CSC148 winter 2018

binary trees week 8

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```

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Outline

general trees continued...

binary trees

traversals

binary search trees

queues, stacks, recursion

You may have noticed in the last slide there were no recursive calls, and a queue was used to process a recursive structure in level order.

Careful use of a stack allows you to process a tree in preorder.



tree inheritance issues

one approach to BinaryTree would be to make it a subclass of Tree, but there are some design considerations

- ▶ any client code that uses Tree would be required not to violate the branching factor (2) of BinaryTree
- one way to achieve this would be to make Tree immutable: make sure there is no way to change children or value, and then have subclasses that might be mutable

we will take a different approach: a completely separate BinaryTree class





BTNode

Change our generic Tree design so that we have two named children, left and right, and can represent an empty tree with None

special methods...

We'll want the standard special methods:

contains

True

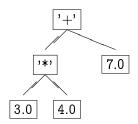
```
you've implemented contains on linked lists, nested Python lists, general
Trees before; implement this function, then modify it to become a method

def contains(node: BTNode, value: object) -> bool:
    """
    Return whether tree rooted at node contains value.

>>> contains(None, 5)
False
>>> contains(BTNode(5, BTNode(7), BTNode(9)), 7)
```

arithmetic expression trees

Binary arithmetic expressions can be represented as binary trees:





evaluating a binary expression tree

- ▶ there are no empty expressions
- ▶ if it's a leaf, just return the value
- otherwise...
 - ▶ evaluate the left tree
 - evaluate the right tree
 - combine left and right with the binary operator

Python built-in eval might be handy.





inorder

A recursive definition:

- visit the left subtree inorder
- visit this node itself
- visit the right subtree inorder

The code is almost identical to the definition.



preorder

- visit this node itself
- ▶ visit the left subtree in preorder
- visit the right subtree in preorder



postorder

- ▶ visit the left subtree in postorder
- visit the rightsubtree in postorder
- visit this node itself



level order

- ▶ visit root
- visit root's children
- ▶ visit root's grandchildren
- ▶ visit root's greatgrandchildren
- **.**..

definition

Add ordering conditions to a binary tree:

- data are comparable
- ▶ data in left subtree are less than node.data
- data in right subtree are more than node.data



why binary search trees?

Searchs that are directed along a single path are efficient:

- ▶ a BST with 1 one has height 1
- ▶ a BST with 3 nodes may have height 2
- ▶ a BST with 7 nodes may have height 3
- a BST with 15 nodes may have height 4
- ▶ a BST with n nodes may have height $\lceil \lg n \rceil$.



bst_contains

If node is the root of a "balanced" BST, then we can check whether an element is present in about $\lg n$ node accesses.

```
def bst_contains(node: BTNode, value: object) -> bool:
    """
    Return whether tree rooted at node contains value.

Assume node is the root of a Binary Search Tree

>>> bst_contains(None, 5)
    False
    >>> bst_contains(BTNode(7, BTNode(5), BTNode(9)), 5)
    True
    """
# use BST property to avoid unnecessary searching
```





mutation: insert

```
.....
Insert data in BST rooted at node if necessary, and return new root
Assume node is the root of a Binary Search Tree.
>>> b = BTNode(8)
>>> b = insert(b, 4)
>>> b = insert(b, 2)
>>> b = insert(b, 6)
>>> b = insert(b, 12)
>>> b = insert(b, 14)
>>> b = insert(b, 10)
>>> print(b)
        14
    12
        10
8
        6
```

def insert(node: BTNode, data: object) -> BTNode: