

HANDIN 1

WORKSHOP 2 - SI725018 - NLA - LORENZO STIULLANO

EXERCISE 1

i) Write an algorithm called algorithm OP that computes the so called outer product $\vec{x}\vec{y}^T \in \mathbb{R}^{n \times n}$ of vectors $\vec{x}, \vec{y} \in \mathbb{R}^n$

Algorithm OP; outer product

Input: $\vec{x}, \vec{y} \in \mathbb{R}^n$

Output: $A = \vec{x}\vec{y}^T$

1. for i, \dots, n do:
2. for j, \dots, n do:
3. $A_{ij} = x_i y_j$
4. end for
5. end for
6. return A

ii) The cost of algorithm OP is $C(n, n) = \sum_{i=1}^n \sum_{j=1}^n 1^* = \sum_{i=1}^n n = n \cdot n = n^2 \therefore C(n) = n^2$
↑
sum over the columns
sum over the rows
*there is a 1 because we only perform 1 computation per for loop ($x_i y_j$).

iii) $(I - \vec{u}\vec{u}^T) \vec{x}$ for $u, x \in \mathbb{R}^n$

$$\begin{aligned} &= (I \vec{x} - \vec{u}\vec{u}^T \vec{x}) \\ &= \vec{x} - \vec{u}\vec{u}^T \vec{x} \\ &= \vec{x} - \lambda \vec{u} \end{aligned}$$

$\leftarrow Ix = x$ with compute cost $O(n)$

$\leftarrow u^T x = \lambda$ for some scalar and compute cost for dot product is $O(n)$

$\leftarrow \vec{x} - \lambda \vec{u} = O(n)$ since we multiply each value by λ and take away

So as a result our compute cost is:

we first re-write our expression as $\vec{I}\vec{x} - \vec{u}\vec{u}^T\vec{x}$ (1)
 then we calculate dot product $\vec{u}\vec{u}^T$ at cost n
 then we calculate $\vec{I}\vec{x}$ again at cost n
 then we are left with $\vec{x} - \lambda\vec{u}$ where multiplying by scalar and subtracting \vec{x} from $\lambda\vec{u}$ is at a cost n

so cost of algorithm = $n + n + n = 3n = O(n)$ \square .

EXERCISE 3

(i) let $x^T = [4, 1, 3, -3, -1]$ $\|x\|_1$, $\|x\|_2$, $\|x\|_\infty$

$$\|x\|_1 = \sum_{i=1}^n |x_i| = 4 + 1 + 3 + 3 + 1 = \boxed{12}$$

$$\|x\|_2 = \left(\sum_{i=1}^n x_i^2 \right)^{1/2} = (4^2 + 1^2 + 3^2 + 3^2 + 1^2)^{1/2} = \boxed{6}$$

$$\|x\|_\infty = \max_{1 \leq i \leq n} |x_i| = \boxed{4}$$

(ii) $A = \begin{bmatrix} 3 & -5 & 0 \\ 2 & 3 & 2 \\ 6 & -7 & 8 \end{bmatrix}$ $\|A\|_1$ and $\|A\|_\infty$

By Theorem 1.5 $\|A\|_1 = \max_{1 \leq j \leq n} \sum_{i=1}^m |a_{ij}|$, max column sum.

$$\therefore \|A\|_1 = \max (|3| + |2| + |6|, |5| + |3| + |7|, |0| + |2| + |8|) = \max(11, 15, 10) = \boxed{15}$$

$$\|A\|_\infty = \max_{1 \leq i \leq m} \sum_{j=1}^n |a_{ij}|, \text{ max row sum by 1.5}$$

$$\therefore \|A\|_\infty = \max (|3| + |-5| + |0|, |2| + |3| + |2|, |6| + |-7| + |8|) = \max(8, 7, 21) = \boxed{21}$$