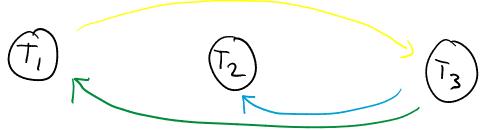
Part 2: Conflict Serializability (20 points)

Consider the following three transactions and schedule (time goes from top to bottom). Is this schedule conflict-serializable? Show why or why not.

| T1 | T2 | Т3 |
|--------|--------|--------|
| R(A) | | |
| W(A) | | |
| | | Ρ(Λ) |
| | | R(A) |
| | | W(A) |
| | R(A) | |
| R(B) | | |
| | | |
| | | R(B) |
| W(B) | | |
| | | W(B) |
| | R(B) | |
| | commit | |
| commit | | |
| | | |
| | | commit |



Ti3 T3 crease a loop, so it is not a conflict serializable schedule