## Problem Set 3

- 1. (a) Let A be a DFA accepting the language L. Is the reverse of all the strings accepted by L(A) regular?
  - (b) Let L be a regular language over  $\{a, b, c\}$ . Define the projection of L with respect to  $\{b, c\}$  denoted  $L \downarrow \{b, c\}$  as the language

 $\{w' \mid w' \text{ is obtained from } w \in L \text{ after deleting all occurrences of symbol } a\}$ 

Is  $L \downarrow \{b, c\}$  regular?

- (c) Show that every NFA can be converted into an equivalent one with a single accepting state.
- (d) Let  $N = (Q, \Sigma, \delta, q_0, F)$  be a DFA. Construct an automaton  $N_1 = (Q, \Sigma, \delta_1, q_0, F_1)$  as follows:

•  $F_1 = F \cup \{q_0\}$ 

• Define  $\delta_1$  such that for any  $q \in Q$  and  $a \in \Sigma \cup \{\epsilon\}$ ,

$$\delta_1(q,a) = \delta(q,a)$$
 for  $q \notin F$  or  $a \neq \epsilon$ 

$$\delta_1(q, a) = \delta(q, a) \cup \{q_0\} \text{ for } q \in F \text{ and } a = \epsilon$$

Is  $L(N_1) = (L(N))^*$ ?

(e) Let L be a regular language. Is the language  $L_{\frac{1}{2}}$ , the set of first halves of strings in L regular? Formally,

$$L_{\frac{1}{2}} = \{x \mid \exists y, |x| = |y|, xy \in L\}$$

- 2. Let  $L_n = \{x \mid x \text{ is a binary number that is a multiple of } n\}$ . Show that for all  $n \geq 1$ ,  $L_n$  is regular.
  - 3. Recall that we defined an angelic acceptance condition for NFAs in class: a word w is accepted whenever it has at least one accepting run. Under this, we showed that the languages accepted by NFAs are regular. Consider the following devilish acceptance condition, which says that an NFA M accepts a word x iff every possible computation of M on x ends in an accept state. Show that NFAs with the devilish acceptance condition recognize the class of regular languages.
- 4. Let L be a regular language. Consider the language L' defined as

$$\{w \in L \mid \text{no proper prefix of } w \text{ is in } L\}$$

Show that L' is regular.

5. Show that any DFA accepting the language

$$L_n = \{x \in \{0,1\}^* \mid \text{ the } n \text{ th bit from the right is a } 1 \}$$

has  $\geq 2^{n-1}$  states.