

# **CSCI 48** *Intro. to Computer Science*

## **Lecture 9: BST** (*insert, delete*)

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<http://www.cs.toronto.edu/~ahchanaei/teaching/20165/csci48/>

# Last week

- ❖ Binary trees (branch factor =2)
- ❖ Depth-first traversal
  - inorder, preorder, and postorder
- ❖ Breadth-first traversal
  - level-order
- ❖ Binary Search Trees
- ❖ **Today**
  - More on BST
    - insert
    - delete

# Binary Search Trees

- ❖ Add ordering conditions to a binary tree:
  - data are comparable
  - data in left subtree are less than node.data
  - data in right subtree are more than node.data

# Binary Search Trees

# Binary Search Trees

- ❖ a BST with 1 node has height 1
  - ❖ a BST with 3 nodes may have height 2
  - ❖ a BST with 7 nodes may have height 3
  - ❖ a BST with 15 nodes may have height 4
  - ❖ a BST with  $n$  nodes may have height  $\lceil \lg n \rceil$
- 
- ❖ if the BST is “balanced”, then we can check whether an element is present in about  $\lg n$  node accesses
    - This is significantly faster than a linear search:  $O(n)$

# bst\_contains

```
def bst_contains(node, value):
```

```
    """  
    Return whether tree rooted at node contains value.
```

```
    Assume node is the root of a Binary Search Tree
```

```
    :param node: node of a Binary Search Tree
```

```
    :type node: BinaryTree|None
```

```
    :param value: value to search for
```

```
    :type value: object
```

```
    :rtype: bool
```

```
>>> bst_contains(None, 5)
```

```
False
```

```
>>> bst_contains(BinaryTree(7, BinaryTree(5), BinaryTree(9)), 5)
```

```
True
```

```
    """
```

# bst\_contains

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def bst_contains(node, value):
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    Return whether tree rooted at node contains value.
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```
>>> bst_contains(None, 5)
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```
False
```

```
>>> bst_contains(BinaryTree(7, BinaryTree(5), BinaryTree(9)), 5)
```

```
True
```

```
"""
```

```
if node is None:
```

```
    return False
```

```
elif value < node.data:
```

```
    return bst_contains(node.left, value)
```

```
elif value > node.data:
```

```
    return bst_contains(node.right, value)
```

```
else:
```

```
    return True
```

# bst\_insert

```
def insert(node, data):
    """Insert data in BST rooted at node if necessary, and return new root.
    Assume node is the root of a Binary Search Tree.
    :param node: root of a binary search tree.
    :type node: BinaryTree
    :param data: data to insert into BST, if necessary.
    :type data: object

    >>> b = BinaryTree(5)
    >>> b1 = insert(b, 3)
    >>> print(b1)
    5
        3
    <BLANKLINE>
    """

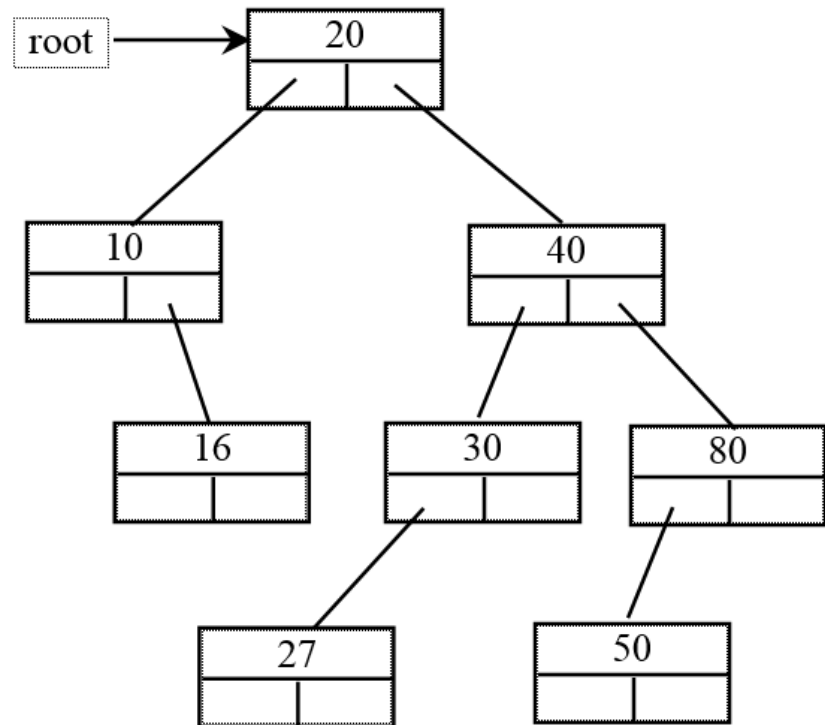
    return_node = node
    if not node:
        return_node = BinaryTree(data)
    elif data < node.data:
        node.left = insert(node.left, data)
    elif data > node.data:
        node.right = insert(node.right, data)
    else: # nothing to do
        pass
    return return_node
```



# bst\_insert

- ❖ **Let's trace it for a few examples:**

# bst\_delete

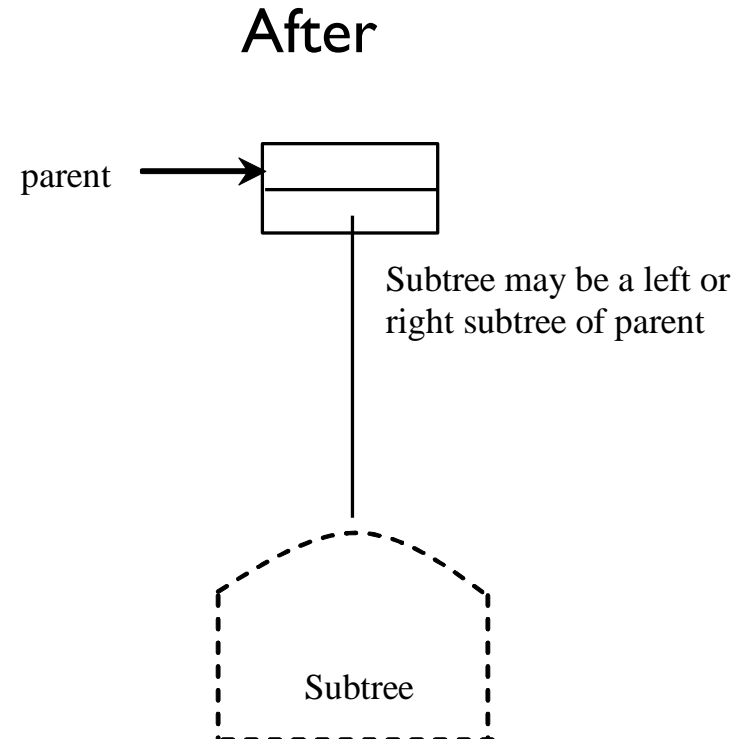
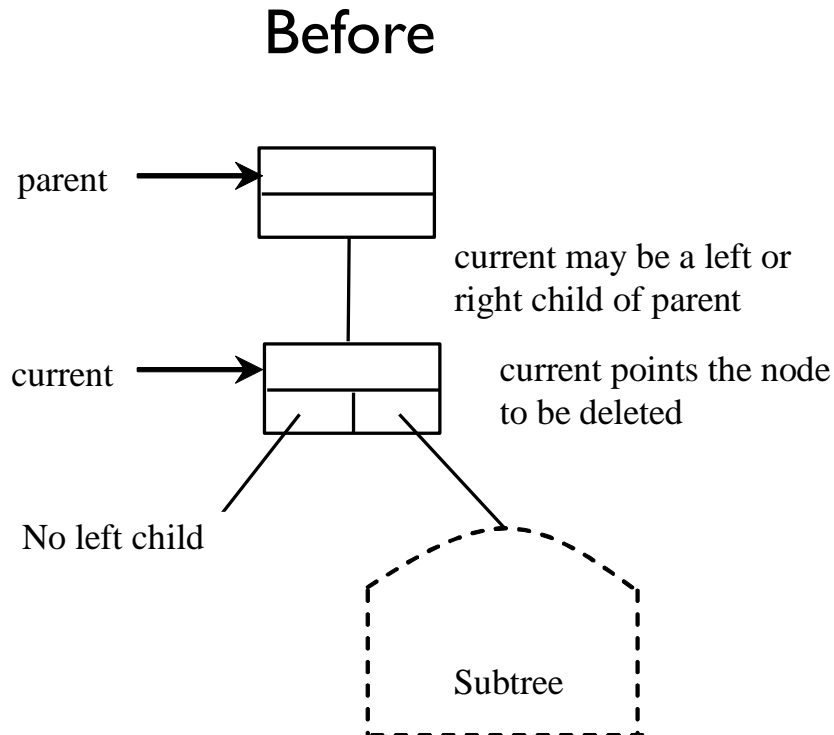


# bst\_delete

- ❖ First locate the node that contains the element and also its parent node.
- ❖ Let current point to the node that contains the element in the tree and parent point to the parent of the current node.
- ❖ There are two cases to consider ...

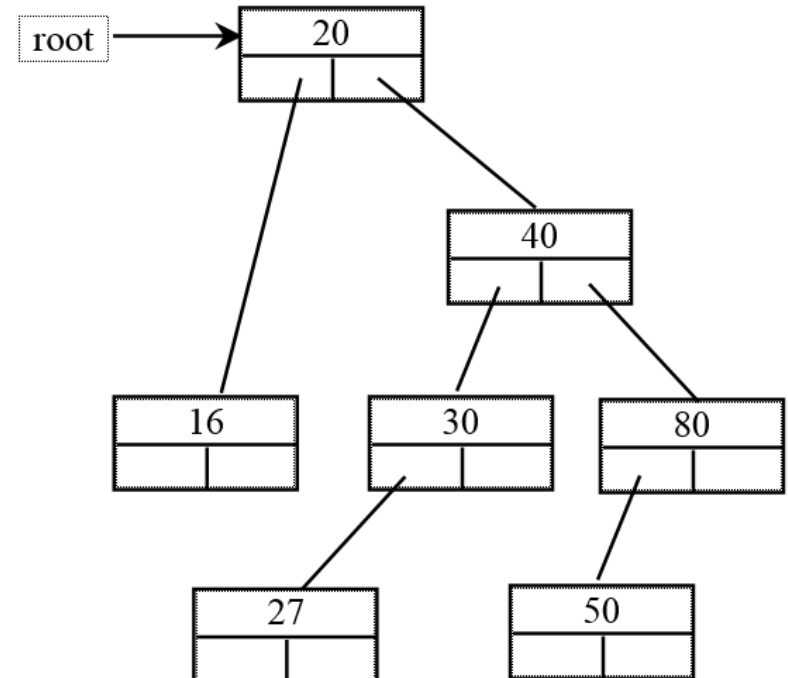
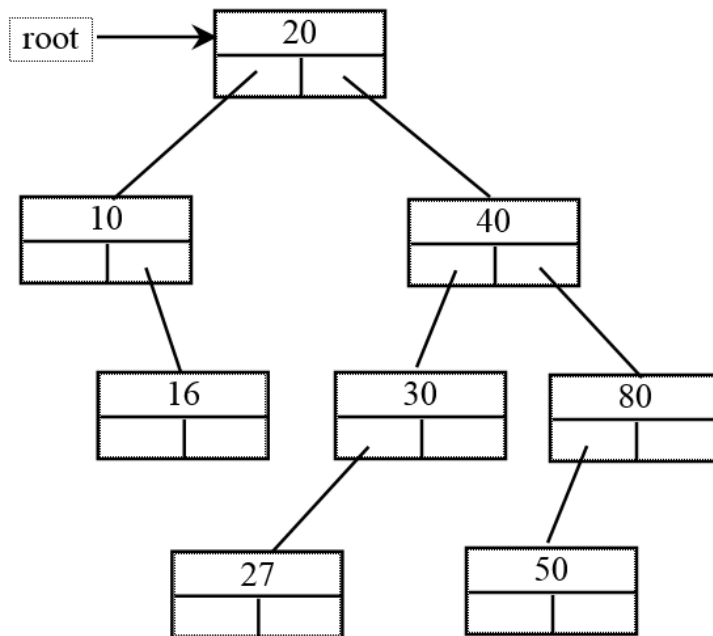
## Case I: The current node has no left child

- ❖ Simply connect the parent with the right child of the current node.



# Example for **Case I**. Deleting node 10

Connect the parent of node 10 with the right child of node 10.



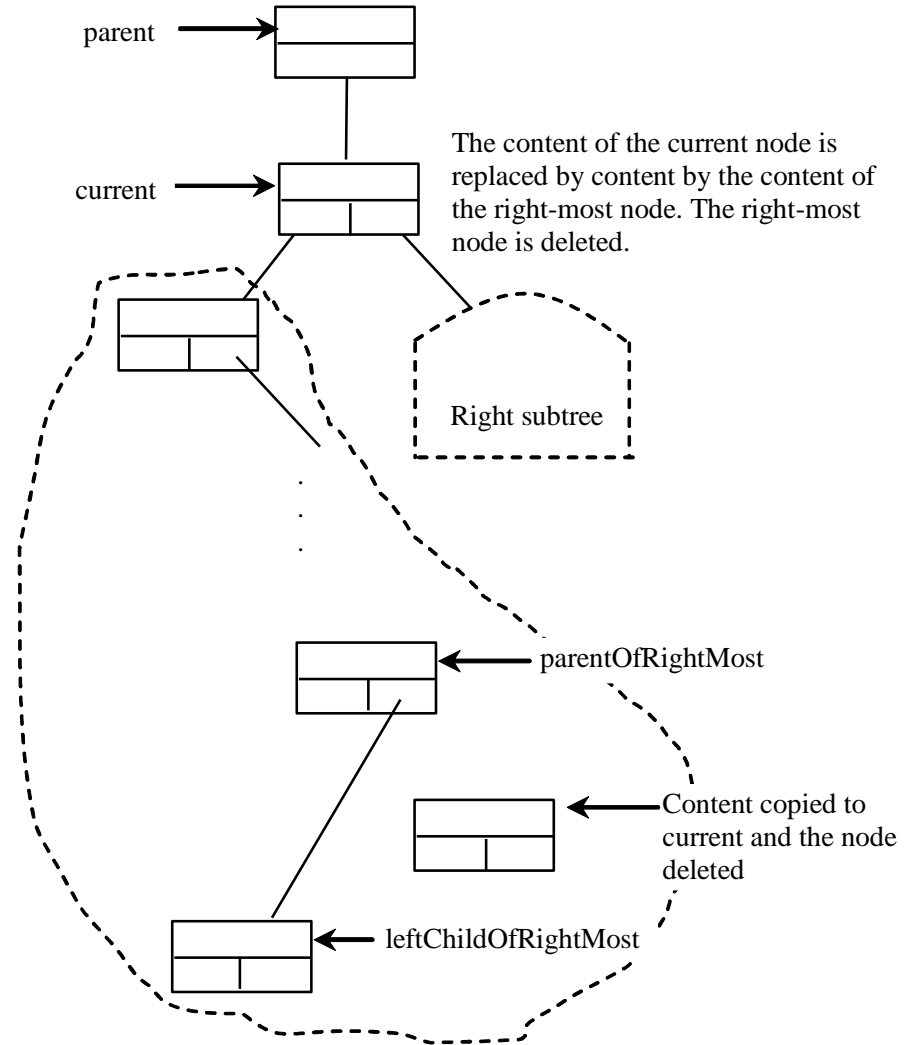
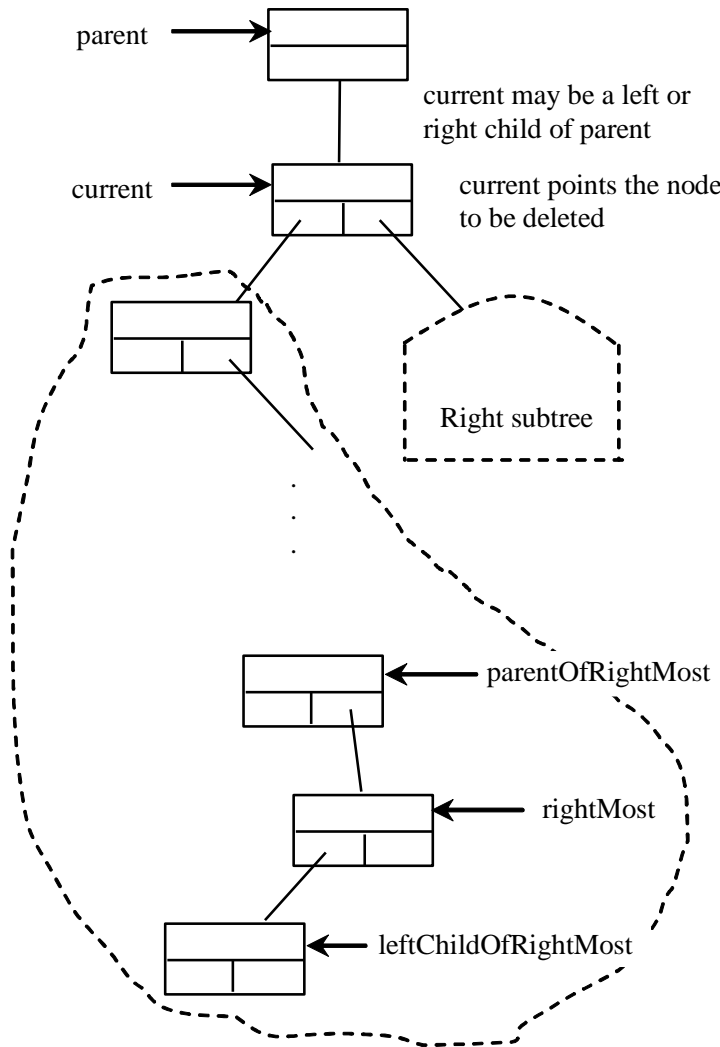
## **Case 2: The current node has a left child.**

- ❖ Let right\_most point to the node that contains the largest element in the left subtree of the current node.
- ❖ Let parent\_of\_right\_most point to the parent node of the right\_most node.

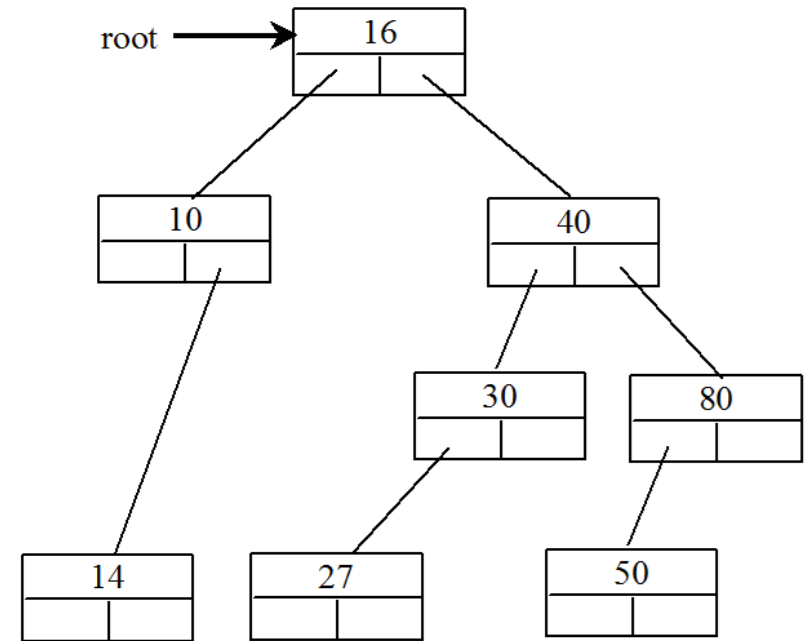
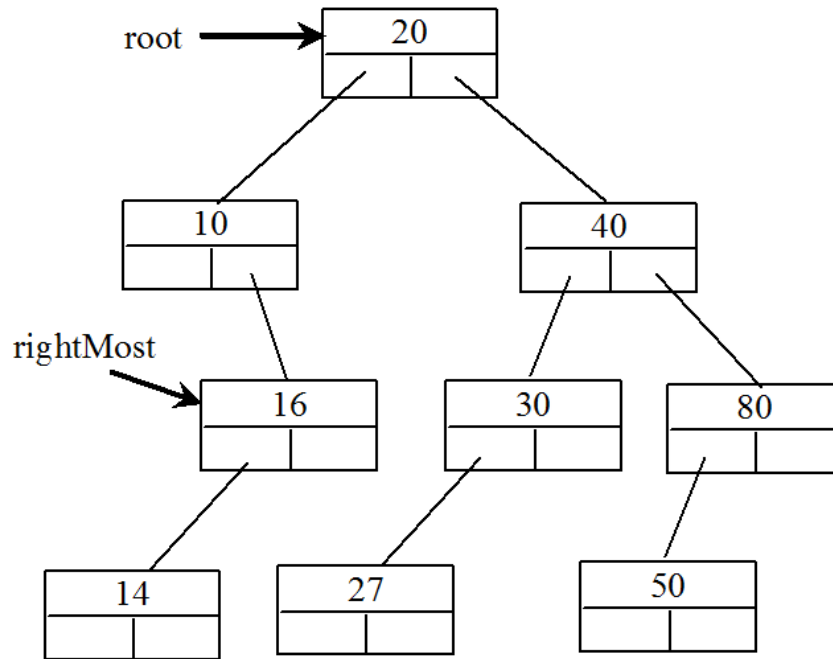
Then:

1. Replace the element value in the current node with the one in the right\_most node,
2. Connect the parent\_of\_right\_most node with the left child of the right\_most node.

# Case 2 (diagram)

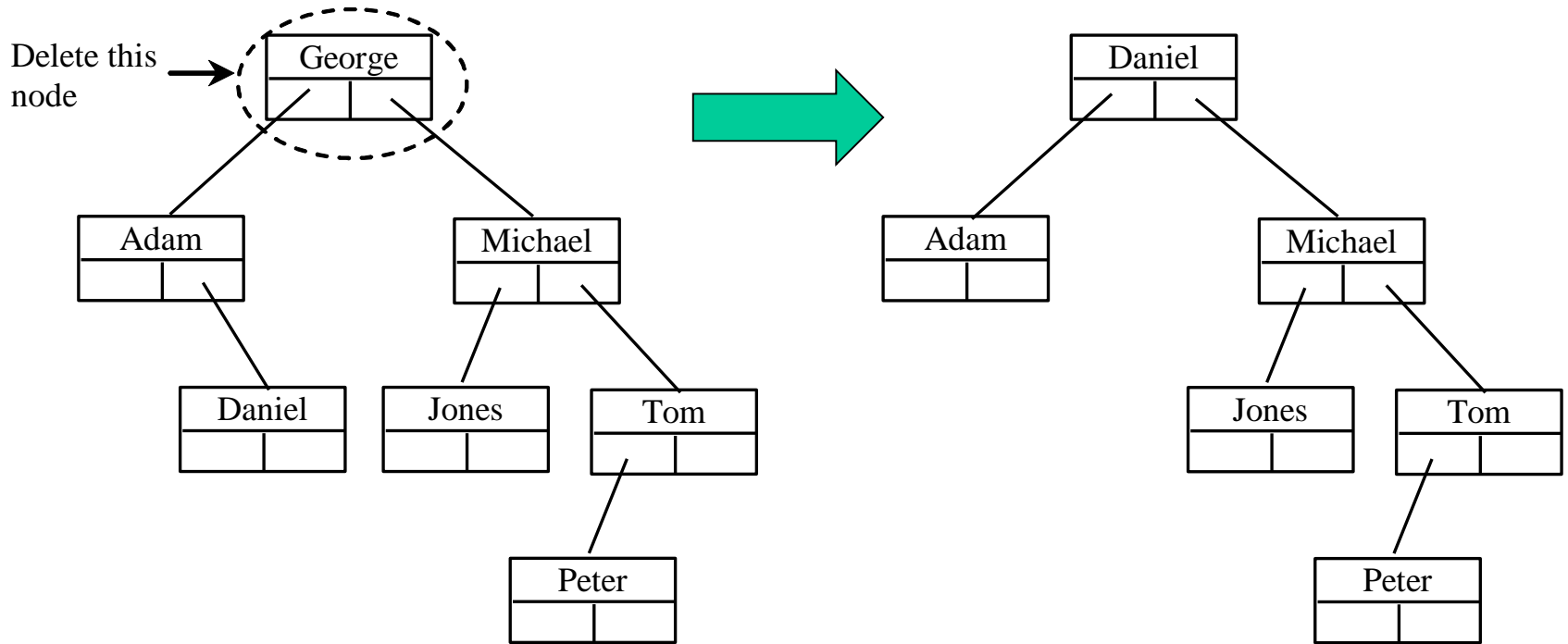


# Example for **Case 2. Deleting node 20**



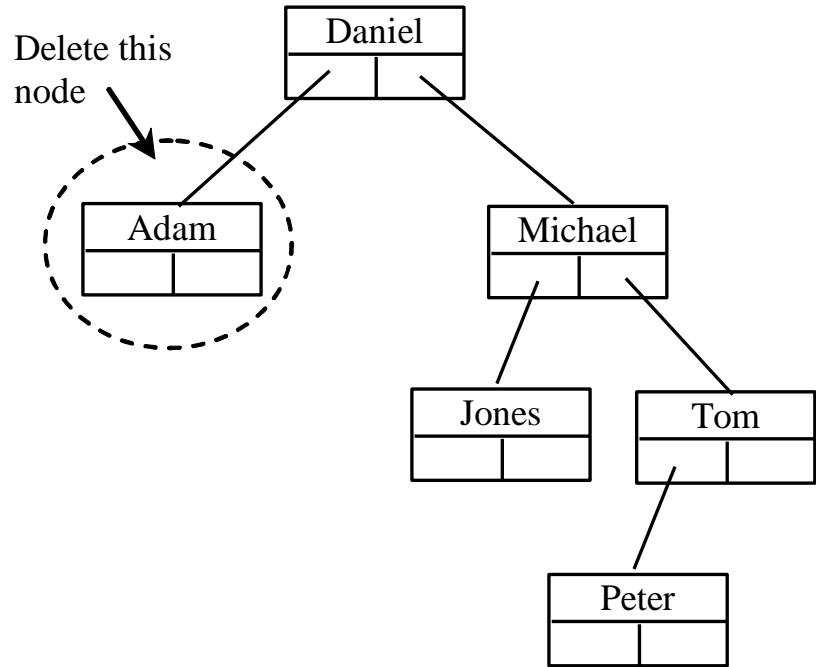
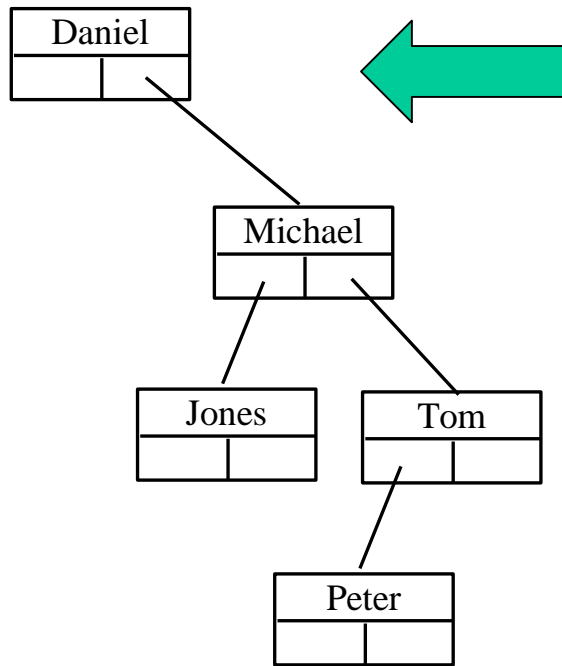


# More Examples



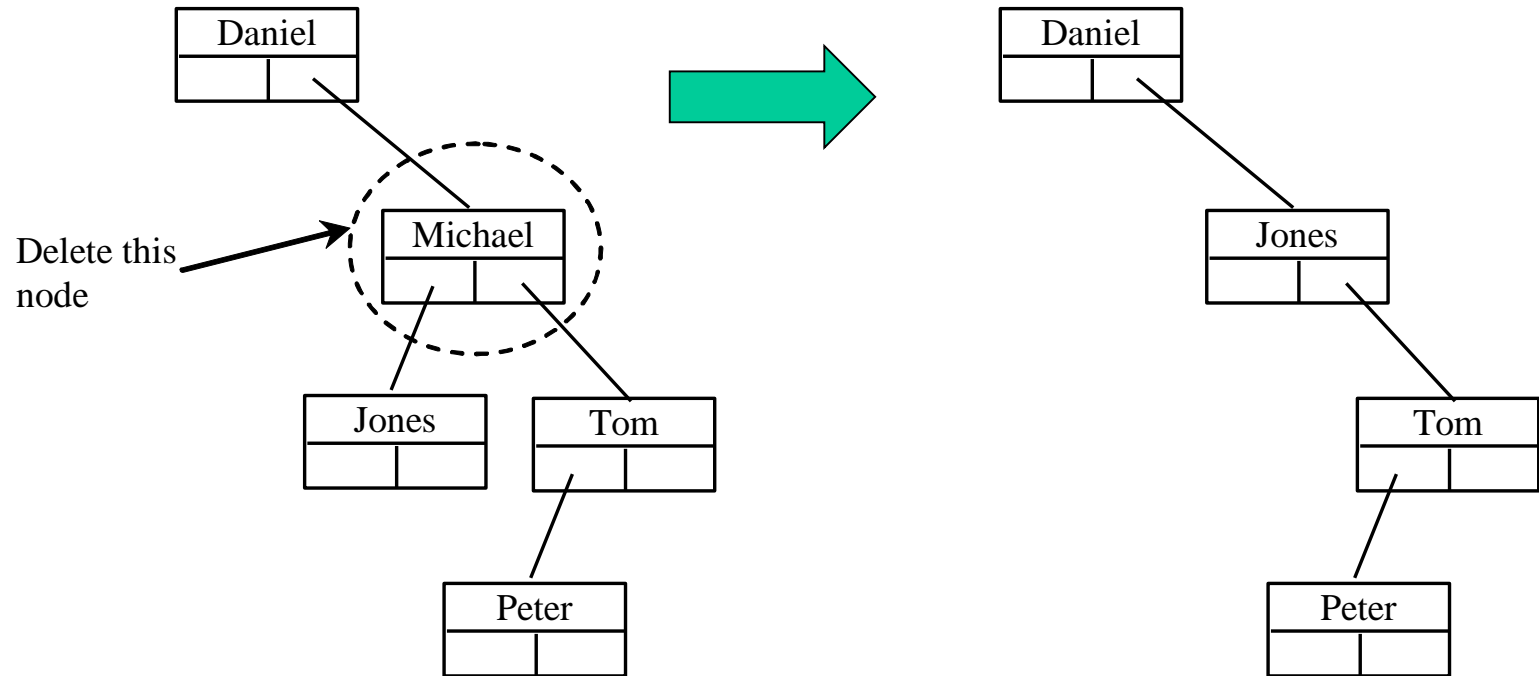
**Case 1 or 2?**    2

# More Examples



**Case 1 or 2?** |

# More Examples



**Case 1 or 2?** 2

## bst\_delete

- First locate the nodes that contain the element and its parent. Call them current and parent.

```
parent = None  
current = root
```

```
while current is not None and current.data != data:
```

```
    if data < current.data:  
        parent = current  
        current = current.left
```

```
    elif data > current.data:  
        parent = current  
        current = current.right
```

```
    else: pass # Element is in the tree pointed at by current
```

```
if current is None: return False # Element is not in the tree
```

# Case 1: bst\_delete

*# Case 1: current has no left child*

**if** current.left **is None**:

*# Connect the parent with the right child of the  
current node*

*# Special case, assume the node being deleted is at  
root*

**if** parent **is None**:

current = current.right

**else**:

*# Identify if parent left or parent right should  
be connected*

**if** data < parent.data:

parent.left = current.right

**else**:

parent.right = current.right

**else**:

*# Case 2: The current node has a left child*

## Case II: bst\_delete

```
# Locate the rightmost node in the left subtree of  
# the current node and also its parent
```

```
parent_of_right_most = current  
right_most = current.left
```

```
while right_most.right is not None:
```

```
    parent_of_right_most = right_most
```

```
    right_most = right_most.right # Keep going to the right
```

```
# Replace the element in current by the element in rightMost
```

```
current.element = right_most.element
```

```
# Eliminate rightmost node
```

```
    if parent_of_right_most.right == right_most:
```

```
        parent_of_right_most.right = right_most.left
```

```
else:
```

```
    # Special case: parent_of_right_most == current
```

```
    parent_of_right_most.left = right_most.left
```

```
return True # Element deleted successfully
```

# Summary

## ❖ Homework:

- In Slides 12 and 14,
  - replace every *left* with *right*, every *right* with *left*, and also *largest* with *smallest*.
- And, implement the method.

## ❖ Next Week:

- How `bst_delete` can be written recursively?