{n-1} \end{array} }$$

$$s \leftarrow A$$

#### First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & & ^{n-2} & ^{n-1} \end{array} }$$

$$s \leftarrow A$$

Last In First Out

LIFO =

First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & & n-2 & n-1 \end{array} }$$

$$s \leftarrow A$$

Last In First Out

$$\mathsf{LIFO} = \boxed{ \begin{picture}(40,0) \put(0,0){\line(1,0){100}} \put(0,0){\$$

ADD(FIFO, A)

First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & & n-2 & n-1 \end{array} }$$

$$s \leftarrow A$$

Last In First Out

$$\mathsf{LIFO} = \boxed{ \begin{tabular}{c|c} A & B \\ \hline $\uparrow$ & $\uparrow$ \\ \hline $1$ & $2$ \\ \end{tabular} }$$

ADD(FIFO, B)

#### First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & & ^{n-2} & ^{n-1} \end{array} }$$

$$s \leftarrow A$$

$$\mathsf{LIFO} = \underbrace{ \begin{bmatrix} \mathsf{A} & \mathsf{B} & \cdots & \mathsf{Y} \\ \uparrow & \uparrow & \uparrow \\ 1 & 2 & \mathit{n-1} \end{bmatrix} }_{}$$

#### First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & ^{n-2} & ^{n-1} \end{array} }$$

$$s \leftarrow A$$

#### First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & ^{n-2} & ^{n-1} \end{array} }$$

$$s \leftarrow A$$

$$\mathsf{LIFO} = \boxed{ \begin{tabular}{c|c} A & B & \cdots & Y & Z \\ \hline $\uparrow$ & $\uparrow$ & $\uparrow$ & $\uparrow$ \\ \hline $1$ & $2$ & $n-1$ & $n$ \\ \hline \end{tabular} }$$

$$s \leftarrow \mathsf{POP}(\mathsf{LIFO})$$

#### First In First Out

$$\mathsf{FIFO} = \boxed{ \begin{array}{c|c} \mathsf{B} & \cdots & \mathsf{Y} & \mathsf{Z} \\ \uparrow & \uparrow & \uparrow \\ 1 & ^{n-2} & ^{n-1} \end{array} }$$

$$s \leftarrow A$$

$$\mathsf{LIFO} = \boxed{ \begin{tabular}{c|c} A & B & \cdots & Y \\ \hline \uparrow & \uparrow & \uparrow \\ 1 & 2 & \textit{n-1} \end{tabular} }$$

$$s \leftarrow Z$$

## Presentación

- 1 Elementos de la implementación
- 2 Algoritmo primero en anchura
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- 4 Algoritmo backtracking

# Pseudocódigo

```
función primero_anchura(nodo):
 estado ← nodo
 Si estado es hoja cerrada
          retornar None
Si no, si estado es hoja abierta
          retornar estado.interp()
 frontera ← lista FIFO [estado]
 Mientras frontera no vacía
          estado ← POP(frontera)
          Para cada hijo en estado.expandir()
                     Si hijo es hoja abierta
                        retornar hijo.interp()
                     Si no, si hijo no es hoja
                        ADD(frontera, hijo)
 retornar None
```

A:  $\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$ 

Frontera

FIFO =

A:  $\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$ 

#### Frontera

$$FIFO = A$$

A: 
$$\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$$

#### Frontera

$$FIFO = A$$

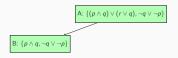
 $estado \leftarrow POP(Frontera)$ 

A:  $\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$ 

Frontera

 $\mathsf{FIFO} =$ 

 $\textit{estado} \leftarrow \mathsf{A}$ 



Frontera

 $\mathsf{FIFO} =$ 

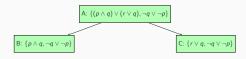
 $\textit{estado} \leftarrow \mathsf{A}$ 

 $\textit{hijo} \leftarrow \mathsf{B} \; (\mathsf{hoja} \; \mathsf{abierta?} \; \mathsf{No})$ 

Frontera

 $\textit{estado} \leftarrow \mathsf{A}$ 

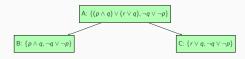
ADD(frontera, B)



Frontera

 $\textit{estado} \leftarrow \mathsf{A}$ 

 $\textit{hijo} \leftarrow C \; (\text{hoja abierta? No})$ 

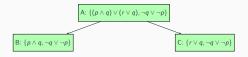


Frontera

$$FIFO = B C$$

 $estado \leftarrow A$ 

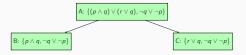
ADD(frontera, C)



#### Frontera

$$FIFO = B C$$

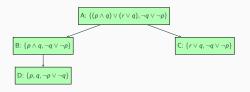
 $estado \leftarrow POP(Frontera)$ 



### Frontera

$$FIFO = C$$

 $\textit{estado} \leftarrow \mathsf{B}$ 

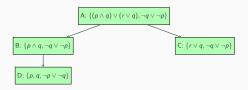


Frontera

$$FIFO = C$$

 $estado \leftarrow B$ 

 $\textit{hijo} \leftarrow \mathsf{D} \; (\mathsf{hoja} \; \mathsf{abierta?} \; \mathsf{No})$ 

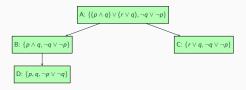


Frontera

$$FIFO = CD$$

 $estado \leftarrow B$ 

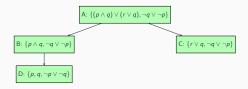
ADD(frontera, D)



### Frontera

$$FIFO = CD$$

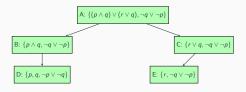
 $estado \leftarrow POP(Frontera)$ 



### Frontera

$$FIFO = D$$

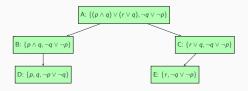
$$\textit{estado} \leftarrow C$$



Frontera

 $\textit{estado} \leftarrow C$ 

 $\textit{hijo} \leftarrow \mathsf{E} \; (\mathsf{hoja} \; \mathsf{abierta?} \; \mathsf{No})$ 

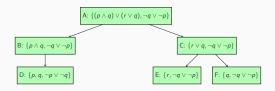


Frontera

$$FIFO = D E$$

 $\textit{estado} \leftarrow \mathsf{C}$ 

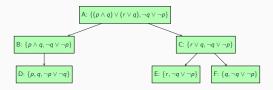
ADD(frontera, E)



#### Frontera

$$FIFO = D E$$

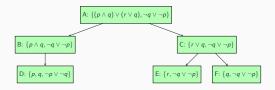
hijo ← F (hoja abierta? No)



#### Frontera

$$\mathsf{FIFO} = \boxed{\mathsf{D} \ \mathsf{E} \ \mathsf{F}}$$

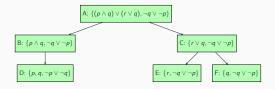
ADD(frontera, F)



#### Frontera

$$\mathsf{FIFO} = \boxed{\mathsf{D} \ \mathsf{E} \ \mathsf{F}}$$

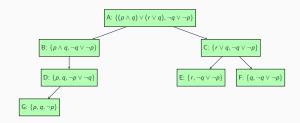
*estado* ← POP(Frontera)



### Frontera

$$\mathsf{FIFO} = \boxed{\mathsf{E} \ \mathsf{F}}$$

 $estado \leftarrow D$ 

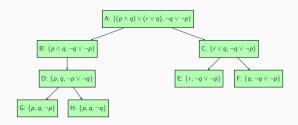


### Frontera

$$\mathsf{FIFO} = \boxed{\mathsf{E} \mathsf{F}}$$

 $estado \leftarrow D$ 

 $hijo \leftarrow G$  (hoja cerrada)

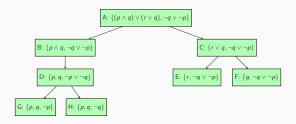


### Frontera

$$\mathsf{FIFO} = \boxed{\mathsf{E} \mathsf{F}}$$

 $estado \leftarrow D$ 

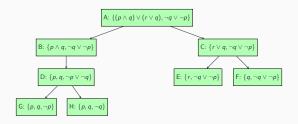
hijo ← H (hoja cerrada)



### Frontera

$$FIFO = \boxed{E} \boxed{F}$$

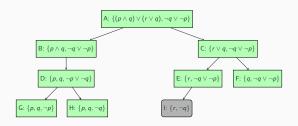
*estado* ← POP(Frontera)



### Frontera

$$FIFO = \boxed{F}$$

 $\textit{estado} \leftarrow \mathsf{E}$ 



Frontera

 $estado \leftarrow E$ 

hijo ← I (¡hoja abierta!)

### Presentación

- 1 Elementos de la implementación
- 2 Algoritmo primero en anchura
- 3 Algoritmo primero en profundidad
- 4 Algoritmo backtracking

# Pseudocódigo

```
función primero_profundidad(nodo):
 estado \leftarrow nodo
 Si estado es hoja cerrada
           retornar None
 Si no, si estado es hoja abierta
           retornar estado.interp()
 frontera ← lista LIFO [estado]
 Mientras frontera no vacía
           estado ← POP(frontera)
           Para cada hijo en estado.expandir()
                      Si hijo es hoja abierta
                        retornar hijo.interp()
                      Si no ADD(frontera, hijo)
 retornar None
```

A:  $\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$ 

Frontera

LIFO =

A:  $\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$ 

### Frontera

$$\mathsf{LIFO} = \boxed{\mathsf{A}}$$

A: 
$$\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$$

Frontera

$$LIFO = A$$

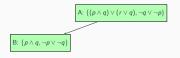
 $estado \leftarrow POP(Frontera)$ 

A:  $\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$ 

Frontera

 $\mathsf{LIFO} =$ 

 $\textit{estado} \leftarrow \mathsf{A}$ 

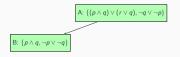


Frontera

 $\mathsf{LIFO} =$ 

 $\textit{estado} \leftarrow \mathsf{A}$ 

 $\textit{hijo} \leftarrow \mathsf{B} \; (\mathsf{hoja} \; \mathsf{abierta?} \; \mathsf{No})$ 

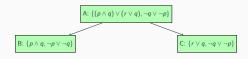


Frontera

$$LIFO = B$$

 $\textit{estado} \leftarrow \mathsf{A}$ 

ADD(frontera, B)

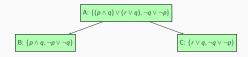


Frontera

$$LIFO = B$$

 $\textit{estado} \leftarrow \mathsf{A}$ 

 $\textit{hijo} \leftarrow C \; (\text{hoja abierta? No})$ 

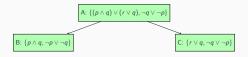


Frontera

$$LIFO = B C$$

 $estado \leftarrow A$ 

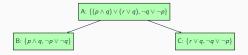
ADD(frontera, C)



### Frontera

$$LIFO = B C$$

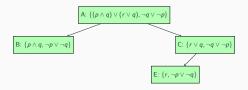
*estado* ← POP(Frontera)



Frontera

$$LIFO = B$$

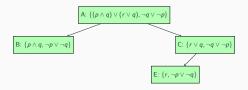
 $\textit{estado} \leftarrow C$ 



Frontera

 $\textit{estado} \leftarrow C$ 

 $\textit{hijo} \leftarrow \mathsf{E} \; (\mathsf{hoja} \; \mathsf{abierta?} \; \mathsf{No})$ 

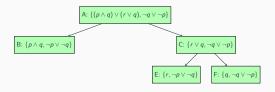


Frontera

$$LIFO = B E$$

 $\textit{estado} \leftarrow \mathsf{C}$ 

ADD(frontera, E)

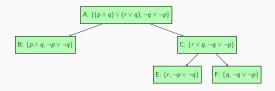


Frontera

$$LIFO = B E$$

 $\textit{estado} \leftarrow \mathsf{C}$ 

 $hijo \leftarrow F \text{ (hoja abierta? No)}$ 

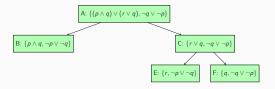


Frontera

$$LIFO = B E F$$

 $estado \leftarrow C$ 

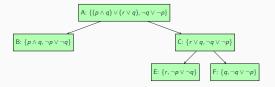
ADD(frontera, F)



#### Frontera

$$\mathsf{LIFO} = \boxed{\mathsf{B} \ \mathsf{E} \ \mathsf{F}}$$

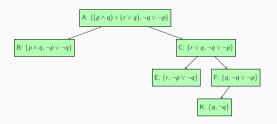
*estado* ← POP(Frontera)



#### Frontera

$$\mathsf{LIFO} = \boxed{\mathsf{B}} \boxed{\mathsf{E}}$$

 $estado \leftarrow F$ 

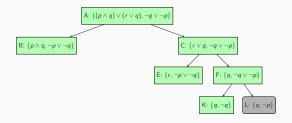


Frontera

$$LIFO = B E$$

 $\textit{estado} \leftarrow \mathsf{F}$ 

 $hijo \leftarrow K \text{ (hoja cerrada)}$ 



#### Frontera

$$\mathsf{LIFO} = \boxed{\mathsf{B}} \ \mathsf{E}$$

*hijo* ← L (¡hoja abierta!)

#### Presentación

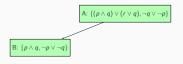
- 1 Elementos de la implementación
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- 4 Algoritmo backtracking

#### Pseudocódigo

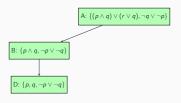
```
función backtracking(nodo):
 estado \leftarrow nodo
 Si estado es hoja cerrada
           retornar None
 Si no, si estado es hoja abierta
           retornar estado.interp()
 Para cada hijo en estado.expandir()
           resultado ← backtracking(hijo)
           Si resultado no es None
             retornar resultado
 retornar None
```

A: 
$$\{(p \land q) \lor (r \lor q), \neg q \lor \neg p\}$$

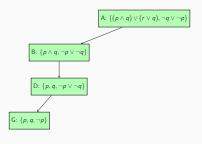
$$backtracking(A) = ?$$



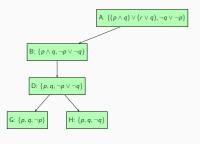
$$\begin{aligned} \mathsf{backtracking}(\mathsf{A}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{B}) &= ? \end{aligned}$$



$$\begin{aligned} backtracking(A) &= ? \\ backtracking(B) &= ? \\ backtracking(D) &= ? \end{aligned}$$



$$\begin{aligned} \mathsf{backtracking}(\mathsf{A}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{B}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{D}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{G}) &= \mathsf{None} \end{aligned}$$

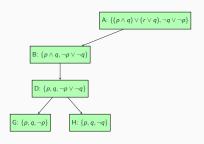


$$\mathsf{backtracking}(\mathsf{A}) = ?$$
 
$$\mathsf{backtracking}(\mathsf{B}) = ?$$

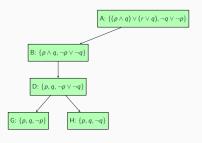
$$backtracking(D) = ?$$

backtracking(G) = None

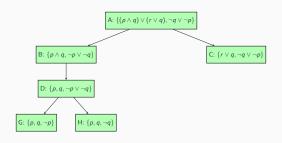
backtracking(H) = None



$$\begin{aligned} \mathsf{backtracking}(\mathsf{A}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{B}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{D}) &= \mathsf{None} \end{aligned}$$



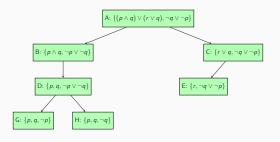
$$\begin{aligned} \mathsf{backtracking}(\mathsf{A}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{B}) &= \mathsf{None} \end{aligned}$$



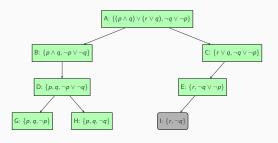
backtracking(A) = ?

backtracking(B) = None

backtracking(C) = ?



$$\begin{aligned} backtracking(A) &= ? \\ backtracking(B) &= None \\ backtracking(C) &= ? \\ backtracking(E) &= ? \end{aligned}$$

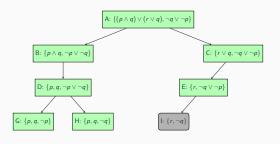


$$\mathsf{backtracking}(\mathsf{A}) = ?$$
 
$$\mathsf{backtracking}(\mathsf{B}) = \mathsf{None}$$

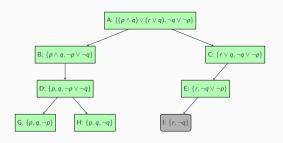
$$backtracking(C) = ?$$

$$backtracking(E) = ?$$

$$backtracking(I) = \{'q' : False, 'r' : True\}$$



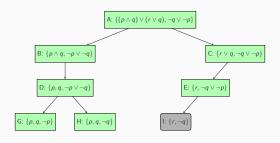
$$\begin{aligned} \mathsf{backtracking}(\mathsf{A}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{B}) &= \mathsf{None} \\ \\ \mathsf{backtracking}(\mathsf{C}) &= ? \\ \\ \mathsf{backtracking}(\mathsf{E}) &= \{'q' : \mathsf{False}, 'r' : \mathsf{True}\} \end{aligned}$$



$$backtracking(A) = ?$$

backtracking(B) = None

 $backtracking(C) = \{'q' : False, 'r' : True\}$ 



 $backtracking(A) = \{'q' : False, 'r' : True\}$ 

#### Fin de la sesión

En esta sesión usted ha aprendido a:

- 1. Visualizar el algoritmo de generación de tableaux mediante tres algoritmos: búsqueda en anchura, búsqueda en profundidad y backtracking.
- 2. Comparar empíricamente los tiempos de ejecución de los tres algoritmos de búsqueda.