

Cost analysis of the algorithm (c_3 → problem_2.py)

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class Kruskall:
    def __init__(self, size):
        self.size = size
        self.edges = []
        self.vertex = ['']*size

    def createEdgeList(self, weight, u, v):
        if 0 <= u < self.size and 0 <= v < self.size:
            self.edges.append((u, v, weight))

    def createVertexList(self, vertex, d):
        if 0 <= vertex < self.size:
            self.vertex[vertex] = d

    def find(self, parent, i):
        if parent[i]==i:
            return i
        return self.find(parent, parent[i])

    def union(self, parent, rank, x, y):
        xRoot = self.find(parent, x)
        yRoot = self.find(parent, y)
        if rank[xRoot] < rank[yRoot]:
            parent[xRoot]=yRoot
        elif rank[xRoot]>rank[yRoot]:
            parent[yRoot]= xRoot
        else:
            parent[yRoot]=xRoot
            rank[xRoot] +=1

    def kruskalProblem(self):
        self.edges=sorted(self.edges, key=lambda item:
item[2])
        parent = [i for i in range(self.size)]
        rank = [0]*self.size

        result = []

        for u, v, weight in self.edges:
            x = self.find(parent, u)
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        y = self.find(parent,v)                                #c26=E.log n
        if x != y:                                             #c27=E
            result.append((u,v,weight))                        #c28=E
            self.union(parent,rank,x,y)                        #c29=E.log n

    mstResult=[]                                              #c30=1

    for(u,v,weight) in result:                                #c31=V-1
        mstResult.append((self.vertex[u],self.vertex[v],weight))
#c32=V-1

    return mstResult                                          #c33=1

```

*n representa o número de vértices no grafo

*E representa as arestas

- Basic operation:

$$C_{11}, C_{12}, C_{13} = \log n$$

$$C_{25}, C_{26}, C_{29} = E \cdot \log n$$

- Time complexity calculation:

$$T(n) = (c_{11} + c_{12} + c_{13}) \cdot (\log n) + (c_{25} + c_{26} + c_{29}) \cdot (E \cdot \log n)$$

$$T(n) = 3 \log n + 3(E \cdot \log n)$$

$$T(n) = E \cdot \log n$$

$$T(n) \in O(E \cdot \log n)$$

- Solving the recurrence:

$$T(n) = E \cdot \log E + E \cdot \log n, \quad n \geq 1, \quad T(0) = O(E \cdot \log E)$$

$$T(1) = E \cdot \log E + E \cdot \log$$