

Satisfiability Problem

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CS-226: Digital Logic Design



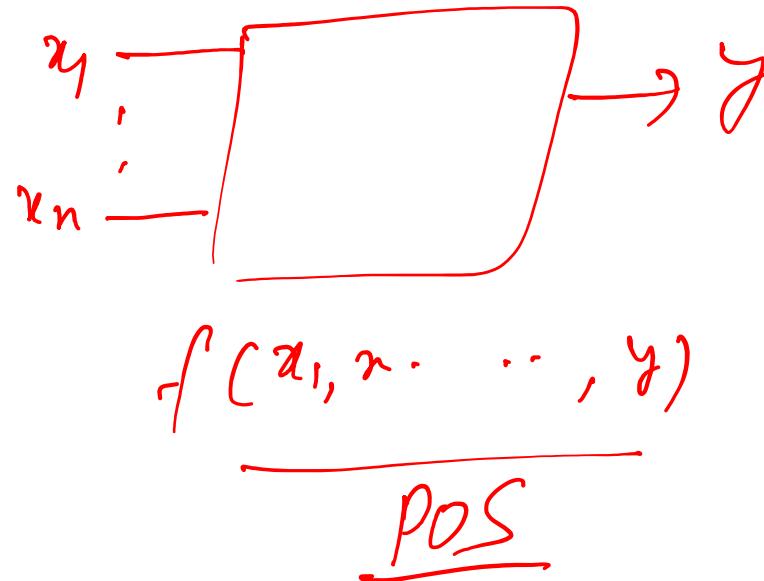
Lecture 11-A: 09 February 2021

CADSL

Function Representation

{
Truth Table
SOM
DOM
ROBDD
RM

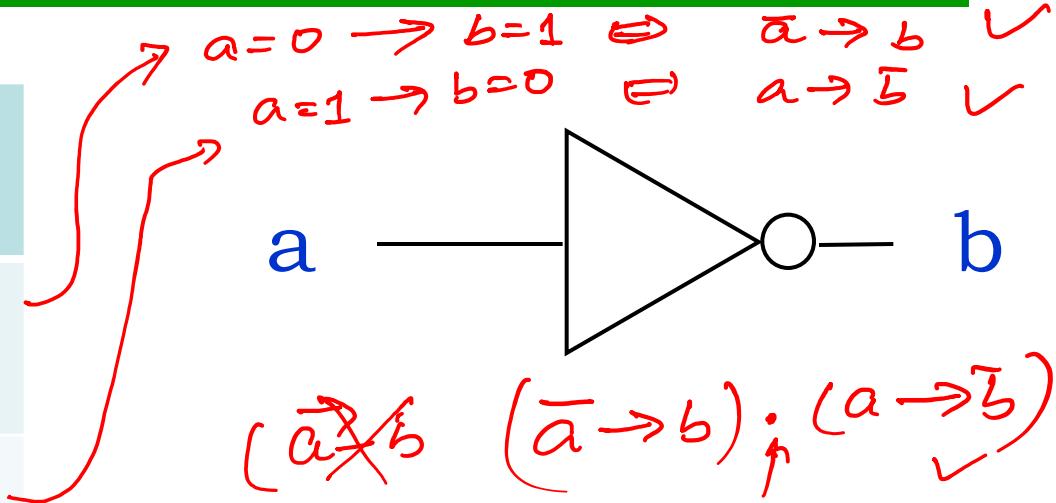
AIG.
SOP
POS



Function Representation

Truth Table

a	b
0	1
1	0



$$f(a,b) = \underbrace{(a+b) \cdot (\bar{a}+\bar{b})}_{\text{OR}} \quad \checkmark$$

$$f(a,b) = (a+b)(\bar{a}+\bar{b})$$

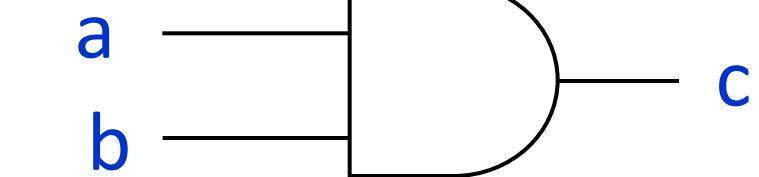


Function Representation

TRUTH TABLE

a	b	c
0	0	0
0	1	0
1	0	0
1	1	1

$$\begin{aligned} & a=0 \& b=0 \rightarrow c=0 \Leftrightarrow \bar{a} \cdot \bar{b} \rightarrow \bar{c} \\ & a=0 \& b=1 \rightarrow c=0 \Leftrightarrow \bar{a} \cdot b \rightarrow \bar{c} \end{aligned}$$



$$\begin{aligned} & a=1 \& b=0 \rightarrow c=0 \Leftrightarrow a \cdot \bar{b} \rightarrow \bar{c} \\ & a=1 \& b=1 \rightarrow c=1 \Leftrightarrow a \cdot b \rightarrow c \end{aligned}$$

$$\begin{aligned} & (\bar{a} \cdot \bar{b} \rightarrow \bar{c}) \cdot (\bar{a} \cdot b \rightarrow \bar{c}) \cdot (a \cdot \bar{b} \rightarrow \bar{c}) \cdot (a \cdot b \rightarrow c) \\ & = (\bar{a} \cdot \bar{b} + \bar{c}) \cdot (\bar{a} \cdot b + \bar{c}) \cdot (\bar{a} \cdot \bar{b} + c) \cdot (\bar{a} \cdot b + c) \\ & = (a + b + \bar{c}) \cdot (a + \bar{b} + \bar{c}) \cdot (\bar{a} + b + \bar{c}) \cdot (\bar{a} + \bar{b} + c) \end{aligned}$$

$$f(a,b,c) = \underbrace{(a+c)}_{\text{POS}} \underbrace{(b+\bar{c})}_{\text{POS}} \underbrace{(\bar{a}+\bar{b}+c)}_{\text{POS}}$$

POS



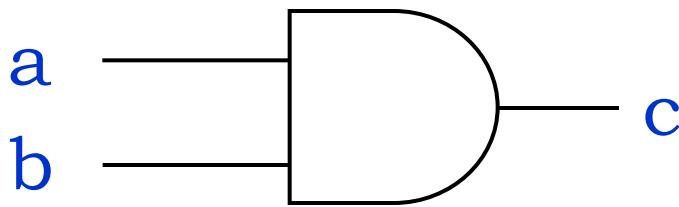
Function Representation

a	b	c	
0	0	0	$\bar{a} \cdot \bar{b} \rightarrow \bar{c}$ ✓ $\bar{a} \cdot b \rightarrow c$ ✓
0	1	1	$a \cdot \bar{b} \rightarrow c$ ✓ $a \cdot b \rightarrow c$ ✓
1	0	1	$(\bar{a} \cdot \bar{b} \rightarrow \bar{c}) \cdot (\bar{a} \cdot b \rightarrow c) (a \cdot \bar{b} \rightarrow c) (a \cdot b \rightarrow c)$ $= (\bar{a} \cdot \bar{b} + \bar{c}) (\bar{a} \cdot b + c) (\overline{a \cdot b}) (\overline{a \cdot b} + c)$ $= (a+b+\bar{c}) (a+\bar{b}+c) (\bar{a}+\bar{b}+c) (\bar{a}+\bar{b}+c+c)$
1	1	1	$= (a+b+\bar{c}) (a+\bar{b}+c) (\bar{a}+\bar{b}+c) (\bar{a}+\bar{b}+c+c)$ $= (a+b+\bar{c}) (a+\bar{b}+c) (\bar{a}+\bar{b}+c) (\bar{a}+\bar{b}+c+c)$ $= (a+b+\bar{c}) (\bar{b}+c) (\bar{a}+c)$

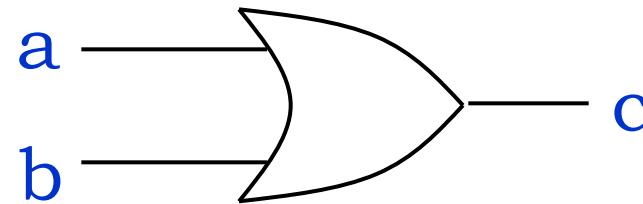
$f(a, b, c)$ $\underline{\text{clause}}$ $\underline{\text{3 clauses}}$ SOP



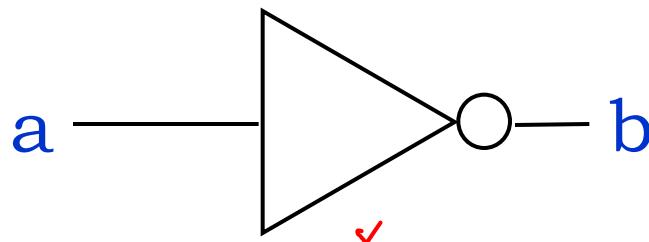
SAT Formulas for Simple Gates



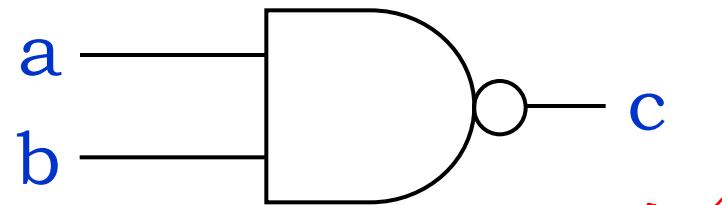
✓ $(\bar{c} + a)(\bar{c} + b)(c + \bar{a} + \bar{b})$



$(c + \bar{a})(c + \bar{b})(\bar{c} + a + b)$ ✓



✓ $(a + b)(\bar{a} + \bar{b})$

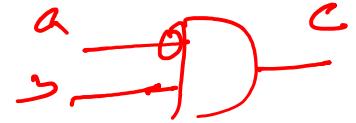


$(c + a)(c + b)(\bar{c} + \bar{a} + \bar{b})$

Function Representation

TRUTH TABLE

a	b	c
0	0	0
0	1	1
1	0	0
1	1	0

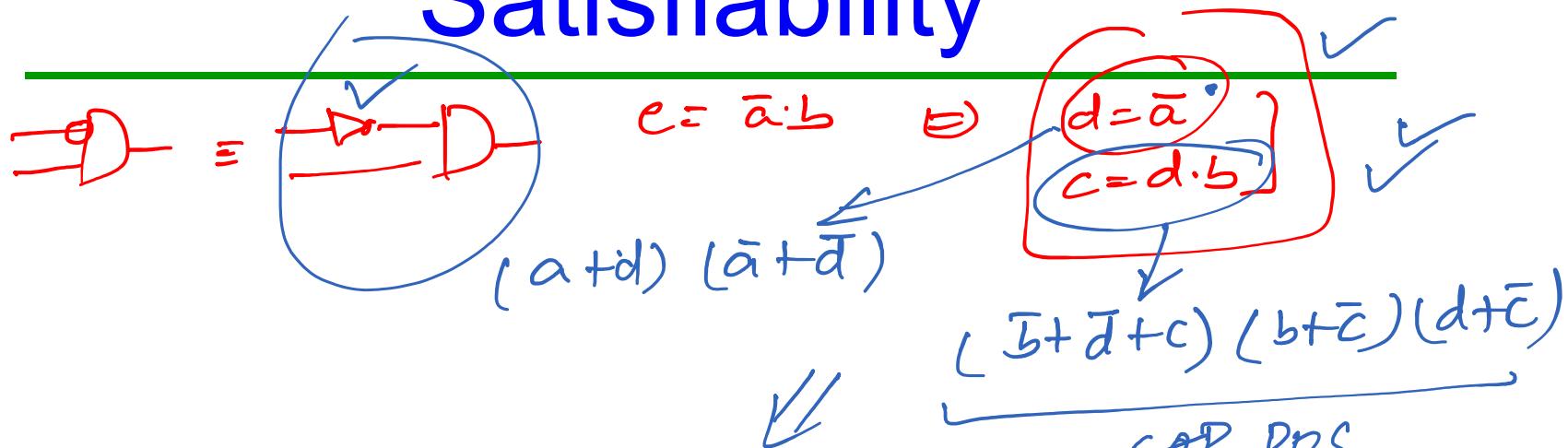


✓

$$\begin{aligned} & \bar{a} \cdot \bar{b} \rightarrow \bar{c} \\ & \bar{a} \cdot b \rightarrow c \\ & a \cdot \bar{b} \rightarrow \bar{c} \\ & a \cdot b \rightarrow \bar{c} \\ & = (\bar{a} \cdot \bar{b} \rightarrow \bar{c}) \cdot (\bar{a} \cdot b \rightarrow c) (a \cdot \bar{b} \rightarrow \bar{c}) \\ & = (\bar{a} \cdot \bar{b} + \bar{c}) (\bar{a} \cdot b + c) (\bar{a} \cdot \bar{b} + \bar{c}) (\bar{a} \cdot b + \bar{c}) \\ & = (a + b + \bar{c}) (a + \bar{b} + c) (\bar{a} + \bar{b} + \bar{c}) (\bar{a} + \bar{b} + c) \\ & = (a + b + \bar{c}) (a + \bar{b} + c) (\bar{a} + c) \\ f(a, b, c) &= \underline{(a + b + \bar{c}) (\bar{a} + \bar{b} + c) (\bar{a} + c)} \end{aligned}$$



Satisfiability



$$f(a,b,c,d) = (a+d)(\bar{a}+\bar{d}) (\bar{b}+\bar{d}+c) (\bar{b}+\bar{c})(d+\bar{c})$$

PDS 5 clauses

what can be done using this representation

$$f: (\underline{a+\bar{c}}) (\underline{b+\bar{c}}) (\overline{\bar{a}+\bar{b}+c}) \leftarrow \text{AND oper.}$$

$$f = \begin{cases} a=0, b=0, c=0 & \text{can satisfy or not (SAT)} \\ a=0, b=0, c=1 & \rightarrow \text{Not satisfied.} \end{cases}$$



Satisfiability

$$f(a, b, c) = \frac{(a + \bar{b})}{\checkmark} (b + \bar{c}) (\bar{a} + \bar{b} + c)$$

if $a = 0, b = 0$ $c = ?$

$$\cdot \bar{c} \cdot \bar{c} \cdot 1 = \bar{c} = 1$$

Hence = " $\boxed{c=0}$ "

For what value of $a \underline{\text{or}} b$ $\underline{\text{or}} c = 1$

$\underline{\text{c=1}}$

$a=?$ $b=?$

$$\therefore a \cdot b \cdot 1 = a \cdot b$$

$$\underline{a \cdot b = 1} \Rightarrow \begin{cases} a = 1 \\ b = 1 \end{cases} \checkmark$$

Satisfiability



SAT Problem Definition

Given a CNF formula, f :

- A set of variables, V
- Conjunction of clauses
- Each clause: disjunction of literals over \underline{V}
AN? \leftarrow
OR.

$$\begin{array}{l} (a, b, c) \\ \underline{(C_1, C_2, C_3)} \end{array}$$



$$\frac{c_1 \cdot c_2 \cdot c_3 = 1}{\begin{cases} c_1 = 1 \\ c_2 = 1 \\ c_3 = 1 \end{cases}}$$

Does there exist an assignment of Boolean values to the variables, V which sets at least one literal in each clause to '1' ?

Example :

$$(a \checkmark + b \# + \bar{c}) (\bar{a} + \checkmark c) (a \# + \bar{b} + \checkmark c)$$

$\underbrace{\phantom{a \# + b \# + \bar{c}}}_{C_1}$ $\underbrace{\phantom{\bar{a} + c}}_{C_2}$ $\underbrace{\phantom{a \# + \bar{b} + c}}_{C_3}$

$$a = b = c = \checkmark$$



SAT Problem

$f(a, b, c \dots)$

POS



Thank You

