

LAB #2 on DISTRIBUTED HASH TABLES

Advancing our DHT implementation in Python

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- Previous LAB we started implementing a basic DHT by implementing:
 - hash-based function to compute keys and nodeIDs
 - clockwise distance metric
 - Node class with successor and predecessor to be setup so to achieve a "DHT ring" aka "double linked-list"
 - Recursive JOIN for placing new nodes dynamically in the DHT
 - Recursive STORE and LOOKUP methods

What we (YOU!) do today

1. JOIN/LEAVE protocol:

- JOIN -> not just placing nodes, but also getting initial items from predecessor
- LEAVE -> graceful goodbye, passing own keys to new responsible node

2. FINGER TABLE! To improve efficiency

- Learn how to initialize/update a finger table
- then reimplement efficiently Lookup, this time based on Finger table
- 3. Make everything more realistic... with Flask:)

JOIN/LEAVE protocol

- JOIN: "the predecessor of JOINER has some keys that now should belong to JOINER"
 - Identify these keys and store them in JOINER
 - Then remove them from predecessor
- LEAVE: "all keys of LEAVER should go to its predecessor"
 - Store all LEAVER items at predecessor
 - Then leave by informing pred and succ to rewire so to close the gap

Towards the Finger Table

- Our JOIN/LEAVE will make the insertion and removal of nodes independent of the insertion and removal of data! :)
- However, the performance is terrible
 - O(n) with an expected performance of n/2
 - Consider a DHT with 1000 nodes & the need of setting up a TCP/IP connection for each request forwarding... traversing n/2=500 nodes can be quite slow! (e.g. $20ms \times 500 = 10s$)
- Add Finger Table to to access O(log n) performance!

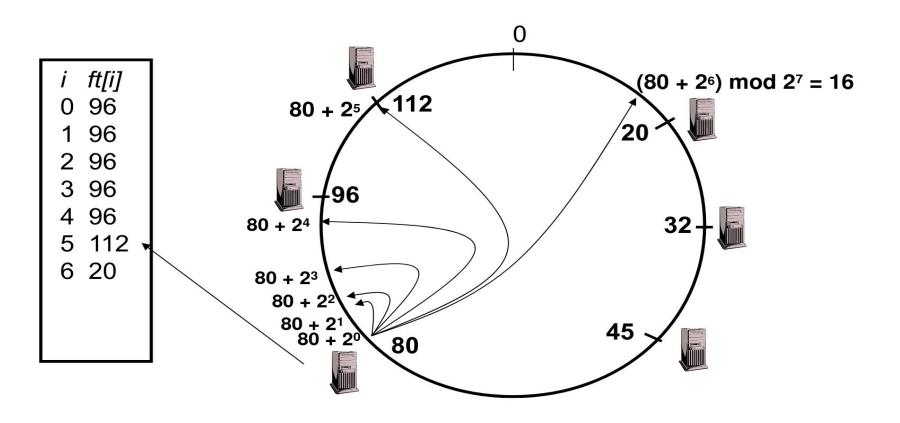
Finger Table idea

- Instead of storing a pointer to the succ node, each node stores a "finger table" containing the addresses of **k** nodes
- The distance between the current node's ID and the IDs of the nodes in the finger table increases exponentially
- So each node on the path to a given key is logarithmically closer than the last ⇒ O(log n) nodes traversed worst-case :)
- Updating a finger table requires that a node address is found for each of the **k** slots in the table...

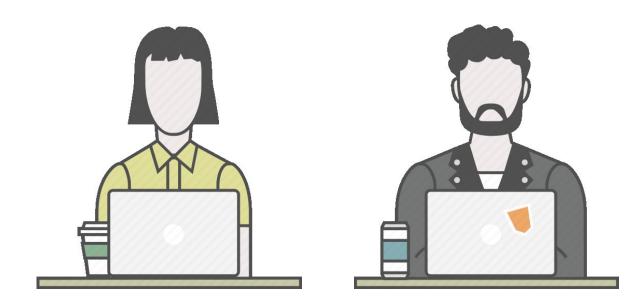
Lookup with Finger Table

- For any slot x, where x is 1 to k, finger[x] is determined by taking the current node's ID and looking up the node responsible for the key (id + $2^{(x-1)}$) % 2^k
- When doing lookups, you now have k nodes to choose from at each hop, instead of only one at each
- [SHORTCUT] When a node receive a Lookup request, it will now forward it to the node in his finger-table which has the shortest distance to the key

Finger Table Visualization



Coding Time!



Errata Corrige and hints

- OK, now you are already working on implementing JOIN/LEAVE protocol, the finger table and fingerLookup but... wait few more minutes please!
- Some corrections from the last time and further recommendations for you

compute_key(string, bitlength=k)

- We wanted a function that
 - 1. Gets the binary digest computed by some hash function (sha256)
 - 2. Extract only the k rightmost bits
 - 3. Returns the int number given by considering these k bits as an uint
- check out <u>bitstring</u>
 - pure Python module designed to help make the creation and analysis of binary data as simple and natural as possible

```
digest = sha256(bytes(string, 'utf-8')).hexdigest()
bindigest = BitArray(hex=digest).bin
subbin = bindigest[:bitlength]
return BitArray(bin=subbin).uint
```

- Separate the main file, responsible for interacting with the DHT, from the DHT code
 - e.g., LAB solution consists of 2 files, main.py and advancedDHT.py
- Define a **findNode(startnode, key)** function which...
 - "Recursively find the node whose ID is the greatest but smaller ID in the DHT compared to key"
 - it embeds the JOIN criterion of last lab, reused in lookup and store
- Define an **update()** method for the node class
 - compute there the finger table

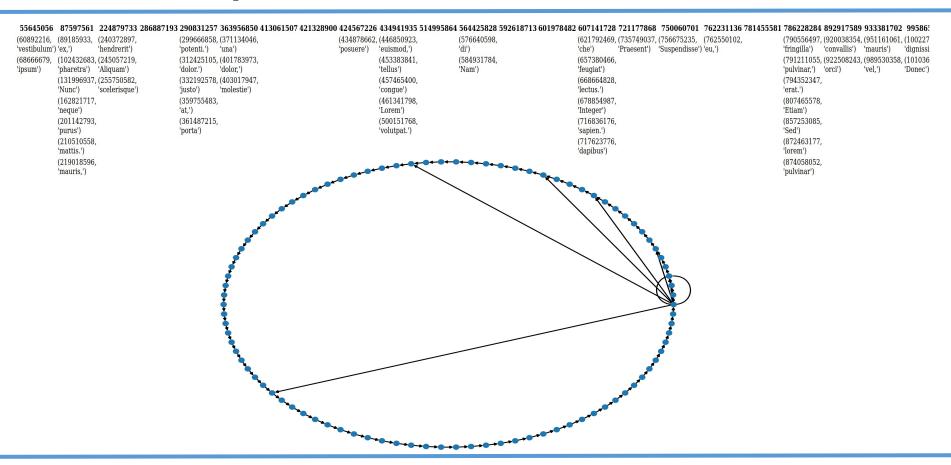
Further recommendations

- At some point it's a good idea to inject more nodes and contents to test the finger table initialization
 - printing a large DHT becomes almost impossible... use tabulate to print the DHT in html and open it in a browser (NB: tablefmt = html)
 - Draw the DHT as circular graph with networkx!!!

 nx.draw(G, pos=nx.circular_layout(G), with_labels=False, node_size=0.1)
 - For some node, draw also the finger edges, check they make "exponential jumps"

Wishful output

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Further recommendations

- In the proposed solution for the main file
- (initDHT) -> printDHT -> Create new node and JOIN -> printDHT -> make one node LEAVE -> printDHT; so you check content passing works properly
- 2. Add up to 100 nodes and 1000 items -> printDHThtml -> drawCircularGraph **NB:** BITLENGTH >> 8 or may have hash conflicts!!!
 - Check graph is truly circular ^ check from html that content is fairly distributed
- 3. Compute finger table for each node -> drawGraphWithFingerLinks4someNode
- 4. Issue many fingerLookups and standard lookups
 - keep track of recursionLevel (how many forwardings)
 - compare the recursionLevel for same key looked up with/without fingertable... who performs better???

Wishful output

Comparison of recursion level for same key looked-up at the same node first

WITH then WITHOUT finger-table

content	fingerSteps	standardSteps
nel	0	+ 17
mezzo	4	77
del	2	65
cammin	3	23
di	3	36
nostra	4	91
vita	0	6
mi	1	19
ritrovai	1	93
per	4	89
una	1	30
selva	3	87
oscura	3	24
che	1	39
la	4	91
diritta	1	70
via	0	57
era	0	57
smarrita	1	93

Questions?

