

*Circles within circles...*

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# Operating System Structures

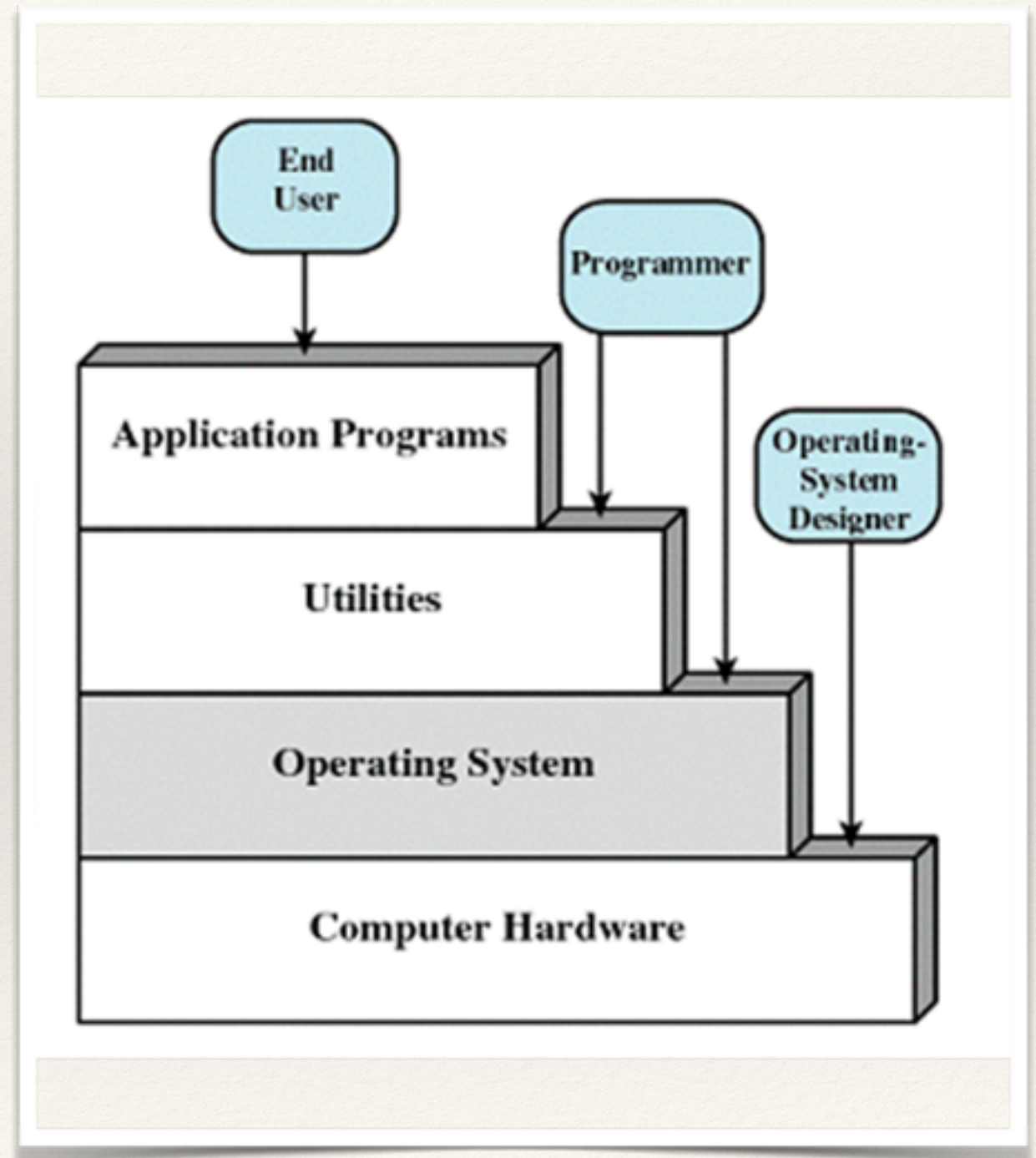
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- Layered Structure
- Multi-kernels for Multi-core processors
- Microkernel Structure
- Virtual Machine Structure



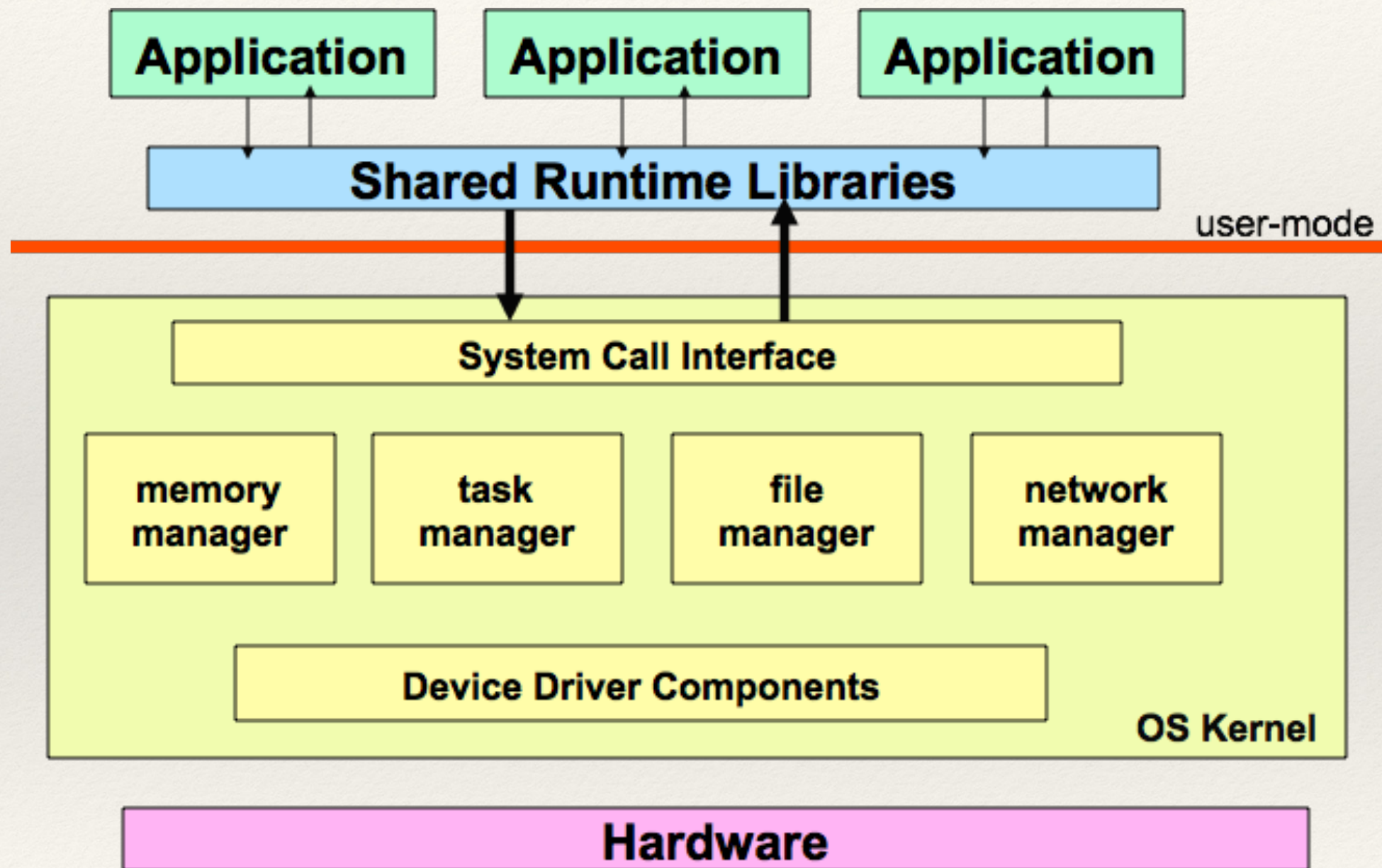
# Layered System Structure

- ❖ System structured as a series of levels - each level performs a related subset of functions.
- ❖ Each level relies on the next lower level for more primitive functions.
- ❖ Layer Concepts
  - ❖ Kernel
  - ❖ Basic services and I/O drivers
  - ❖ Resource management





# A Modern OS Design





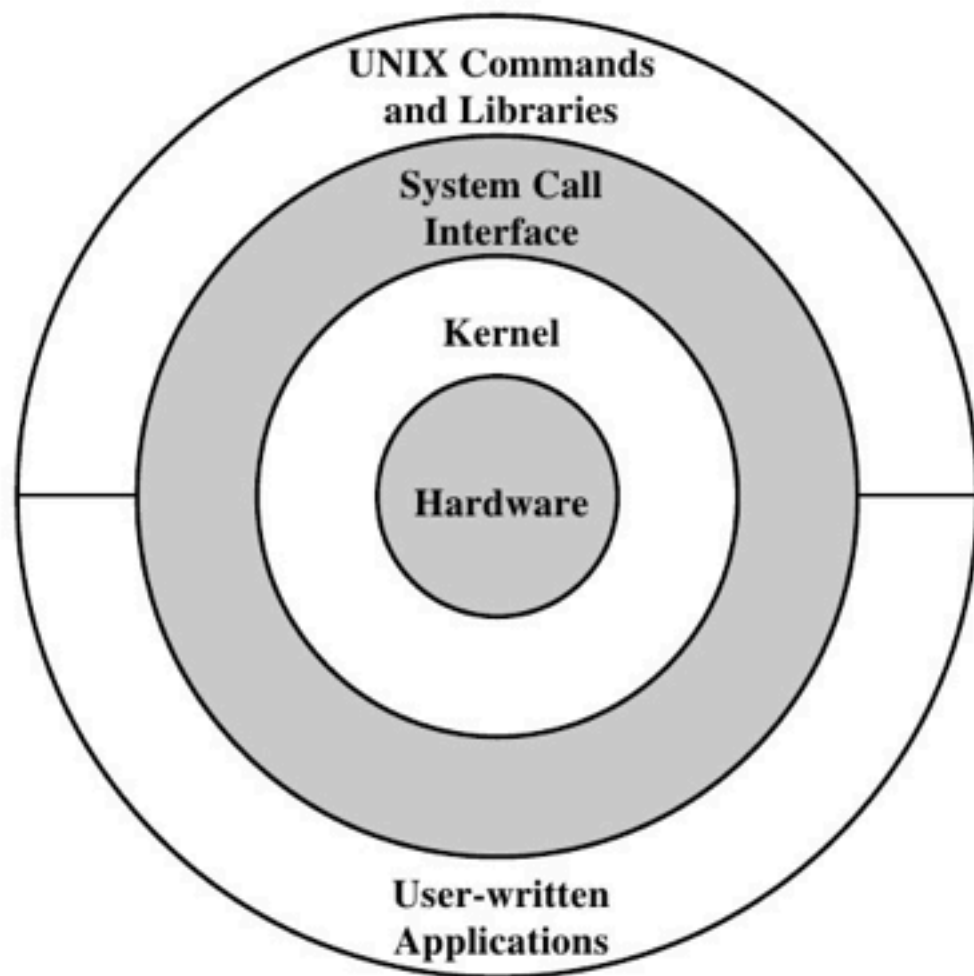
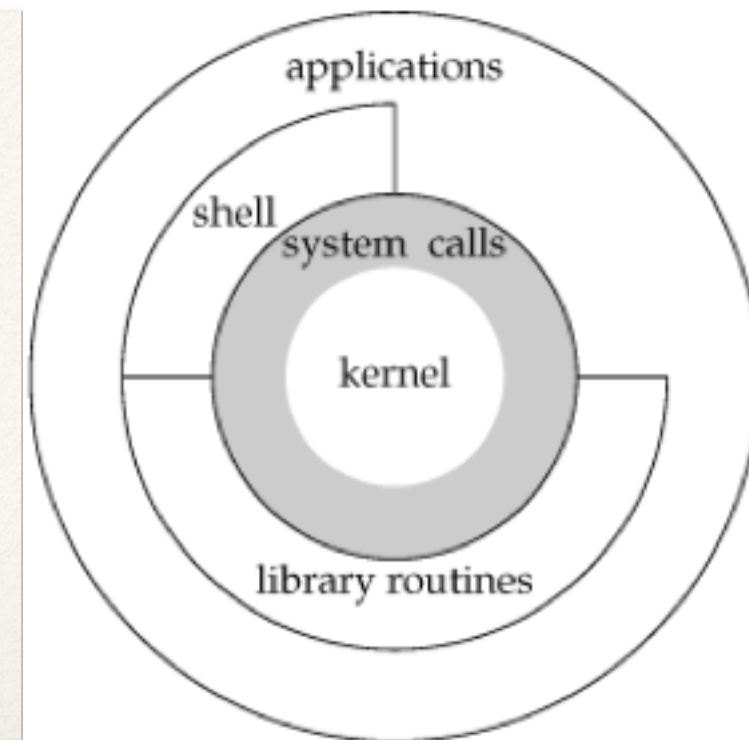
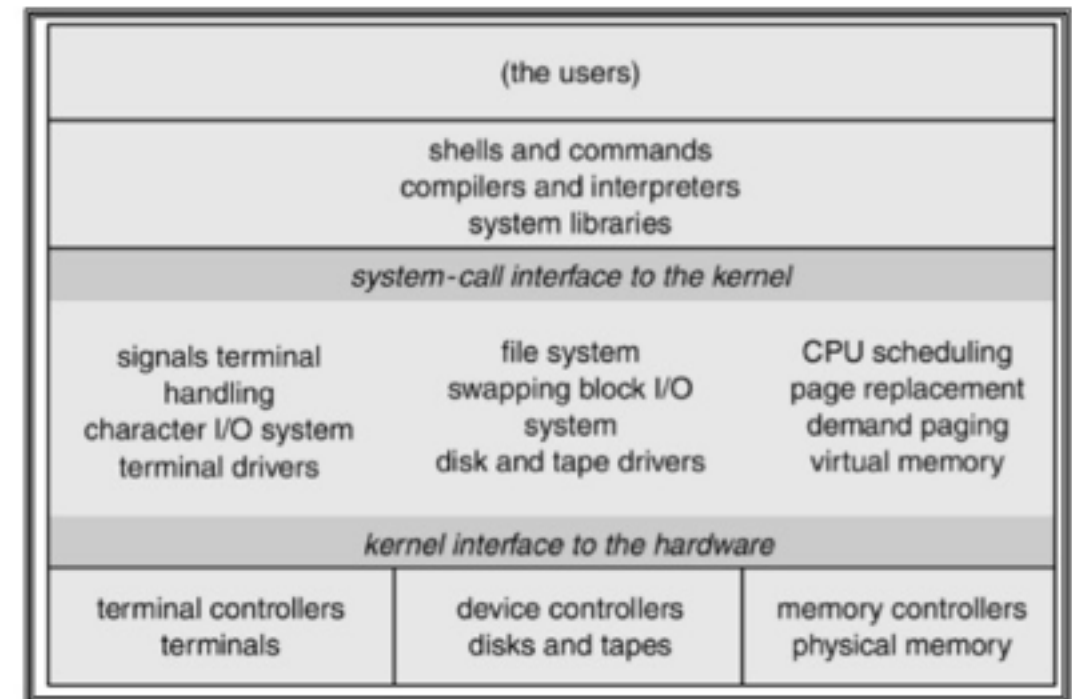


Figure 2.15 General UNIX Architecture



## UNIX System Structure





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# Layered Approach

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- ❖ The operating system is divided into a number of layers (levels), each built on top of lower layers. The bottom layer (layer 0), is the hardware; the highest (layer N) is the user interface.
- ❖ With modularity, layers are selected such that each uses functions (operations) and services of only lower-level layers.



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# Multi-Processors

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- ❖ Typically each processor does self-scheduling from the pool of available process or threads.
  - ❖ Timer interrupt
  - ❖ Ready queue
- ❖ OS Support
  - ❖ Any thread (including kernel threads) can run on any processor.
  - ❖ **Soft affinity** (close relation) – Try to reschedule a thread on the same processor.
  - ❖ **Hard affinity** – Restrict a thread to certain processors.



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# Multi-Core Processors

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- ❖ In a multi-core environment, the control over which core to run a specific thread or application on is essential.
- ❖ Without this control, the threads / applications may get assigned to the wrong processors or cores and cause unnecessary performance degradation.
- ❖ Generally each processor has its own cache, shared cache, shared memory and I/O



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# Multi-Core OS Design Issues

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- ❖ **Single kernel or Multi-kernel**
  - ❖ Kernel routines must be reentrant to allow multiple threads to execute them
- ❖ **Scheduling**
  - ❖ Must avoid conflicts
  - ❖ May be able to run processes or threads concurrently
- ❖ **Synchronization**
  - ❖ Mutual exclusion, event ordering
- ❖ **Memory Management**
  - ❖ Deal with multi-port memory
  - ❖ Have a unified paging scheme
- ❖ **Reliability and Fault Tolerance**



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# Microkernel Structure

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- ❖ Popularized by use in Mach OS
- ❖ **Monitors** (1960s):
  - ❖ Built as a single large program, any routine can call any other routine
  - ❖ Used in most early systems
- ❖ **Layered OS** (1970s – now):
  - ❖ Based on modular programming
  - ❖ Major changes still had wide-spread effects on other layers
- ❖ **Microkernel** (1980s – now):
  - ❖ Only essential functions in the kernel
  - ❖ File System, device drivers, *etc.*, are now external subsystems / processes
  - ❖ Processes interact through (IPC) messages passed through the kernel



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# Microkernel

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- ❖ Identify and isolate a small operating system core that contains only essential OS functions.
- ❖ Move many services included in the traditional kernel OS to external subsystems
  - ❖ device drivers
  - ❖ file systems
  - ❖ virtual memory manager
  - ❖ windowing system and security services



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# Microkernel Design

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- ❖ **Primitive Memory Management**

- ❖ Kernel handles virtual ~~→~~ physical mapping, rest is a user mode process
  - ❖ VM module can decide what pages to move to/from disk
  - ❖ Module can allocate memory
- ❖ Three microkernel memory operations
  - ❖ *Grant* – Grant pages to someone
  - ❖ *Map* – Map pages in another space
  - ❖ *Flush* – Reclaim pages granted or mapped

- ❖ **Inter-process Communication (IPC)**

- ❖ Based on messages

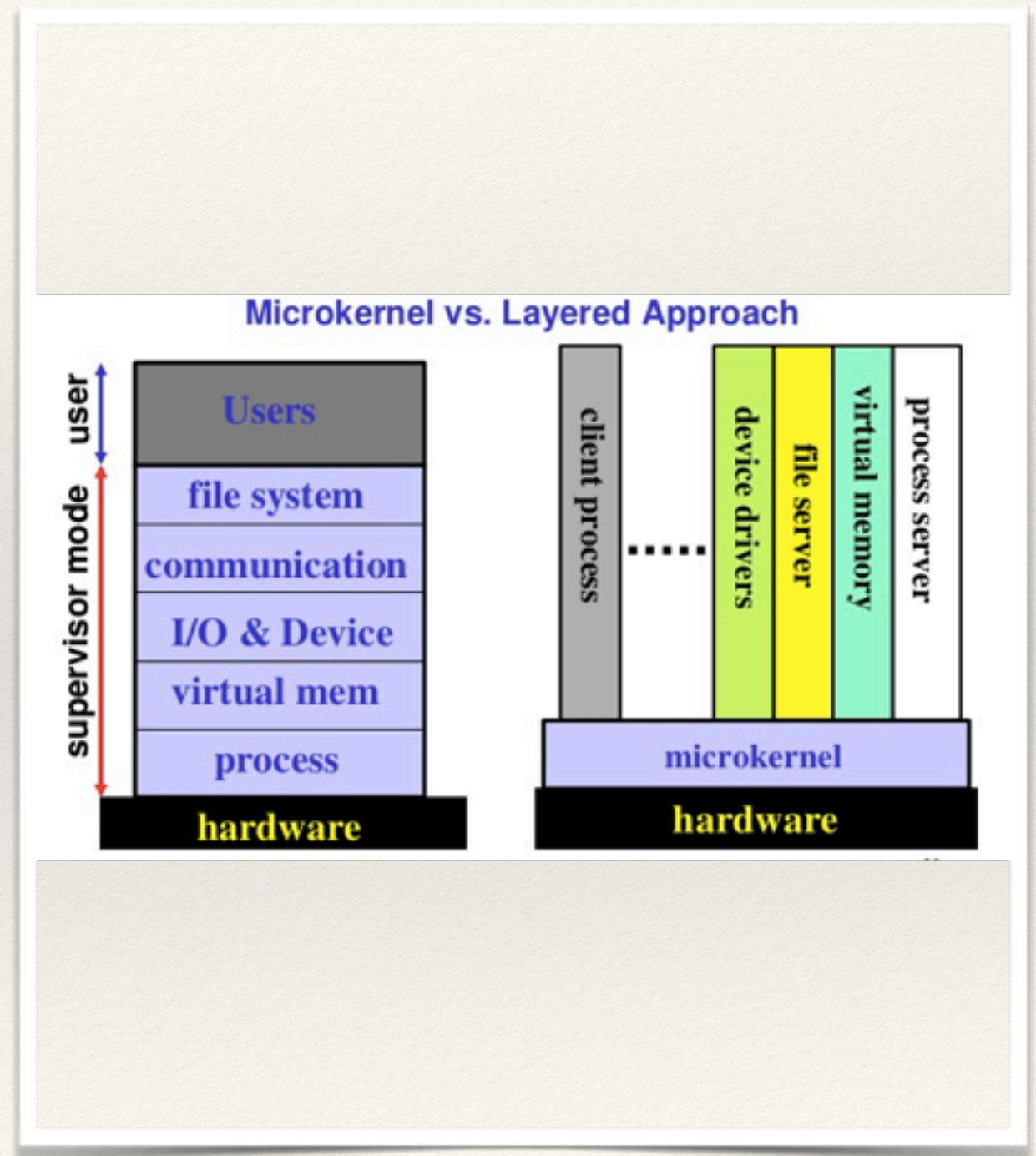
- ❖ **I/O and Interrupts**

- ❖ Handle interrupts as messages



# Microkernel Benefits

- ❖ **Uniform Interface**
  - ❖ Same message for user / system services
- ❖ **Extensibility**
  - ❖ Easy to add new services
  - ❖ Modifications need only change directly affected components
  - ❖ Could have multiple file servers
- ❖ **Flexibility**
  - ❖ Can customize system by omitting services
- ❖ **Portability**
  - ❖ Isolate nearly all processor-specific code in the kernel





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# Microkernel Benefits

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- ❖ **Reliability**

- ❖ Easy to test kernel
- ❖ Fewer system calls to master
- ❖ Less interaction with other components

- ❖ **Distributed System Support**

- ❖ Just as easy to send a message to another machine as this machine
  - ❖ Need system-wide unique Ids
- ❖ Processes do not have to know where a service resides

- ❖ **Object-Orientated OS**

- ❖ Lends OO disciplines to the kernel



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# Kernel Performance

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- ❖ Sending a message generally slower than simple kernel call (soft interruption)
- ❖ Depends on size of the microkernel
  - ❖ First generation systems slower
  - ❖ Then tried to include critical system items into kernel (Mach)
    - ❖ Fewer user / system mode switches
    - ❖ Lose some microkernel benefits
- ❖ Trying approach of very small kernel
  - ❖ Mach L4 kernel (version 2) - 12K code, 7 system calls. Speed seems to match Unix
  - ❖ Mac OS X adopt this concept



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# Virtual Machine Structure

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- ❖ A virtual machine takes the layered approach to its logical conclusion. It treats hardware and the operating system kernel as though they were all hardware.
- ❖ A virtual machine provides an interface identical to the underlying bare hardware.
- ❖ The operating system creates the illusion of multiple processes, each executing on its own processor with its own (virtual) memory.



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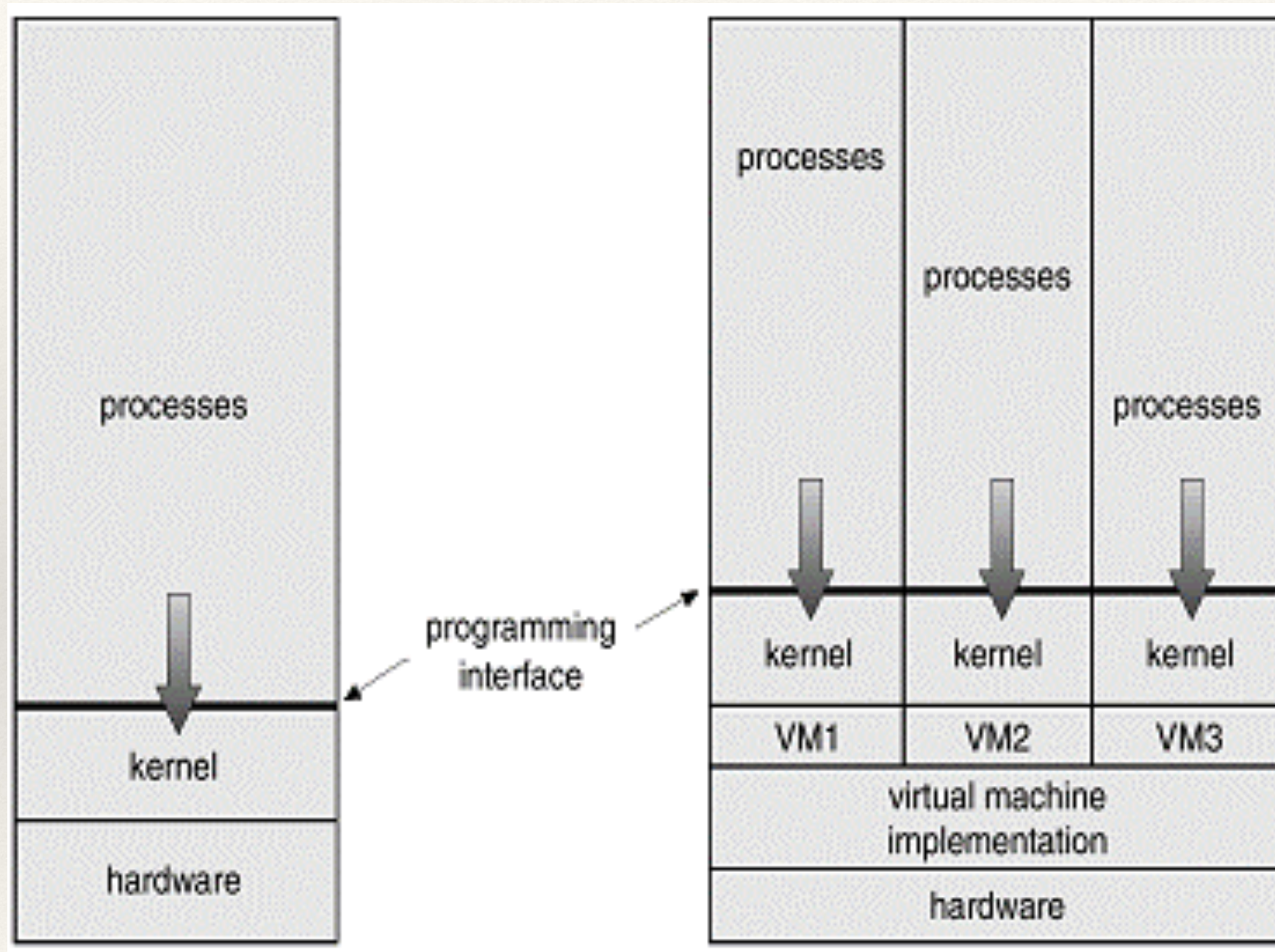
# Virtual Machines

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- ❖ The resources of the physical computer are shared to create the virtual machines.
- ❖ CPU scheduling can create the appearance that users have their own processor.
- ❖ Spooling and a file system can provide virtual hard drivers and virtual printers.
- ❖ A normal user display device serves as the virtual machine operator's console.



# System Models





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# Advantages/Disadvantages of VM's

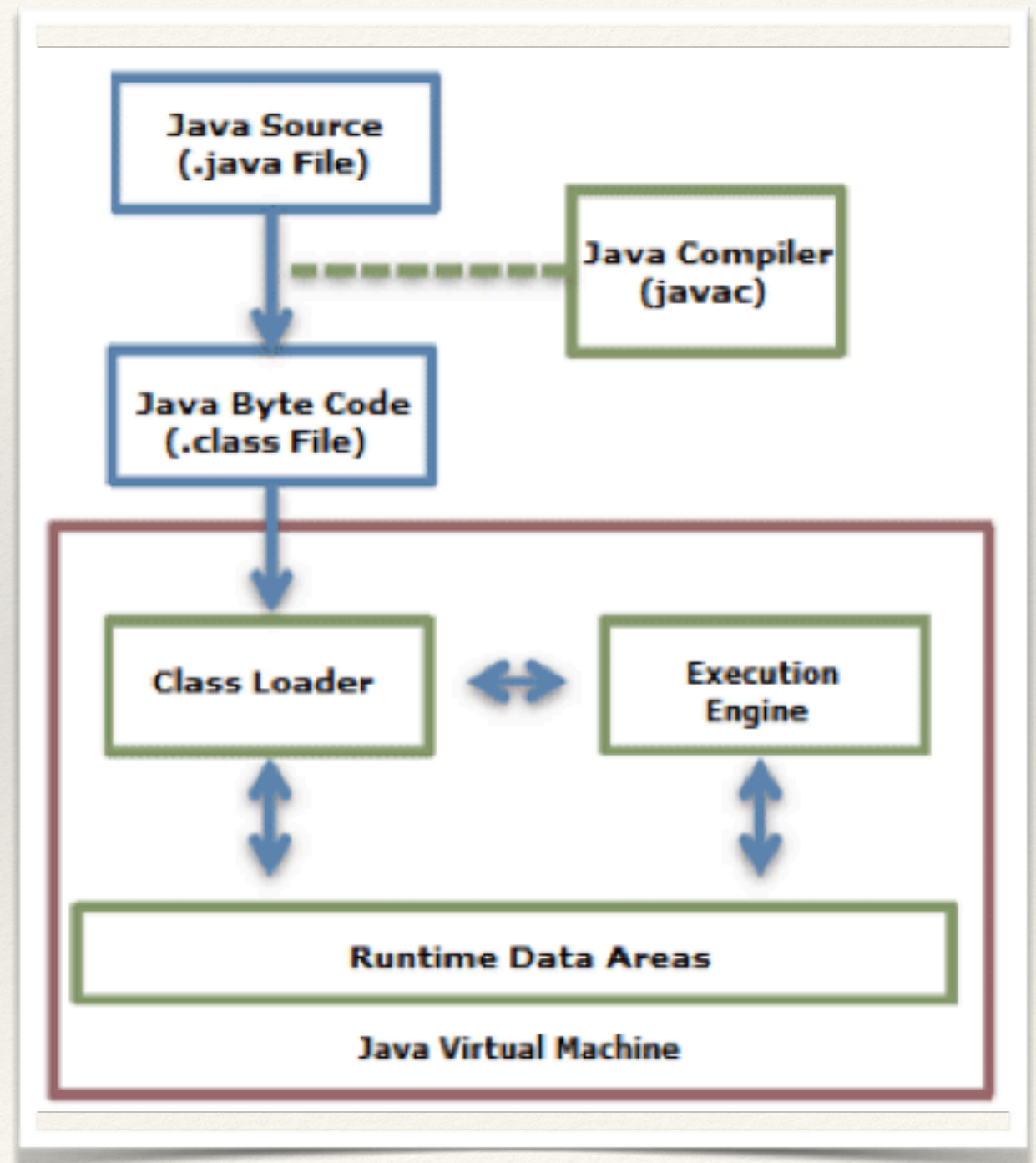
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- ❖ The virtual-machine concept provides complete protection of system resources since each virtual machine is isolated from all other virtual machines. This isolation, however, permits no direct sharing of resources.
- ❖ A virtual-machine system is a perfect vehicle for operating systems research and development.
- ❖ System development is done on the virtual machine, instead of on a physical machine, and so does not disrupt normal system operation.
- ❖ The virtual machine concept is difficult to implement due to the effort required to provide an exact duplicate to the underlying machine.



# Java Virtual Machine

- ❖ Compiled Java programs are platform-neutral bytecodes executed by a Java Virtual Machine (JVM).
- ❖ JVM consists of
  - ❖ class loader
  - ❖ class verifier
  - ❖ runtime interpreter
- ❖ Just-In-Time (JIT) compilers increase performance.





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# System Design Goals

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- ❖ **User goals** – operating system should be convenient to use, easy to learn, reliable, safe, and fast.
- ❖ **System goals** – operating system should be easy to design, implement, and maintain, as well as flexible, reliable, error-free, and efficient.



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# Mechanisms and Policies

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- ❖ **Mechanisms** determine how to do something, **policies** decide what will be done.
- ❖ The separation of policy from mechanism is a very important principle, it allows maximum *flexibility* if policy decisions are to be changed later.



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# System Implementation

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- ❖ Traditionally written in assembly language, operating systems can now be written in higher-level languages.
- ❖ Code written in a high-level language:
  - ❖ can be written faster
  - ❖ is more compact
  - ❖ is easier to understand and debug
- ❖ An operating system is far easier to port (move to some other hardware) if it is written in a high-level language.