

Foundations of Artificial Intelligence

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1 Course introduction

There will be no online streaming. The course topics are basically the same as the previous editions of the course so you are welcome to watch the existing recordings that will be available throughout this course. The list of links to the recordings is available on the course page.

For the topics that are not covered by previous year recordings, or they have substantially updated, we will provide additional recordings. Note that these are not new topics. They were already included in the course plan last year, but we were unable to include them during the previous edition.

The course calendar is available on the instructors' WeBeep pages (<https://webeep.polimi.it/>)

1.1 Evaluation

Evaluation is based on closed-book written exams with open questions and numerical problems. The evaluation assigns up to 32 points. The laude (30 e lode) is assigned when students receive 32 points in the written exam.

Sample exams are available on the WeBeep pages of the course. The course is quite new so older exams are partially representative of the questions you should expect in the exam. The textbook is also a good source of problems and exercises. Students can reject the assigned grade and repeat the exam. All the five exams will be in presence only. There is no other way to pass the exam or increase the grade assigned to the written exam.

2 Lecture 2

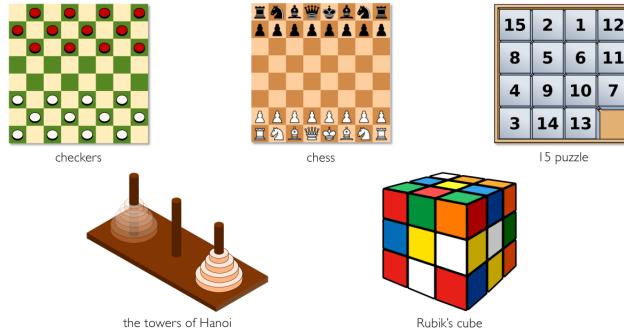
3 Problem solving by search

Problem formulation: Several real-world problems can be formulated as search problems. In a search problem, the solution can be found by exploring different alternatives.

Problem-solving agents are examples of goal-based agents

- Problem formulation
- Searching the solution
- Execute the solution

You have to start from a feasible situation, for example in the 8 problem not all configurations are solvable.



Examples of search problems.

3.1 Eight puzzle



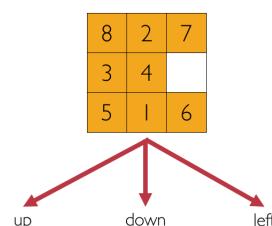
The search process

3.1.1 The states

We define a state as a feasible configuration of the 8 tiles on the 3x3 grid. The search for a solution is represented by a sequence of states in the state space.

3.1.2 The action() function

For the 8-puzzle we can define a function `actions(s)`. The function `actions(s)` returns, given a state s , the actions that are applicable in that state. An action is represented by the movement of the blank position. In this case, `actions(s)` returns `{up, down, left}`



3.1.3 The result() function

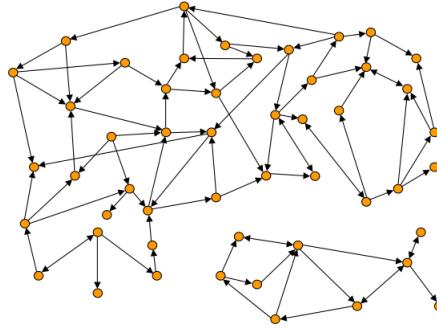
For the 8-puzzle, we can define a function $\text{result}(s,a)$ that given a state s and an action a applicable in s , returns the state s' reached by executing a in s . State s' is a successor of s .

3.2 Search problems

A set of states S and an initial state s_0 . The function $\text{actions}(s)$ that given a state s , returns the set of feasible actions. The function $\text{result}(s,a)$ that given a state s and an action a returns the state reached. A goal test that given a state s return true if the state is a goal state. A step cost $c(s,a,s')$ of an action a from s to s' .

3.2.1 The state space

The state space is a directed graph with nodes representing states, arcs representing actions. There is an arc from s to s' if and only if s' is a successor of s



The solution to a search problem is a path in the state space from the initial state to a state that satisfies the goal test

3.2.2 The optimal solution

The optimal solution is the solution with the lowest cost. The cost of a path (path cost) is the sum of the costs of the arcs that compose the path (step costs) **A problem could have no solution!** Consider, for example, a case in which the starting and the ending point are in two different "island" like in the figure (the graph) above.

How many states has the state space of the n-puzzle?

- The 8-puzzle has $9!$ (362880) states
- The 15-puzzle has $16!$ (1.3×10^{12}) states
- The 24-puzzle has $25!$ (10^{25}) states

It is usually impossible to build an explicit representation of the entire state space. Consider the n-puzzle problem: suppose to generate 100 millions of states per second, it would require 0.036 seconds to generate the 8-puzzle state space. The 15-puzzle would require a little less than 4 hours, the 24-puzzle would require more than 10⁹ years

This is also an issue of memory not just time: the number of states in the game of chess is $10^{43} - 10^{50}$, the game of go has 10^{170} states, the estimated number of atoms in the universe is between 10^{79} and 10^{81} .

A problem-solving agent must find (build) a solution by exploring only a small portion of the state space.

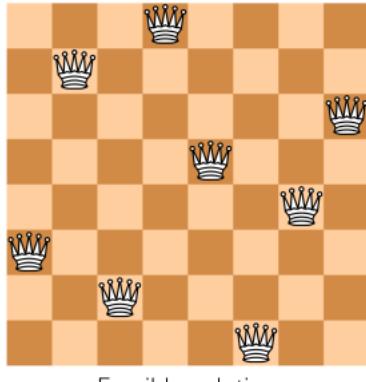
Search problems are typical of agents that operate in environments that are fully observable, static, discrete, and deterministic.

3.3 Abstraction

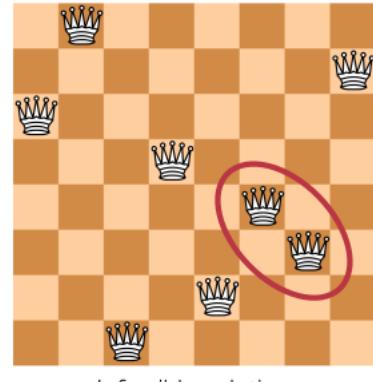
A state is an abstract representation of possible physical situations that share the same fundamental properties and differ in some details. A state of the 8-puzzle represents a set of possible physical situations that share the relative positions of the tiles but might differ in the colors of the tiles, in the materials of the tiles, etc.

3.4 The Eight-Queens Puzzle

The objective is to position 8 queens in a chessboard so that no two queens are in the same row, column, or diagonal.



Feasible solution



Infeasible solution

A first formulation could be:

- States: all arrangements of 0, 1, 2, ..., 8 queens on the board. The state space contains $64 \times 63 \times \dots \times 57 \sim 1.8 \times 10^{14}$ states
- Initial state: 0 queens on the board
- Function actions(): all the possible ways to add one queen in an empty square
- Function result()
- Goal test: 8 queens are on the board, with no queens attacking each other
- Step cost: irrelevant, unitary

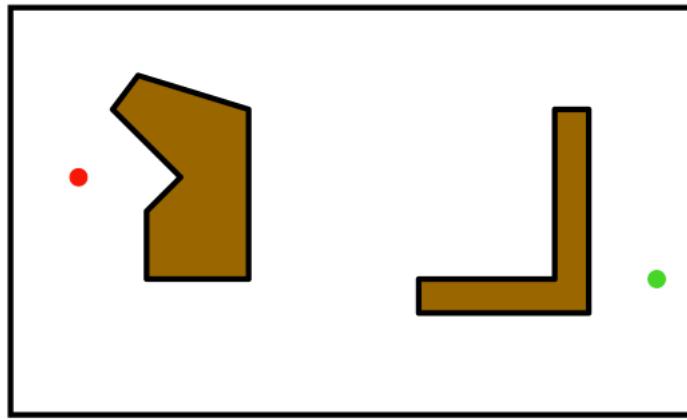
A second formulation could be:

- States: all arrangements of $k = 0, 1, 2, \dots, 8$ queens in the k leftmost columns with no two queens attacking each other
- The state space is made of 2057 states
- Initial state: 0 queens on the board
- Function actions(): all the possible ways to add one queen in any square that is not attacked by any queen already in the board, in the leftmost empty column
- Function result():
- Goal test: 8 queens are on the board
- Step cost: irrelevant, unitary

In **automatic assembly problems**, the aim is to find an order in which to assemble the parts of some object. A good automatic assembly sequence should allow to add parts later in the sequence without undoing some of the work already done.

3.5 Path planning problems

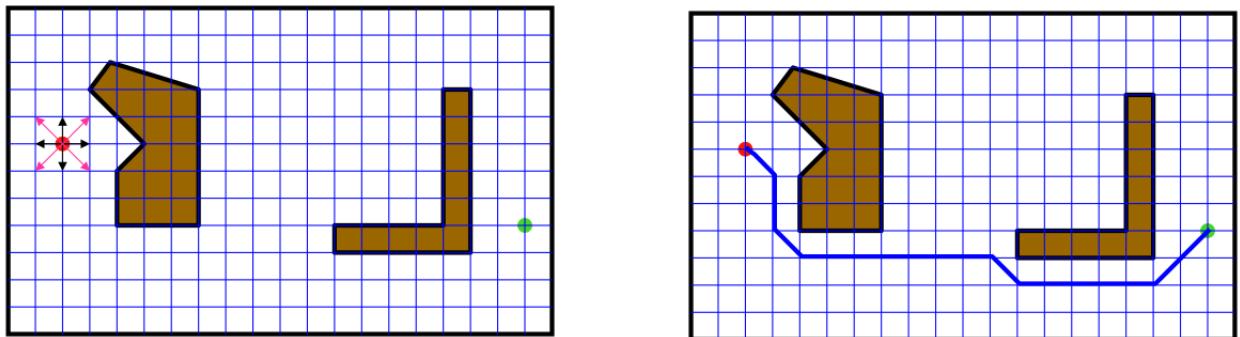
An example of path planning problem is the cleaner robot which, from a starting point, has to reach another point while avoiding obstacles.



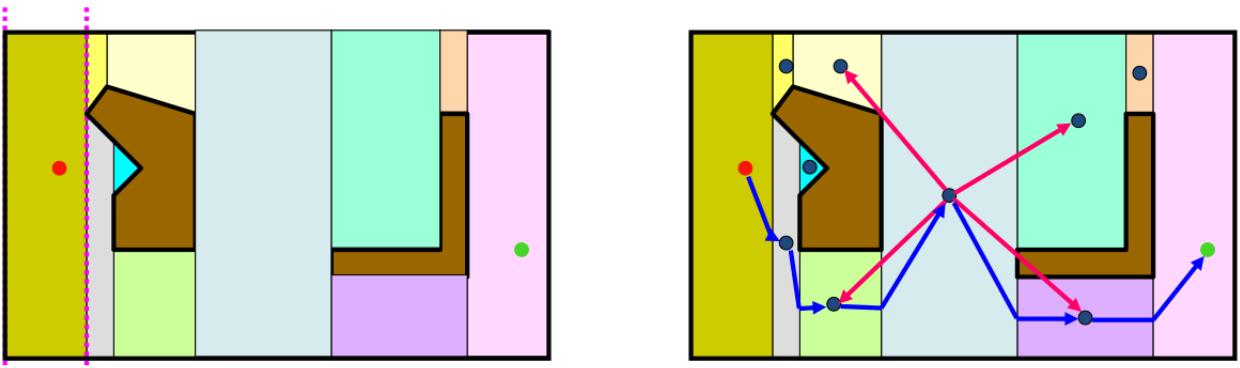
What are the possible states? How is the function actions(s) defined? And the function result(s,a)? And the goal test?

3.5.1 First formulation

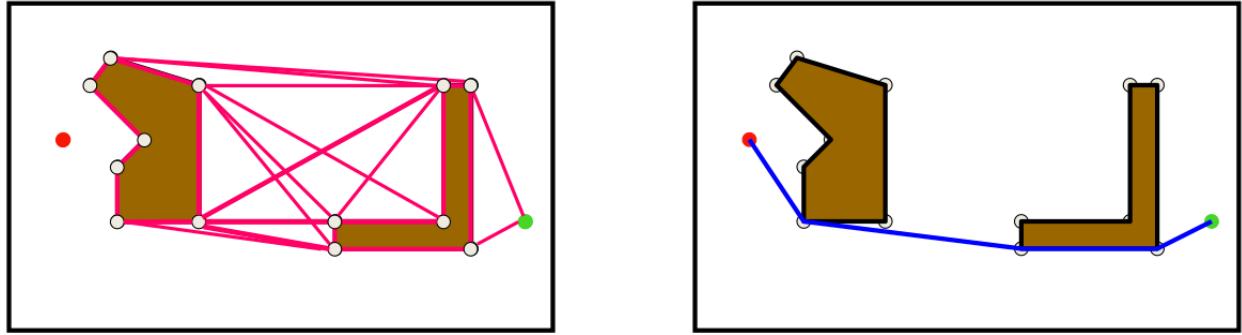
First formulation of path planning. The cost for a horizontal or vertical movement is 1; 1.41 for a diagonal movement.



3.5.2 Second formulation

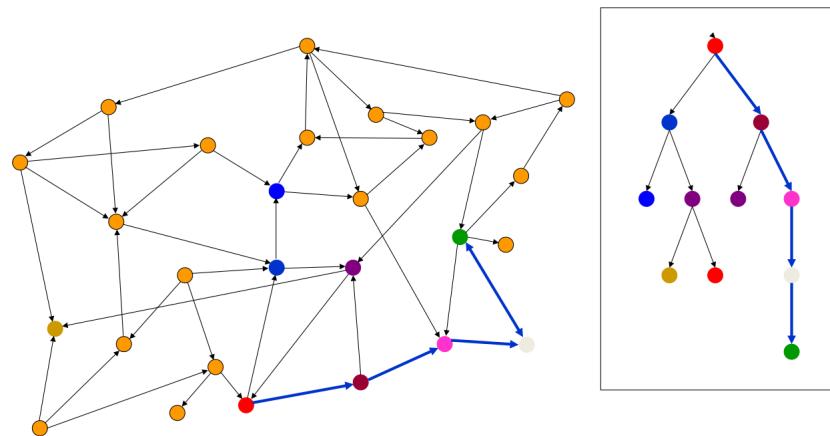


3.5.3 Third formulation



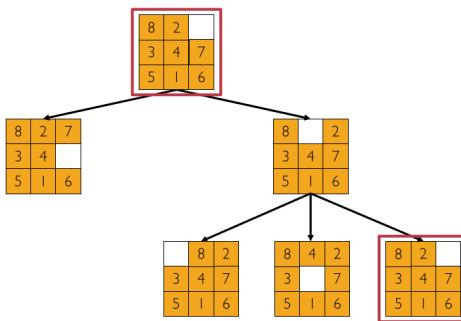
3.6 Searching for solution

Solutions to a search problem are found by building a search tree from the space state graph.



A **search tree** is composed of search nodes. Each node corresponds to a path from the initial state to a state in the state space. Each state of the space state can correspond to multiple nodes, when a state can be reached, from the initial state, following multiple paths.

If the states can be visited multiple times, then the search tree can be infinite even if the state space is finite!



In the figure you can see that the starting state is obtained again during the search. This would end up in a infinite loop of the search tree.

3.6.1 Nodes data structure

```

1 class Node:
2     def __init__(self, state, parent=None, action=None, path_cost=0):
3         """Create a search tree Node, derived from a parent by an action."""
4         self.state = state
5         self.parent = parent

```

```

6     self.action = action
7     self.path_cost = path_cost
8     self.depth = 0
9     if parent:
10        self.depth = parent.depth + 1

```

3.6.2 Tree search algorithm

Initialize the root of the tree with the initial state of problem

LOOP DO

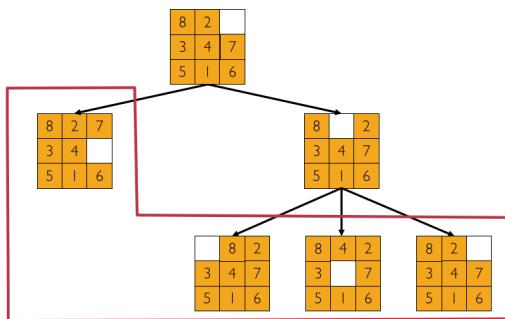
```

    IF there are no more nodes candidate for expansion
        THEN RETURN failure
    select a node not yet expanded
    IF the node corresponds to a goal state
        THEN RETURN the corresponding solution
    ELSE expand the chosen node, adding the resulting nodes to the tree

```

3.6.3 Frontier

The frontier is the set of nodes of the search tree that have been generated, but not yet chosen. The frontier is implemented as a priority queue.



- Search Strategies: A search strategy determines the ordering of nodes in the frontier
- Node Selection: The first node in the frontier is usually selected

The expansion of a node in the search tree involves two steps:

1. Apply function `actions()` to the state s of node n to compute all the available actions
2. For each action a , compute the successor state s' using `result(s, a)` and generate a child node for each successor state

The new generated nodes are inserted in the frontier. The node expansion code from the textbook's notebook is:

```

1 def expand(self, problem):
2     """List the nodes reachable in one step from this node."""
3     return [self.child_node(problem, action) for action in problem.actions(self.state)]
4 def child_node(self, problem, action):
5     """[Figure 3.10]"""
6     next_state = problem.result(self.state, action)
7     next_node = Node(next_state, self, action, \
8                      problem.path_cost(self.path_cost, self.state, action, next_state))
9     return next_node

```

The search strategies discussed so far can generate many nodes (in the search tree) corresponding to the same state (in the state space). This is unavoidable in problems with reversible actions. Such “repeated states” (revisited states) can generate infinite search trees even if the state space is finite. The search thus becomes inefficient.

To avoid repeated states, we need to be able to compare the states corresponding to nodes. We use a closed list containing the states from nodes already selected from the frontier. When a node is chosen for expansion, its state is checked against the closed list. When a node is expanded, its state is added to the closed list.

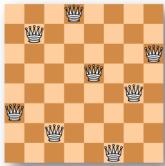
3.6.4 Graph search algorithm

```

1 initialize the root of the tree with the initial state of problem
2 initialize the closed list to the empty set
3 LOOP DO
4     IF there are no more nodes candidate for expansion
5         THEN RETURN failure
6     choose a node not yet expanded
7     IF the node corresponds to a goal state
8         THEN RETURN the corresponding solution
9     ELSE IF the state corresponding to the node is not in the list
10        THEN add the corresponding state to the closed list
11        expand the chosen node, adding the resulting nodes to the tree

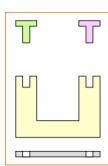
```

no repeated states few repeated states



8 queens

few repeated states

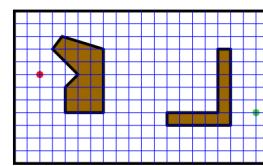


assembly planning

many repeated states



8-puzzle



robot navigation

search tree is finite

search tree is infinite

3.6.5 Best first search

```

function BEST-FIRST-SEARCH(problem,f) returns a solution node or failure
    node  $\leftarrow$  NODE(STATE=problem.INITIAL)
    frontier  $\leftarrow$  a priority queue ordered by f, with node as an element
    reached  $\leftarrow$  a lookup table, with one entry with key problem.INITIAL and value node
    while not IS-EMPTY(frontier) do
        node  $\leftarrow$  POP(frontier)
        if problem.IS-GOAL(node.STATE) then return node
        for each child in EXPAND(problem, node) do
            s  $\leftarrow$  child.STATE
            if s is not in reached or child.PATH-COST < reached[s].PATH-COST then
                reached[s]  $\leftarrow$  child
                add child to frontier
    return failure

function EXPAND(problem, node) yields nodes
    s  $\leftarrow$  node.STATE
    for each action in problem.ACTIONS(s) do
        s'  $\leftarrow$  problem.RESULT(s, action)
        cost  $\leftarrow$  node.PATH-COST + problem.ACTION-COST(s, action, s')
        yield NODE(STATE=s', PARENT=node, ACTION=action, PATH-COST=cost)

```

Figure 3.7 The best-first search algorithm, and the function for expanding a node. The data structures used here are described in Section 3.3.2. See Appendix B for **yield**.

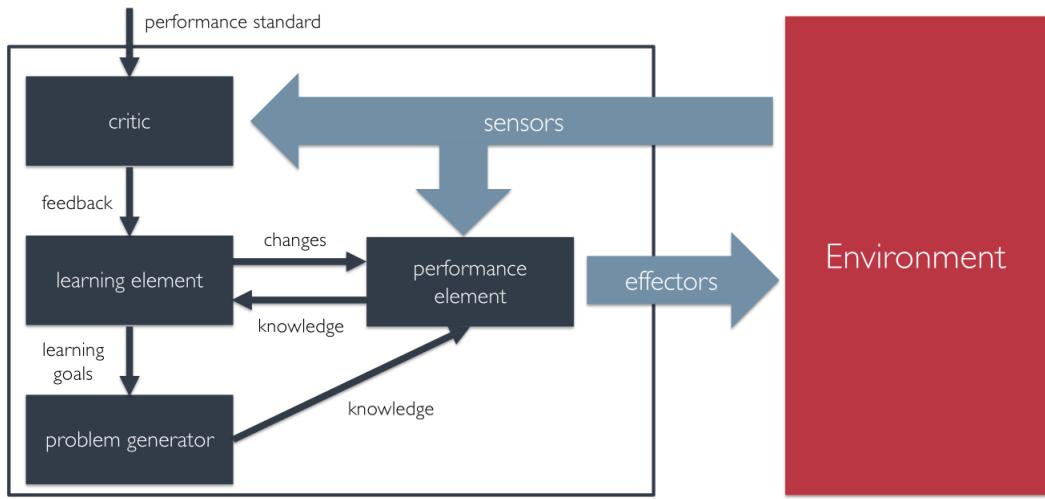
4 Learning Agents

So far we have described the following agents:

- simple reflex agents
- model-based reflex agents
- goal-based agents
- utility-based agents

Check the lecture 3 for more details. All these agents seen so far can improve their performance with learning. Every component of the agent's decisional process can be modified in order to perform better.

Learning allows the agent to operate in initially unknown environments and to become more competent than its initial knowledge alone might allow. "What kind of performance element will my agent use to do this once it has learned how?"



The agents seen till now are inside the "performance element" block.

Components:

- Learning Element: it is responsible for making improvements.
- Performance Element: it selects external actions.
- Critic: it provides feedback on the agent is doing and determines how the performance element should be modified to improve future performance
- Problem Generator: it suggests exploratory actions that will lead to new and informative experiences

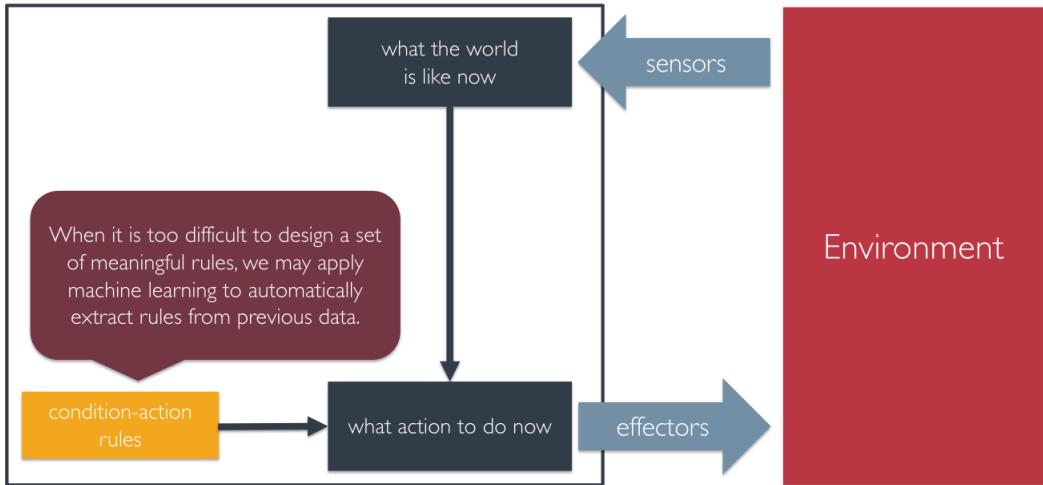
The performance element analyses the sensor inputs and decide what action to perform. Thus it is what we have previously considered to be the entire agent.

4.1 Machine learning

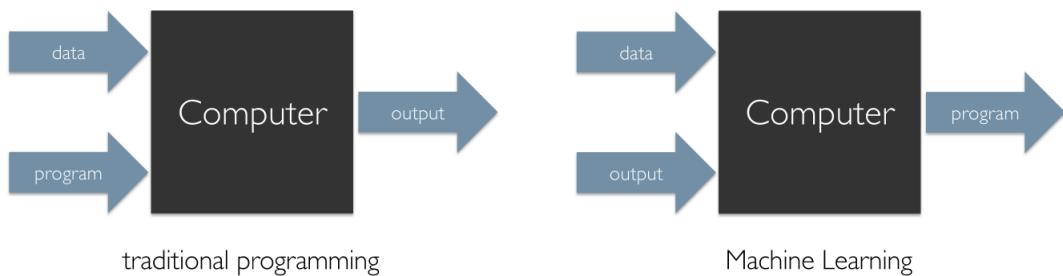
"A computer program is said to learn from experience E with respect to some class of tasks T and performance measure P, improves with experience E" Mitchell (1997).

It is an area of Artificial Intelligence focused on building algorithms capable of learning, extracting knowledge from experience. Machine Learning algorithms extract knowledge, they cannot create it. The goal is to build programs that can make informed decisions on new unseen data.

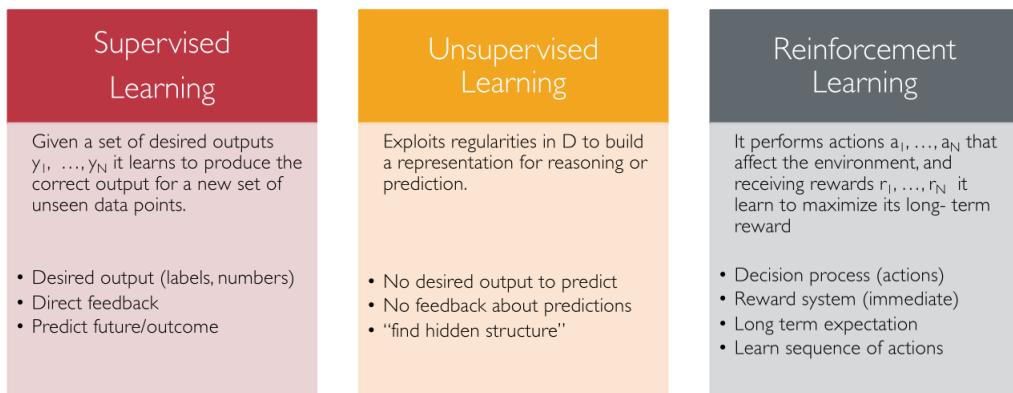
Example of application of machine learning to extract rules for a simple reflex agent:



Differences between programming and machine learning:



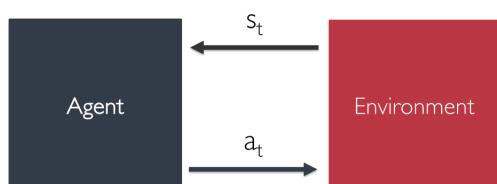
Suppose we have the experience we collected encoded as a dataset $D = x_1, \dots, x_N$:

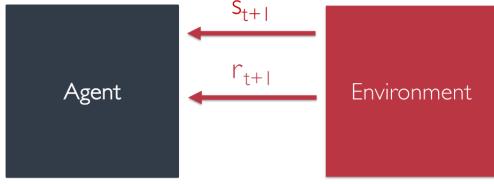


4.1.1 Reinforcement learning

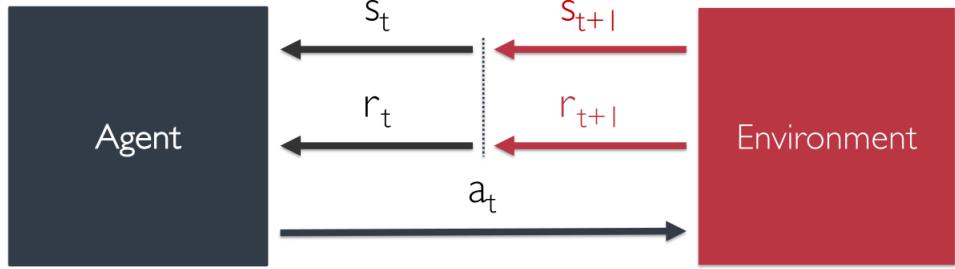
In sequential decision making:

- we take a sequence of decisions (or actions) to reach the goal
- the optimal actions are context-dependent
- we generally do not have examples of correct actions
- actions may have long-term consequences
- short-term consequences of optimal actions might seem negative





At time t , the agent perceives the environment to be in state s_t and decides to perform action a_t . As a result, in the next time step $t + 1$ the environment state changes to s_{t+1} and the agent receives a reward r_{t+1} . This is the generic agent-environment interaction in reinforcement learning:



$$\dots \rightarrow (S_t, A_t, R_{t+1}, S_{t+1}, A_{t+1}, R_{t+2}, S_{t+2}, A_{t+2}, R_{t+3}, S_{t+3}, A_{t+3}, \dots)$$

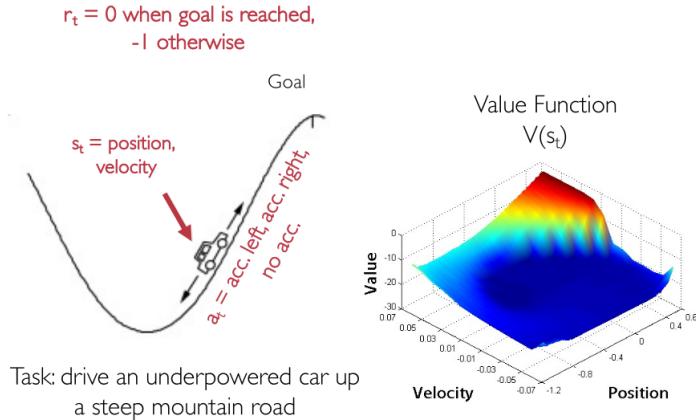
The agent's goal is to maximize the total amount of reward received. How much future reward will it get when it performs a_t in s_t and then continues to do its best from there on? What is the expected payoff from s_t and a_t ? The agent needs to compute an **action-value function** mapping state-action pairs to expected payoffs:

$$Q(a_s, s_t) \rightarrow \text{payoff}$$

or a **state-value functions** mapping to expected payoffs

$$V(s_t) \rightarrow \text{payoff}$$

Reinforcement learning assumes that $Q(s_t, a_t)$ is represented as a table. But the real world is complex, the number of possible inputs can be huge! We cannot afford to compute an exact $Q(s_t, a_t)$ (more about this later).



What the car should do is to swing back and forth in order to achieve enough speed to climb the mountain.

Action selection

At each time step, the agent must decide what action to take in step t based on its current evaluation of the expected payoff in s_t using a policy function. At any given point in time, a policy $\pi(s_t)$ select what actions the agent should perform. The policy defines the behavior of an agent based on its payoff evaluation. The policy can be deterministic or stochastic.

- Deterministic policy

- In the simplest case the policy can be modeled as a function $\pi : \mathcal{S} \rightarrow \mathcal{A}$.
- For example, the policy might simply select the action with the largest expected payoff
- This type of policy can be conveniently represented using a table
- Stochastic policy
 - It maps each state to a probability distribution over the actions $\pi : \mathcal{S}, \mathcal{A} \rightarrow \mathcal{R}$
 - $\pi(s, a)$ returns the probability of selecting a in s
 - Since $\pi(s, a)$ is a probability distribution, it always return value greater or equal to zero and the sum over all the actions is 1
 - A stochastic policy can be used to represent also a deterministic policy

To obtain a lot of reward, the agent must prefer actions that it has tried in the past and found to lead to high payoff. However, to discover such actions, it has to try actions that it has not selected before. The agent needs to find a trade-off between the exploration of new actions and the exploitation promising actions. This is called exploration-exploitation dilemma.

- Greedy Policy: for each state, it deterministically selects an action with maximal value
- ϵ -Greedy Policy: with probability ϵ it performs a random action, with probability $1 - \epsilon$ it performs the action promising the highest payoff

The environment

The environment must satisfy the Markov property. The next state s_{t+1} and reward r_{t+1} depend only on the current state s_t and action a_t . The environment can thus be modeled as a **Markov Decision Process (MDP)** that has a one-step dynamic described by the probability distribution $p(s_{t+1}, r_{t+1}|s_t, a_t)$.

- $p : \mathcal{S} \times \mathcal{R} \times \mathcal{S} \times \mathcal{A} \rightarrow [0, 1]$
- $\sum_{s \in \mathcal{S}} \sum_{r \in \mathcal{R}} p(s', r|s, a) = 1 \quad \forall s \in \mathcal{A}, \forall a \in \mathcal{A}(s)$

Expected payoff

In reinforcement learning, the agent has to maximize the reward it receives in the long run

$$G_t \doteq r_{t+1} + r_{t+2} + r_{t+3} + \dots r_{t+k} + \dots \stackrel{?}{=} \infty$$

To provide an upper bound to the payoff, we introduce a discount the future rewards by a factor $\gamma \in (0, 1)$:

$$G_t \doteq r_{t+1} + \gamma r_{t+2} + \gamma^2 r_{t+3} + \dots \gamma^{k-1} r_{t+k} + \dots < \infty$$

Thus, the expected reward to maximize will be defined as:

$$\mathbb{E}[G_t] = \mathbb{E} \left[\sum_{k=0}^{\infty} \gamma^k R_{t+k+1} \right] \leq R_{max} \frac{1}{1-\gamma}$$

The reward hypothesis

“That all of what we mean by goals and purposes can be well thought of as maximization of the expected value of the cumulative sum of a received scalar signal (reward).” (Rich Sutton)

In reinforcement learning, the agent learns how to maximize the expected future payoff. We must design a reward function that adequately represents our learning goal. Examples

- Goal-reward representation returns 1 when the goal is reached, 0 otherwise
- Action-penalty representation returns -1 when the state is not the goal, 0 once the goal is reached

Challenges to reward hypothesis: how to represent risk-sensitive behavior? How to capture diversity in behavior?

The value function The action-value function $Q(s_t, a_t)$ estimates the expected future payoff when performing action a_t in state s_t . The state-value function $V(s_t)$ estimates the expected future payoff starting from s_t (in the former case) They can be both decomposed as the sum of the immediate reward received r_{t+1} and the future rewards.

The state-value function can again be decomposed into immediate reward plus discounted value of successor state (Bellman Expectation Equation):

$$V(s) = \mathbb{E}[r_{t+1} + \gamma V(s_{t+1})|s_t = s]$$

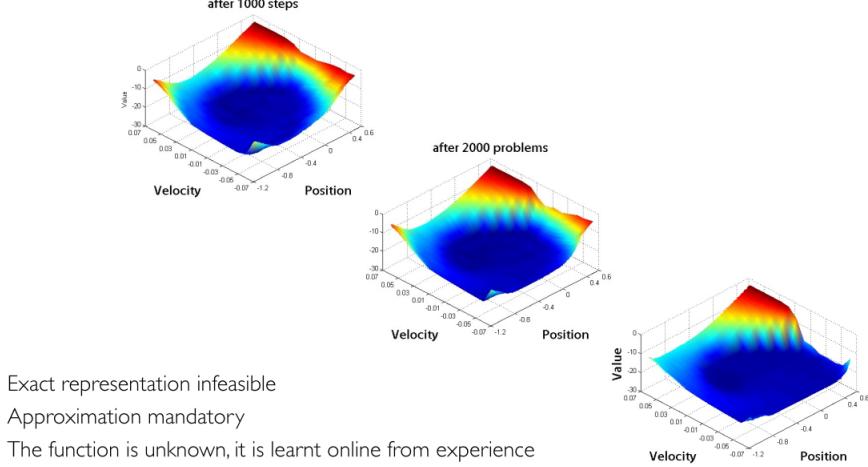
The action-value function can be similarly decomposed:

$$Q(s, a) = \mathbb{E}[r_{t+1} + \gamma V(s_{t+1})|s_t = s, a_t = a]$$

At the beginning the table $Q(\cdot, \cdot)$ is initialized with random values. At time t ,

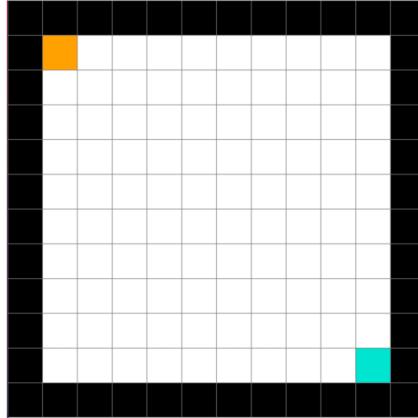
$$Q(s_t, a_t) \leftarrow Q(s_t, a_t) + \beta \left[r_{t+1} + \gamma \max_{a \in \mathcal{A}} Q(s_{t+1}, a) - Q(s_t, a_t) \right]$$

The parameters are the discount factor γ , the learning rate β , the action selection strategy $\pi(s_t, a_t)$; ϵ -Greedy is the most common choice during learning but sufficient exploration must be guaranteed to tackle the exploration-exploitation dilemma.

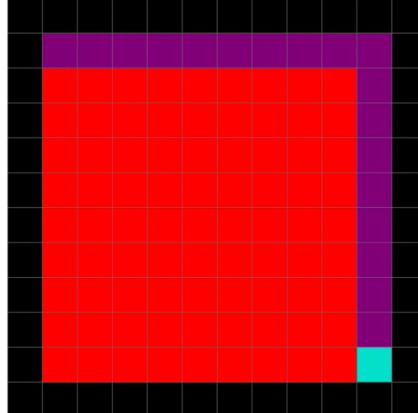


Tabular representation is infeasible in practice and approximators must be used for interesting problems. Reinforcement learning computes an unknown value function while also trying to approximate it. Approximator works on intermediate estimates while also providing information for the learning. Convergence is not guaranteed.

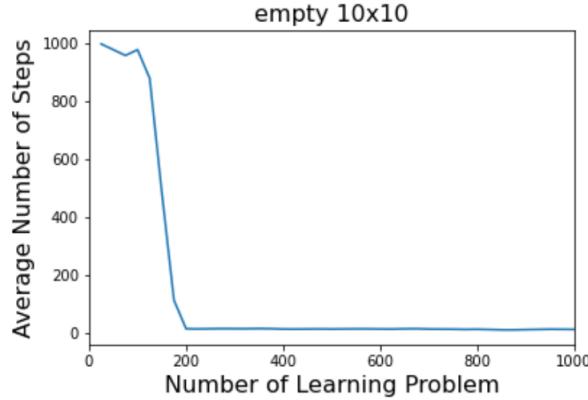
Let's now consider this simple empty environment with one start position (yellow) and one goal position (blue):



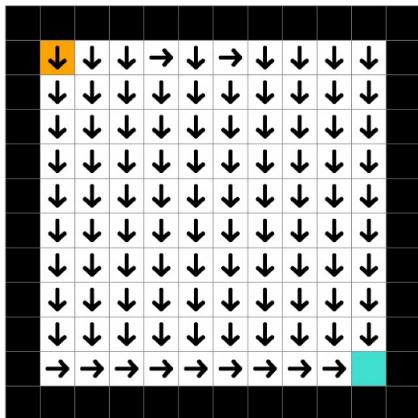
Here is the solution using a search algorithm:



When applying reinforcement learning we need to select a reward function. Let's keep it simple, zero everywhere, except when we reach the goal when we reach the goal, we receive one. Let's measure the performance as the average number of steps in the last 100 problems.

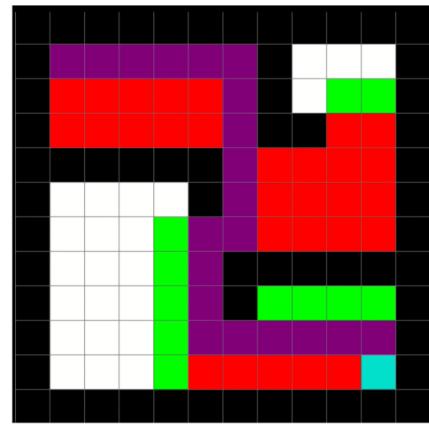
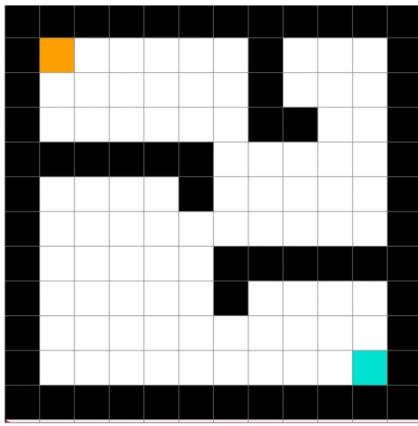


Search finds a solution, reinforcement learning an action value function. Here in the figure below is represented the best action for every position based on the action-value function on the left, while on the right the computed action value function:

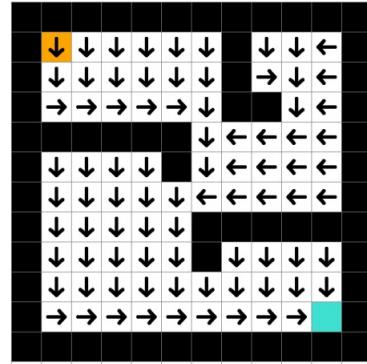
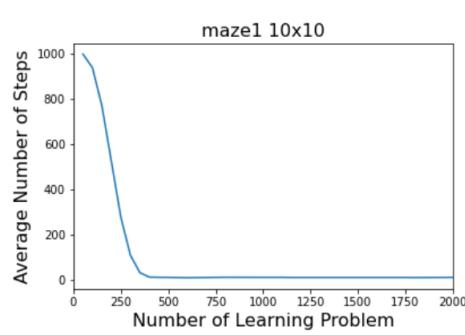


State	Up	Right	Down	Left
(1, 2)	0.397	0.440	0.440	0.000
(1, 3)	0.418	0.463	0.463	0.000
(1, 4)	0.440	0.488	0.488	0.000
(1, 5)	0.463	0.513	0.513	0.000
(1, 6)	0.488	0.540	0.540	0.000
(1, 7)	0.513	0.569	0.569	0.000
(1, 8)	0.540	0.599	0.599	0.000
(1, 9)	0.569	0.630	0.630	0.000
(1, 10)	0.599	0.663	0.000	0.000
(2, 1)	0.000	0.440	0.440	0.397
(2, 2)	0.418	0.463	0.463	0.418
(2, 3)	0.440	0.488	0.488	0.440
(2, 4)	0.463	0.513	0.513	0.463
(2, 5)	0.488	0.540	0.540	0.488
(2, 6)	0.513	0.569	0.569	0.513
...

Let's consider another example. On the left is the new problem to solve while on the right there is the solution found by using the A* search algorithm:



Down below the solution obtained using reinforcement learning:



On the slides there is also the computed action value function for this last example problem.

4.1.2 Supervised learning

Given a set of inputs-output pairs $(x_1, y_1), \dots, (x_N, y_N)$ it learns to produce the correct output for a new set of unseen data points. When the target values are real values, we have a regression problem. When they are discrete or symbolic, we have a classification problem.

5 Uninformed search strategies

Search Strategies: a search strategy determines the ordering of nodes in the frontier. Selecting a node amounts to decide the ordering of the frontier and to select the first node.

Uninformed Search Strategies: they use only the information contained in the problem formulation.

The Five Elements of Uninformed Search Strategies

1. The set of states (and an initial state)
2. Function actions()
3. Function result()
4. Goal test
5. Step cost

5.1 Tree-search algorithm

```
1 Initialize the root of the tree with the initial state of problem
2 LOOP DO
3     IF there are no more nodes candidate for expansion
4         THEN RETURN failure
5     select a node not yet expanded
6     IF the node corresponds to a goal state
7         THEN RETURN the corresponding solution
8     ELSE expand the chosen node, adding the resulting nodes to the tree
```

5.2 Graph-search algorithm

Here is introduced the closed list in order to avoid visiting the same node multiple times.

```
1 initialize the root of the tree with the initial state of problem
2 initialize the closed list to the empty set
3 LOOP DO
4     IF there are no more nodes candidate for expansion
5         THEN RETURN failure
6     choose a node not yet expanded
7     IF the node corresponds to a goal state
8         THEN RETURN the corresponding solution
9     ELSE IF the state corresponding to the node is not in the closed list
10        THEN add the corresponding state to the closed list expand the chosen node, adding the
11        resulting nodes to the tree
```

5.3 Evaluation of search strategies

- **Completeness:** is the search strategy guaranteed to find a solution when there is one?
- **(Cost) Optimality:** does the search strategy find the optimal solution (minimum cost solution)?
- **(Time and Space) Complexity:** how much time and memory does the search strategy take to perform the search?
- **Parameters:**
 - **b**, the branching factor of the search tree (the maximum number of successors of any node)
 - **d**, the depth d of the shallowest goal node

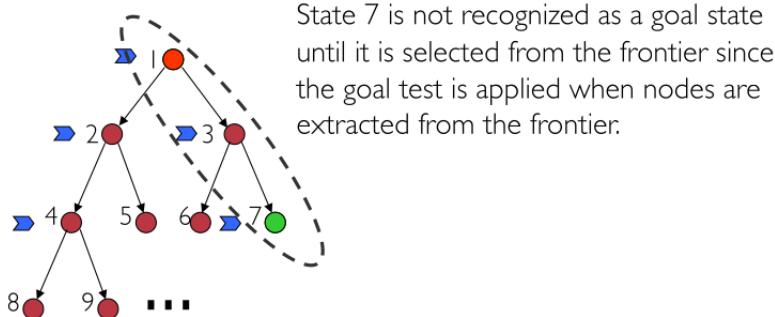
When the branching factor is large, the problem is difficult, the number of nodes will grow quickly. When the value of d is large the closest solution will require more exploration of the search tree. Sometimes, the depth d is fixed for any solution (like for example, in the eight-queen puzzle)

5.4 Breadth-first search

First, it selects the root, then all the root's successors, then all their successors, and so on. In breadth-first search, all nodes at level k of the search tree must be selected before selecting any node at level $k + 1$.

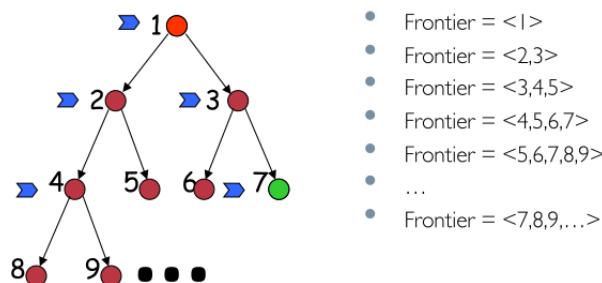
Breadth-first search always selects the node with minimum depth (the shallowest one). When more nodes are at the same depth, a tie-breaker is applied determined by the insertion order.

Example:



The node "2" is expanded before the node "3" so in a certain instant the frontier is: 3,4,5 as it is written also in the figure below. The node "7" is not immediately recognized as a goal, it is inserted in the frontier before. Once the goal node is reached, the path to the root is selected as solution.

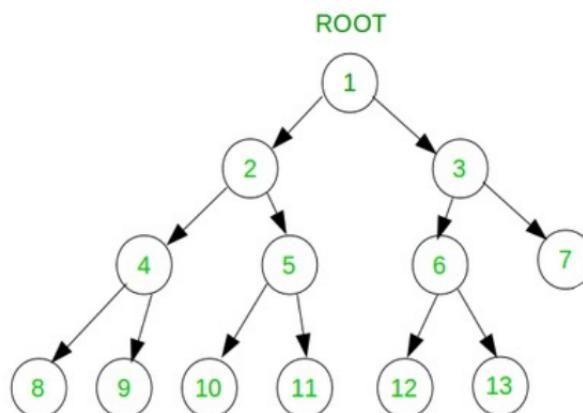
Breadth-first search can be implemented by considering the frontier a FIFO (First In First Out) queue. The new nodes are appended at the end of the frontier. Thus, the nodes generated first are selected first.



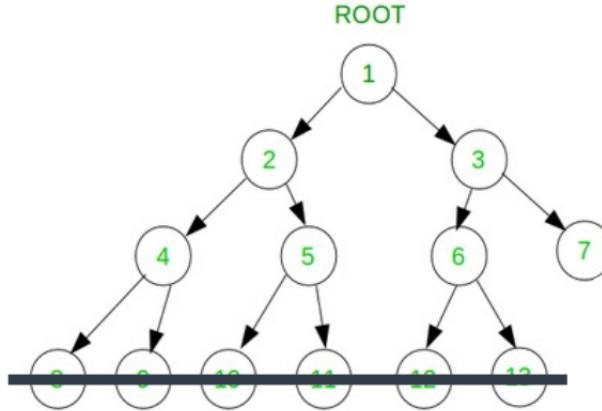
5.4.1 Evaluation of Breadth-First Search

- Breadth-first search is complete: assuming b is finite
- Breadth-first search is optimal: it finds the shallowest solution when step costs are all equal (e.g., they are unitary). When the path cost to any node is a non-decreasing function of the depth of the node. The cost for the second level is greater than the first, and so on.
- Complexity: breadth-first search has temporal and spatial (worst-case) complexity of $O(bd + 1)$ nodes. The root of the search tree generates b nodes at level 1. Each one of them generates b nodes, totaling b^2 nodes at level 2. Thus, at level $d+1$ there will be $b^{d+1} - b$ nodes

Worst case of Breadth-First Search:



In breadth-first search, we can perform the goal test when a node is generated, instead of when it is selected from the frontier. This variant is still complete and optimal (when the path cost to any node is a non-decreasing function of the depth of the node). This variant has temporal and spatial (worst-case) complexity of $O(b^d)$ nodes since we avoid expanding a node before adding it to the frontier.



Here is the breadth-first search algorithm from the textbook:

```

function BREADTH-FIRST-SEARCH(problem) returns a solution node or failure
    node  $\leftarrow$  NODE(problem.INITIAL)
    if problem.IS-GOAL(node.STATE) then return node
    frontier  $\leftarrow$  a FIFO queue, with node as an element
    reached  $\leftarrow$  {problem.INITIAL}
    while not IS-EMPTY(frontier) do
        node  $\leftarrow$  POP(frontier)
        for each child in EXPAND(problem, node) do
            s  $\leftarrow$  child.STATE
            if problem.IS-GOAL(s) then return child
            if s is not in reached then
                add s to reached
                add child to frontier
    return failure

```

And the python version with the goal test when a node is selected from the frontier is:

```

1 def breadth_first_search(problem, verbose=False):
2     node = Node(problem.initial)
3     frontier = FIFOQueue([node])
4     reached = {problem.initial}
5     while frontier:
6         node = frontier.pop()
7         if problem.is_goal(node.state):
8             return node
9         for child in expand(problem, node):
10            s = child.state
11            if s not in reached:
12                reached.add(s)
13                frontier.appendleft(child)
14    return failure, visited_states

```

A similar python version with the goal test when a node is generated:

```

1 def breadth_first_search_early_goal_test(problem):
2     node = Node(problem.initial)
3     if problem.is_goal(problem.initial):
4         return node
5     frontier = FIFOQueue([node])

```

```

6     reached = {problem.initial}
7     while frontier:
8         node = frontier.pop()
9         for child in expand(problem, node):
10            s = child.state
11            if problem.is_goal(s):
12                return child
13            if s not in reached:
14                reached.add(s)
15                frontier.appendleft(child)
16    return failure

```

We could implement breadth-first search as a call to Best-First-Search where the evaluation function $f(n)$ is the depth of the node (the number of actions it takes to reach the node). However, the native implementation is more efficient since:

- A first-in-first-out queue will be faster than a priority queue
- Reached can be a set of states rather than a mapping from states to nodes, because once we've reached a state, we can never find a better path to the state
- Thus, breadth-first search can perform an early goal test, checking whether a node is a solution as soon as it is generated, rather than the late goal test that best-first search uses, waiting until a node is popped off the queue

5.5 Uniform cost search

It generalizes breadth-first search: it sorts the nodes in the frontier according to their increasing path cost from the root and it selects the node n with the smallest path cost from the root. It does not necessarily choose the shallowest node! It is implemented as best-first-search where the evaluation function $f(n)$ returns the path cost.

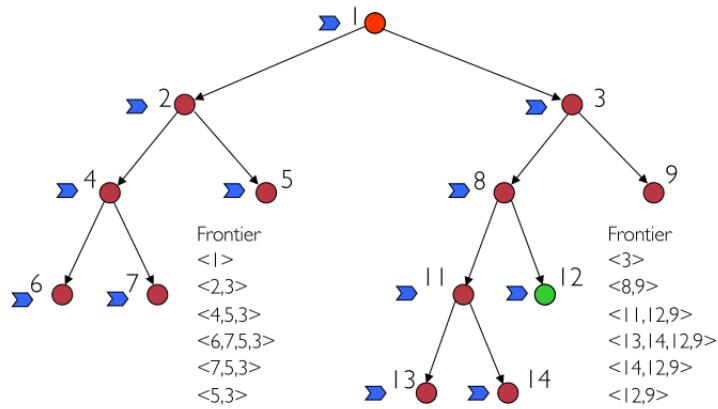
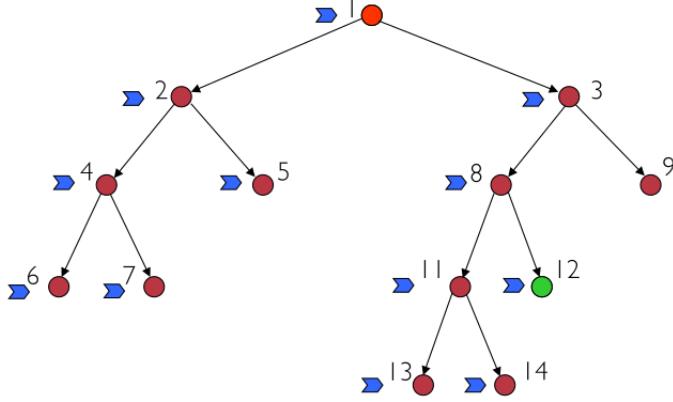
function UNIFORM-COST-SEARCH(*problem*) returns a solution node, or *failure*
return BEST-FIRST-SEARCH(*problem*, PATH-COST)

5.5.1 Evaluation of Uniform-Cost Search

- Uniform-cost search is complete: assuming b finite and step costs larger than a small positive ϵ
- Uniform-cost search is optimal: uniform-cost search expands nodes in order of their optimal path costs
- Complexity: Uniform-cost search has temporal and spatial worst-case complexity of $O(b^{(1+\lfloor C^*/\epsilon \rfloor)})$ nodes, where C^* is the cost of the optimal solution. The fraction at the exponent represents the depth of the tree since ϵ is very small will always be chosen by the algorithm.

5.6 Depth-first search

First, it selects the root node, then, it expands one of its successors, then one of its successors, and so on. Depth-first search always selects the node at maximum depth (the deepest one). When it reaches a node without successors, it “backtracks” and chooses one of the deepest nodes not yet chosen.



The python *recursive implementation* of depth-first search is:

```

1 def depth_first_recursive_search(problem, node=None):
2     if node is None:
3         node = Node(problem.initial)
4     if problem.is_goal(node.state):
5         return node
6     elif is_cycle(node):
7         return failure
8     else:
9         for child in expand(problem, node):
10            result = depth_first_recursive_search(problem, child)
11            if result:
12                return result
13    return failure

```

The python *iterative implementation* of depth-first search is:

```

1 def depth_first_search_iterative(problem):
2     "Search deepest nodes in the search tree first."
3     frontier = LIFOQueue([Node(problem.initial)])
4     result = failure
5     while frontier:
6         node = frontier.pop()
7         if problem.is_goal(node.state):
8             return node
9         elif not is_cycle(node):

```

```

10         for child in expand(problem, node):
11             frontier.append(child)
12     return result
13
14 def is_cycle(node, k=30):
15     "Does this node form a cycle of length k or less?"
16     def find_cycle(ancestor, k):
17         return (ancestor is not None and k > 0 and
18                 (ancestor.state == node.state or find_cycle(ancestor.parent, k - 1)))
19     return find_cycle(node.parent, k)

```

5.6.1 Evaluation of Depth-First Search

- Depth-first search is not complete: it can follow path of infinite length
- Depth-first search is not optimal: it does not find the shallowest solution, might go very deeply in a wrong direction.
- Complexity: given the maximum depth of the search tree m , depth-first has spatial (worse-case) complexity of $O(bm)$ nodes. It has temporal (worst-case) complexity of $O(b^m)$ nodes; m could be infinite; $m \geq d$ (often $m \gg d$). At each step of the search, only a single path (from the root to a leaf node) must be stored in memory; also, the nodes not yet expanded at each level should be stored

Also depth-first search could be implemented as a call to best-first search where the evaluation function $f(n)$ is the negative of the depth. However, it is usually implemented not as a graph search but as a tree-like search that does not keep a table of reached states.

5.7 Depth-limited search

Depth-limited search is a depth-first search with a predetermined depth limit L . Nodes at level L are assumed to have no successors. Depth-limited search is complete when $L \geq d$, however, d is often unknown.

Depth-limited search is not optimal. Complexity: Depth-limited search has spatial (worst-case) complexity of $O(bL)$ nodes, it has temporal (worst-case) complexity of $O(b^L)$ nodes.

Python recursive implementation of depth-limited search:

```

1 def depth_limited_search(problem, limit=40):
2     "Search deepest nodes in the search tree first."
3     frontier = LIFOQueue([Node(problem.initial)])
4     result = failure
5     while frontier:
6         node = frontier.pop()
7         if problem.is_goal(node.state):
8             return node
9         elif len(node) >= limit:
10            result = cutoff
11        elif not is_cycle(node):
12            for child in expand(problem, node):
13                frontier.append(child)
14    return result

```

5.8 Iterative deepening search

It performs repeated depth-limited searches, increasing the depth limit $L : L = 0, L = 1, L = 2, \dots$. Recursive python implementation of iterative deepening search:

```

1 def iterative_deepening_search(problem):
2     "Do depth-limited search with increasing depth limits."
3     for limit in range(1, sys.maxsize):
4         result = depth_limited_search(problem, limit)

```

```

5     if result != cutoff:
6         return result

```

```

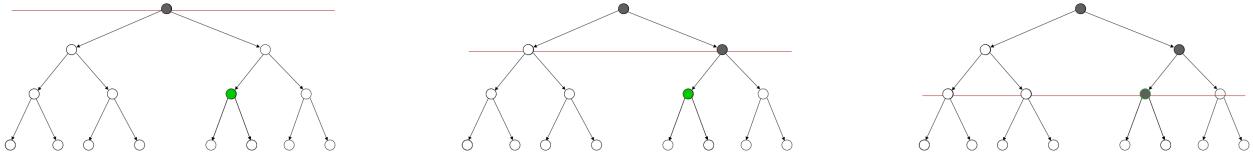
function ITERATIVE-DEEPENING-SEARCH(problem) returns a solution node or failure
    for depth = 0 to  $\infty$  do
        result  $\leftarrow$  DEPTH-LIMITED-SEARCH(problem, depth)
        if result  $\neq$  cutoff then return result

function DEPTH-LIMITED-SEARCH(problem, ℓ) returns a node or failure or cutoff
    frontier  $\leftarrow$  a LIFO queue (stack) with NODE(problem.INITIAL) as an element
    result  $\leftarrow$  failure
    while not Is-EMPTY(frontier) do
        node  $\leftarrow$  POP(frontier)
        if problem.IS-GOAL(node.STATE) then return node
        if DEPTH(node)  $>$  ℓ then
            result  $\leftarrow$  cutoff
        else if not IS-CYCLE(node) do
            for each child in EXPAND(problem, node) do
                add child to frontier
    return result

```

Figure 3.12 Iterative deepening and depth-limited tree-like search. Iterative deepening repeatedly applies depth-limited search with increasing limits. It returns one of three different types of values: either a solution node; or *failure*, when it has exhausted all nodes and proved there is no solution at any depth; or *cutoff*, to mean there might be a solution at a deeper depth than ℓ . This is a tree-like search algorithm that does not keep track of *reached* states, and thus uses much less memory than best-first search, but runs the risk of visiting the same state multiple times on different paths. Also, if the IS-CYCLE check does not check *all* cycles, then the algorithm may get caught in a loop.

Example of iterative deepening search:



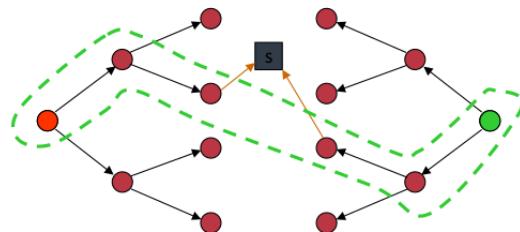
Every time the depth is increased, the entire tree is regenerated (optimization exists)

5.8.1 Evaluation of Iterative Deepening Search

- Iterative deepening search is complete (assuming b finite): if a solution exists, eventually $L = d$
- Iterative deepening search is optimal: when the step costs are all equal (as in breadth-first search)
- Complexity: iterative deepening search has spatial (worst-case) complexity of $O(bd)$ nodes. It has temporal (worst-case) complexity of $O(b^d)$ nodes

5.9 Bidirectional search

It performs two searches in parallel a “forward” search from the initial state to a goal state a “backward” search from a goal state to the initial state. Example of bidirectional search:



5.10 Tree-search Vs. Graph-search algorithms

Tree-search algorithms work well when the state space is a tree. Graph-search algorithms work well when the state space is a graph (general case).

Several variants are possible. For example, repeated states can be checked when a node is generated (not when a node is expanded). Trade-off between memory usage for storing the closed list and time spent to check the closed list.

5.11 Summary of uninformed search strategies

Search strategy	with repeated states		without repeated states		$S(n)$	$T(n)$
	Complete	Optimal	Complete	Optimal		
Breadth-First Search	yes	yes	yes	yes	$O(b^d)$	$O(b^{d+1})$
Uniform-Cost Search	yes	yes	yes	yes	$O(b^{1+\lceil \frac{C^*}{\varepsilon} \rceil})$	$O(b^{1+\lceil \frac{C^*}{\varepsilon} \rceil})$
Depth-First Search	no	no	yes	no	$O(bm)$	$O(b^m)$
Limited-Depth Search	no	no	no	no	$O(bL)$	$O(b^L)$
Iterative Deepening	yes	yes	yes	yes	$O(bd)$	$O(b^d)$

6 Informed Search Strategies

They exploit specific knowledge that is not contained in problem formulation. They select a node from the frontier according to an evaluation function $f(n)$. The evaluation function $f(n)$ provides an estimate of how much a node is “promising”. There are different ways to calculate the evaluation function, we will focus on: Greedy best-first search and A* search.

Conventionally, best (most promising) nodes have small values of evaluation function $f(n)$. Informed search strategies implement the frontier as a priority queue ordered in increasing order of $f(n)$. Thus, nodes with minimum $f(n)$ are chosen first.

```

function BEST-FIRST-SEARCH(problem, f) returns a solution node or failure
  node ← NODE(STATE, problem.INITIAL)
  frontier ← a priority queue ordered by f, with node as an element
  reached ← a lookup table, with one entry with key problem.INITIAL and value node
  while not IS-EMPTY(frontier) do
    node ← POP(frontier)
    if problem.IS-GOAL(node.STATE) then return node
    for each child in EXPAND(problem, node) do
      s ← child.STATE
      if s is not in reached or child.PATH-COST < reached[s].PATH-COST then
        reached[s] ← child
        add child to frontier
    return failure

function EXPAND(problem, node) yields nodes
  s ← node.STATE
  for each action in problem.ACTIONS(s) do
    s' ← problem.RESULTS(s, action)
    cost ← node.PATH-COST + problem.ACTION-COST(s, action, s')
    yield NODE(STATE=s', PARENT=node, ACTION=action, PATH-COST=cost)

```

Figure 3.7 The best-first search algorithm, and the function for expanding a node. The data structures used here are described in Section 3.3.2. See Appendix B for *yield*.

6.1 Best-first greedy search

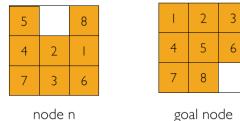
Best-first greedy search uses an evaluation function that is equal to a heuristic function $h(n)$. Thus, the evaluation $f(n)$ corresponds to the heuristic function $h(n)$ that evaluates the estimated cost of the shortest path from a node n to a goal node. To apply best-first greedy search, $h(n)$ must be known. For example, it is sufficient to know the heuristic value for each state.

The heuristic $h(n)$ is an estimate, not the actual cost. If we knew the actual cost, the problem would have already been solved.

6.1.1 Examples of best-first greedy search

Let's consider two examples of heuristics for the 8-puzzle:

- $h_1(n)$ is the number of misplaced tiles
- $h_2(n)$ is computed as the sum of Manhattan distances of each tile to its final position (computed only considering horizontal and vertical movements)



Considering the node N we have:

- $h_1(n) = 6$ because only 4 and 7 are in the right position
- $h_2(n) = 2$ (for tile 5) + 3 (for tile 8) + 0 (for tile 4) + 1 (for tile 2) + 3 (for tile 1) + 0 (for tile 7) + 3 (for tile 3) + 1 (for tile 6) = 13

What are the underlying intuitions behind $h_1(n)$ and $h_2(n)$?

- $h_1(n)$: if a tile is misplaced, I will need at least one move to place it in its correct position
- $h_2(n)$: i am counting how many moves each tile would require if no other tiles were placed on the board

$h_1(n)$ and $h_2(n)$ are defined considering simpler problems (that is, by relaxing some of the original constraints).

6.2 Heuristic function accuracy

Given two (consistent) heuristic functions h_1 and h_2 such that $h_1(n) \leq h_2(n)$ for any node n , h_2 **dominates** h_1 since each node expanded by A^* search using h_1 is also expanded by h_2 . In the 8-puzzle h_2 dominates h_1 . When h_2 dominates h_1 , h_2 we say it is **more accurate or more informed than h_1** .

Best-first Greedy Tree Search python code (it still checks for cycles but has no reached table):

```

1 def best_first_tree_search(problem, f):
2     "A version of best_first_search without the 'reached' table."
3     frontier = PriorityQueue([Node(problem.initial)], key=f)
4     while frontier:
5         node = frontier.pop()
6         if problem.is_goal(node.state):
7             return node
8         for child in expand(problem, node):
9             if not is_cycle(child):
10                frontier.add(child)
11    return failure

```

If $f(n)=g(n)$ then it is uninformed search algorithm since $g(n)$ is a given data corresponding to the path cost from root to the node.

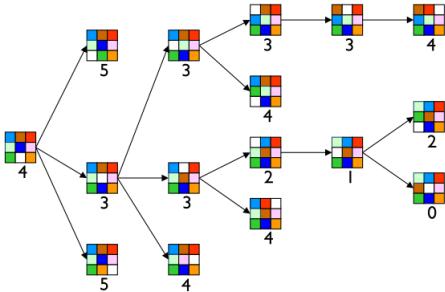
Best-first Greedy Search python code (with reached table)

```

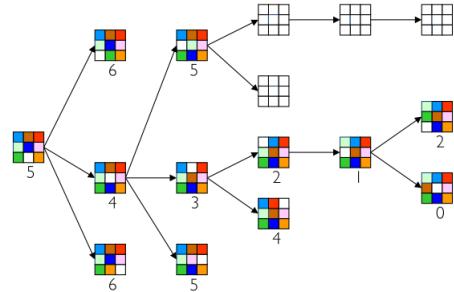
1 def best_first_search(problem, f):
2     "Search nodes with minimum f(node) value first."
3     node = Node(problem.initial)
4     frontier = PriorityQueue([node], key=f)
5     reached = {problem.initial: node}
6     while frontier:
7         node = frontier.pop()
8         if problem.is_goal(node.state):
9             return node
10        for child in expand(problem, node):
11            s = child.state
12            if s not in reached or child.path_cost < reached[s].path_cost:
13                reached[s] = child
14                frontier.add(child)
15    return failure

```

Example of best-first greedy search using $h_1(n)$ and graph-search (because of repeated states):

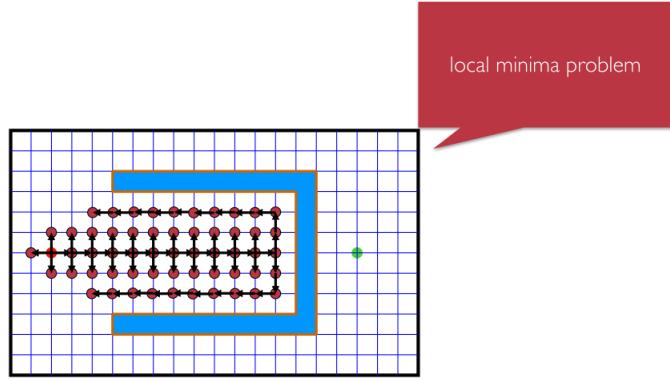


Example of best-first greedy search using $h_2(n)$ and graph-search (because of repeated states)



$h_1(n)$ estimated that we would need four moves to reach the goal $h_2(n)$ estimated that we needed five moves. They are both estimates, but some estimates might be more accurate than others (more later).

Another example of best-first greedy search. In this case $h(n)$ is computed as the straight distance from the goal. We apply graph-search since we might have repeated nodes.



6.2.1 Evaluation of Best-First Greedy Search

- Best-first greedy search is not complete: since it can get easily stuck in local optima (and cannot backtrack)
- Best-first greedy search, in general, is not optimal: for the same reason
- Complexity: it has temporal and spatial (worst-case) complexity of $O(b^m)$ nodes, m is the maximum depth of the search tree (and could be infinite)

6.3 A* search

The evaluation function $f(n)$ of a node n is computed as the sum of two components: the cost to reach n from the root $g(n)$ and an heuristic function $h(n)$.

$$f(n) = g(n) + h(n)$$

$f(n)$ estimates the cost of a solution that passes through node n :

- $g(n)$ is the path cost from the root to n (what we know for sure)
- $h(n)$ estimates the cost of the shortest path from n to a goal node (what we are guessing)

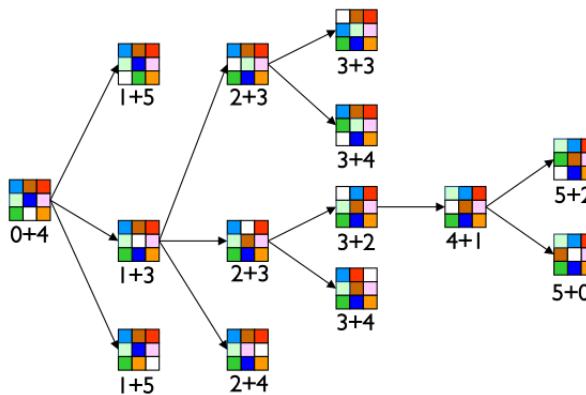
A* Search python code (with reached table):

```

1 def g(n): return n.path_cost
2
3 def astar_search(problem, h=None):
4     """Search nodes with minimum f(n) = g(n) + h(n)."""
5     h = h or problem.h
6     return best_first_search(problem, f=lambda n: g(n) + h(n))

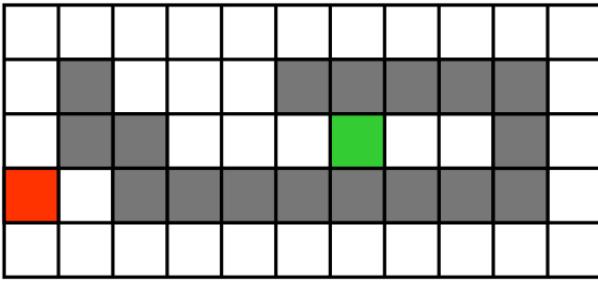
```

We must know $h(n)$ to apply A* search. Example of A* search using graph-search. $h(n)$ is computed as the number of misplaced tiles:



6.3.1 Advantages of A* search

Navigation example: the red square is the starting position, the green square is the goal, grey squares are obstacles. Best-first greedy search using $h(n)$ computed as the Manhattan distance from the goal (without considering obstacles) on the right.



8	7	6	5	4	3	2	3	4	5	6
7			5	4	3					5
6					3	2	1	0	1	2
7	6									5
8	7	6	5	4	3	2	3	4	5	6

The values of the distance can be stored or computed on the fly depending on the problem. Best-first greedy search using $h(n)$ computed as the Manhattan distance from the goal (without considering obstacles) on the left while on the right A* search using $h(n)$ computed as the Manhattan distance from the goal (without considering obstacles).

8	7	6	5	4	3	2	3	4	5	6
7		5	4	3						5
6			3	2	1	0	1	2		4
7	6									5
8	7	6	5	4	3	2	3	4	5	6

8+3	7+4	6+3	5+6	4+7	3+8	2+9	3+10	4	5	6
7+2		5+6	4+7	3+8						5
6+1			3	2+9	1+10	0+11	1	2		4
7+0	6+1									5
8+1	7+2	6+3	5+4	4+5	3+6	2+7	3+8	4	5	6

By using both the cost so far $g(n)$ and the estimate of the cost of reaching the goal $h(n)$, A* can lead the agent away from suboptimal paths.

6.3.2 Evaluation of A* search

- A* search is complete and optimal for tree-search: when $h(n)$ is admissible
- A* search is complete and optimal for graph-search: when $h(n)$ is consistent
- Complexity: it has a temporal and spatial (worst-case) complexity that is exponential in the length of the solution

6.4 Admissible heuristic functions

A heuristic function $h(n)$ is admissible when, for each node n

$$0 \leq h(n) \leq h^*(n)$$

where $h^*(n)$ represents the actual cost from node n to the solution. When n is a goal node, $h(n)$ should be zero. $h(n)$ is “optimistic” since it always underestimates the true cost to reach the goal.

Example

- $h(n)$ computed as the estimated cost of the shortest path from n to a goal node
- $h^*(n)$ returning the actual cost of the shortest path from n to a goal node

How can I define an admissible $h(n)$ without knowing $h^*(n)$?

Considering

5		8
4	2	1
7	3	6

node n

1	2	3
4	5	6
7	8	

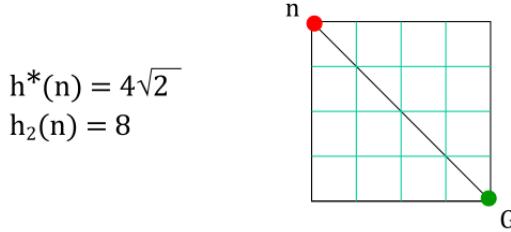
goal node

Is $h_1(n)$ admissible? (the number of misplaced tiles). Is $h_2(n)$ admissible? (the sum of Manhattan distances of each tile to its final position).

Step cost of horizontal/vertical actions is 1 and $\sqrt{2}$ for a diagonal actions. (x_G, y_G) is the goal position.



Why is $h_2(n)$ not admissible, when moving diagonally?



When moving diagonally, the actual $h^*(n)$ is 5.66 while $h_2(n)$ is 8 (thus, it overestimates the cost). How to create an admissible heuristic function? **An admissible heuristic can usually be seen as the cost of an optimal solution to a relaxed problem, obtained by removing constraints.**

In the robot navigation example,

- The Manhattan distance corresponds to the optimal solution when we remove the obstacles but keep the constraints of moving in four directions
- The Euclidean distance corresponds to the optimal solution when we remove both the obstacles and allow diagonal movements

In the 8-puzzle,

- $h_1(n)$ was based on removing the constraint of moving only by sliding the tiles
- $h_2(n)$ considered the moves needed to position each tile as if it was alone on the board

6.5 First optimality theorem for A^* search

A^* search with tree-search is complete and optimal when $h(n)$ is admissible (when all step costs are larger than a small positive ϵ).

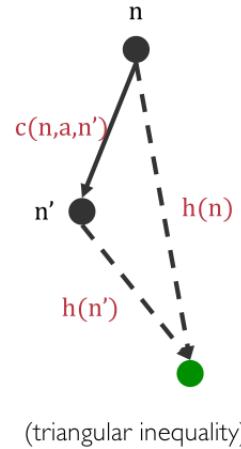
Completeness: if a solution exists, it will have a finite cost C^* . Since the step costs are at least ϵ , then $g(n)$ will always increase and exceed the cost of the solution. Thus, A^* will expand forever the tree. **Thus, if a solution exists, A^* search terminates with the solution.**

Optimality: call C^* the cost of the optimal solution. Consider a sub-optimal goal node G' in the frontier $f(G') = g(n) + h(n) \leq C^*$. Consider a node n in the frontier, which is on the path corresponding to the optimal solution; thus, $f(n) = g(n) + h(n) \leq C^*$. From the two equations we derive that $f(n) \leq C^* \leq f(G')$ and thus the suboptimal goal G' will not be chosen from the frontier. **Thus, when A^* search selects a node from the frontier; the path to the corresponding state is optimal.**

6.6 Consistent Heuristic Functions

A heuristic function $h(n)$ is consistent when, for each node n and each one of its successors n' :

$$h(n) \leq c(n, a, n') \leq h(n')$$



A consistent heuristic function is also admissible; vice versa is not true in general. If n is a goal node, then $h(n) = 0$.

When is a heuristic function not consistent?

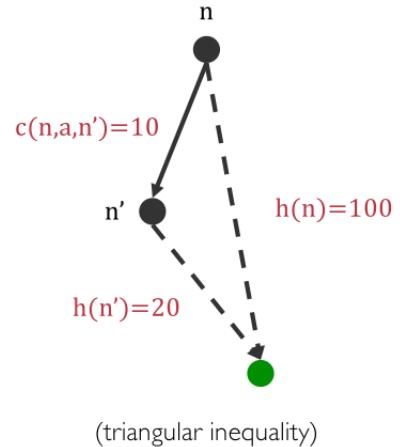
Consider a heuristic function for which:

- the estimated cost to the goal node from n is 100
- the cost of moving from n to n' is 10
- the estimated cost to the goal node from n' is 20

In this case, the triangular inequality does not hold:

$$h(n) \leq c(n, a, n') \leq h(n')$$

$$(100 \not\leq 10 + 20)$$



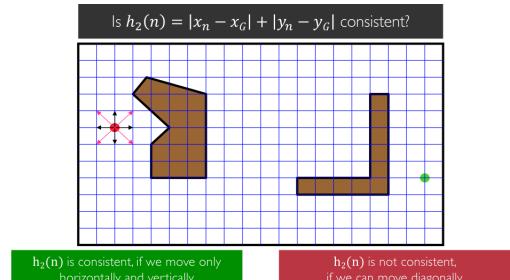
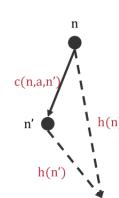
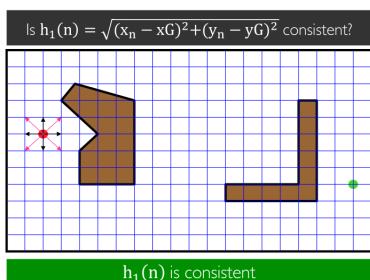
Let's consider two examples of heuristics for the 8-puzzle:

- $h_1(n)$ is the number of misplaced tiles
- $h_2(n)$ is computed as the sum of Manhattan distances of each tile to its final position (computed only considering horizontal and vertical movements)

Considering the node N we have:

- $h_1(n) = 6$ because only 4 and 7 are in the right position
- $h_2(n) = 2$ (for tile 5) + 3 (for tile 8) + 0 (for tile 4) + 1 (for tile 2) + 3 (for tile 1) + 0 (for tile 7) + 3 (for tile 3) + 1 (for tile 6) = 13

Both heuristics are consistent. Now let's get back on the path problem with obstacles:



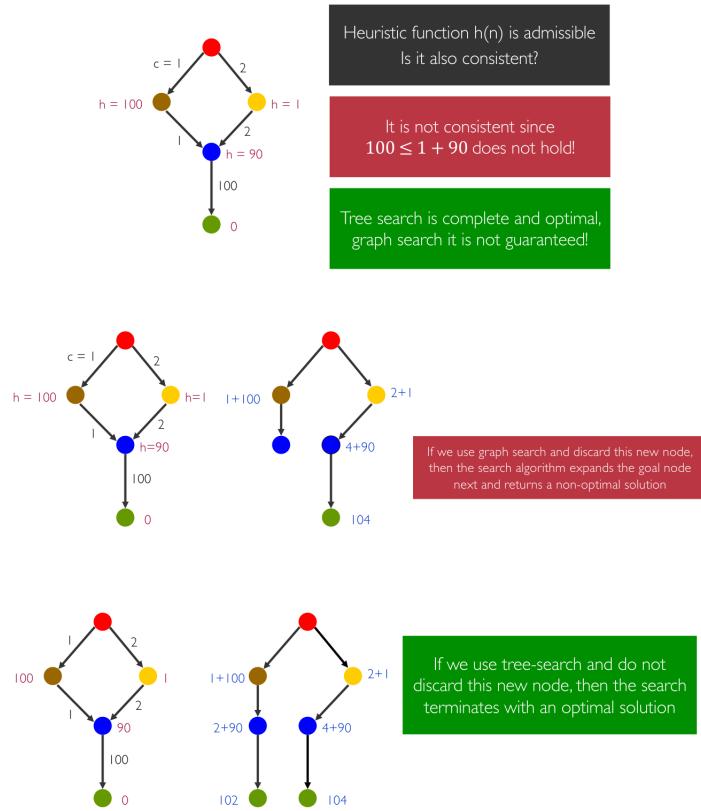
6.7 Second optimality theorem for A^* search

A^* search with graph-search is complete and optimal when $h(n)$ is consistent (when all step costs are larger than a small positive ϵ).

Completeness: it can be proved that in the same way we did for tree search with an admissible $h(n)$.

Optimality: given n' successor of n , $f(n') = g(n') + h(n') = g(n) + c(n, a, n') + h(n') \geq g(n) + h(n) = f(n)$. Thus, A^* search chooses (and expands) nodes in non-decreasing order of $f(n)$. When A^* selects a node from the frontier, the current path to the corresponding state is optimal. Thus, when A^* selects the first goal node from the frontier, this is the optimal solution.

Consistency is a stronger property than admissibility. In fact, there are more admissible than consistent functions Graph-search is more constrained since it does not consider nodes corresponding to the same state. Thus, it requires a stronger property to guarantee completeness and optimality.



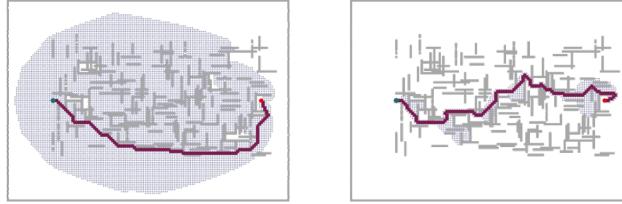
Other properties of A^* search:

- A^* search expands all nodes with $f(n) < C^*$
- A^* search expands some nodes with $f(n) = C^*$
- A^* search expands no nodes with $f(n) > C^*$
- A^* search is **optimally efficient**: there exists no other optimal search algorithm that, for any heuristic function, expands less nodes than A^* search (except some nodes with $f(n) = C^*$)

So, if we have a consistent heuristic, we are ready to apply A^* search, correct? Not really... There are very dumb consistent heuristic functions The heuristic $h_0(n)$ that returns 0 for all nodes n is consistent and thus admissible, but uninteresting. A^* search with $h_0(n)$ implements uniform-cost search. Uniform-cost search is thus a special case of A^* search.

6.8 Weighted A*

It introduces a weight factor w ($1 \leq w < \infty$) over the heuristics so that, $f(n) = g(n) + wh(n)$. The theoretical properties do not apply anymore, but this extension of A* can be more efficient. A* search (left) compared to weighted A* search with $w = 2$ (right):



6.9 Iterative deepening A* (IDA*) search

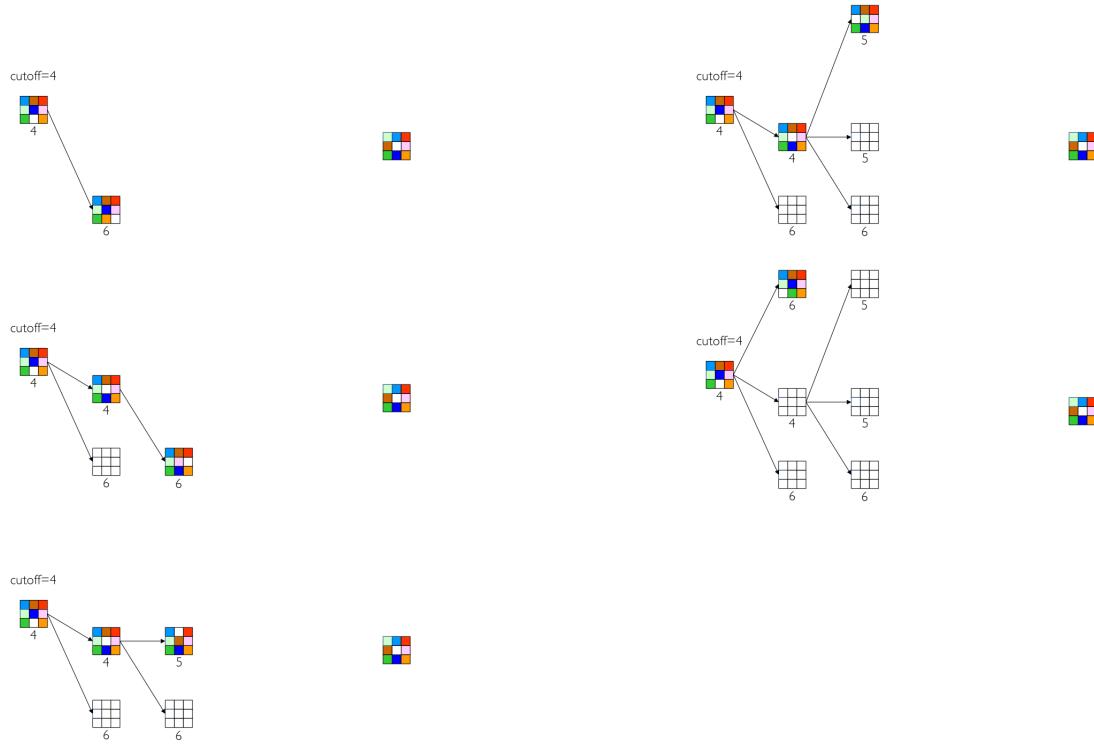
IDA* reduces memory requirement of A* search by applying a limit to the values of $f(n)$. It assumes to have a consistent heuristic function $h(n)$. Iterative Deepening A* (IDA*) Search:

```

1 cutoff = f(root-node)
2 repeat
3     1. perform depth-first search by expanding all nodes n such that f(n) <= cutoff
4         2. cutoff = smallest value f() of non-expanded (leaf) nodes until
5             (a solution is found) or (a termination condition on time or memory is met)

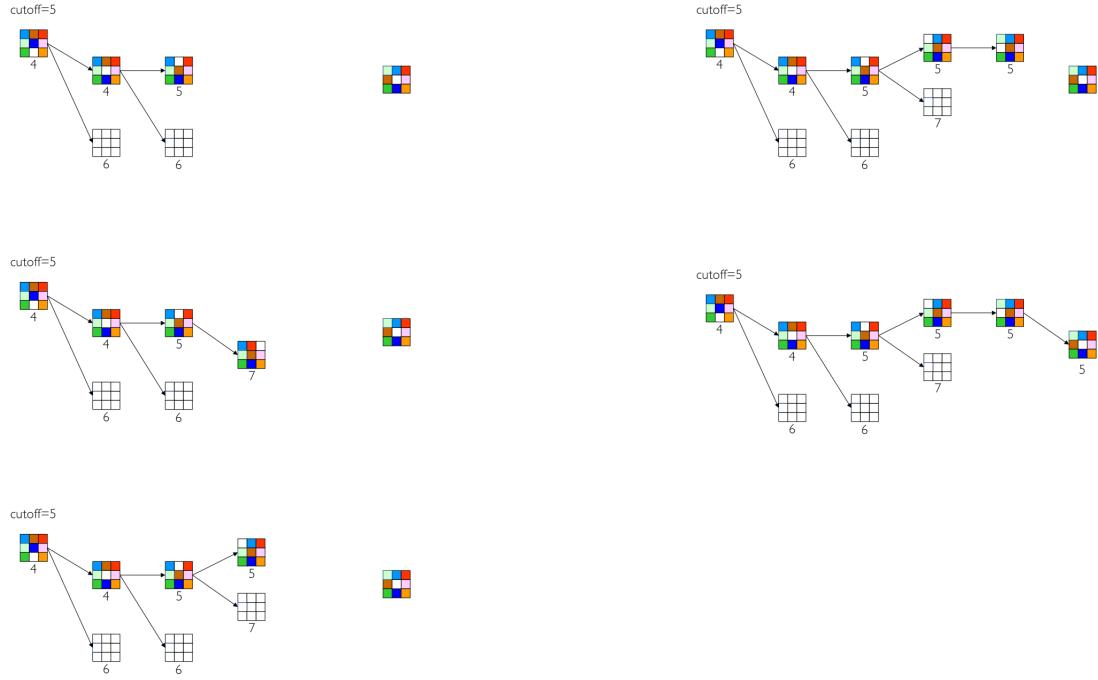
```

Example of IDA* search. 8-puzzle with $h(n)$ computed as the number of misplaced tiles and cutoff equal to 4:



Cutoff set to 5:





6.9.1 Evaluation of IDA* search

- Completeness: IDA* search is complete when $h(n)$ is admissible
- Optimality: IDA* search is optimal when $h(n)$ is admissible
- Complexity: IDA* search requires less memory than A* search and avoids to sort the frontier. However, IDA* cannot avoid to revisit states not on the current path, because it uses too little memory (it just remembers the current cutoff). Thus, it is very efficient memory-wise but not time-wise since it revisits states
- Several other variants of A* search exist and are widely used

7 Adversarial search strategies

Adversarial search considers multiagent environments in which agents compete. Board games are a typical scenario for adversarial search strategies.

We consider games with:

- two players
- turn-taking
- zero sum
- perfect information
- deterministic and stochastic (they include elements of chance)

Zero sum means that if a player wins (+1) the other player loses (-1). Thus sum, thus, is zero. Perfect information means that the agent has perfect information about the game state which is fully observable (like chess, unlike poker).

Games can be formulated as search problems with an element of uncertainty due to the action (moves) of the opponent. In general, there is no sequence of actions that makes a player win independently of the actions of the opponent. At each turn, a player must choose the action to play without the possibility to explore the whole state space.

7.1 Formulation of search problem for a game

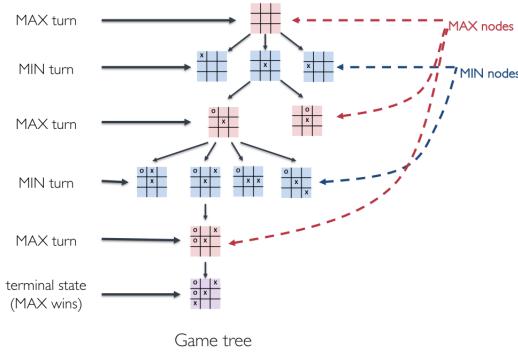
- **Players:** two players, MAX and MIN, that take turn. MAX moves first
- **Initial state:** initial set up of the game (initial configuration of the board)
- **Function Actions(s):** given a state s, returns the set of legal moves in s
- **Function Result(s,a):** defines the state resulting from taking action a in state s
- **Goal test:** given a state, returns true if the game is over
- **Utility(s,p):** returns the numerical value to player p in terminal state s (e.g., +1, 0, or -1)

Python formulation:

```
1  class Game:
2      def actions(self, state):
3          """Return a collection of the allowable moves from this state."""
4          raise NotImplementedError
5      def result(self, state, move):
6          """Return the state that results from making a move from a state."""
7          raise NotImplementedError
8      def is_terminal(self, state):
9          """Return True if this is a final state for the game."""
10         return not self.actions(state)
11     def utility(self, state, player):
12         """Return the value of this final state to player."""
13         raise NotImplementedError
```

Playing a game

```
1  def play_game(game, strategies: dict, verbose=False):
2      """Play a turn-taking game. 'strategies' is a {player_name: function} dict,
3      where function(state, game) is used to get the player's move."""
4      state = game.initial
5      while not game.is_terminal(state):
6          player = state.to_move
7          move = strategies[player](game, state)
8          state = game.result(state, move)
9          if verbose:
10              print('Player', player, 'move:', move)
11              print(state)
12      return state
```



Solving the game corresponds to finding the best action for MAX (or for MIN) in a state.

7.2 Action selection - Minimax Search

It considers the game from MAX's perspective. Starting from the current state (corresponding to the root), MAX builds the game tree. The game tree can be built completely or up to a maximum depth h (which could depend on the available time to perform a move). The game tree is built using a depth-first search. MAX evaluates the leaf nodes of the game tree using Utility(s) or by an evaluation function $e(s)$. MAX “backs up” the minimax values from the leaves to the root. MAX chooses the best move (the move that maximizes the minimax value).

7.3 The evaluation function: $e(s)$

When the state is not terminal, we cannot compute utility. In non-terminal state, we can compute an evaluation function $e(s)$ that, given a state s and a player p , returns a real number estimating the expected utility of s to p . Evaluation function:

- $e(s) > 0$ means that s is good for MAX
- $e(s) < 0$ means that s is good for MIN
- $e(s) = 0$ means that s is neutral

7.4 The utility function

The Utility function returns the numerical value to player p in terminal state s (e.g., +1, 0, or -1) Utility:

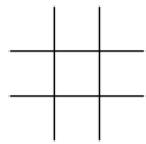
- Utility(s) > 0 means that MAX wins
- Utility(s) < 0 means that MIN wins
- Utility(s) $= 0$ means a tie

Convention Utility(s) and $e(s)$ return values relative to MAX. The values for MIN are obtained by taking the opposite (i.e., multiplying by -1) because the games are zero sum. Utility(s) and $e(s)$ are positive when they are good for MAX, negative when they are good for MIN.

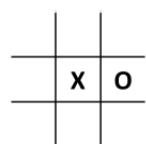
Example of evaluation function for Tic-Tac-Toe $e(s)$ is computed as the difference between:

- The number of rows, columns and diagonals that are open for MAX and
- The number of rows, columns, and diagonals open for MIN

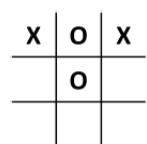
$e(s)$ could be normalized between -1 and +1 to keep values consistent with utility.



$$e(s) = 8 - 8 = 0$$



$$e(s) = 6 - 4 = 2$$



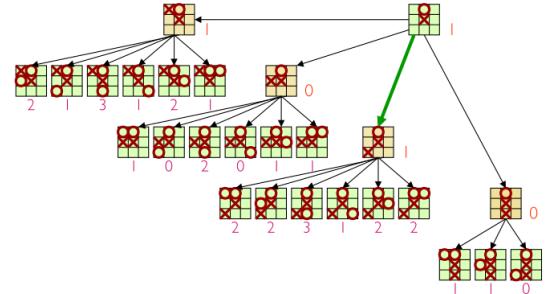
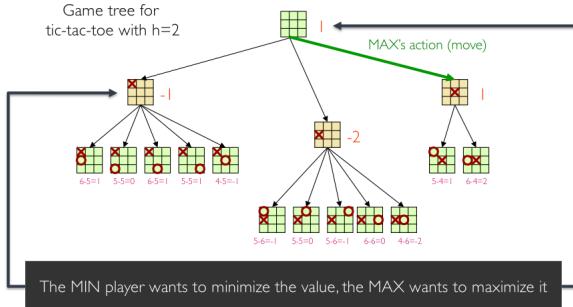
$$e(s) = 3 - 3 = 0$$

Starting from the values assigned (by function Utility() or by the evaluation function) to the leaves and going up to the root, a minimax value is assigned to each node of the game tree. The minimax value of a node is the estimate of the expected utility for a player in the state corresponding to that node, if both players play optimally. The minimax value of a leaf node is equal to the value returned by:

- The function Utility(s), if the leaf node is terminal
- The evaluation function e(s), if the leaf node is not terminal

The minimax value of a MAX node is the maximum of the minimax values of its children. The minimax value of a MIN node is the minimum of the minimax values of its children.

Examples of Minimax Search:



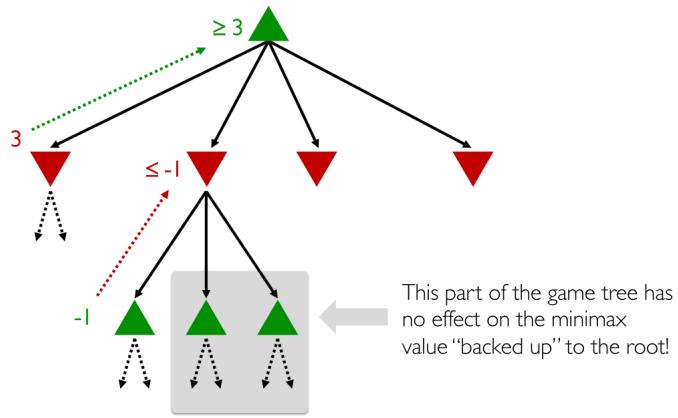
Minimax search python implementation:

```

1 def minimax_search(game, state):
2     """Search game tree to determine best move; return (value, move) pair."""
3     player = state.to_move
4     def max_value(state):
5         if game.is_terminal(state):
6             return game.utility(state, player), None
7         v, move = -infinity, None
8         for a in game.actions(state):
9             v2, _ = min_value(game.result(state, a))
10            if v2 > v:
11                v, move = v2, a
12        return v, move
13
14    def min_value(state):
15        if game.is_terminal(state):
16            return game.utility(state, player), None
17        v, move = +infinity, None
18        for a in game.actions(state):
19            v2, _ = max_value(game.result(state, a))
20            if v2 < v:
21                v, move = v2, a
22        return v, move
23    # the Minimax search plays for MAX
24    return max_value(state)

```

Minimax is optimal when using the utility function (when the entire game tree is built). However, Minimax search builds a game tree in which nodes grow exponentially with the maximum depth h . Can MAX choose the optimal action without examining all the nodes in the game tree?

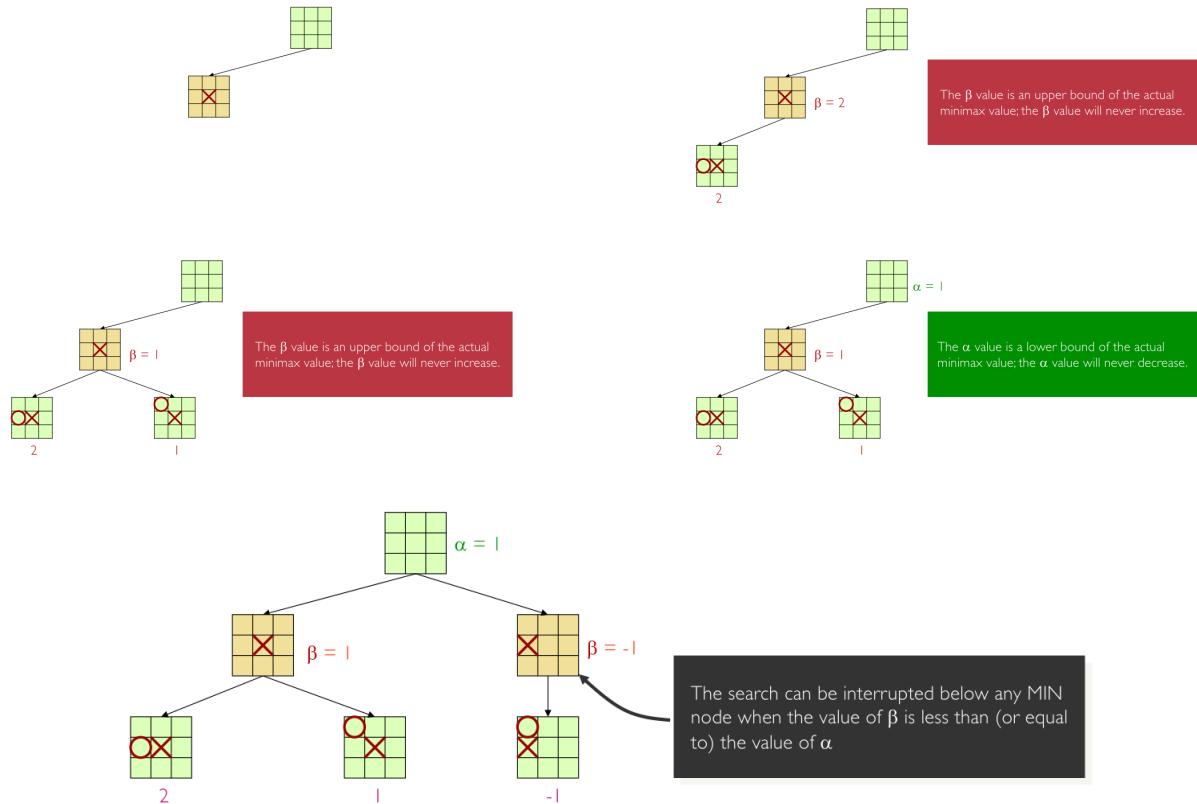


7.5 Alpha-Beta Pruning

Reduces the complexity of the minimax search by not considering branches of the game tree that cannot influence the final decision. Alpha-beta pruning returns the same outcome of the “vanilla” minimax search. Alpha-beta pruning uses two parameters α and β :

- Parameter α : it represents the value of the best move found so far for MAX (i.e. with the highest minimax value). α can be interpreted as “at least”
- Parameter β : it represents the value of the best move found so far for MIN (i.e. with the lowest minimax value). β can be interpreted as “at most”

Alpha-beta pruning for Tic-Tac-Toe



Here is a sketch of the algorithm of Alpha-Beta pruning:

1. Update the α and β values as soon as possible
2. Interrupt the search below a MAX node when the α value is larger or equal to the β value
3. Interrupt the search below a MIN node when the β value is smaller or equal to the α value

7.5.1 Evaluation of Minimax Search and Alpha-Beta pruning

Given the maximum number of legal moves in a state is b , minimax search will expand $O(b^h)$ nodes. The efficiency of alpha-beta pruning depends on the order in which nodes are examined. In the worst case, alpha-beta pruning expands $O(b^h)$ nodes.

In the best case, alpha-beta pruning expands $O(b^{h/2})$ nodes

- MIN nodes that are children of a MAX node should be considered in increasing order of minimax values
- MAX nodes that are children of a MIN node should be considered in decreasing order of minimax values
- This cannot be guaranteed in practice

In the average case (nodes considered in random order), alpha-beta pruning expands $O(b^{3h/4})$ nodes.

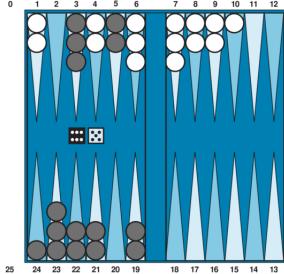
Move ordering example:



Move ordering example.

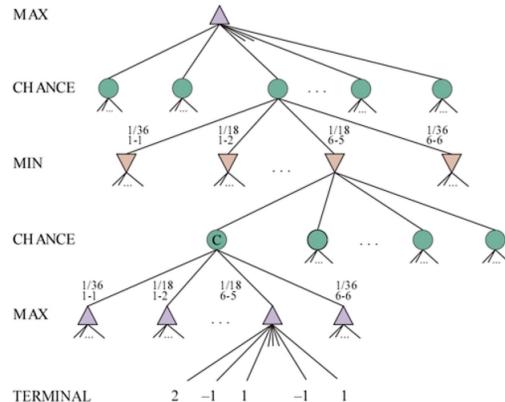
8 Stochastic games

Many games include unpredictable stochastic events, like throwing of dice in backgammon. We can apply an approach like minimax search using a game tree with chance nodes and expectiminimax values.

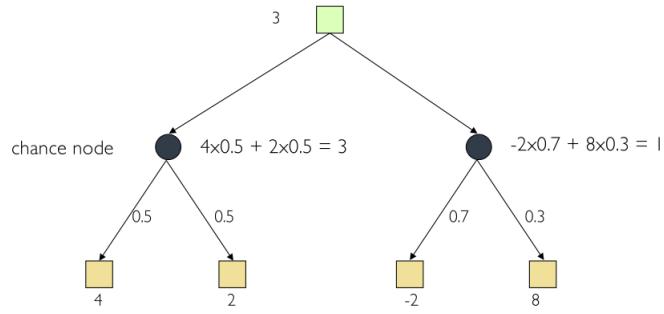


Black has rolled and must choose among four legal moves: (5-11,5-10), (5-11,19-24), (5-10,10-16), and (5-11,11-16), where the notation (5-11,11-16) means move one piece from position 5 to 11, and then move a piece from 11 to 16.

A typical backgammon position. The goal of the game is to move all one's pieces off the board. Here is a schematic game tree for a backgammon position:

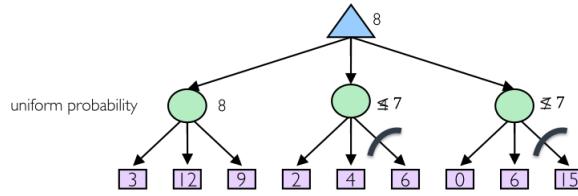


Stochastic games example for chance nodes with only two possible outcomes:

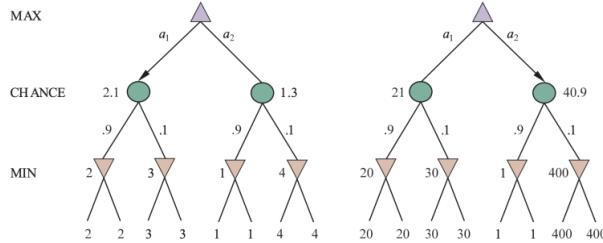


8.1 Alpha-Beta pruning for stochastic games

Alpha-beta pruning can be applied to stochastic games. Consider the following game tree:



What happens if you don't know the range of $\text{Utility}()$ and $e()$? What happens if, instead, you know the range $[0, 15]$? An order-preserving transformation of leaves values can change the best move at the root. Do values matter also in deterministic games (minimax)? Why?



9 Monte Carlo Tree Search

Alpha-Beta pruning is effective in many games, however it cannot play some games, like Go. Why Go is difficult?

- Its branching factor is very large
- The game comprises more than 200 moves (about 250 moves on average)
- Order of magnitude greater than the branching factor of 20 for Chess
- It lacks a good evaluation function
- It is difficult to model, similar looking positions can have completely different outcomes

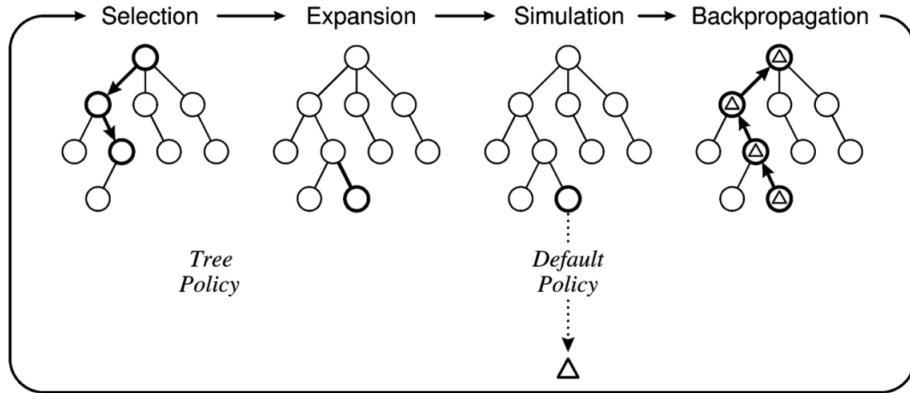
Alpha-beta pruning reached a proficiency level weak to intermediate.

Monte Carlo Tree Search (MCTS) builds a “sparse” game tree to choose the best move at the root. Monte Carlo for random simulation MCTS runs random simulations and builds a search tree from the results. Markovian Decision Problem (MDP) with sequence of decisions, problem are phrased as state, action pairs.

Monte Carlo tree search has different phases:

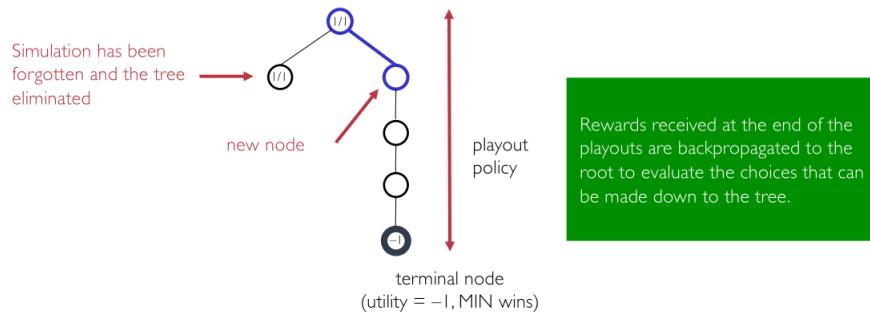
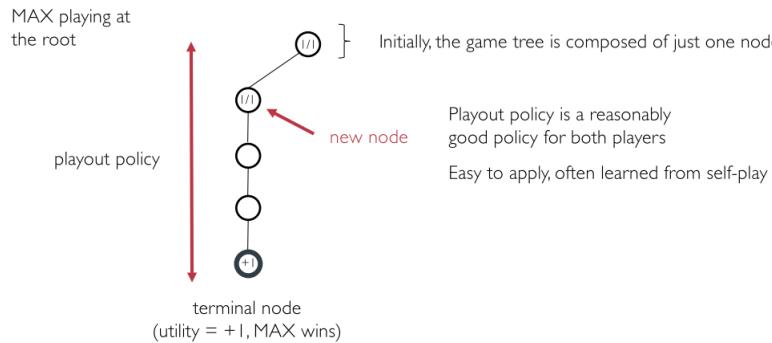
1. Selection: it descends the game tree selecting successors using a selection policy until either a node with an unexpanded branch or a leaf is reached
2. Expansion: if the node is not a goal, the tree is expanded by generating one or more children nodes.

3. Simulation: a simulation (a playout or roll-out) is run from the current node to evaluate the game outcome by following a default playout policy (typically random); This results in a reward computed according to a reward function.
4. Backpropagation: the reward is used to update all the nodes visited during the previous steps.

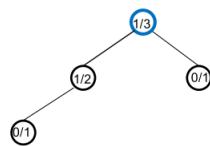
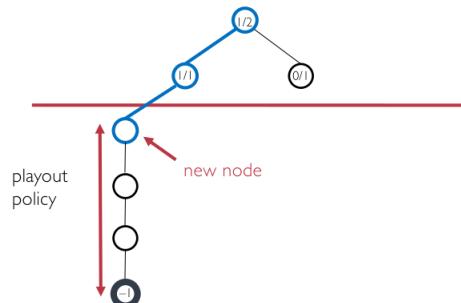
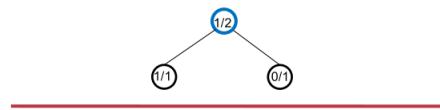


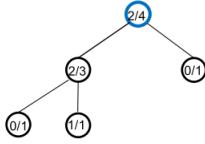
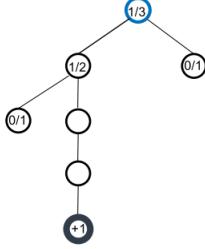
MCTS builds a “sparse” game tree by sampling the game tree focusing on the most promising paths. MCTS have more difficulty in problems in which sampling might not be effective. For example, problems with very few good moves with respect of the number of available ones.

When does MCTS stop? It is an **anytime method**. Whenever it stops, it can provide an answer. Typically, it can be stopped when a time limit has been reached or a number of rollouts have been performed.



- Each node n stores a pair of values
 - Number of wins for the player at the root
 - Number of performed simulations that involve n
- Alternatively, each node n stores a pair of values
 - Number of wins for the player at the node
 - Number of performed simulations that involve n
- In general, the utilities $U(n)$ are stored in nodes





How do we select the successors nodes? Selection policy should balance exploration and exploitation. I wish to visit nodes that promise the highest chance of winning. I wish to try less explored nodes to learn what might happen if I take those actions. **UCB1(n)** is a common solution to balance exploration and exploitation in MCTS.

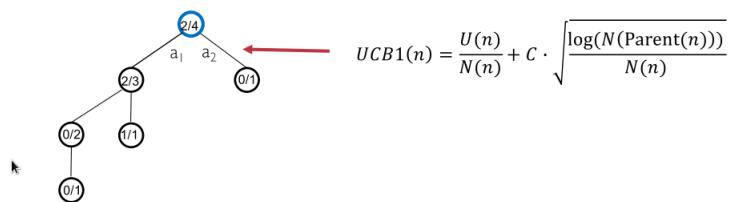
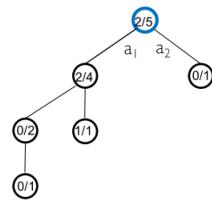
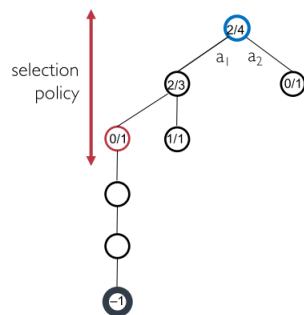
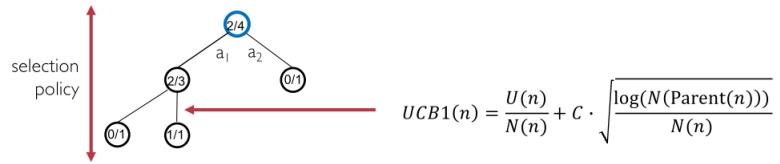
$$UCB1(n) = \frac{U(n)}{N(n)} + C \cdot \sqrt{\frac{\log(N(Parent(n)))}{N(n)}}$$

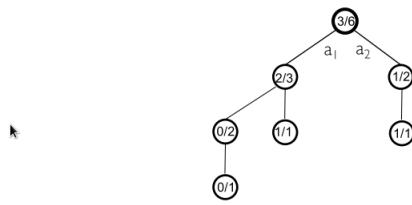
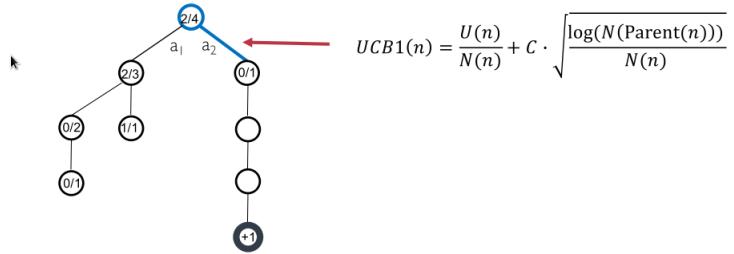
The node with the largest UCB1(n) is selected. Where:

- $U(n)$ is the utility of node n (number of wins)
- $N(n)$ is the number of simulations involving node n
- C is the exploration parameter
- $N(Parent(n))$ is the number of simulations involving the parent of node n

And, in general, the first term is the exploitation term and the second term is the exploration term. When C is infinity, then we have pure exploration. When C is zero, we have pure exploitation.







MCTS does not need an explicit evaluation function. Random layouts are enough to explore the search space. Thus, it can be employed an evaluation function, or a substantial corpus of knowledge, is unavailable.