

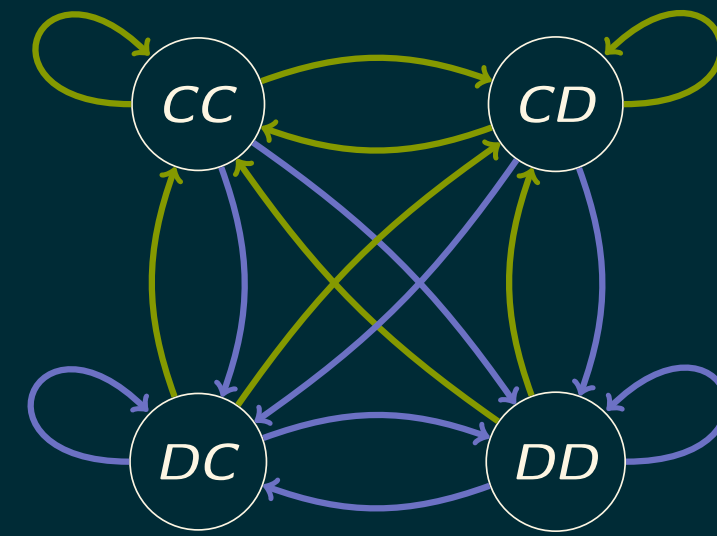
# THE POWER OF MEMORY

Is memory size advantageous in interactions (social, biological, ... ) ?

	C	D
C	3, 3	0, 5
D	5, 3	1, 1

	1	2	3	...	$n-1$	$n$
player $p$	C	C	C	...	D	...
player $q$	C	C	D	...	D	...

memory size  $n-1$



$$\begin{bmatrix} p_1 q_1 & p_1(-q_1+1) & q_1(-p_1+1) & (-p_1+1)(-q_1+1) \\ p_2 q_3 & p_2(-q_3+1) & q_3(-p_2+1) & (-p_2+1)(-q_3+1) \\ p_3 q_2 & p_3(-q_2+1) & q_2(-p_3+1) & (-p_3+1)(-q_2+1) \\ p_4 q_4 & p_4(-q_4+1) & q_4(-p_4+1) & (-p_4+1)(-q_4+1) \end{bmatrix}$$

W. H. Press and F. J. Dyson. **Iterated Prisoner's Dilemma contains strategies that dominate any evolutionary opponent** PNAS 2012.

$$p^* \rightarrow \text{manipulates} \rightarrow q$$

This work considers an optimisation approach to identify:

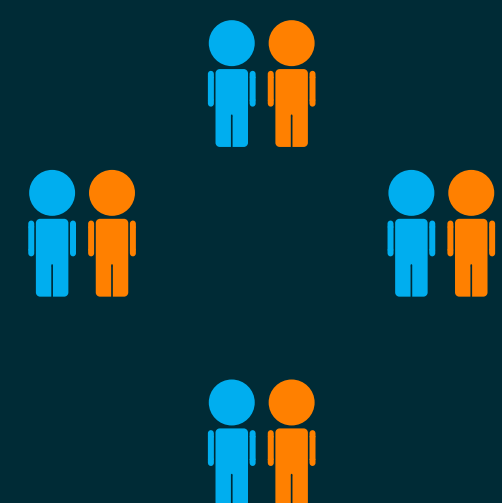
$$p^* \rightarrow \text{best response} \rightarrow q$$

$$1. \max_q : \frac{\frac{1}{2} p Q p^T + c^T p + a}{\frac{1}{2} p \bar{Q} p^T + \bar{c}^T p + \bar{a}}$$

$$st : p \in \mathbb{R}_{[0,1]}^4$$

## PURELY RANDOM STRATEGIES $p = (p, p, p, p)$

### 2. AGAINST A SINGLE OPPONENT

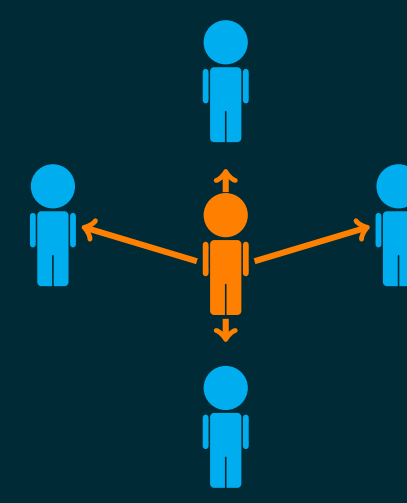


$$p^* = \operatorname{argmax}(u_q(p)), p \in S_q,$$

where the set  $S_q$  is defined as:

$$S_q = \left\{ 0, p_{\pm}, 1 \mid \begin{array}{l} 0 < p_{\pm} < 1, \\ p_{\pm} \neq \frac{-d_0}{d_1} \end{array} \right\}$$

### 3. AGAINST MULTIPLE OPPONENTS



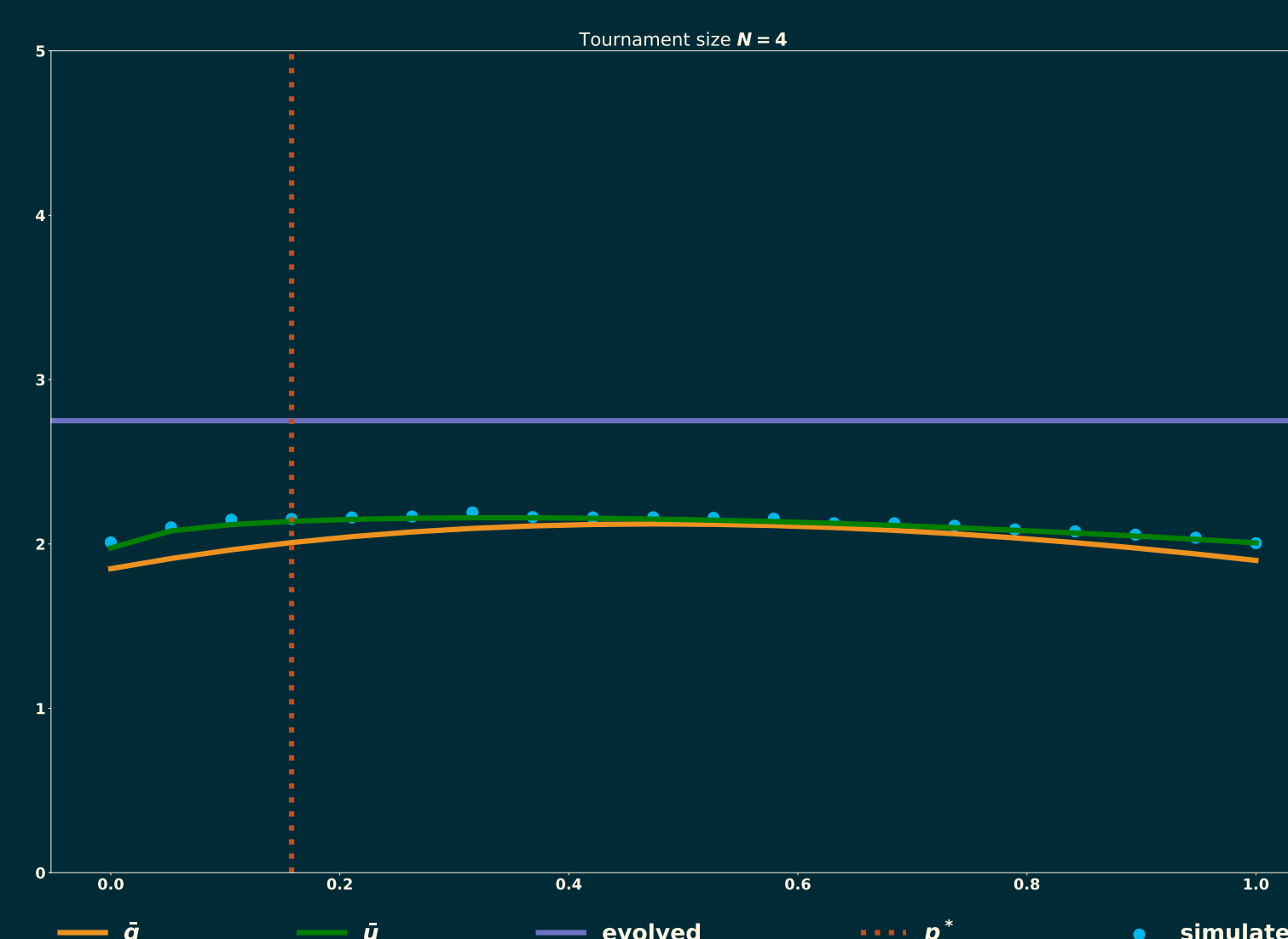
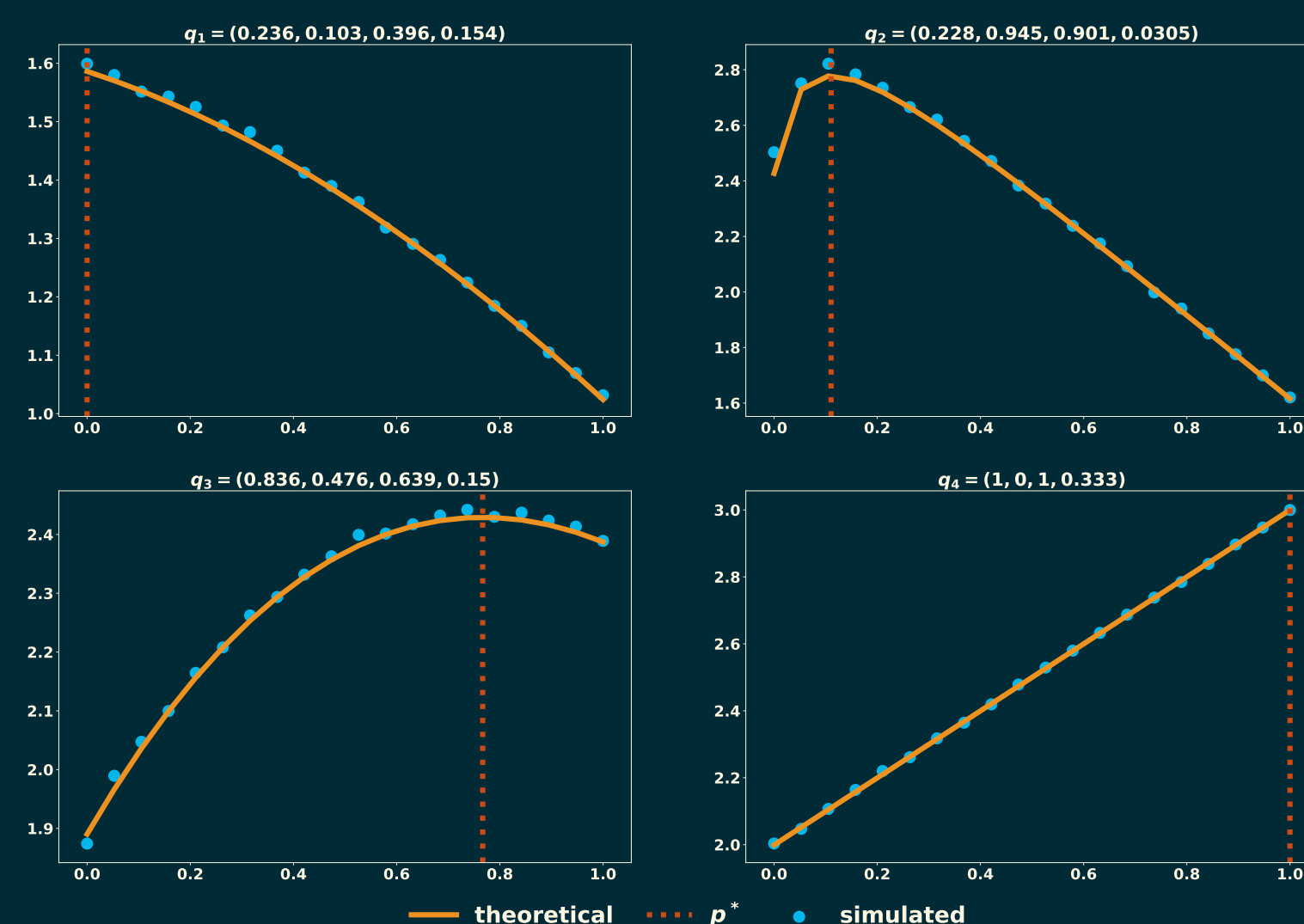
$$p^* = \operatorname{argmax} \left( \sum_{i=1}^N u_{q(i)}^{(i)}(p) \right), p \in S_{q(i)},$$

where the set  $S_{q(i)}$  is defined as:

$$S_{q(i)} = \left\{ \frac{2N}{U} \lambda_i \cup \{0, 1\} \mid \lambda_i \neq \frac{d_{0i}}{d_{1i}} \right\}$$

## RESULTS

- The utility of a given player  $p$  against a given opponent  $q$  is written in a compact way.
- Defining the optimal random behaviour  $p^*$  is reduced to a search over a small finite set.
- Optimising against the mean utility can not be captured by optimising against the mean opponent.



## FUTURE WORK RESULTANT THEORY

