

# CSC240 Problem Set 2

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1.

We will use proof by contraposition.

Assume NOT  $2CR(C)$ .

NOT  $2CR(C)$  IMPLIES  $\exists \alpha, \beta, \gamma \in \{1, \dots, N\}$ .

$\alpha \neq \beta$  AND  $\alpha \neq \gamma$  AND  $(\forall e \in C_\alpha (e \in C_\beta \text{ OR } e \in C_\gamma$

By definition,  $i \in C_j$  IFF  $M_{[i,j]} = 1$  IFF  $j \in R_i$ .

Let  $\alpha \in \{1, \dots, N\}$  such that  $\exists \beta \in \{1, \dots, N\}. \exists \gamma \in \{1, \dots, N\}$ .

$\alpha \neq \beta$  AND  $\alpha \neq \gamma$  AND  $(\forall e \in C_\alpha (e \in C_\beta \text{ OR } e \in C_\gamma$

Let  $e := \alpha \in \{a, b, c\}$

Let  $i \in \{1, \dots, N\}$  be arbitrary.

Either  $\alpha \in R_i$  OR NOT  $\alpha \in R_i$ .

Case 1: NOT  $\alpha \in R_i$ .

NOT  $e \in R_i$ .

Case 2:  $\alpha \in R_i$ .

By above  $\alpha \in R_i$  IMPLIES  $\beta \in R_i$  OR  $\gamma \in R_i$

$\exists d \in S. d \in T$  AND NOT  $d = e$ , namely  $\gamma$  or  $\beta$ .

Thus either  $e \notin R_i$  OR  $\exists d \in S. d \in T$  AND NOT  $d = e$

$i$  was arbitrary so it holds for all  $R_i$ .

$\exists S \in \mathcal{P}_3(\{1, \dots, N\}). \exists e \in S \forall T \in R. e \notin T$  OR  $\exists d \in S. d \in T$  AND NOT  $d = e$

NOT SEL(3, R).

2.

1 Assume NOT 2CR( $C$ ).

2 To obtain a contradiction, assume  $\forall \alpha \in \{1, \dots, N\}. \forall \beta \in \{1, \dots, N\}. \forall \gamma \in \{1, \dots, N\}.$

$[(\text{NOT } (\alpha = \beta)) \text{ AND } (\text{NOT } (\alpha = \gamma)) \text{ AND } (\exists e \in C_\alpha (\text{NOT MEMBER}(e, C_\beta) \text{ AND NOT MEMBER}(e, C_\gamma)))]$

3  $\exists S \in C. \exists X \in C. \exists Y \in C.$

$[(\text{NOT } (S = X)) \text{ AND } (\text{NOT } (S = Y)) \text{ AND } (\exists e \in S (\text{MEMBER}(e, X) \text{ OR MEMBER}(e, Y)))]$   
modus ponens: 1.

4 Let  $\alpha \in \{1, \dots, N\}$  such that  $C_\alpha = S$ .

5 let  $\beta \in \{1, \dots, N\}$  such that  $C_\beta = X$ .

6 Let  $\gamma \in \{1, \dots, N\}$  such that  $C_\gamma = Y$ .

7 This is a contradiction: 3,4,5,6.

8 Therefore  $\exists \alpha \in \{1, \dots, N\}. \exists \beta \in \{1, \dots, N\}. \exists \gamma \in \{1, \dots, N\}.$

$[(\text{NOT } (\alpha = \beta)) \text{ AND } (\text{NOT } (\alpha = \gamma)) \text{ AND } (\forall e \in C_\alpha (\text{MEMBER}(e, C_\beta) \text{ OR MEMBER}(e, C_\gamma)))]$

9 Let  $i \in \{1, \dots, n\}$  be arbitrary.

10  $\text{MEMBER}(i, C_n) \text{ IMPLIES } M_{[i,n]} = 1$ : modus ponens: by assumption.

11  $M[i, n] = 1 \text{ IMPLIES MEMBER}(n, R_i)$ : modus ponens: by assumption.

12 Since  $i$  is an arbitrary element of  $\{1, \dots, N\}$ .

13  $\forall i \in \{1, \dots, N\}. \text{MEMBER}(i, C_n) \text{ IMPLIES MEMBER}(n, R_i)$ : generalization 9.

14  $\exists \alpha \in \{1, \dots, N\}. \exists \beta \in \{1, \dots, N\}. \exists \gamma \in \{1, \dots, N\}.$

$[(\text{NOT } (\alpha = \beta)) \text{ AND } (\text{NOT } (\alpha = \gamma)) (\exists e \in C_\alpha (\text{MEMBER}(e, C_\beta) \text{ OR MEMBER}(e, C_\gamma)))]$ ;  
modus ponens: 8.

15 let  $a \in \{1, \dots, N\}$  be such that  $\forall b \in \{1, \dots, N\}. \exists c \in \{1, \dots, N\}.$

$[(\text{NOT } (a = b)) \text{ AND } (\text{NOT } (a = c)) \text{ AND } (\forall e \in C_a (\text{MEMBER}(e, C_b) \text{ OR MEMBER}(e, C_c)))]$

Instantiation: 8

16 Let  $\{a, b, c\} \in \mathcal{P}_3(\{1, \dots, N\})$ .

17 Let  $e := a \in \{a, b, c\}$

18 Let  $i \in \{1, \dots, n\}$  be arbitrary.

19 Either  $\text{MEMBER}(a, R_i) \text{ OR NOT MEMBER}(a, R_i)$ .

20 Case 1: Assume  $\text{MEMBER}(a, R_i)$

21  $\text{MEMBER}(a, R_i) \text{ IMPLIES MEMBER}(i, C_a)$ ; modus ponens: by assumption.

22  $\text{MEMBER}(i, C_a) \text{ IMPLIES } \exists b \in \{1, \dots, n\}. \exists c \in \{1, \dots, n\}.$

$[(\text{NOT } (a = b)) \text{ AND } (\text{NOT } (a = c)) \text{ AND } (\text{MEMBER}(i, C_b) \text{ OR MEMBER}(i, C_c))]$ ; modus ponens: 15.

23 Either  $\text{MEMBER}(i, C_b) \text{ OR NOT MEMBER}(i, C_b)$ .

24 Case 1.a: Assume  $\text{MEMBER}(i, C_b)$ .

25  $\text{MEMBER}(i, C_b) \text{ IMPLIES MEMBER}(b, R_i)$ ; modus ponens: 13.

26 Thus  $\exists d \in R_i. \text{MEMBER}(d, \{a, b, c\}) \text{ AND NOT } d = e$ .

27  $\text{MEMBER}(i, C_b) \text{ IMPLIES } \exists d \in R_i. \text{MEMBER}(d, \{a, b, c\}) \text{ AND NOT } d = e$   
direct proof: 25, 26.

28 Case 1.b: Assume NOT  $\text{MEMBER}(i, C_b)$

29 NOT  $\text{MEMBER}(i, C_b) \text{ IMPLIES MEMBER}(c, R_i)$ ; modus ponens: 13.

30 Thus  $\exists d \in R_i. \text{MEMBER}(d, \{a, b, c\}) \text{ AND NOT } d = e$ .

31 NOT  $\text{MEMBER}(i, C_b) \text{ IMPLIES } \exists d \in R_i. \text{MEMBER}(d, \{a, b, c\}) \text{ AND NOT } d = e$   
direct proof: 29, 30.

32 MEMBER( $a, R_i$ ) IMPLIES  $\exists d \in R_i$ .MEMBER( $d, \{a, b, c\}$ ) AND NOT  $d = e$ ;  
 direct proof: 24, 29.  
 33 Case 2: Assume NOT MEMBER( $a, R_i$ ).  
 34 NOT MEMBER( $a, R_i$ ) IMPLIES NOT MEMBER( $e, R_i$ ); modus ponens: 17.  
 35 NOT MEMBER( $a, R_i$ ) IMPLIES NOT MEMBER( $e, R_i$ ); direct proof: 34  
 36 Since  $i$  is an arbitrary element of  $\{1, \dots, n\}$   
 37  $\forall i \in \{1, \dots, n\}.$  $\exists d \in R_i$ .MEMBER( $d, \{a, b, c\}$ ) AND NOT  $d = e$  OR NOT MEMBER( $e, R_i$ ).  
 generalization: 18.  
 38  $\exists S \in \mathcal{P}_3(R).$  $\exists e \in S.$  $\forall T \in R.$  $\exists d \in S$ .MEMBER( $d, T$ ) AND NOT  $d = e$  OR NOT MEMBER( $e, T$ ).  
 Construction: 15, 16, 17.  
 39 NOT SEL( $R$ )  
 40 SEL( $3, R$ ) IMPLIES 2CF( $C$ ); indirect proof: 1, 40.