PUMP II Problem Set 1

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lemma. 1.0 $\binom{n}{k} = 0$

Proof.

By Pascal's inequality:

$$\binom{n}{k-1} = \binom{n+1}{k} - \binom{n}{k}$$
$$\binom{n}{-1} = \binom{n+1}{0} - \binom{n}{0}$$
$$= \frac{(n+1)!}{0!(n+1-0)!} - \frac{n!}{0!(n-0)!}$$
$$= 1-1$$
$$= 0$$

1. a)

 $\mathcal{P}(A \cup B) = \{A \cup B, A, B, \{1\}, \{2\}, \{cat\}, \{dog\}, \{1, dog\}, \{1, cat\}, \{2, dog\}, \{2, cat\}, \{1, 2, cat\}, \{1, 2, dog\}, \{1, cat, dog\}, \{2, cat, dog\}, \emptyset\}$

 $\mathcal{P}(\mathcal{P}(B)) = \{\mathcal{P}(B), \{B\}, \{cat\}, \{dog\}, \{\emptyset\}, \{B, \{cat\}\}, \{B, \{dog\}\}, \{B, \emptyset\}, \{B\{cat\}, \{dog\}\}, \{B, \{cat\}, \{\emptyset\}\}, \{\{cat\}, \emptyset\}, \{\{cat\}, \{dog\}\}, \{\{cat\}, \{dog\}, \emptyset\}, \{\{cat\}, \{dog\}, \emptyset\}, \{\{cat\}, \{dog\}, \emptyset\}, \emptyset\}$

b)

Proof.

The total number of ways that n objects in a set can be arranged into $s = \{1, 2, 3, ..., n\}$ sized subsets is equal to $\sum_{k=0}^{n} \binom{n}{k}$. Thus it is sufficient to prove that $\sum_{k=0}^{n} \binom{n}{k} = 2^n$. We prove this statement through induction on n.

Base Case: n = 0

$$\sum_{k=0}^{0} {0 \choose k} = \frac{0!}{0!(0-0)!} = 1 = 2^{0}$$

Thus the base case is true. Inductive Step: j = n + 1

$$\sum_{k=0}^{n+1} \binom{n+1}{k} = \sum_{k=0}^{n+1} \binom{n}{k} + \binom{n}{k-1}, \text{ via Pascal's inequality}$$

$$= \sum_{k=0}^{n+1} \binom{n}{k} + \sum_{k=0}^{n} \binom{n}{k-1}$$

$$= \sum_{k=0}^{n} \binom{n}{k} + \binom{n}{n+1} + \sum_{k=0}^{n} \binom{n}{k-1}, \text{ via reindexing}$$

$$= 2^n + 2^n + 0 + 0, \text{ by lemma } 1.1$$

$$= 2^{n+1}$$

c)

Proof.

We will induct on n.

$$|\mathcal{P}(S)| = 2^n$$

Base case: $n = 0, S = \emptyset$

$$|\mathcal{P}(S)| = 1 = 2^0$$

Thus the base case is true.

Inductive Step: j=n+1 Let A be a set s.t. $A\cap S=\emptyset, |A|=1$ We can construct some set V s.t. $V=\{x: \forall y\in \mathcal{P}(S), x=y\cup A\}$ V and $\mathcal{P}(S)$ are equal in size but disjoint.

$$\mathcal{P}(A \cup S) = \{T : T \in \mathcal{P}(S) \text{ or } T \in V\}$$
$$|V| = |\mathcal{P}(S)| = 2^n$$

Thus $|\mathcal{P}(A \cup S)| = |V| + |\mathcal{P}(S)| = 2^n + 2^n = 2^{n+1}$

lemma. 2.0 If x^2 is a multiple of 7, then x is a multiple of 7.

Proof.

Take the contrapositive.

$$x = 7n + a, n \in \mathbb{Z}, a \in \{1, 2, 3, 4, 5, 6\}$$
$$x^{2} = 49k^{2} + 14ak + a^{2}$$
$$= 7(7k^{2} + 2ak) + a^{2}$$

 a^2 can never be a multiple of 7, and thus x^2 is not a multiple of 7.

lemma. 2.1 If $x \in \mathbb{Q}$ and $y \notin \mathbb{Q}$, then $x + y \notin \mathbb{Q}$.

Proof.

Assume $x + y \in \mathbb{Q}$.

$$\frac{p}{q} + x = \frac{m}{n}, p, q, m, n \in \mathbb{Z}$$
$$x = \frac{m}{n} - \frac{p}{q}$$

Thus x must be in \mathbb{Q} , a contradiction.

lemma. 2.3 0 is the only rational that when multiplied with an irrational results in a rational.

Proof.

Assume there is a rational, $x \in \mathbb{Q}$ s.t. $x \neq 0$ and x multiplied by an irrational is rational.

$$yx = \frac{m}{n}, y \notin \mathbb{Q}, x \in \mathbb{Q}, m, n \in \mathbb{Z}$$
$$y\frac{p}{q} = \frac{m}{n}, p, q \in \mathbb{Z}$$
$$y = \frac{mq}{np}$$

Thus, y would be rational, a contradiction.

2. a)

Proof.

Assume that $x \in \mathbb{Q}$

$$x^{2} = 7$$

$$\frac{p^{2}}{q^{2}} = 7, p \in \mathbb{N}, q \in \mathbb{Z}$$

$$p^{2} = 7q^{2}$$

From lemma 2.0 if p^2 is a multiple of 7, then p is a multiple of 7. Thus we can write $p = 7n, n \in \mathbb{Z}, n < p$.

$$(7n)^2 = 7q^2$$
$$49n^2 = 7q^2$$
$$7n^2 = q^2$$

Thus we can do the same thing we did to q as we did to p, where $k < q, k \in \mathbb{Z}$. Thus we get to the stage where $n^2 = 7k^2$. We then get to the equality that we started, and we can repeat the steps we've taken; however, since n < p then we need to be able to descend infinitely on the naturals, which we cannot, and thus we reach a contradiction.

Proof.

Assume that $x \in \mathbb{Q}$

$$7\frac{p^2}{q^2} - 2 = 0, p, q \in \mathbb{Q}$$
$$\frac{p^2}{q^2} = \frac{2}{7}$$

Therefore $p^2 = 2$ and $q^2 = 7$, but we have already proven that these equalities do no have integer solutions; thus, we have reached a contradiction.

b)

Proof.

$$u = \frac{x}{x^2 - 7y^2}, v = \frac{-y}{x^2 - 7y^2}$$

Because $x, y, 7 \in \mathbb{Q}$, and \mathbb{Q} being a field, $u, v \in \mathbb{Q}$.

c)

Proof.

Assume $\sqrt{2} \in \mathbb{Q}(\sqrt{7})$

$$\sqrt{2} = (u + \sqrt{7}v), u, v \in \mathbb{Q}$$
$$2 = u^2 + 2\sqrt{7}uv + 7v^2$$

Because 2 is rational, and the sum of a rational and an irrational is always irrational by lemma 2.2, either u=0 or v=0 for $2\sqrt{7}uv=0$ (by lemma 2.3).

Case 1: u = 0

$$7v^2 = 2$$
$$v^2 = \frac{2}{7}$$

We have already proven that this equation has no solution in \mathbb{Q} , thus a contradiction.

Case 2: v = 0

$$u^2 = 2$$

We have already proven that this equation has no solutions in \mathbb{Q} , thus a contradiction.

Case 3: u = v = 0

$$2 = 0$$

This is clearly a contradiction.

3. a)

The fallacy that the student is employing is that every subset of students of size n have the same grade. It is only true that the $first\ n$ students have the same grade.

- b) i) We need to prove P(1).
- ii) We need to prove P(2) and P(3).
- iii) We have to prove that $P(n) \Rightarrow P(n+1)$.

4. a)

$$\neg(\forall \epsilon > 0, \exists \delta > 0: 0 < |x - a| < \delta \Rightarrow 0 < |f(x) - f(a)| < \epsilon)$$
 Is the same as:
$$\exists \epsilon > 0, \forall \delta > 0: 0 < |x - a| < \delta \wedge |f(x) - f(a)| \geq \epsilon$$

b) i)

$$\emptyset$$
, because $|f(x) - f(a)| \ge 0$

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$$\{f: f \text{ is continuous at } x\},$$
 because if $|f(x)-f(a)|=\frac{\epsilon}{2}$ then $|f(x)-f(a)|<\epsilon$

iii)

 $\{f:f \text{ is a function}\}$, because if $\delta=0$, then the condition is vacuously satisfied.

iv)

$$f: D \to C$$
 $\{f: \exists \delta > 0, \exists x \in D \text{ s.t. } f \text{ is constant on the interval } (x - \delta, x + \delta)\}$