## Problem Set 1

- 1 Consider the formula  $\varphi = \forall x \exists y R(x,y) \land \exists y \forall x \neg R(x,y)$ . Show that  $\varphi$  is satisfiable over a structure whose universe is infinite and countable.
- 2. Let  $\tau$  be a signature consisting of a binary relation P and a unary relation F. Let  $\mathcal{F}$  be a structure consisting of a universe of people, P(x,y) is interpreted on  $\mathcal{F}$  as "x is a parent of y" and F(x) is interpreted as "x is female". Given the  $\tau$ -structure  $\mathcal{F}$ ,
  - Define a formula  $\varphi_B(x,y)$  which says x is a brother of y
  - Define a formula  $\varphi_A(x,y)$  which says x is an aunt of y
  - Define a formula  $\varphi_C(x,y)$  which says x and y are cousins
  - Define a formula  $\varphi_O(x)$  which says x is an only child
  - Give an example of a family relationship that cannot be defined by a formula
- 3. Consider the signature  $\tau$  that has the binary functions  $+, \times$ . Let  $\mathcal N$  be the structure over  $\tau$  having as universe the set  $\mathbb N$  of natural numbers and which interprets  $+, \times$  in the usual way. Construct FO formulae Zero(x), One(x), Even(x), Odd(x) and Prime(x) using  $\tau$  such that
  - For any  $a \in \mathbb{N}$ ,  $\mathcal{N} \models Zero(a)$  iff a is zero.
  - For any  $a \in \mathbb{N}$ ,  $\mathcal{N} \models One(a)$  iff a is one.
  - For any  $a \in \mathbb{N}$ ,  $\mathcal{N} \models Even(a)$  iff a is even.
  - For any  $a \in \mathbb{N}$ ,  $\mathcal{N} \models Odd(a)$  iff a is odd.
  - For any  $a \in \mathbb{N}$ ,  $\mathcal{N} \models Prime(a)$  iff a is prime.
  - Goldbach's conjecture says that every even integer greater than 2 is the sum of two primes. Whether or not this is true is an open question in number theory. State Goldbach's conjecture as a FO-sentence over  $\tau$ .
- 4. A group is a structure (G,+,0) where G is a set,  $0 \in G$  is a special element called the identity and  $+: G \times G \to G$  is a binary operation such that
  - (A1) The operation + is associative
  - (A2) The constant 0 is a right-identity for the operation +
  - (A3) Every element in G has a right inverse: for each  $x \in G$ , we can find  $y \in G$  such that x + y = 0
  - (A4) For any three elements  $x, y, z \in G$ , if x + z = y + z, then x = yUsing a signature  $\tau = (c, op)$  where c is a constant and op is a binary function symbol write (A1)-(A4) in FO.

5. Let  $\tau$  be a signature consisting of the binary function symbol + and a constant 0. We denote by x + y the function +(x, y). Consider the following sentences:

$$\varphi_1 : \forall x \forall y \forall z [(x + (y + z)) = ((x + y) + z)]$$
$$\varphi_2 : \forall x [(x + 0) = x \land (0 + x) = x]$$
$$\varphi_3 : \forall x [\exists y (x + y = 0) \land \exists z (z + x) = 0]$$

Let  $\psi$  be the conjunction of the three sentences.

- Show that  $\psi$  is satisfiable by exhibiting a  $\tau$ -structure.
- Show that  $\psi$  is not valid.
- Let  $\alpha$  be the sentence  $\forall x \forall y ((x+y) = (y+x))$ . Does  $\alpha$  follow as a consequence of  $\psi$ ? That is, is it the case that  $\psi \to \alpha$ ?
- Show that  $\psi$  is not equivalent to any of  $\varphi_1 \wedge \varphi_2$ ,  $\varphi_2 \wedge \varphi_3$  and  $\varphi_1 \wedge \varphi_3$ .
- 6. Explain the difference between the first order prefixes  $\exists x \forall y \exists z$  and  $\forall x \exists y \forall z$ .
- 7. Show that the sentences  $\forall x \exists y \forall z (E(x,y) \land E(x,z) \land E(y,z))$  and  $\exists x \forall y \exists z (E(x,y) \land E(x,z) \land E(y,z))$  are not equivalent by exhibiting a graph which satisfies one but not both of the sentences.
- 8 For each  $n \in \mathbb{N}$ ,  $\exists^{\geq n}$  denotes a counting quantifier. Intuitively,  $\exists^{\geq n}$  means that "there exist at least n such that". FO with counting quantifiers is the logic obtained by adding these quantifiers (for each  $n \in \mathbb{N}$ ) to the fixed symbols of FO. The syntax and semantics are as follows:

Syntax : For any formula  $\varphi$  of FO with counting quantifiers,  $\exists^{\geq n} x \varphi$  is also a formula.

Semantics:  $\mathcal{A} \models \exists^{\geq n} x \varphi$  iff  $\mathcal{A} \models \varphi(a_i)$  for each of n distinct elements  $a_1, a_2, \ldots, a_n$  from the universe  $u(\mathcal{A})$ .

- Using counting quantifiers, define a sentence  $\varphi_{45}$  such that  $\mathcal{A} \models \varphi_{45}$  iff  $|u(\mathcal{A})| = 45$ .
- Define a FO sentence  $\varphi$  (not using counting quantifiers) that is  $\wedge$  equivalent to the sentence  $\exists^{\geq n} x (x = x)$ .
- 9. Write an FO formula that will evaluate to true only over a structure that has at least n elements and at most m elements.