

Canonical Representation

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CADSL

Canonical Forms

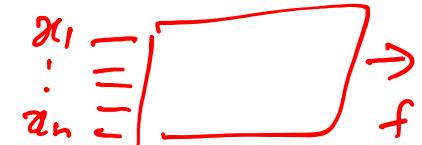
- Canonical Forms in common usage:

- ✓ – Truth Table $\rightarrow 2^n$
- Sum of Minterms (SOM) ✓ 2^n .
- Product of Maxterms (POM) ✓
- Binary Decision Diagram (BDD) ✓ graphical.
↑
compact
- Reed Muller Representation



Graphical Method

- ❖ BDD is canonical form of representation



- ❖ Shanon's expansion theorem

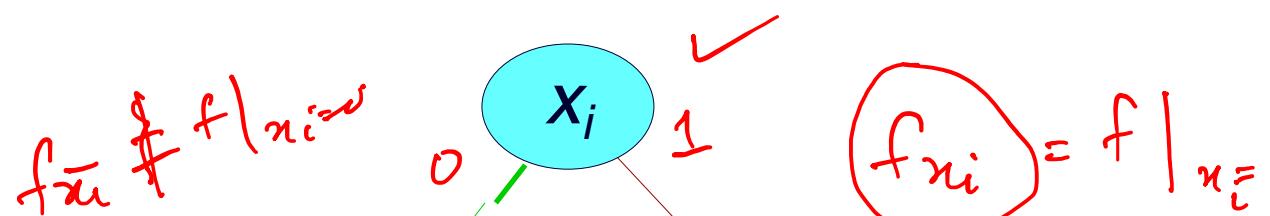
- ❖ $f(x_1, x_2, \dots, x_i, \dots, x_n) =$

$$x_i \cdot f(x_1, x_2, \dots, x_i=1, \dots, x_n) +$$

$$x_i' \cdot f(x_1, x_2, \dots, x_i=0, \dots, x_n)$$

$$\begin{aligned} f(x_1, x_2, \dots, x_n) &= \underline{x_i \cdot f_{x_i}} + \underline{\bar{x}_i \cdot f_{\bar{x}_i}} \\ &= x_i \cdot \underline{f_{x_i}} \oplus \bar{x}_i \cdot \underline{f_{\bar{x}_i}} \end{aligned}$$

$f|_{x_i=1}$ $f|_{x_i=0}$



$$f(x_1, x_2, \dots, x_i=1, \dots, x_n)$$

$$f(x_1, x_2, \dots, x_i=1, \dots, x_n)$$



Decision Structures

Truth Table

x_1	x_2	x_3	f
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

2^n

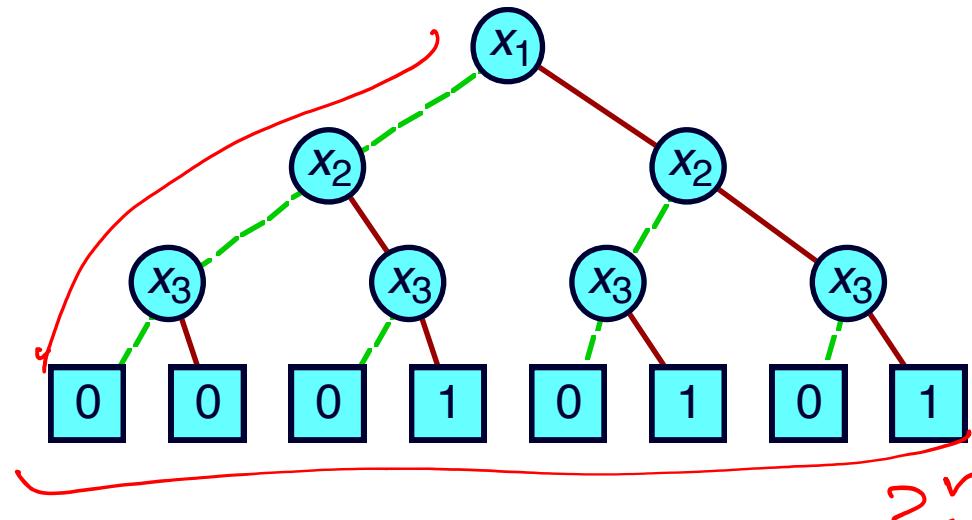
$$\underline{x_1 < x_2 < x_3}$$

order-

Binary

Decision Tree

$$2^{n+1} - 1$$

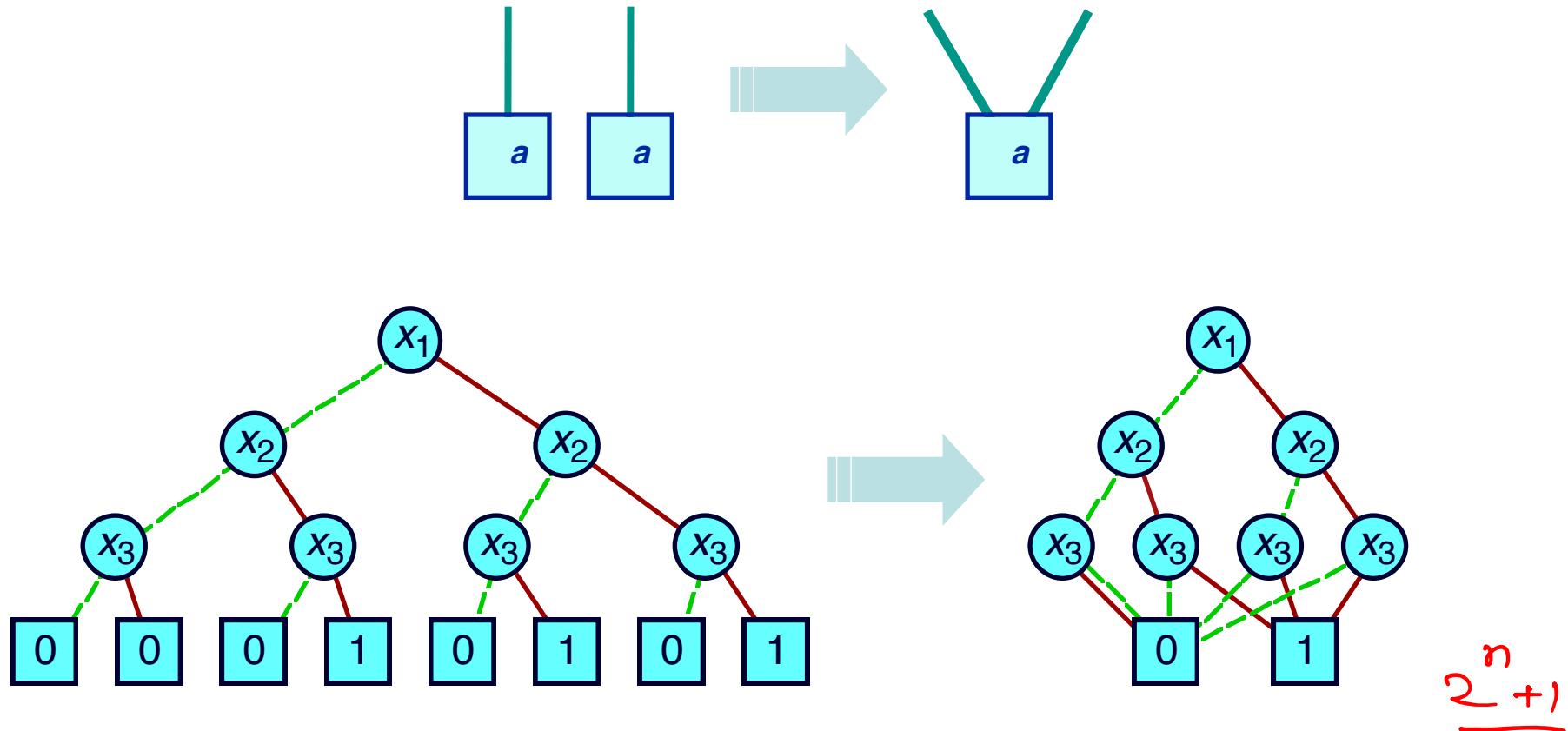


- Vertex represents decision
- Follow green (dashed) line for value 0
- Follow red (solid) line for value 1
- Function value determined by leaf value.



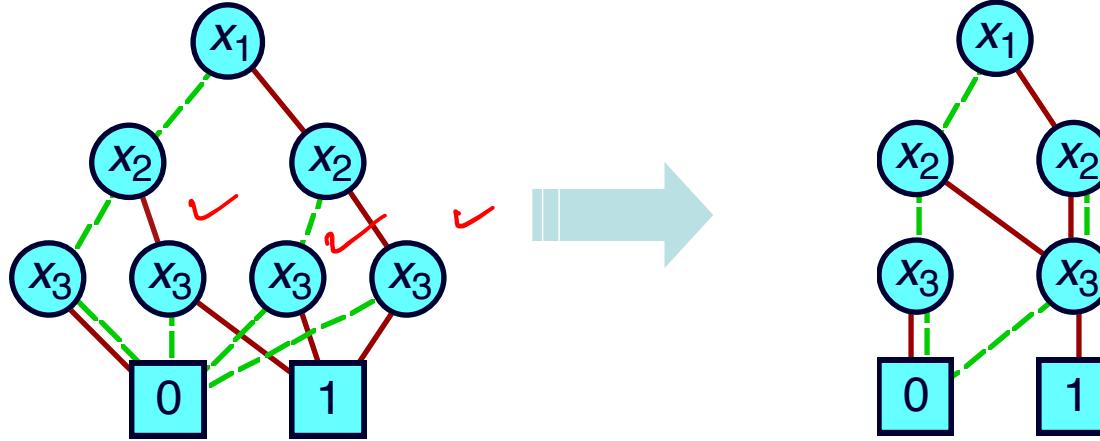
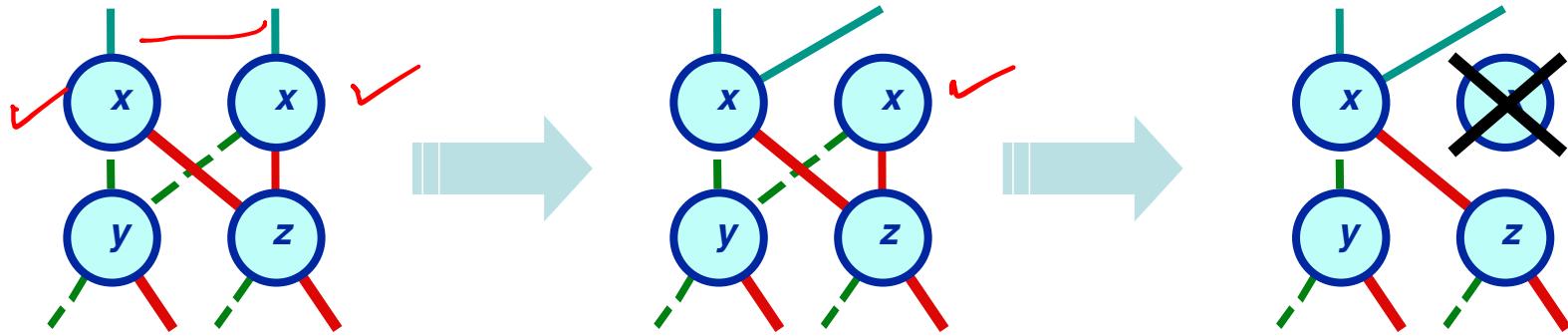
Reduction Rule #1

Merge equivalent leaves



Reduction Rule #2

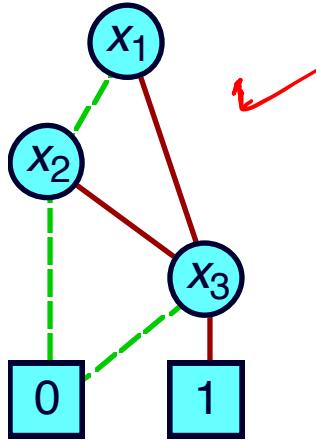
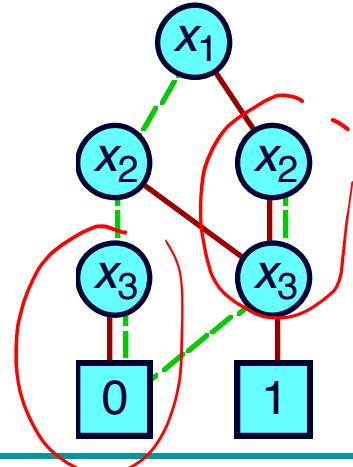
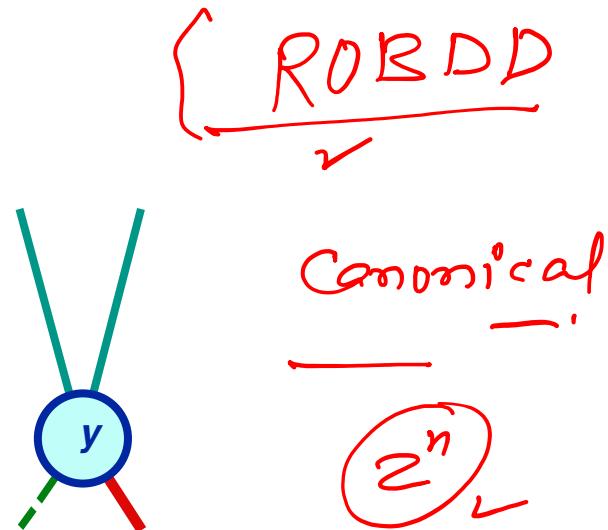
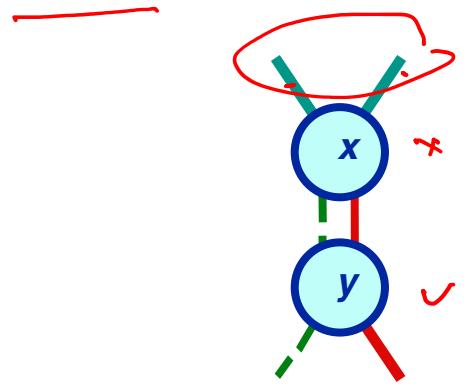
Merge isomorphic nodes



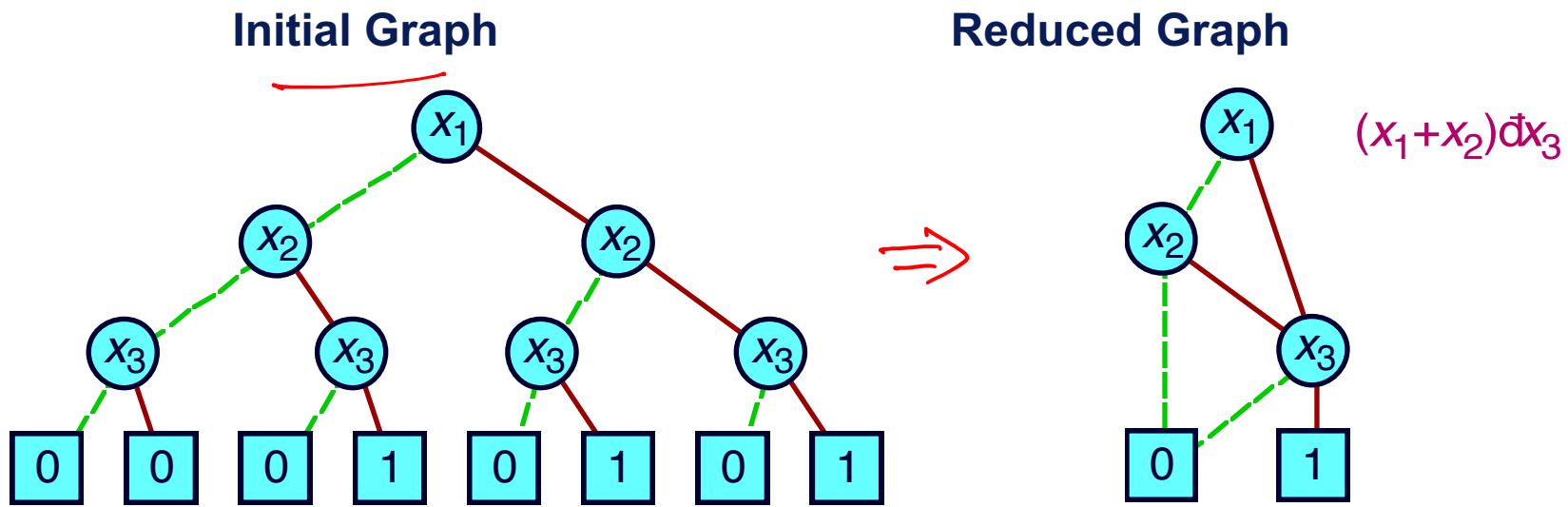
Reduction Rule #3

$x_1 < \underline{x_2 < x_3}$

Eliminate Redundant Tests



Example OBDD

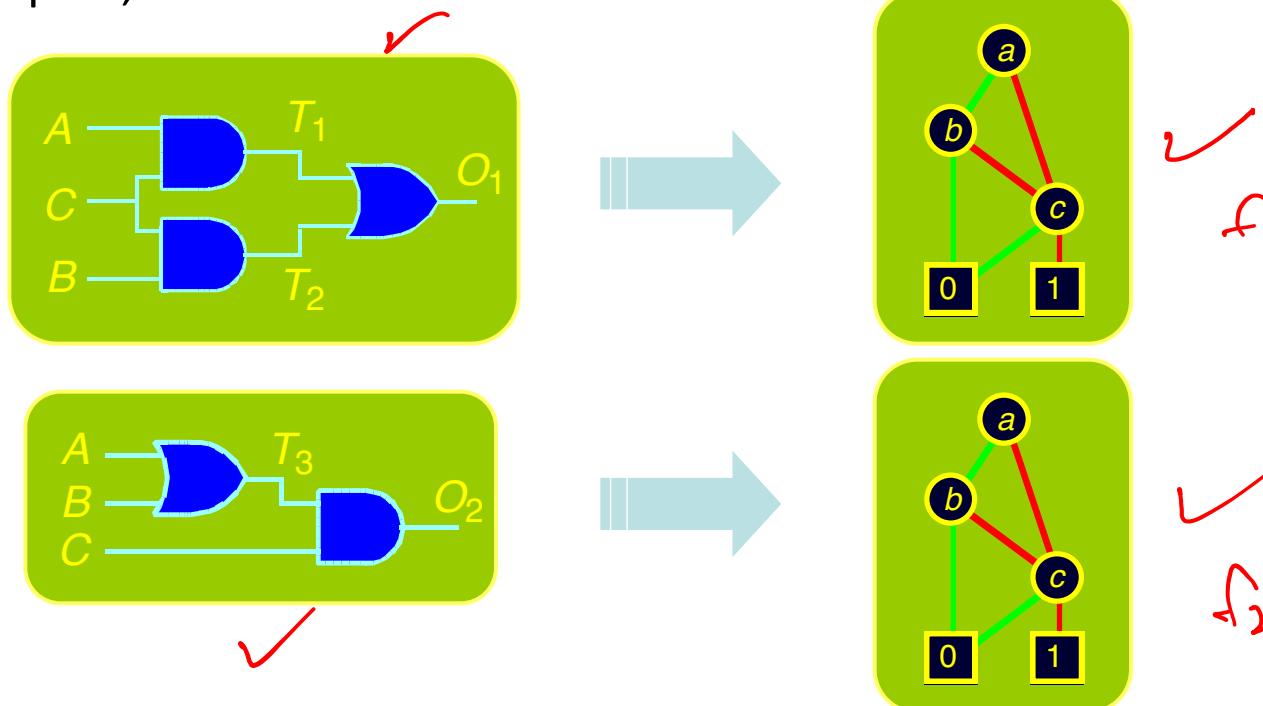


- Canonical representation of Boolean function
 - ❖ For given variable ordering f_1 f_2
 - Two functions equivalent if and only if graphs isomorphic
 - o Can be tested in linear time
 - Desirable property: *simplest form is canonical.*

Binary Decision Diagram

ROBDD

- Generate Complete Representation of Circuit Function
 - Compact, canonical form



- Functions equal if and only if representations identical
- Never enumerate explicit function values
- Exploit structure & regularity of circuit functions

$$f_1 \equiv f_2$$

Isomorphic

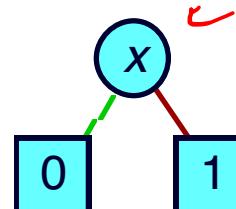


Example Functions

Constants

- ✓ 0 Unique unsatisfiable function
- 1 Unique tautology

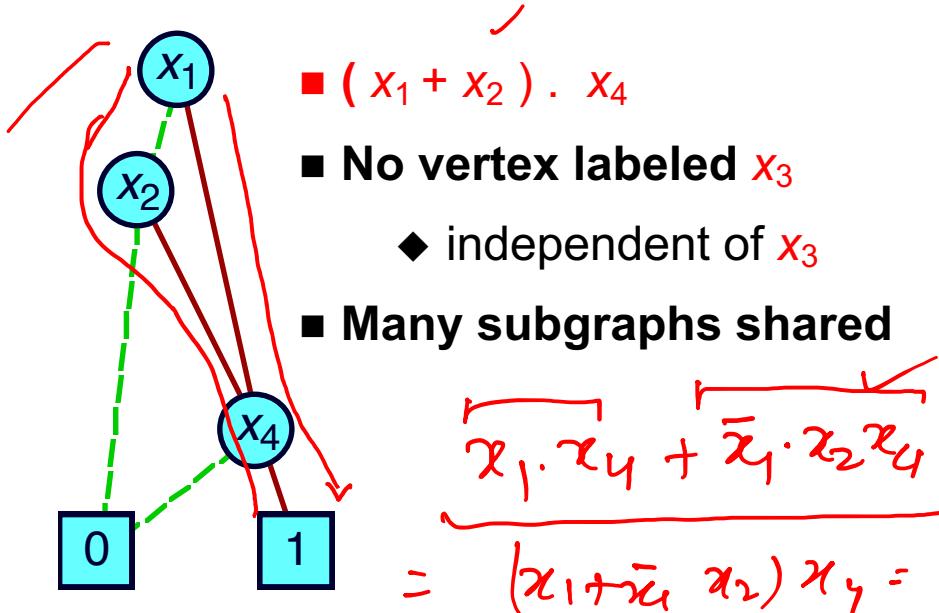
Variable



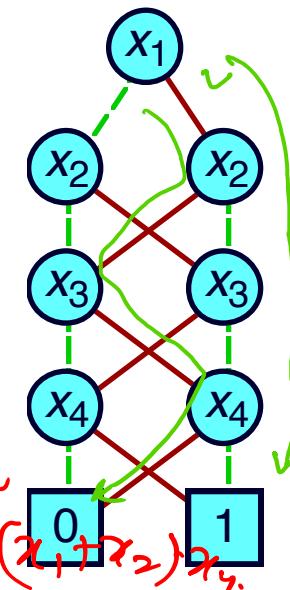
ROBDD

Treat variable
as function

Typical Function



Odd Parity



$f(x_1, x_2, x_3, x_4)$

Linear
representation



Effect of Variable Ordering

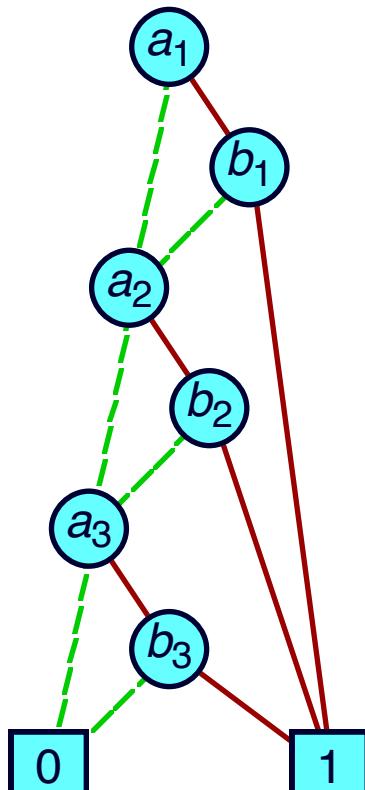
$$(a_1 \vee b_1) \wedge (a_2 \vee b_2) \wedge (a_3 \vee b_3)$$

$a_1 < b_1 < a_2 < b_2 < a_3 < b_3$

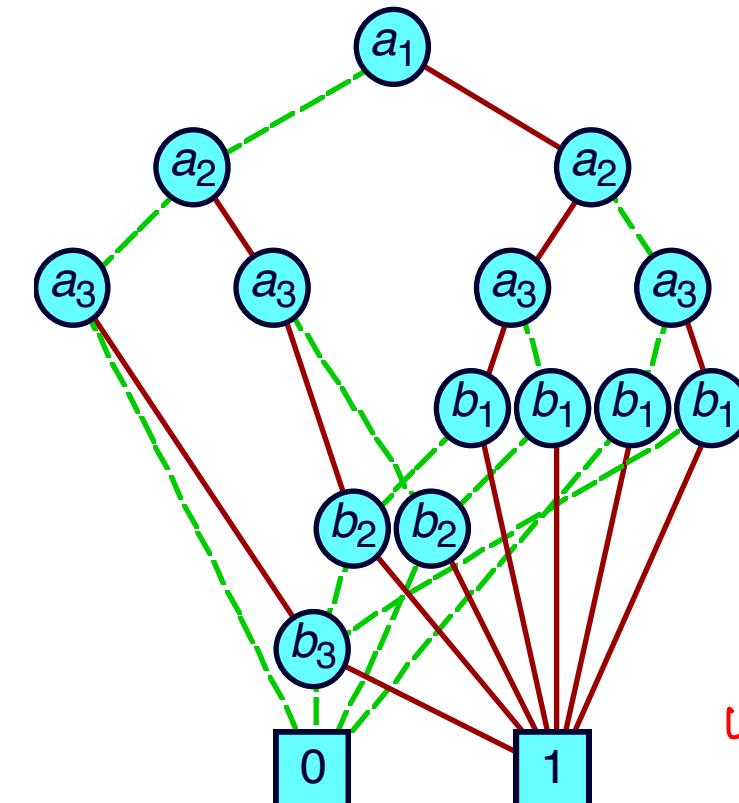
$a_1 < a_2 < a_3 < b_1 < b_2 < b_3$

Good Ordering

Bad Ordering



Linear Growth



Exponential Growth



Selecting Good Variable Ordering

- Intractable Problem
 - Even when problem represented as OBDD
 - i.e., to find optimum improvement to current ordering
- Application-Based Heuristics
 - Exploit characteristics of application
 - e.g., Ordering for functions of combinational circuit
 - Traverse circuit graph depth-first from outputs to inputs
 - Assign variables to primary inputs in order encountered



Selecting Good Variable Ordering

➊ Static Ordering ✓

- Fan In Heuristic
- Weight Heuristic

➋ Dynamic Ordering

- Variable Swap
- Window Permutation
- Sifting



Dynamic Variable Reordering



Richard Rudell, Synopsys



Periodically Attempt to Improve Ordering for All BDDs

❖ Move each variable through ordering to find its best location



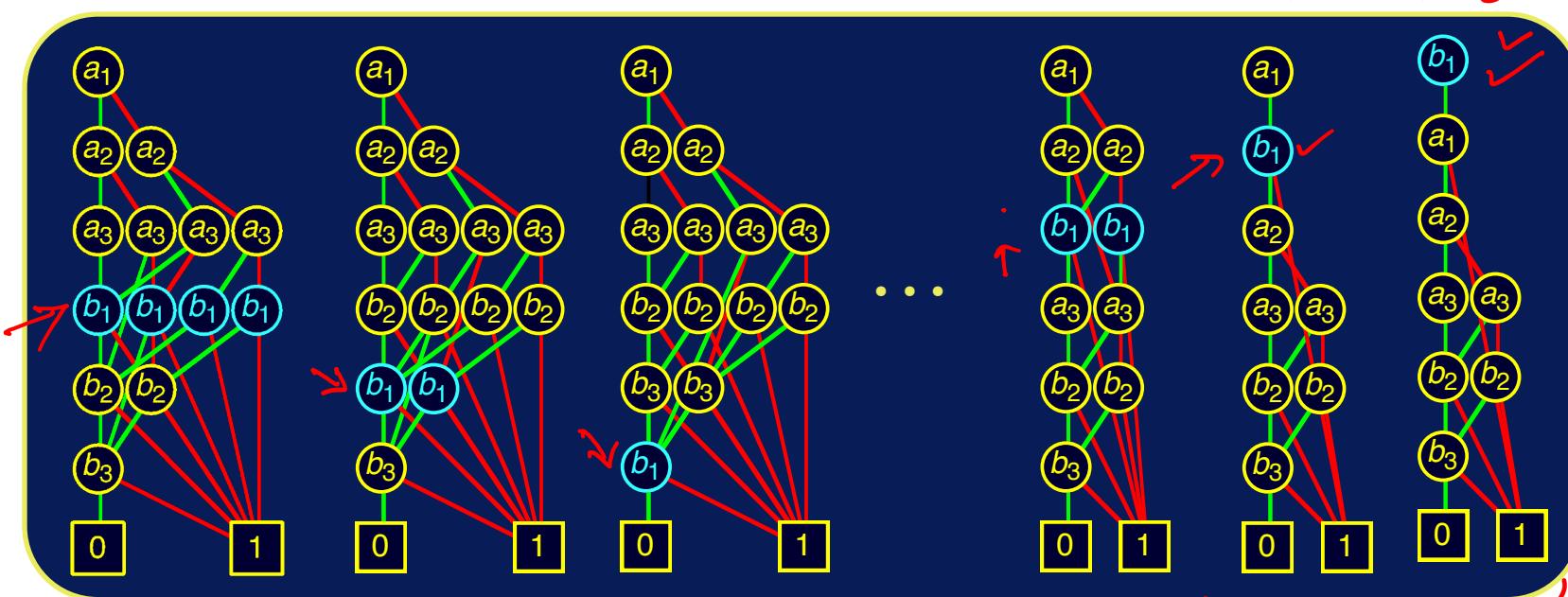
Has Proved Very Successful

❖ Time consuming but effective

Dynamic Reordering By Sifting

- Choose candidate variable
- Try all positions in variable ordering
 - 🟡 Repeatedly swap with adjacent variable
- Move to best position found

Best Choices



ROBDD Sizes & Variable Ordering

- Bad News 💣
 - Finding optimal variable ordering NP-Hard
 - Some functions have exponential BDD size for all orders e.g. multiplier
- Good News 😊
 - Many functions/tasks have reasonable size ROBDDs
 - Algorithms remain practical up to 1,000,000 node OBDDs
 - Heuristic ordering methods generally satisfactory
- What works in Practice 🤝
 - Application-specific heuristics e.g. DFS-based ordering for combinational circuits
 - Dynamic ordering based on variable sifting (*R. Rudell*)



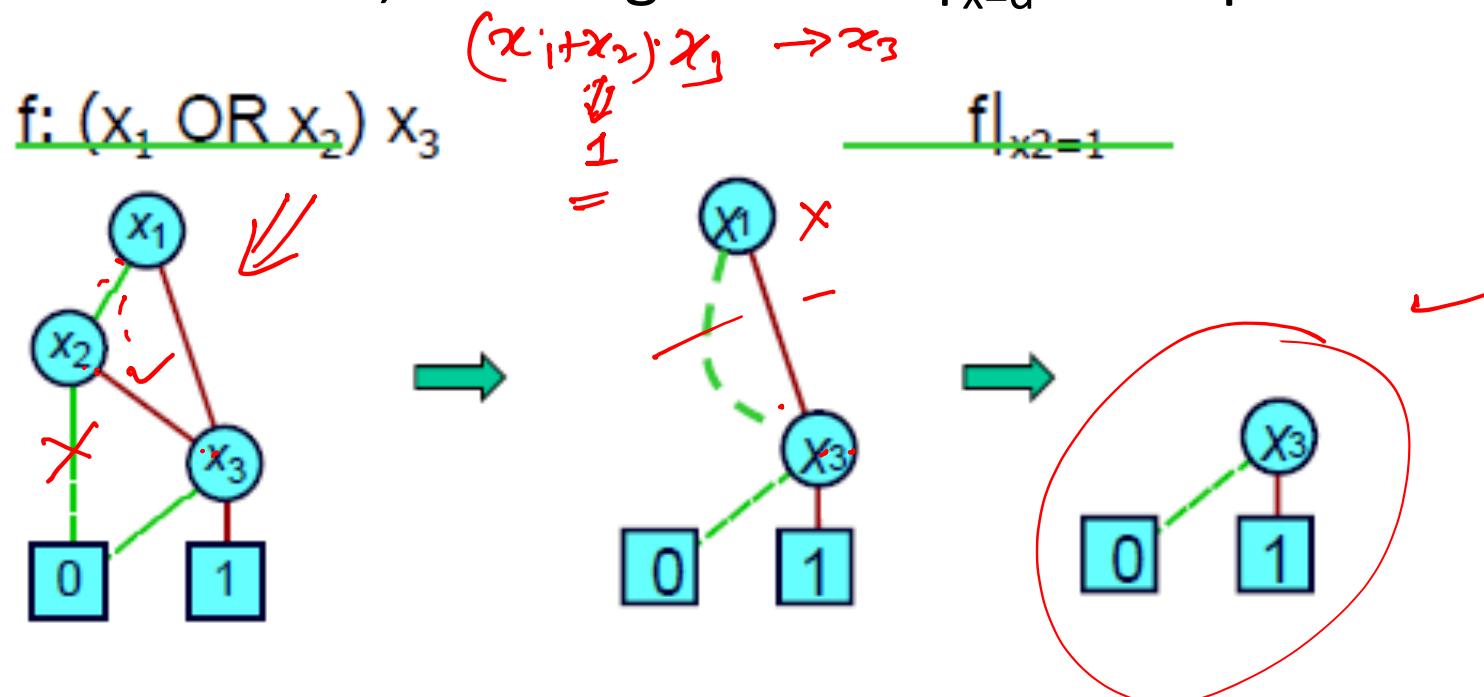
$$f_1 = a \cdot b + b \cdot c$$

16

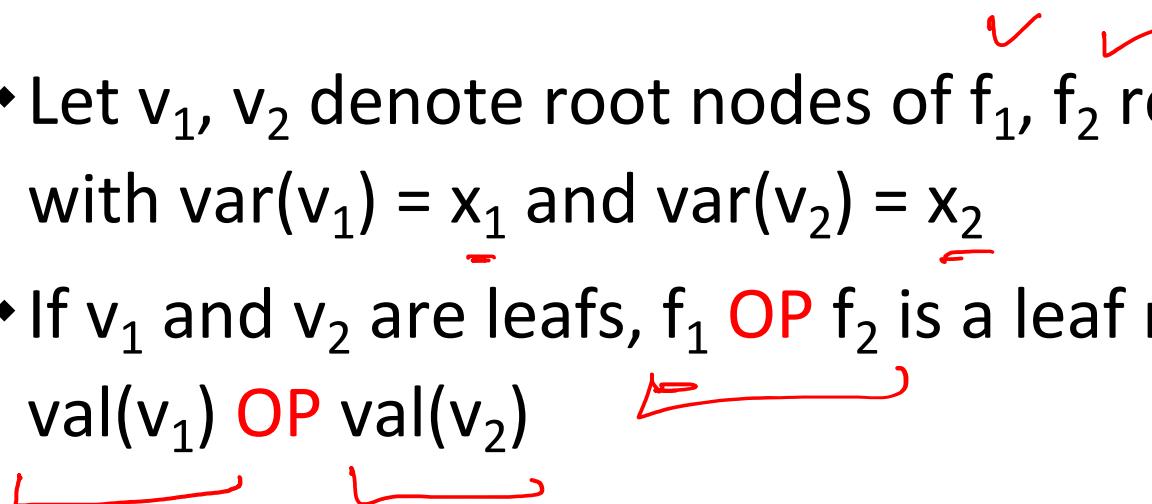
$$\begin{array}{c} a \\ \square \\ \square \end{array}$$
$$\begin{array}{c} b \\ \square \\ \square \end{array}$$
$$\begin{array}{c} c \\ \square \\ \square \end{array}$$
$$f = a \cdot b$$

Operations with BDD

- ❖ **Restriction:** A restriction to a function to $x=d$, denoted $f|_{x=d}$, where $x \in \text{var}(f)$, and $d \in \{0,1\}$, is equal to f after assigning $x = d$.
- ❖ Given BDD of f , deriving BDD of $f|_{x=d}$ is simple



Operations with BDD

- ❖ Let v_1, v_2 denote root nodes of f_1, f_2 respectively ,
with $\text{var}(v_1) = \underline{x_1}$ and $\text{var}(v_2) = \underline{x_2}$
 - ❖ If v_1 and v_2 are leafs, $f_1 \text{ OP } f_2$ is a leaf node with value
 $\text{val}(v_1) \text{ OP } \text{val}(v_2)$
- 



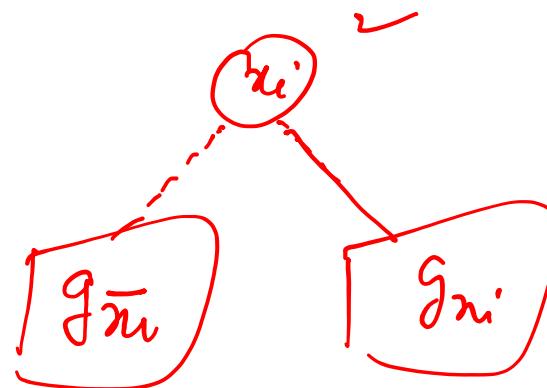
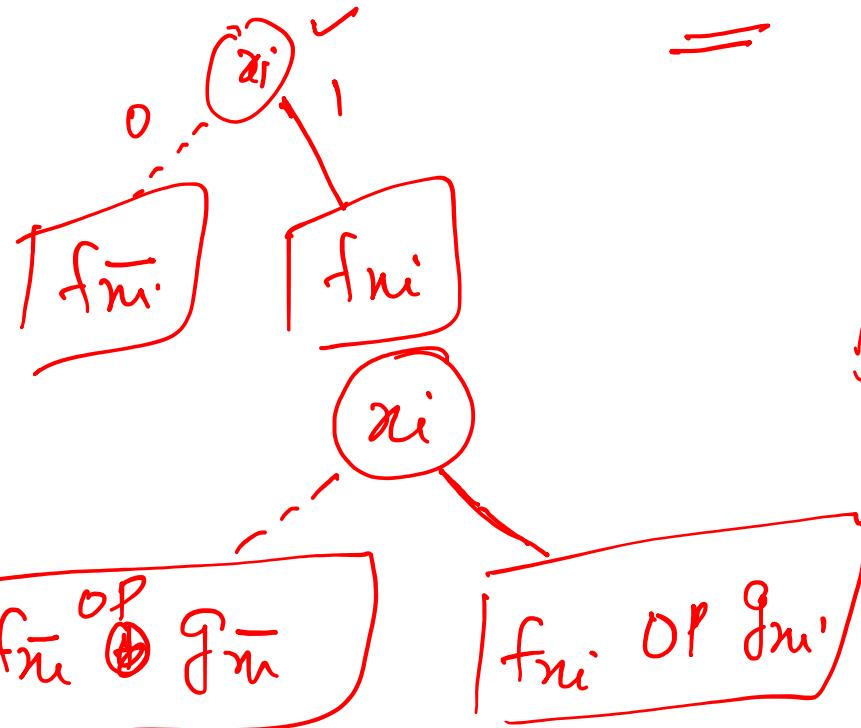
Operations with BDD

$$f(x_1, x_2, \dots, x_i, \dots, x_n)$$

$$x_i f_{xi} \oplus \bar{x}_i \bar{f}_{\bar{x}_i}$$

(OP)

$$g(x_1, x_2, \dots, x_i, \dots, x_n)$$

$$x_i g_{xi} \oplus \bar{x}_i \bar{g}_{\bar{x}_i}$$


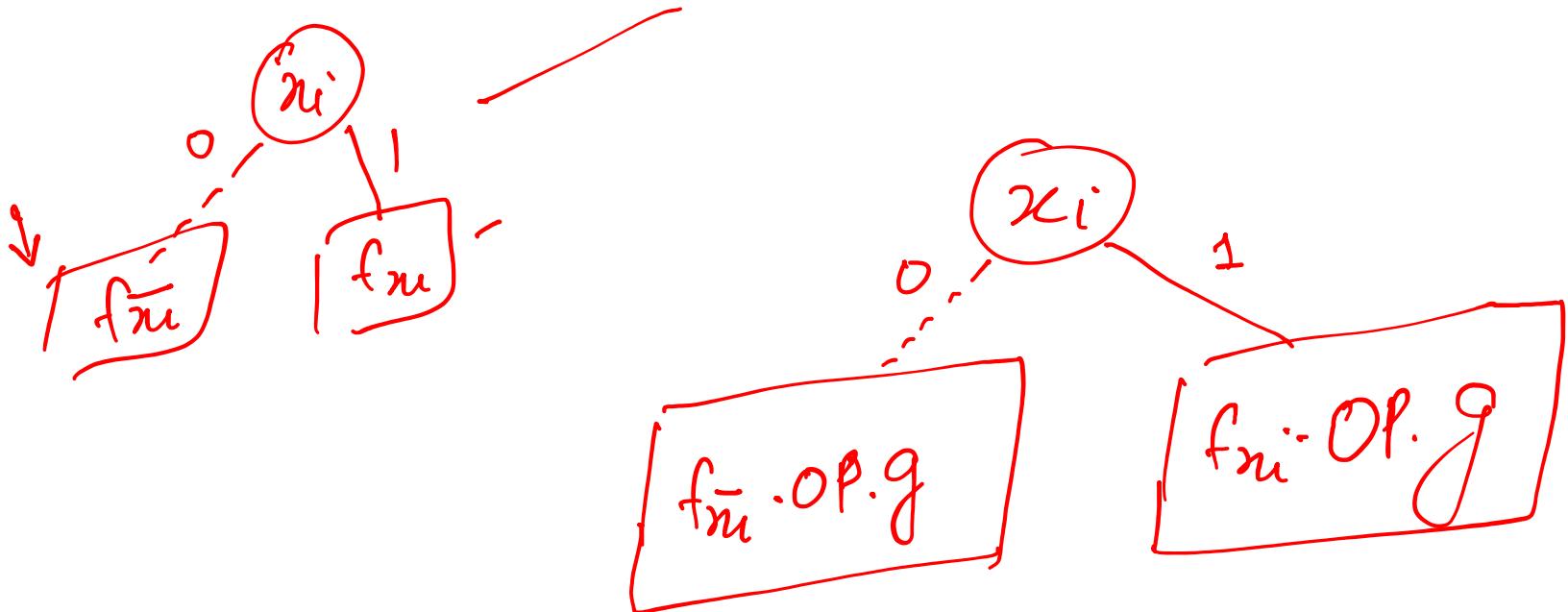
$f \cdot \text{OP} \cdot g$



Operations with BDD

$$f(x_1, x_2, \dots, x_i, \dots, x_n) \oplus g(x_1, x_2, \dots, x_j, \dots, x_n)$$

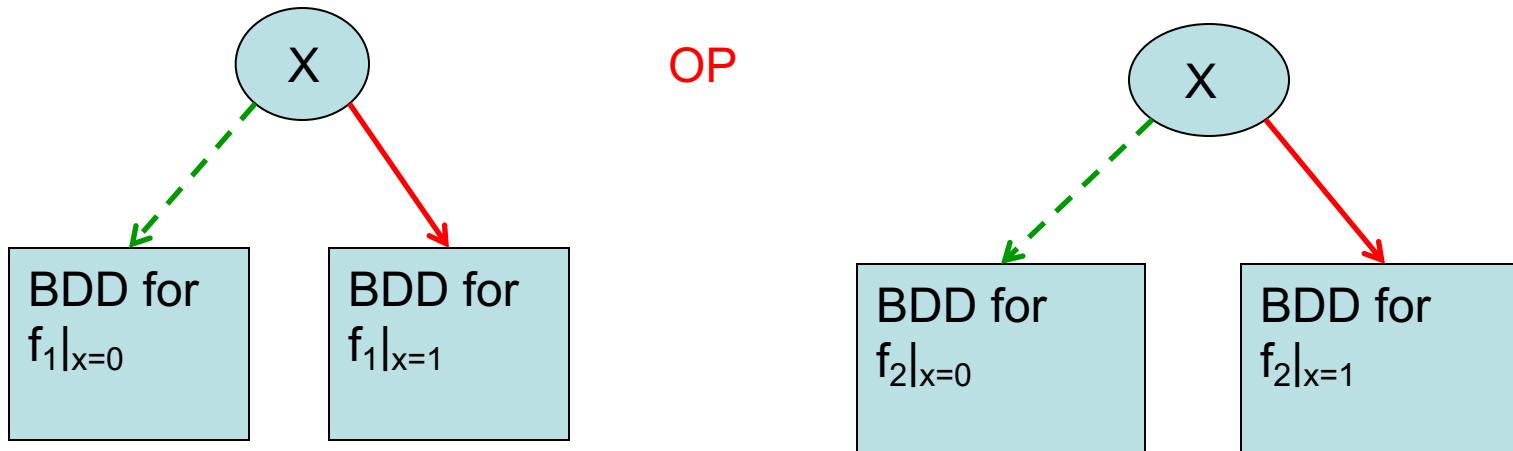
\uparrow does not have
 x_i



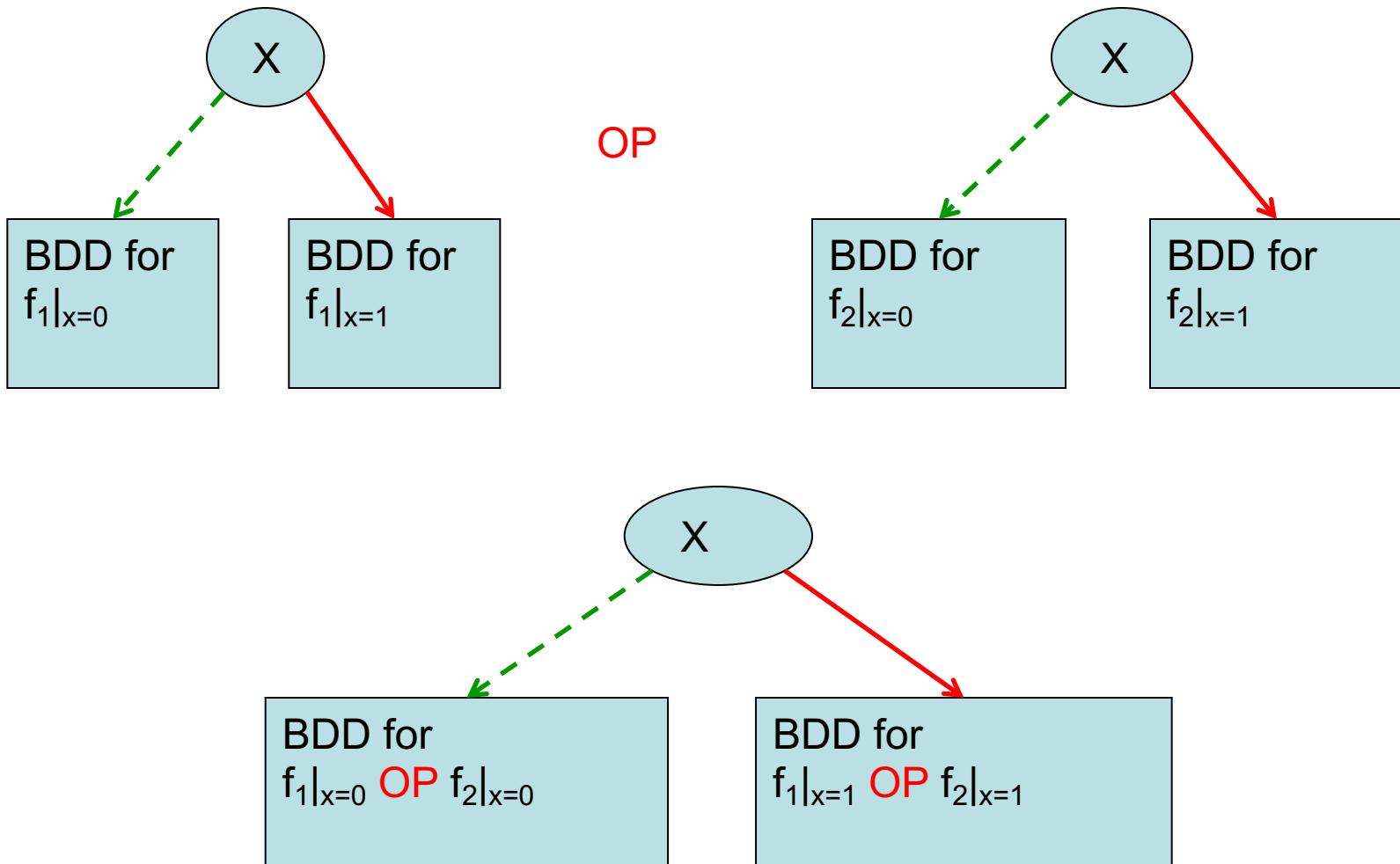
Operations with BDD

- ❖ If $x_1 = x_2 = x$, apply Shannon's expansion

$$f_1 \text{ OP } f_2 = x' \cdot (f_1|_{x=0} \text{ OP } f_2|_{x=0}) + x \cdot (f_1|_{x=1} \text{ OP } f_2|_{x=1})$$



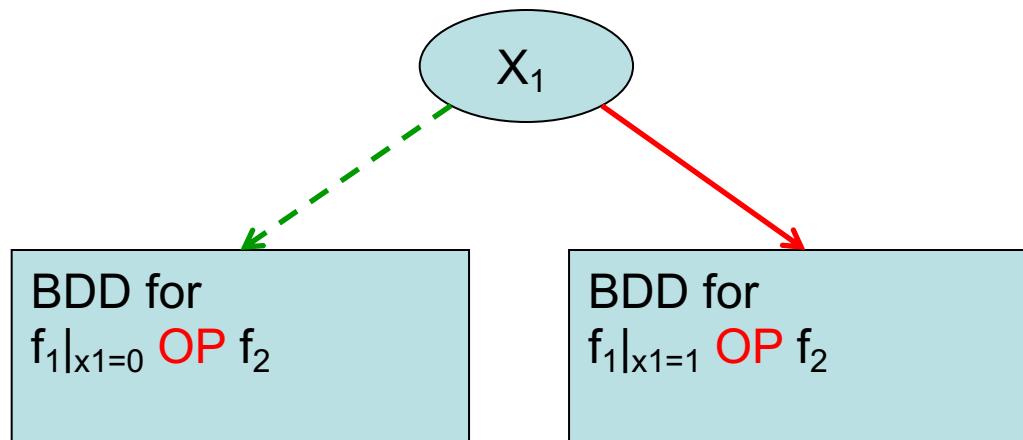
Operations with BDD



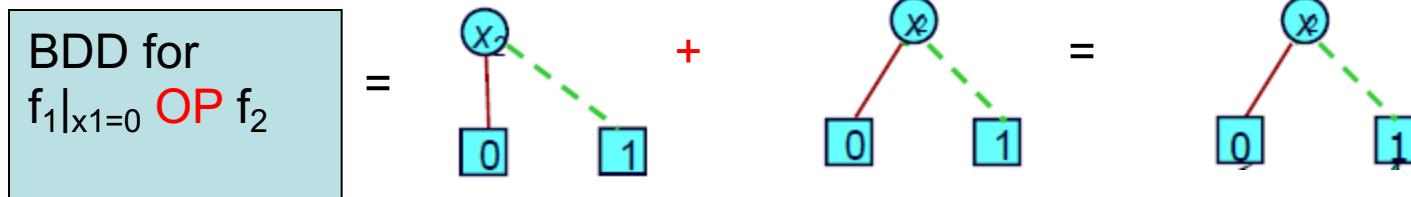
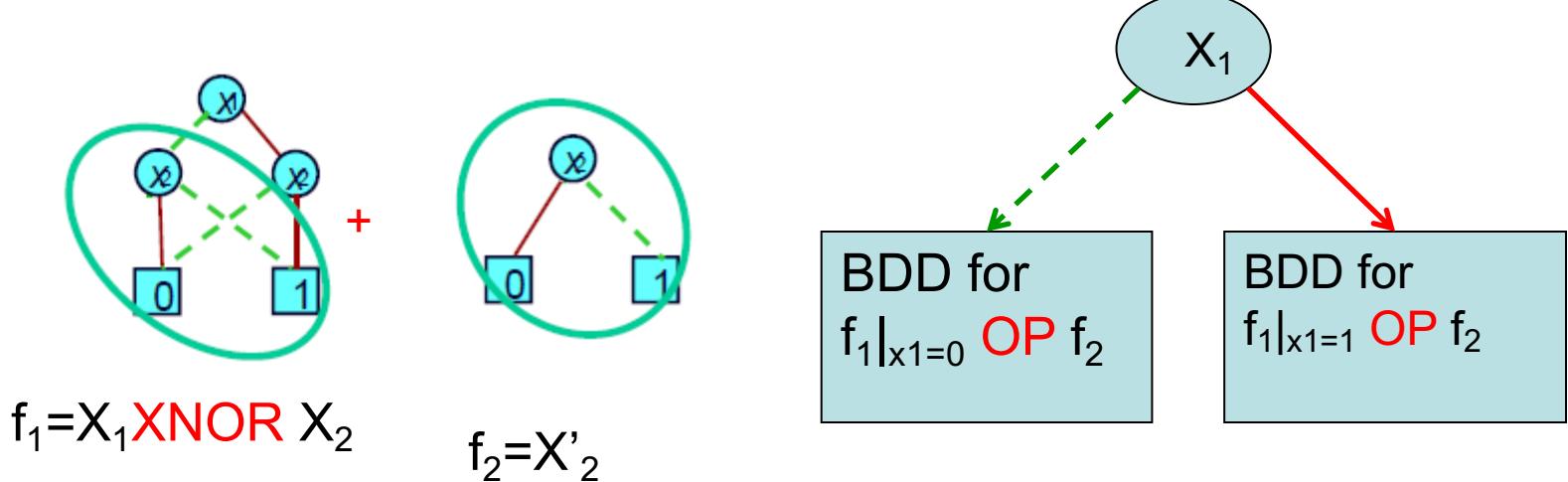
Operations with BDD

- ❖ Else suppose $x_1 < x_2 = x$, in variable order

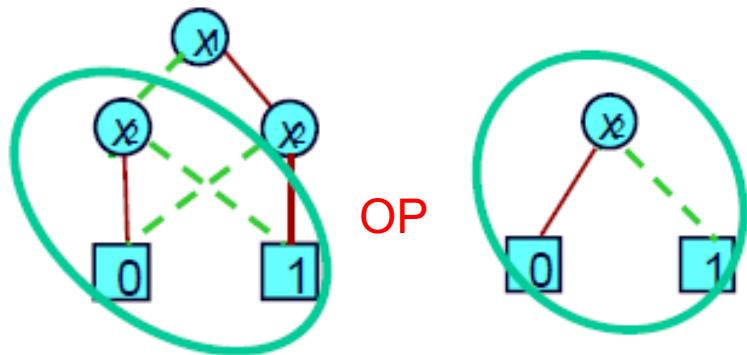
$$f_1 \text{ OP } f_2 = x'_1 (f_1|_{x_1=0} \text{ OP } f_2) + x_1 (f_1|_{x_1=1} \text{ OP } f_2)$$



Operations with BDD: Example

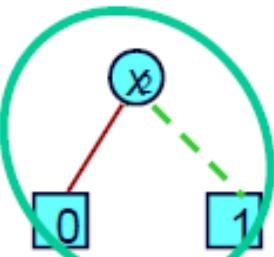


Operations with BDD: Example



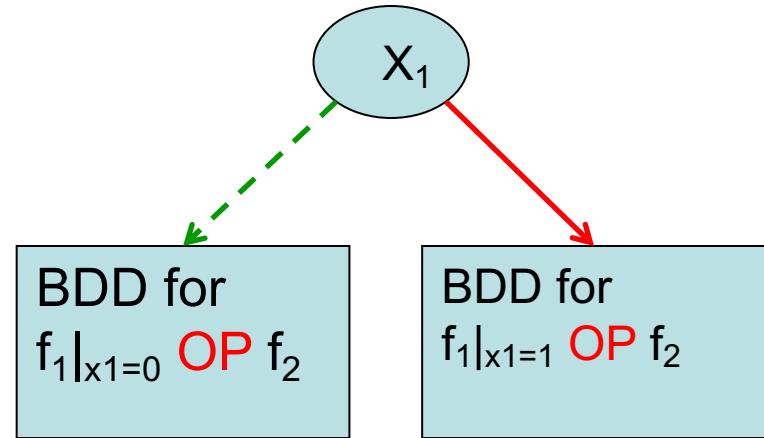
$$f_1 = X_1 \text{XNOR} X_2$$

OP



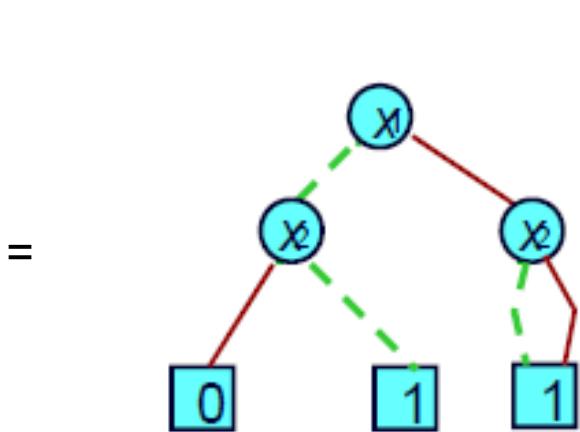
$$f_2 = X'_2$$

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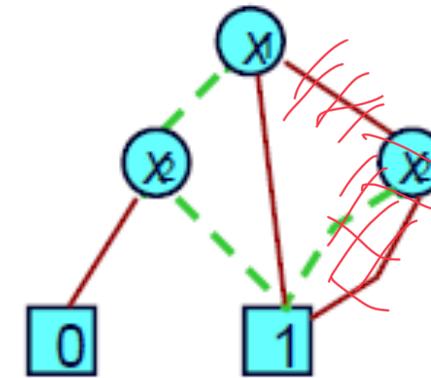
$$\text{BDD for } f_1|_{x_1=0} \text{ OP } f_2$$

$$\text{BDD for } f_1|_{x_1=1} \text{ OP } f_2$$

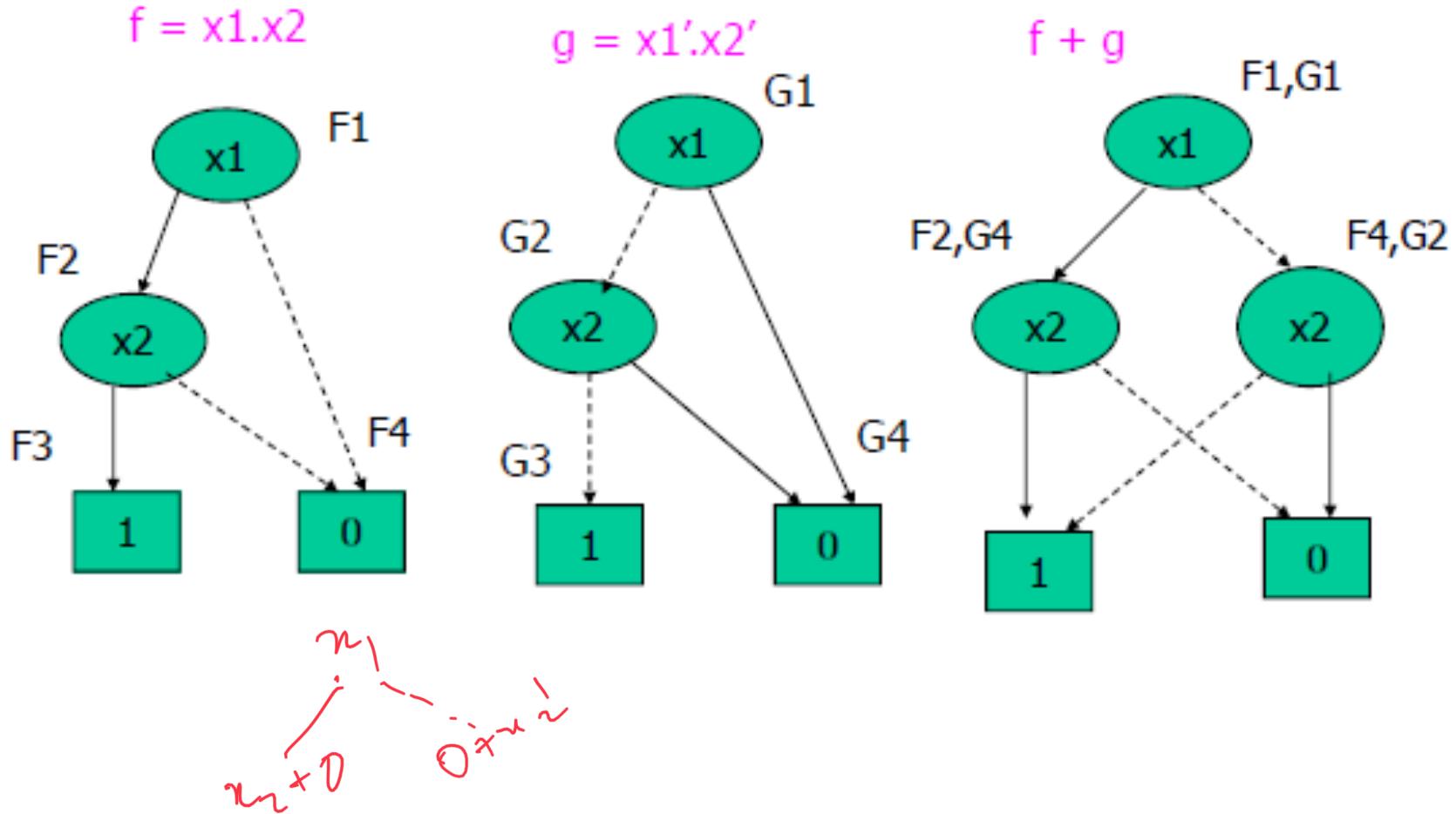


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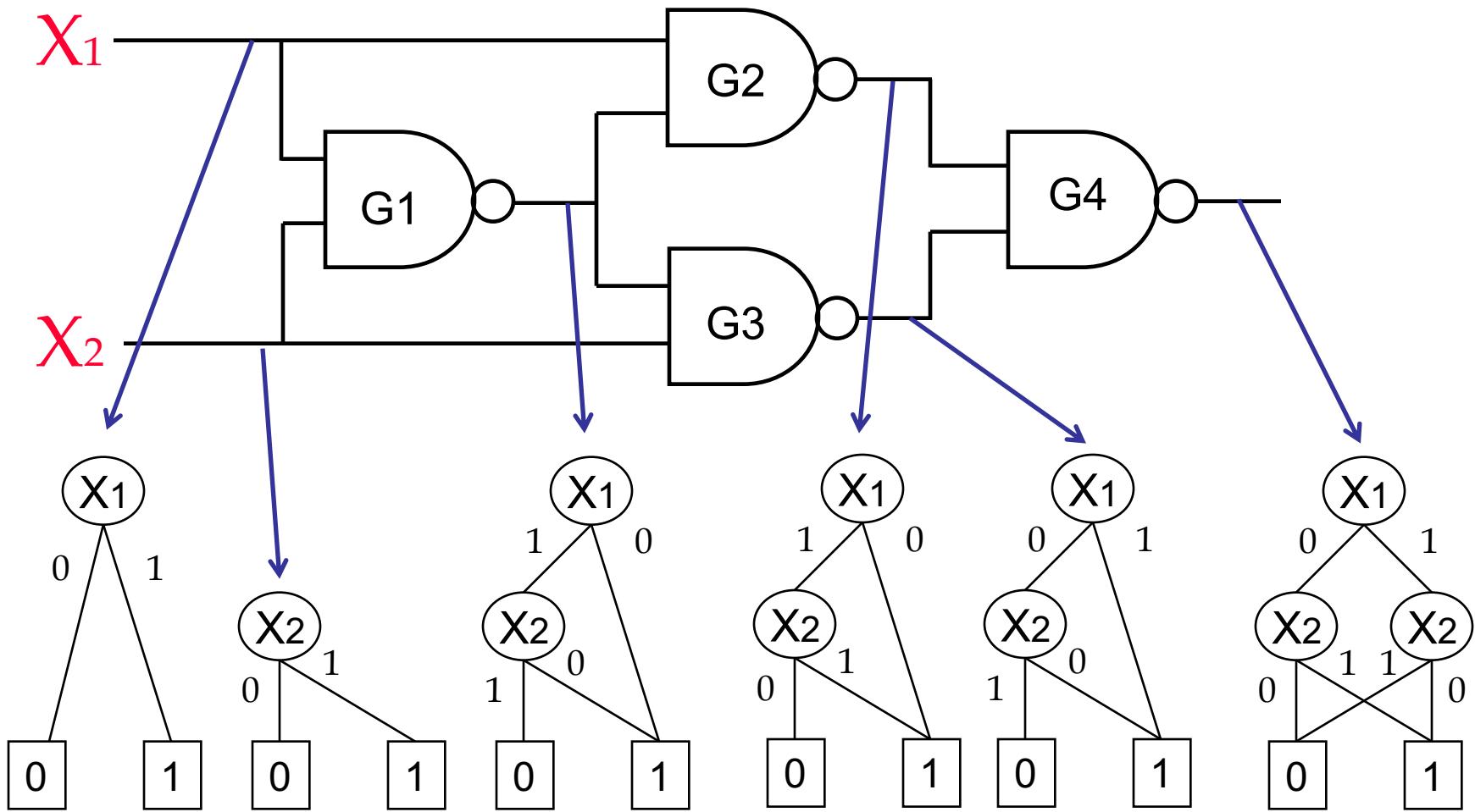
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Operations with BDD: Example



From Circuits to BDD



Thank You

