Passion-In-Action Class:	GPU101
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Assigned Project:	Breadth-First Search (BFS) Algorithm

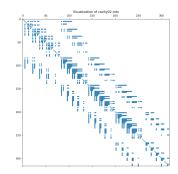
PROJECT GOAL

This Project Goal is to implement and optimise in *CUDA* programming language, for *GPU* devices, the *Breadth-First Search (BFS)* Algorithm, starting from a sequential *BFS* script wrote in *C++* programming language.

STARTING MATERIAL

The starting material is the following one:

- .mtx file rappresenting a Tree Structure as a matrix (in this case a square and simmetrical matrix of size <u>317x317</u>).
- script wrote in C++ programming language which performs the following:
 - reading and representation of the previously mentioned .mtx file in CSR format (*Compressed Sparse Row*).
 - sequential execution by the host of the BFS Algorithm in order to analyze the Tree Structure.

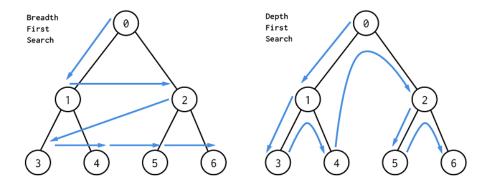


ALGORITHM EXPLANATION

The *Breadth-First Search (BFS)* Algorithm performs a <u>Breadth-Analysis</u> of a Tree Structure which, starting form a root vertex, analyses at the arising vertices by breadth (unlike the *Depth-First Search (DFS)* Algorithm, which performs a <u>Depth Analysis</u>).

The BFS Algorithm analyses first the child vertex from the root (*first frontier*), then the child of the child (*second frontier*) and so on by frontiers, until he has analysed the entire the Tree Structure, or until a vertex satisfying a certain property is found.

In this Project, using the *BFS Algortihm* allows to measure the distances between the vertices and the root, that is given by the user via the variable <u>const_int_host_source</u>.



PROCEDURE

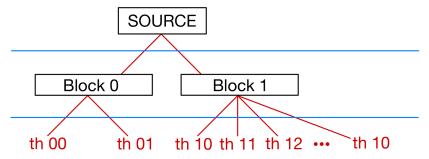
Starting from the given script, the goal is to change the function call **BFS_sequential()** with a <u>GPU-kernell</u> call of a parallelised version of the BFS Algorithm, wich has been named **BFS_gpu()**.

- 1. First of all has to be modified the *main()* function in order to satisfy the *CUDA* requirements:
 - have to be initialised some pointers for the variables to give to the kernell, paying attention to the double cases using the label *host_variable* / *gpu_variable*.
 - has to be allocated the GPU device memory via the function cudaMalloc().
 - have to be copied the data in the direction *HostToDevice* via the function cudaMemcpy().
 - only then can be done the GPU-kernell call to **BFS** gpu() with specific methodology.
 - when the kernell execution has come to an end, has to be free the *GPU* device memory via the function cudaFree().
 - everithing is checked via the macro CHECK(call) for the CUDA function calls, and via the macro CHECK_KERNELLCALL() after the <u>GPU-kernell</u> call.

```
ables allocation
 nt *gpu row ptr:
int *gpu_col_ind;
 nt *gpu_dist;
int *gpu_source;
 / gpu memory allocation
CHECK(cudaMalloc(&gpu_source, sizeof(int)));
CHECK(cudaMalloc(&gpu_row_ptr, host_row_ptr.size() * sizeof(int)));
CHECK(cudaMalloc(&gpu_col_ind, host_col_ind.size() * sizeof(int)));
CHECK(cudaMalloc(&gpu_dist, num_vals * sizeof(int)));
   Copy data from host to device
CHECK(cudaMemcpy(gpu_source, &host_source, sizeof(int), cudaMemcpyHostToDevice));
CHECK(cudaMemcpy(gpu_row_ptr, host_row_ptr.data(), host_row_ptr.size() * sizeof(int), cudaMemcpyHostToDevice));
CHECK(cudaMemcpy(gpu_col_ind, host_col_ind.data(), host_col_ind.size() * sizeof(int), cudaMemcpyHostToDevice));
CHECK(cudaMemcpy(gpu_dist, host_dist.data(), num_vals * sizeof(int), cudaMemcpyHostToDevice));
 '/ Call the gpu_kernel function
BFS_gpu<<<<DIM_GRID,DIM_BLOCK>>>(gpu_source, gpu_row_ptr, gpu_col_ind, gpu_dist);
CHECK_KERNELCALL();
// gpu memory free
CHECK(cudaFree(gpu_source));
CHECK(cudaFree(gpu_row_ptr));
CHECK(cudaFree(gpu_col_ind));
CHECK(cudaFree(gpu_dist));
```

- 2. Since the definition of the *BFS* as a *Breadth Analysis Algorithm* of a Tree Structure, for each frontier have been made the following associations:
 - each vertex corrisponds to a GPU block.
 - in order to analyse one edge of a given vertex, is used a single thread of the corrispective block; a thread may have to analyse more edges, if there are more edges in a single vertex that threads per block
 - operating this way has a significantly less impact on the performance than operating in a completely sequential way.

This associations allow to obtain a complete parallelism, taking advantage of every level that the *GPU* architecture has to offer, being limited only by the *BFS Algorithm* definition of *Breadth Analysis* per layer.



- 3. The GPU-kernell call contains four variables:
 - a pointer to the root vertex from wich the Algorithm begins const int *source_ptr.
 - a pointer <u>const int *rowPointers</u>, which indicates the first edge of a vertex as <u>rowPointers</u>[vertex].
 - a pointer const int *destinations, which indicates the destination vertex of an edge as destinations[edge].
 - a pointer int *distances, wich contains the values of the distances between each vertex and the root; this array has been set at -1 by the *main()* function.

```
// BFS algorithm optimized for GPU
__global__ void BFS_gpu(const int *source_ptr, const int *rowPointers, const int *destinations, int *distances)
{
```

- 4. For each block are initialised four <u>__shared__</u> variables:
 - an array wich represents the previous frontier of the Tree Structure int previousFrontier.
 - an array wich represents the current frontier of the Tree Structure int currentFrontier.
 - two integers used to indicate the size of the corrispective two frontiers, which are establish by the number of outgoing edges from the given frontier int_previousFrontierSize and int_previousFrontierSi

This use of the *GPU* device shared memory allows a big memory optimisation, since all the threads of the block work on the same frontier; in this way less memory is used and the vertex analysis is way more fast and parallelised.

```
// initialize frontiers
__shared__ int currentFrontier[MAX_FRONTIER_SIZE];
__shared__ int currentFrontierSize;
__shared__ int previousFrontier[MAX_FRONTIER_SIZE];
__shared__ int previousFrontierSize;
```

5. Given the pointer to the root <u>const_int_*source_ptr</u>, the next step is to define the first frontier as <u>previousFrontier[_]</u> and the first size as <u>previousFrontierSize</u>, via the <u>__device__</u> function *insertIntoFrontier()*.

Only then is possible to start a <u>while()</u> loop, wich is the core of the *BFS Algorithm*: this loop will go on until the given frontier <u>previousFrontier[]</u> has at least a new child vertex to analyse, this is checked via the condition <u>previousFrontierSize > 0</u>.

- 6. The BFS mechanism has been implemented in the following way:
 - at each iteration is analysed a single frontier, which is represented by the couple previousFrontier[] and prev
 - each frontier use only the necessary number of blocks in order to represent all the vertices, this is done via the condition if (blockldx.x < previousFrontierSize).
 - the threads are managed in a more organised way:
 - first is obtained the addres of the first edge of a given vertex, as int_row_start = rowPointers currentVertex | defining the vertex as int_row_start = previousFrontier | blockldx.x |.
 - in order to establish the limit to the vertex edges, is obtained the addres of the first edge of the next vertex, as int-row_end = rowPointers [currentVertex +1].
 - then is carried out the check <u>row_i < row_end</u> in order to establish if the thread associated to <u>int_row_i = row_start + threadIdx.x</u> has to analyse an edge of the given vertex (a thread may have to analyse from zero to more than one <u>int_row_i</u>, depending on the block size *DIM_BLOCK*).
 - for each <u>int row_i</u>, is done a check on the distance between the root and the destination vertex, if it gives as a result that <u>distances[destinations[row i]] == -1</u>, this means:
 - the destination vertex has not been explored yet, so the first thing to do is to update <u>currentFrontier[]</u> and <u>currentFrontierSize</u>, this is done via the <u>__device__</u> function <u>insertIntoFrontier()</u>.
 - the distance between the new vertex and the root vertex <u>const int source</u> is calculated as <u>distances</u>[destinations[row i]] = <u>distances</u>[currentVertex] +1.
 - once all the vertices of the frontier <u>previousFrontier[]</u>, and all the relative edges <u>rowPointers[blockldx.x]</u>, have been analysed; the vertices distances results updated into the array <u>currentFrontier[]</u>.
 - when all the threads of each block have ended the analysis, the first thread per block calls the <u>__device__</u> function **swap()** between <u>currentFrontier[]</u> and <u>previousFrontier[]</u>; after all the **swap()** executions have come to an end, a new iteration of the <u>while()</u> loop is ready to start, always if the condition <u>previousFrontierSize > 0</u> is satisfied.

CONCLUSIONS

The <u>GPU-kernell</u> call of the **BFS_gpu()** function give as a result the implementation of a completely pararellised and optimized, both at the thread and at the block levels, of the **Breadth-First Search (BFS)** Algorithm, all done in CUDA programming language, for the use of a GPU device.